# Machine-Learning Machines Don't Learn; Al Needs Real Learning, eg Learning Ex Nihilo

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IFLAI2 Oct 6 2022



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Instantly Revealed Fatal Problem for DL:
Representation of Declarative Information —
which logic handles with the ease of driving a
hot knife through a soft stick of butter.

 $\sigma$ : "My best friend's floozerbak makes a bejeeker that's better than anyone else's — I think because it uses some secret ingredient beyond lazerall and sinifer."

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https://arxiv.org/abs/2207.09238

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To represent  $\sigma$  we need to tokenize it. How? We need a vocabulary V that is associated with  $[N_V]$ , a finite set of numbers  $\{1,2,\ldots,N_V\}$ . What is V itself? It's a set composed of sub-words, usually. But without loss of mathematical generality we can just go with words; in that case tokenization gives us

bos\_token, My, best, friend's, floozerbak, makes, a, bejeeker, that's, better, than, anyone, ..., sinifer, eos\_token

which we can then express as a vector composed of the indices; so where  $n_i \in \mathbb{Z}^+$  we have e.g.

$$[n_1, n_2, ..., n_k].$$

...  $\exists x [F(x,I) \land \forall y ((F(y),I \land y \neq x) \rightarrow BF(x,I,y)) \land \exists z (Makes(floozerbak-of(x),z)...$ 

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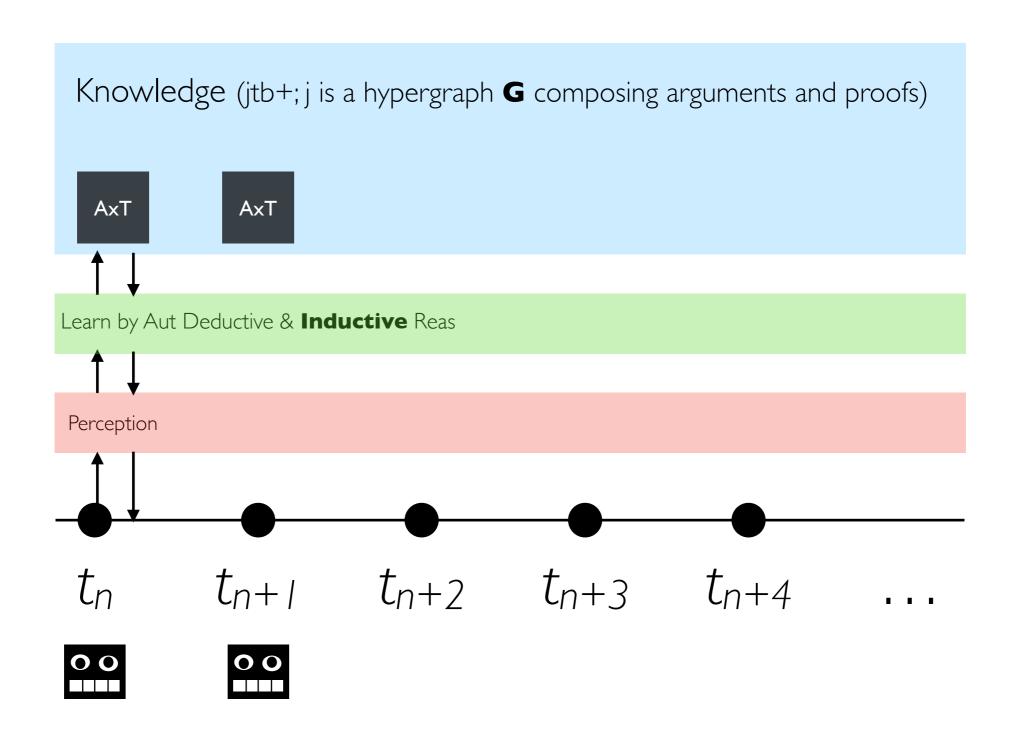
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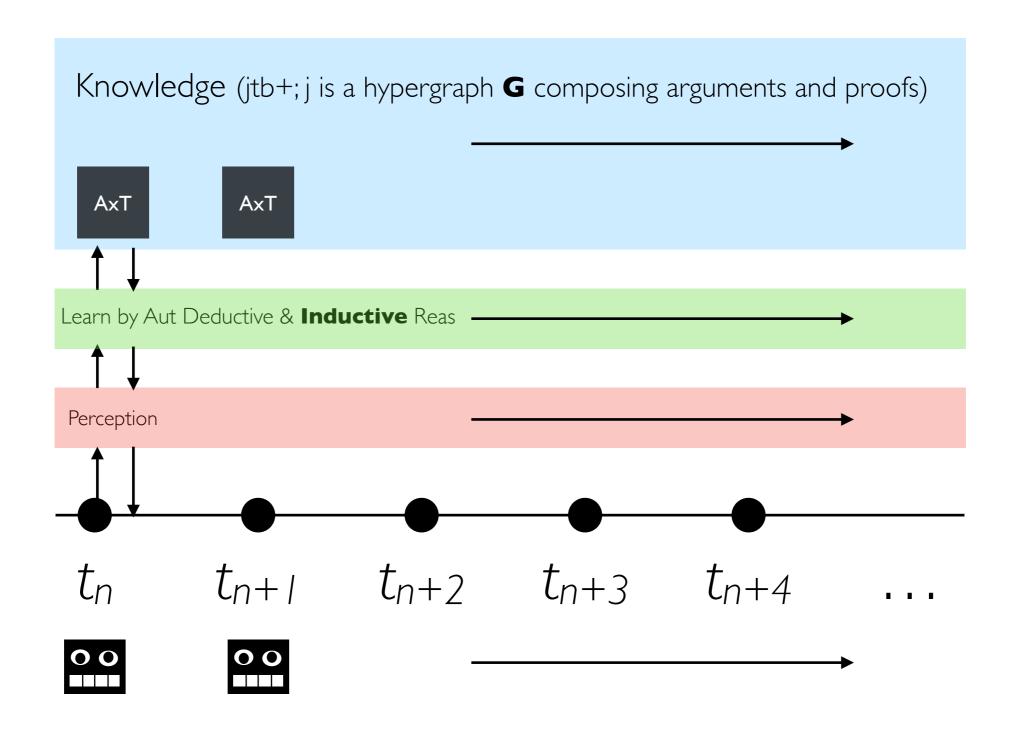
$$[n_1, n_2, ..., n_k].$$

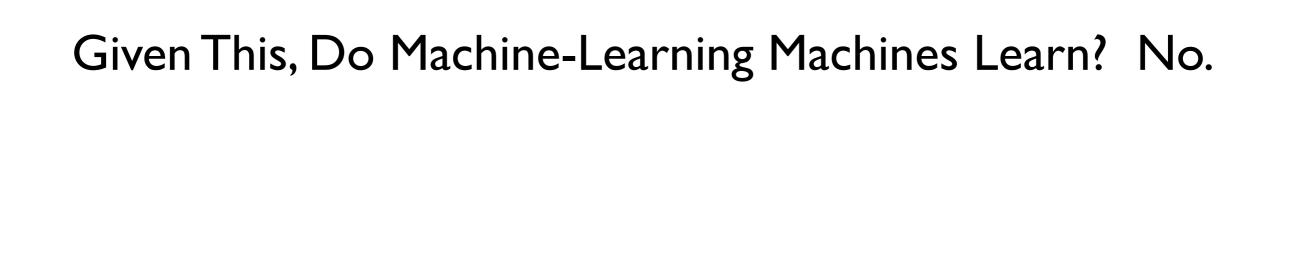
Are there two bejeekers made by two different agents, and believed by the speaker to be singularly good, for reasons beyond their having in them either lazerall or sinifer?

#### Advanced Logicist (Real) Machine Learning



#### Advanced Logicist (Real) Machine Learning





#### Given This, Do Machine-Learning Machines Learn? No.

#### Do Machine-Learning Machines Learn?

Selmer Bringsjord and Naveen Sundar Govindarajulu and Shreya Banerjee and John Hummel

Abstract We answer the present paper's title in the negative. We begin by introducing and characterizing "real learning" ( $\mathcal{RL}$ ) in the formal sciences, a phenomenon that has been firmly in place in homes and schools since at least Euclid. The defense of our negative answer pivots on an integration of *reductio* and proof by cases, and constitutes a general method for showing that any contemporary form of machine learning (ML) isn't real learning. Along the way, we canvass the many different conceptions of "learning" in not only AI, but psychology and its allied disciplines; none of these conceptions (with one exception arising from the view of cognitive development espoused by Piaget), aligns with real learning. We explain in this context by four steps how to broadly characterize and arrive at a focus on  $\mathcal{RL}$ .

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Do Machine-Learning Machines Learn?

#### 17

#### 8 Appendix: The Formal Method

The following deduction uses fonts in an obvious and standard way to sort between functions ( $\mathfrak{f}$ ), agents ( $\mathfrak{a}$ ), and computing machines ( $\mathfrak{m}$ ) in the Arithmetical Hierarchy. Ordinary italicized Roman is used for particulars under these sorts (e.g. f is a particular function). In addition, ' $\mathcal{C}$ ' denotes any collection of conditions constituting jointly necessary-and-sufficient conditions for a form of current ML, which can come from relevant textbooks (e.g. Luger, 2008; Russell and Norvig, 2009) or papers; we leave this quite up to the reader, as no effect upon the validity of the deductive inference chain will be produced by the preferred instantiation of ' $\mathcal{C}$ .' It will perhaps be helpful to the reader to point out that the deduction eventuates in the proposition that no machine in the ML fold that in this style learns a relevant function  $\mathfrak{f}$  thereby also real-learns  $\mathfrak{f}$ . We encode this target as follows:

```
(\star) \neg \exists \mathfrak{m} \ \exists \mathfrak{f} \ [\phi := MLlearns(\mathfrak{m}, \mathfrak{f}) \land \psi := RLlearns(\mathfrak{m}, \mathfrak{f}) \land \mathcal{C}_{\phi}(\mathfrak{m}, \mathfrak{f}) \vdash^{*} (ci') - (ciii)_{\psi}(\mathfrak{m}, \mathfrak{f})]
```

Note that  $(\star)$  employs meta-logical machinery to refer to particular instantiations of  $\mathcal C$  for a particular, arbitrary case of ML  $(\phi)$  is the atomic sub-formula that can be instantiated to make the particular case), and particular instantiations of the triad (ci')–(ciii) for a particular, arbitrary case of  $\mathcal R\mathcal L$   $(\psi)$  is the atomic sub-formula that can be instantiated to make the particular case). Meta-logical machinery also allows us to use a provability predicate to formalize the notion that real learning is produced by the relevant instance of ML. If we "pop"  $\phi/\psi$  to yield  $\phi'/\psi'$  we are dealing with the particular instantiation of the atomic sub-formula.

The deduction, as noted in earlier when the informal argument was given, is indirect proof by cases; accordingly, we first assume  $\neg(\star)$ , and then proceed as follows under this supposition.

```
| (1) | \forall f, a [f : \mathbb{N} \mapsto \mathbb{N} \to (RLlearns(a, f) \to (i) - (iii)) | Def of Real Learning
    (2) MLlearns(m, f) \land RLlearns(m, f) \land f : \mathbb{N} \mapsto \mathbb{N}
                                                                                            supp (for \exists elim on (\star))
    (3) \forall \mathfrak{m}, \mathfrak{f} [\mathfrak{f} : \mathbb{N} \mapsto \mathbb{N} \to (MLlearns(\mathfrak{m}, \mathfrak{f}) \leftrightarrow \mathcal{C}(\mathfrak{m}, \mathfrak{f}))] Def of ML
    (4) \forall f [f : \mathbb{N} \mapsto \mathbb{N} \to (TurComp(f) \lor TurUncomp(f))] theorem
    (5) TurUncomp(f)
                                                                                             supp; Case 1
    (6) \neg \exists \mathfrak{m} \exists \mathfrak{f} [(\mathfrak{f} : \mathbb{N} \mapsto \mathbb{N} \wedge TurUncomp(\mathfrak{f}) \wedge \mathcal{C}(\mathfrak{m}, \mathfrak{f})]
                                                                                           theorem
\therefore (7) \neg \exists \mathfrak{m} MLlearns(\mathfrak{m}, f)
                                                                                            (6), (3)
∴ (8) ⊥
                                                                                            (7), (2)
    (9) |TurComp(f)|
                                                                                             supp; Case 2
|C_{\phi'}(m,f)|
                                                                                            (2), (3)
|(11)|(ci')-(ciii)_{\psi'}(m,f)
                                                                                            from supp for \exists elim on (\star) and provability
|\cdot|(12)|\neg(ci')-(ciii)_{w'}(m,f)
                                                                                            inspection: proofs wholly absent from C
∴|(13)|⊥
                                                                                             (11), (12)
                                                                                            reductio; proof by cases
```

# Let's look @ the paper ...

Step 1: Observe the acute discontinuity of human vs. nonhuman cognition.
(Only humans understand and employ e.g. abstract reasoning schemas unaffected by the physical; layered quantification; recursion; and infinite structures/infinitary reasoning.)

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- Step 4: Real Learning  $(\mathcal{RL})$  is the acquisition of genuine knowledge via **RC**.

But how is this mechanizable?

Well, how about a new form of machine learning? (by reasoning)

# Novel Form of Machine Learning:

# Learning Ex Nihilo

(or Learning Ex Minima)

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#### EPiC Series in Computing

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#### Learning Ex Nihilo

Selmer Bringsjord<sup>1</sup>, Naveen Sundar Govindarajulu<sup>1</sup>, John Licato<sup>2</sup>, and Michael Giancola<sup>1</sup>

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 selmer.bringsjord, naveen.sundar.g, mike.j.giancola, john.licato } @gmail.com

#### Abstract

This paper introduces, philosophically and to a degree formally, the novel concept of learning ex nihilo, intended (obviously) to be analogous to the concept of creation ex nihilo. Learning ex nihilo is an agent's learning "from nothing", by the suitable employment of inference schemata for deductive and inductive reasoning. This reasoning must be in machine-verifiable accord with a formal proof/argument theory in a cognitive calculus (i.e., here, roughly, an intensional higher-order multi-operator quantified logic), and this reasoning is applied to percepts received by the agent, in the context of both some prior knowledge, and some prior and current interests. Learning ex nihilo is a challenge to contemporary forms of ML, indeed a severe one, but the challenge is here offered in the spirit of seeking to stimulate attempts, on the part of non-logicist ML researchers and engineers, to collaborate with those in possession of learning-ex nihilo frameworks, and eventually attempts to integrate directly with such frameworks at the implementation level. Such integration will require, among other things, the symbiotic interoperation of state-of-theart automated reasoners and high-expressivity planners, with statistical/connectionist ML technology.

#### 1 Introduction

This paper introduces, philosophically and to a degree logico-mathematically, the novel concept of learning *ex nihilo*, intended (obviously) to be analogous to the concept of creation *ex nihilo*. Learning *ex nihilo* is an agent's learning "from nothing," by the suitable employment of inference schemata for deductive and inductive<sup>2</sup> (e.g., analogical, enumerative-inductive, abductive, etc.) reasoning. This reasoning must be in machine-verifiable accord with a formal

<sup>&</sup>lt;sup>1</sup>No such assumption as that creation *ex nihilo* is real or even formally respectable is made or needed in the present paper. The concept of creation *ex nihilo* is simply for us an intellectual inspiration — but as a matter of fact, the literature on it in analytic philosophy does provide some surprisingly rigorous accounts. In the present draft of the present paper, we don't seek to mine these accounts.

<sup>&</sup>lt;sup>2</sup>Not to be confused with inductive logic programming (about which more will be said later), or inductive deductive techniques and schemas (e.g. mathematical induction, the induction schema in Peano Arithmetic, etc.). As we explain later, learning *ex nihilo* is in part powered by non-deductive inference schemata seen in inductive logic. An introductory overview of inductive logic is provided in [39].

G. Danoy, J. Pang and G. Sutcliffe (eds.), GCAI 2020 (EPiC Series in Computing, vol. 72), pp. 1–27

Bringsjord, S., Govindarajulu, N.S., Licato, J. & Giancola, M. (2020) "Learning Ex Nihilo" Proceedings of the 6th Global Conference on Artificial Intelligence (GCAI 2020), within International Conferences on Logic and Artificial Intelligence at Zhejiang University (ZJULogAI), in Danoy, G., Pang, J. & Sutcliffe, G., eds., EPiC Series in Computing 72: I-27 (Manchester, UK: EasyChair Ltd), ISSN: 2398-7340.

https://easychair.org/publications/paper/NzWG

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Certain **Evident** Overwhelmingly Likely Beyond Reasonable Doubt Likely More Likely Than Not Counterbalanced More Unlikely Than Not Unlikely Overwhelmingly Unlikely Beyond Reasonable Belief **Evidently False** Certainly False

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Certain Evident

Overwhelmingly Likely

Beyond Reasonable Doubt

Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

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Beyond Reasonable Belief

Evidently False

Certainly False

**Epistemically Positive** 

Certain

**Evident** 

Overwhelmingly Likely

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Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

Overwhelmingly Unlikely

Beyond Reasonable Belief

Evidently False

Certainly False

**Epistemically Positive** 

Certain

**Evident** 

Overwhelmingly Likely

Beyond Reasonable Doubt

Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

Overwhelmingly Unlikely

Beyond Reasonable Belief

Evidently False

Certainly False

**Epistemically Negative** 

**Epistemically Positive** 

Certain

**Evident** 

Overwhelmingly Likely

Beyond Reasonable Doubt

Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

Overwhelmingly Unlikely

Beyond Reasonable Belief

Evidently False

Certainly False

Epistemically Negative

	(6)	Certain
Epistemically Positive	(5)	Evident
	(4)	Overwhelmingly Likely
	(3)	Beyond Reasonable Doubt
	(2)	Likely
	(1)	More Likely Than Not
	(0)	Counterbalanced
		More Unlikely Than Not
	(-2)	Unlikely
	(-3)	Overwhelmingly Unlikely
	(-4)	Beyond Reasonable Belief
	(-5)	Evidently False
Epistemically Negative	(-6)	Certainly False

#### Or ... 11 Strength Factors

**Acceptable** An agent a at time t finds  $\phi$  acceptable *iff* withholding  $\phi$  is not more reasonable than believing in  $\phi$ .

$$\mathbf{B}^{1}(a,t,\phi) \Leftrightarrow \begin{cases} \mathbf{W}(a,t,\phi) \not\succeq_{t}^{a} \mathbf{B}(a,t,\phi); \text{ or } \\ \left(\neg \mathbf{B}(a,t,\phi) \land \neg \mathbf{B}(a,t,\neg\phi)\right) \not\succeq_{t}^{a} \mathbf{B}(a,t,\phi) \end{cases}$$

**Some Presumption in Favor** An agent a at time t has some presumption in favor of  $\phi$  *iff* believing  $\phi$  at t is more reasonable than believing  $\neg \phi$  at time t:

$$\mathbf{B}^{2}(a,t,\phi) \Leftrightarrow \left(\mathbf{B}(a,t,\phi) \succ_{a}^{t} \mathbf{B}(a,t,\neg\phi)\right)$$

**Beyond Reasonable Doubt** An agent a at time t has beyond reasonable doubt in  $\phi$  *iff* believing  $\phi$  at t is more reasonable than withholding  $\phi$  at time t:

$$\mathbf{B}^{3}(a,t,\phi) \Leftrightarrow \begin{cases} \mathbf{B}(a,t,\phi) \succ_{a}^{t} \mathbf{W}(a,t,\phi); \text{ or } \\ \left(\mathbf{B}(a,t,\phi) \succ_{t}^{a} \left(\neg \mathbf{B}(a,t,\phi) \land \neg \mathbf{B}(a,t,\neg\phi)\right) \end{cases}$$

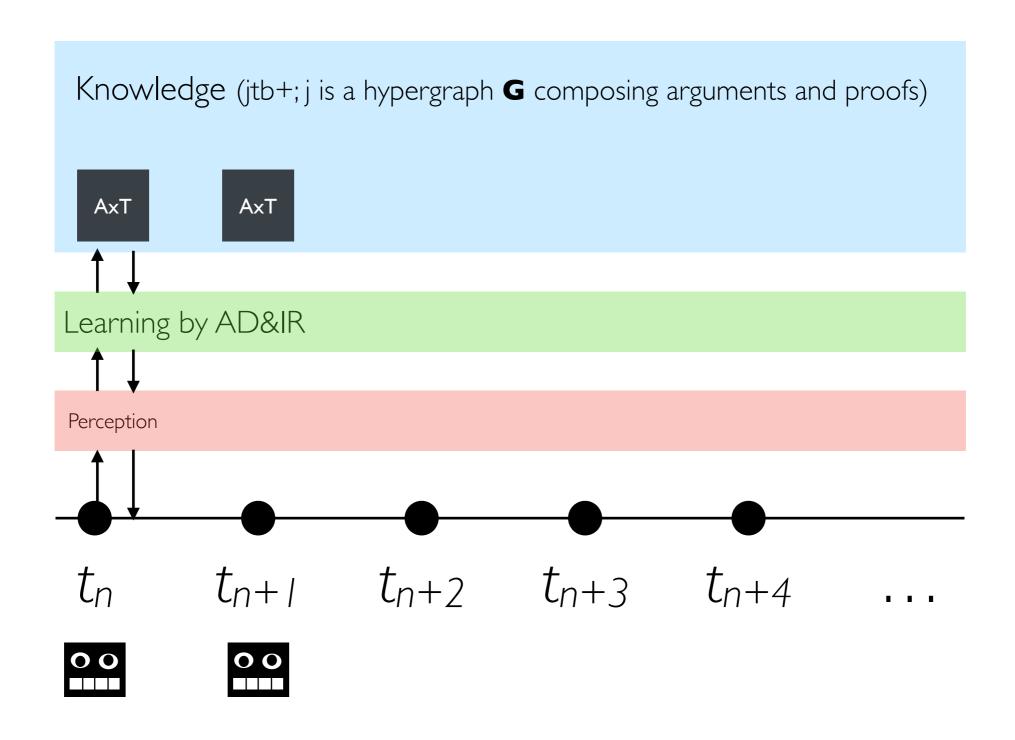
**Evident** A formula  $\phi$  is evident to an agent a at time t iff  $\phi$  is beyond reasonable doubt and if there is a  $\psi$  such that believing  $\psi$  is more reasonable for a at time t than believing  $\phi$ , then a is certain about  $\psi$  at time t.

$$\mathbf{B}^{4}(a,t,\phi) \Leftrightarrow \begin{cases} \mathbf{B}^{3}(a,t,\phi) \land \\ \exists \psi : \begin{bmatrix} \mathbf{B}(a,t,\psi) \succ_{t}^{a} \mathbf{B}(a,t,\phi) \\ \Rightarrow \mathbf{B}^{5}(a,t,\psi) \end{bmatrix} \end{cases}$$

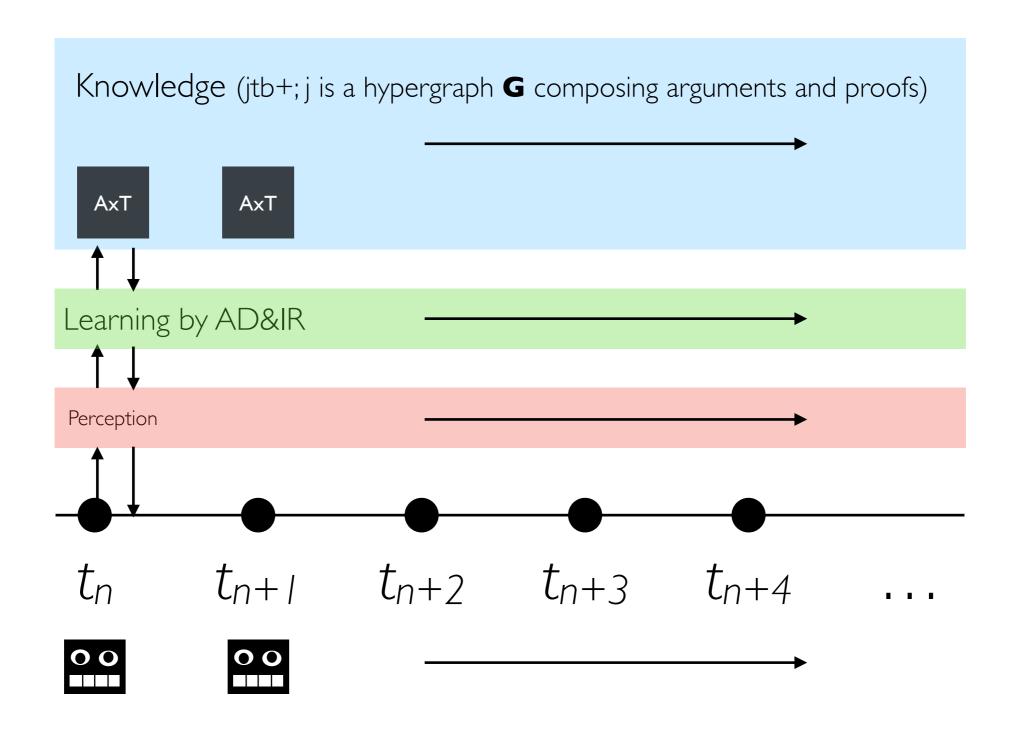
**Certain** An agent a at time t is certain about  $\phi$  iff  $\phi$  is beyond reasonable doubt and there is no  $\psi$  such that believing  $\psi$  is more reasonable for a at time t than believing  $\phi$ .

$$\mathbf{B}^{5}(a,t,\phi) \Leftrightarrow \begin{cases} \mathbf{B}^{3}(a,t,\phi) \wedge \\ \neg \exists \psi : \mathbf{B}(a,t,\psi) \succ_{t}^{a} \mathbf{B}(a,t,\phi) \end{cases}$$

#### Demo: "Sully-esque Al Performs Miracle on the Hudson"



### Demo: "Sully-esque Al Performs Miracle on the Hudson"

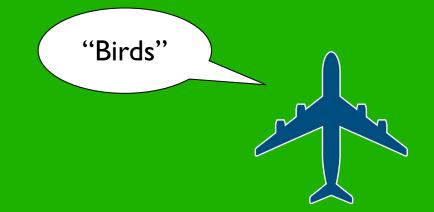


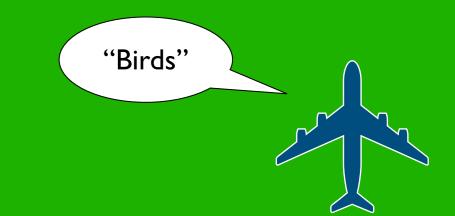


Michael Giancola Graduate Research Assistant (PhD)









ATC

"Birds"



 $\mathbf{S}(atc, capt, t_1, \operatorname{Land}(capt, t_1, lga_{13}))$ 

 $\therefore \mathbf{B}(capt, t_1, \mathbf{B}(atc, t_1, Land(capt, t_1, lga_{13}))) [I_{12}] \checkmark$ 

 $\therefore \mathbf{B}^{1}(capt, t_{1}, \operatorname{Land}(capt, t_{1}, lga_{13})) \qquad [\mathbf{B}^{1}\text{-def}] \checkmark$ 







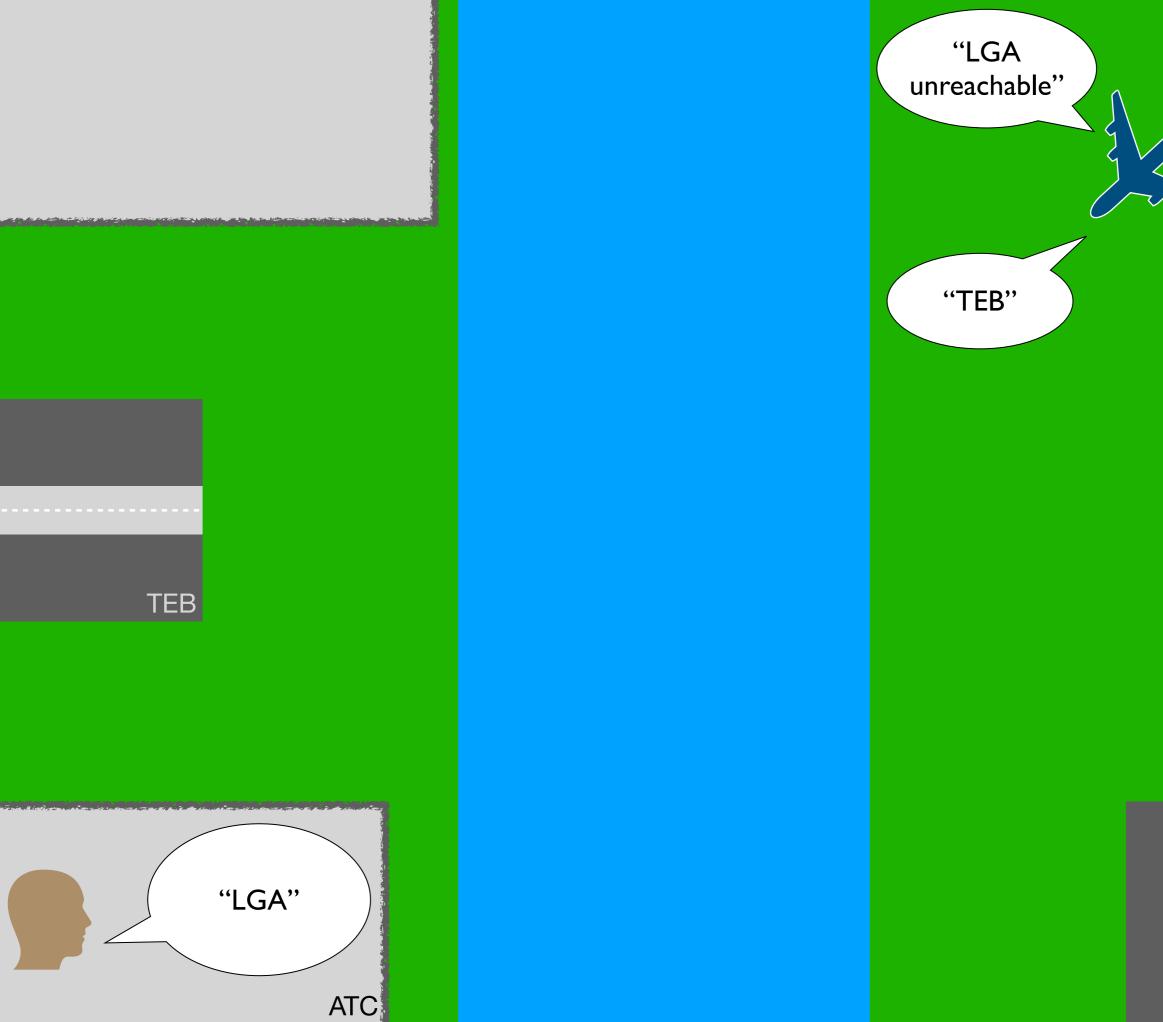


ATC

TEB "LGA" ATC

"LGA unreachable"







$$\therefore \operatorname{Land}(\operatorname{capt}, t_2, \operatorname{teb}) \succ_{t_2}^{\operatorname{capt}} \operatorname{Land}(\operatorname{capt}, t_2, \operatorname{lga}_{13}) \left[ \succ_t^a -\operatorname{def} \right] \checkmark$$

$$\therefore \mathbf{B}^2(\operatorname{capt}, t_2, \operatorname{Land}(\operatorname{capt}, t_2, \operatorname{teb})) \qquad [\mathbf{B}^2 - \operatorname{def}] \checkmark$$



"LGA unreachable"



"TEB"



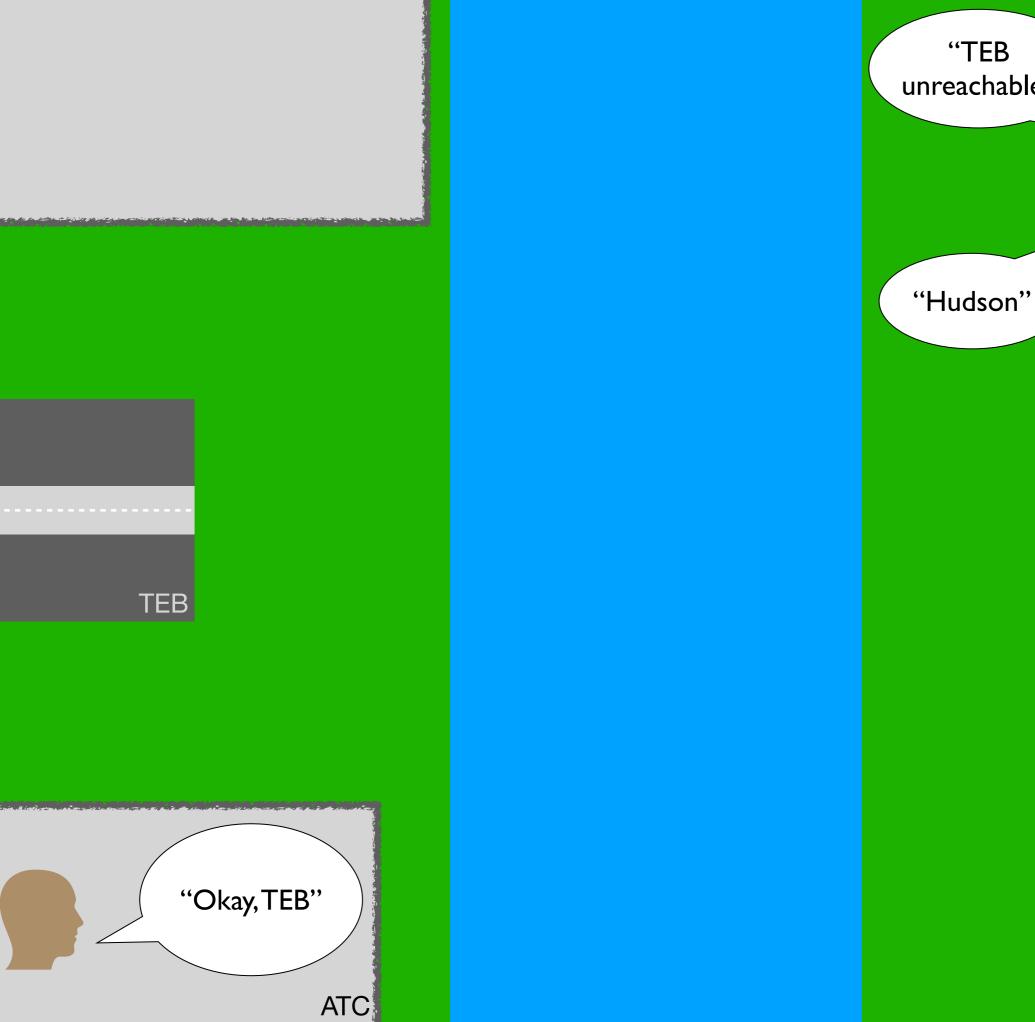


ATC

TEB "Okay, TEB" ATC

"TEB unreachable"







 $\mathbf{S}(atc, capt, t_3, \operatorname{Land}(capt, t_3, teb_1))$ 

- $\therefore \mathbf{B}(capt, t_3, \mathbf{B}(atc, t_3, \operatorname{Land}(capt, t_3, teb_1))) \ [I_{12}] \checkmark$
- $\therefore \mathbf{B}^{1}(capt, t_{3}, \operatorname{Land}(capt, t_{3}, teb_{1})) \qquad [\mathbf{B}^{1}\text{-def}] \checkmark$
- $\therefore \operatorname{Land}(\operatorname{capt}, t_3, \operatorname{hud}) \succ_{t_3}^{\operatorname{capt}} \operatorname{Land}(\operatorname{capt}, t_3, \operatorname{teb}_1) \ [ \succ_t^a -\operatorname{def} ] \checkmark$
- $\therefore \mathbf{B}^2(capt, t_3, \operatorname{Land}(capt, t_3, hud))$  [ $\mathbf{B}^2$ -def]  $\checkmark$

TEB

"Okay, TEB"

"TEB unreachable"



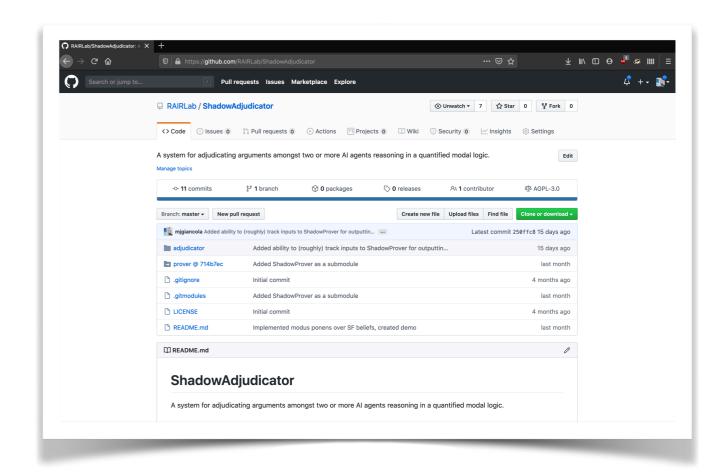
"Hudson"

 $\mathbf{S}(atc, capt, t_3, \operatorname{Land}(capt, t_3, teb_1))$ 

- $\therefore \mathbf{B}(capt, t_3, \mathbf{B}(atc, t_3, Land(capt, t_3, teb_1))) [I_{12}] \checkmark$
- $\therefore \mathbf{B}^{1}(capt, t_{3}, \operatorname{Land}(capt, t_{3}, teb_{1})) \qquad [\mathbf{B}^{1}\text{-def}] \checkmark$
- $\therefore \operatorname{Land}(\operatorname{capt}, t_3, \operatorname{hud}) \succ_{t_3}^{\operatorname{capt}} \operatorname{Land}(\operatorname{capt}, t_3, \operatorname{teb}_1) \ [ \succ_t^a -\operatorname{def} ] \checkmark$
- $\therefore \mathbf{B}^{2}(capt, t_{3}, \operatorname{Land}(capt, t_{3}, hud)) \qquad [\mathbf{B}^{2} \operatorname{def}] \checkmark$

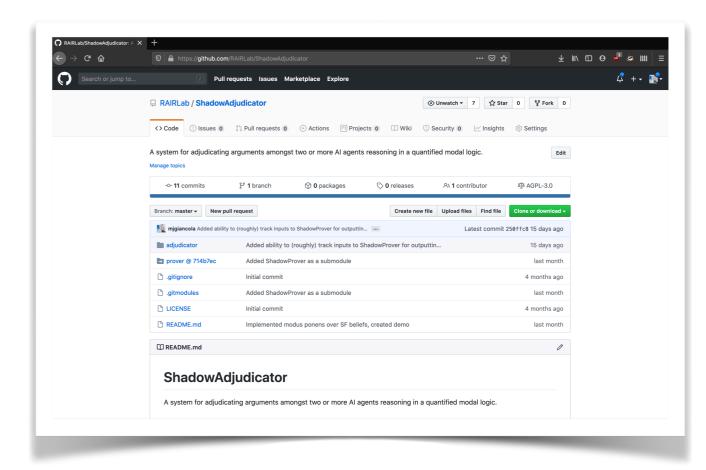






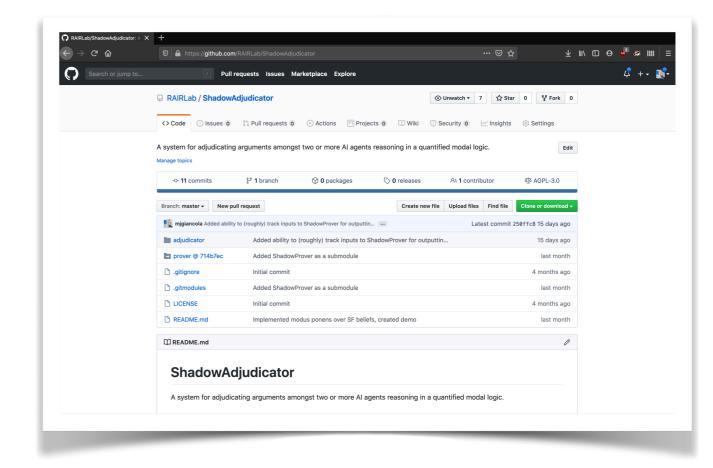


 A nascent automated reasoner for generating and adjudicating arguments



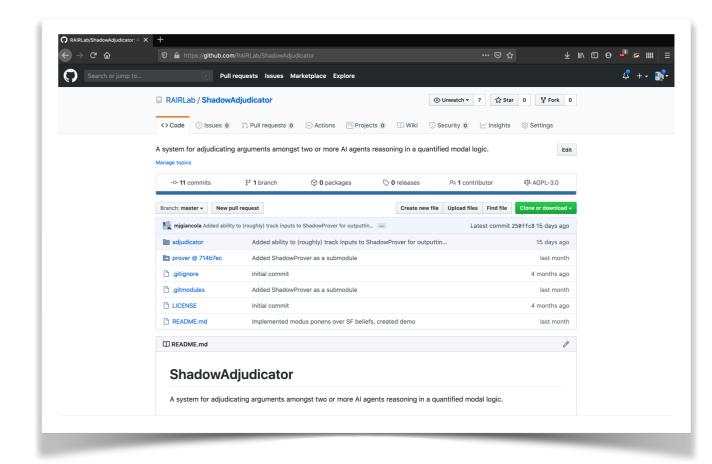


- A nascent automated reasoner for generating and adjudicating arguments
- Builds upon ShadowProver



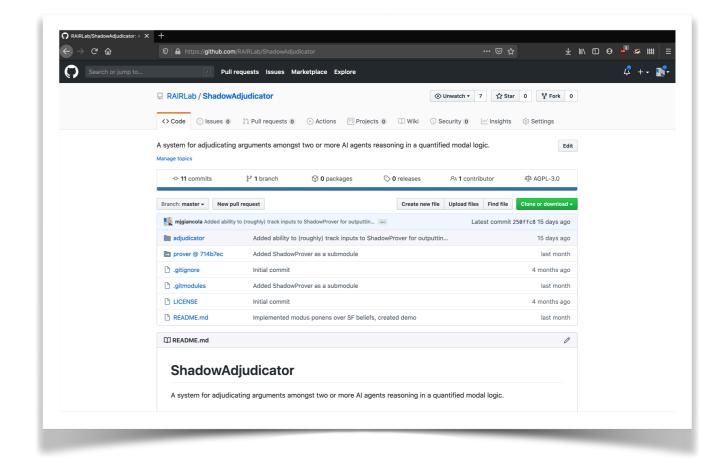


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- A nascent automated reasoner for generating and adjudicating arguments
- Builds upon ShadowProver
  - Uses ShadowProver for subproofs of modal/FOL/PL formulae
  - Implements an algorithm and inference schemata for generating arguments with strength factors





## Simulation



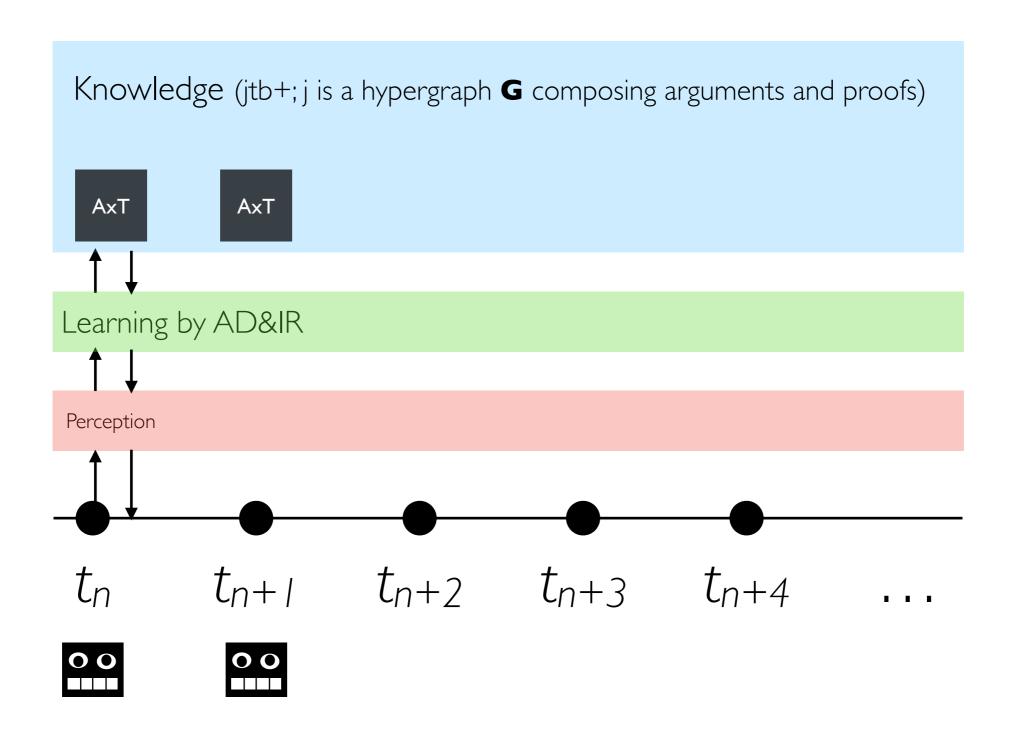


## Simulation

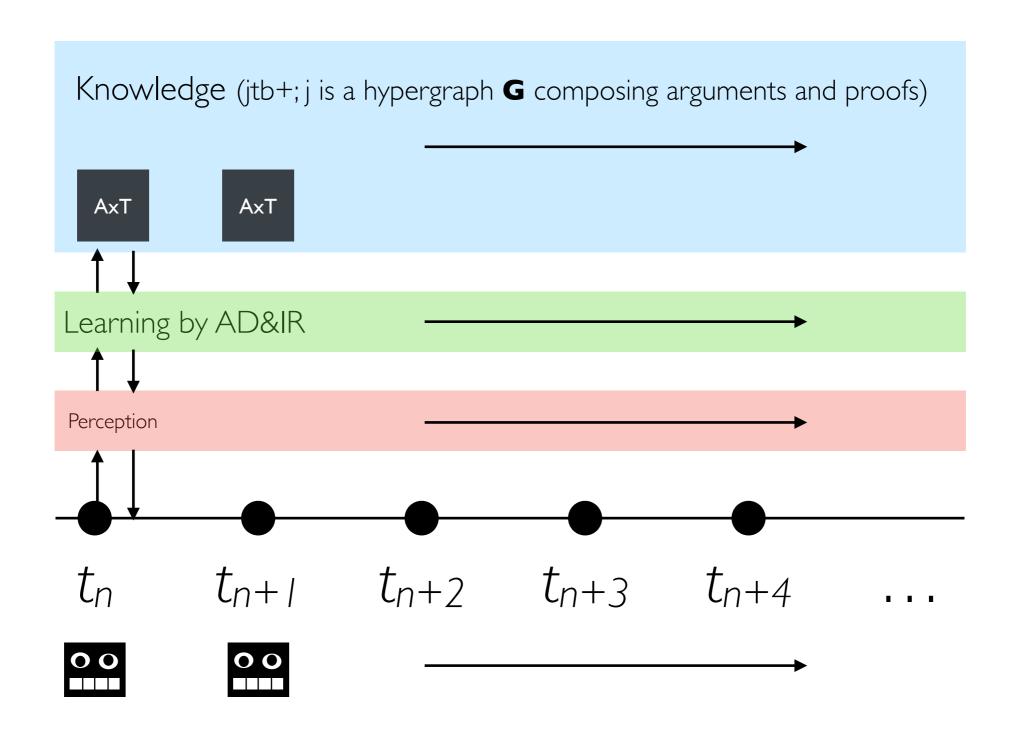




### Formally Shown: "Uber Fatality Avoided"



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Michael Giancola Graduate Research Assistant (PhD)







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At the start of the event, the vehicle was in autonomous mode in the rightmost of four lanes traveling in the same direction, and the pedestrian was walking her bicycle across the street starting on the leftmost side of the roadway.



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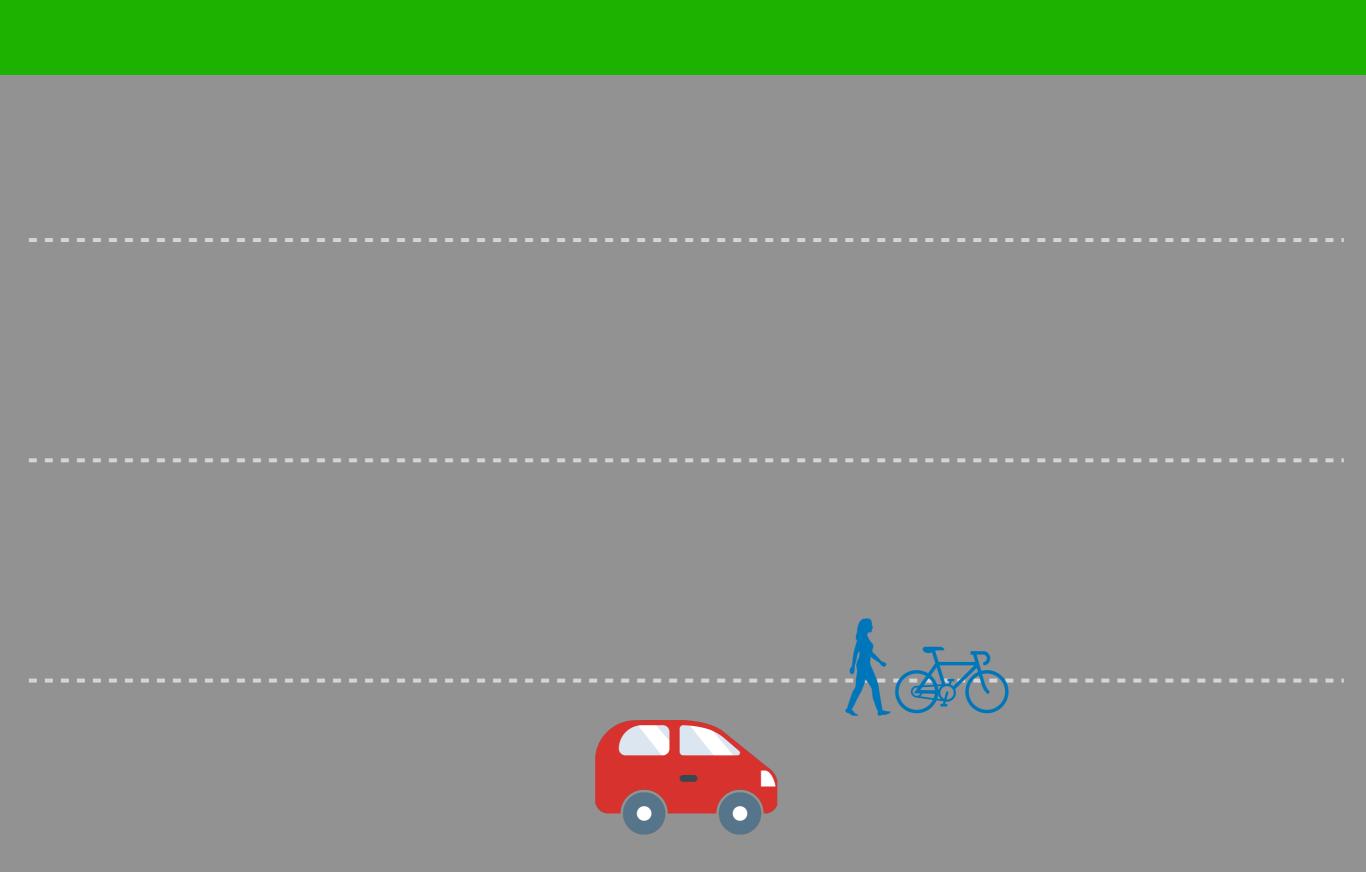
The vehicle's radar first detected the pedestrian **5.6 seconds** before the fatal collision. Less than half a second later, the lidar detected the pedestrian but classified her as **"Other"**.



For the next 2.5 seconds, the lidar re-classified her several times, alternating between "Vehicle" and "Other". The vehicle's automated-driving system (ADS) attempted to predict her direction of travel several times, but **discarded** any previous information about her trajectory every time it reclassified her.



With 2.6 seconds until collision, the lidar classified her as a bicycle but, as it was yet again changing her classification, discarded any past trajectory information, and hence determined that she was not moving. Up to this point, **the car** had not taken any evasive or corrective action.



With 1.5 seconds left, the lidar re-classified her yet again, this time as "Unknown". **The system once again loses all of its tracking history.** However, since at this point the pedestrian had entered the vehicle's lane, the ADS generated a plan to turn the car to the right to avoid her.



Three hundred milliseconds later, the lidar re-classified her as a bicycle, and determined that it would be impossible at this point to maneuver around her. With just 200 ms until collision, the ADS began braking the vehicle, pitifully too late to stop in time.

Now, how our Al would've handled it...



Three hundred milliseconds later, the lidar re-classified her as a bicycle, and determined that it would be impossible at this point to maneuver around her. With just 200 ms until collision, the ADS began braking the vehicle, pitifully too late to stop in time.





Label the time point at which the radar first detected the pedestrian  $\mathbf{t_0}$ , and the location of the pedestrian at that time  $\ell_1$ .

### **Argument I**

### **Argument 2**





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### **Argument I**

#### **Argument 2**

$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$





Label the time point at which the radar first detected the pedestrian  $\mathbf{t_0}$ , and the location of the pedestrian at that time  $\ell_1$ .

## $\mathbf{P}(\mathbf{r}_1, t_0, \text{At}(o^*, t_0, \ell_1)) = \mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$

## **Argument 2**

$$\mathbf{P}(\mathbf{r}_2, t_0, \operatorname{At}(o^*, t_0, \ell_1))$$





Label the time point at which the radar first detected the pedestrian to, and the location of the pedestrian at that time  $\ell_{\rm I}$ .

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## **Argument 2**

$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$

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$$\mathbf{P}(\mathbf{r}_1, t_0, \text{At}(o^*, t_0, \ell_1))$$

•••

$$\mathbf{B}(\mathbf{r}_1, t_1, \neg \text{Moving}(o^*))$$

$$\succ_t^a$$

$$\mathbf{B}(\mathbf{r}_1, t_1, \text{Moving}(o^*))$$

$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$





$$\mathbf{P}(\mathbf{r}_1, t_0, \operatorname{At}(o^*, t_0, \ell_1))$$

• • •

$$\mathbf{B}(\mathfrak{r}_1, t_1, \neg \text{Moving}(o^*))$$

$$\succ_t^a$$

$$\mathbf{B}(\mathfrak{r}_1, t_1, \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \text{NeedToBrake}(c))$$

$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$





$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$

• • •

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$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$

$$\mathbf{P}(\mathbf{r}_2, t_1, \text{At}(o^*, t_0, \ell_2))$$

$$\mathbf{P}(\mathfrak{r}_2,t_1,\ell_1=\ell_2)$$





$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$

•••

$$\mathbf{B}(\mathfrak{r}_1, t_1, \neg \text{Moving}(o^*))$$

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### **Argument 2**

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•••





$$\mathbf{P}(\mathbf{r}_1, t_0, \operatorname{At}(o^*, t_0, \ell_1))$$

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$$\succ_t^a$$

 $\mathbf{B}(\mathbf{r}_1, t_1, \operatorname{Moving}(o^*))$ 

$$\mathbf{B}^2(\mathbf{r}_1, t_1, \neg \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \text{NeedToBrake}(c))$$

### **Argument 2**

$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$

$$\mathbf{P}(\mathbf{r}_2, t_1, \text{At}(o^*, t_0, \ell_2))$$

$$\mathbf{P}(\mathfrak{r}_2,t_1,\ell_1=\ell_2)$$

•••

$$:: \mathbf{B}^5(\mathfrak{r}_2, t_1, \operatorname{Moving}(o^*))$$

$$: \mathbf{B}^5(\mathfrak{r}_2, t_1, \text{NeedToBrake}(c))$$





$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$

$$\mathbf{B}(\mathfrak{r}_1, t_1, \neg \text{Moving}(o^*))$$

$$\succ_t^a$$

 $\mathbf{B}(\mathbf{r}_1, t_1, \operatorname{Moving}(o^*))$ 

$$\mathbf{B}^2(\mathfrak{r}_1,t_1,\neg \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \text{NeedToBrake}(c)) \qquad \therefore \mathbf{B}^5(\mathfrak{r}_2, t_1, \text{NeedToBrake}(c))$$

### **Argument 2**

$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$

$$\mathbf{P}(\mathbf{r}_2, t_1, \text{At}(o^*, t_0, \ell_2))$$

$$\mathbf{P}(\mathfrak{r}_2,t_1,\ell_1=\ell_2)$$

$$\mathbf{B}^5(\mathbf{r}_2, t_1, \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^5(\mathfrak{r}_2, t_1, \text{NeedToBrake}(c))$$



## **Adjudicator**



$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$

$$\mathbf{B}(\mathfrak{r}_1, t_1, \neg \text{Moving}(o^*))$$

$$\succ_t^a$$

$$\mathbf{B}(\mathfrak{r}_1, t_1, \operatorname{Moving}(o^*))$$

$$\mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \operatorname{Moving}(o^*))$$

$$\mathbf{B}^2(\mathfrak{r}_1,t_1,\neg \mathrm{NeedToBrake}(c))$$

### **Argument 2**

$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$

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$$\mathbf{P}(\mathfrak{r}_2, t_1, \ell_1 = \ell_2)$$

$$\mathbf{B}^5(\mathbf{r}_2, t_1, \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \text{NeedToBrake}(c)) \qquad \therefore \mathbf{B}^5(\mathfrak{r}_2, t_1, \text{NeedToBrake}(c))$$



## **Adjudicator**

 $\mathbf{B}(\mathfrak{a}, t_1, \text{NeedToBrake}(c))$ 



$$\mathbf{P}(\mathfrak{r}_1,t_0,\operatorname{At}(o^*,t_0,\ell_1))$$

•••

$$\mathbf{B}(\mathfrak{r}_1, t_1, \neg \text{Moving}(o^*))$$

$$\succ_t^a$$

 $\mathbf{B}(\mathbf{r}_1, t_1, \operatorname{Moving}(o^*))$ 

$$\mathbf{B}^2(\mathbf{r}_1, t_1, \neg \operatorname{Moving}(o^*))$$

$$\therefore \mathbf{B}^2(\mathfrak{r}_1, t_1, \neg \text{NeedToBrake}(c))$$

### **Argument 2**

$$\mathbf{P}(\mathbf{r}_2, t_0, \text{At}(o^*, t_0, \ell_1))$$

$$\mathbf{P}(\mathbf{r}_2, t_1, \text{At}(o^*, t_0, \ell_2))$$

$$\mathbf{P}(\mathfrak{r}_2,t_1,\ell_1=\ell_2)$$

•••

$$\mathbf{B}^5(\mathbf{r}_2, t_1, \operatorname{Moving}(o^*))$$

$$: \mathbf{B}^5(\mathfrak{r}_2, t_1, \text{NeedToBrake}(c))$$

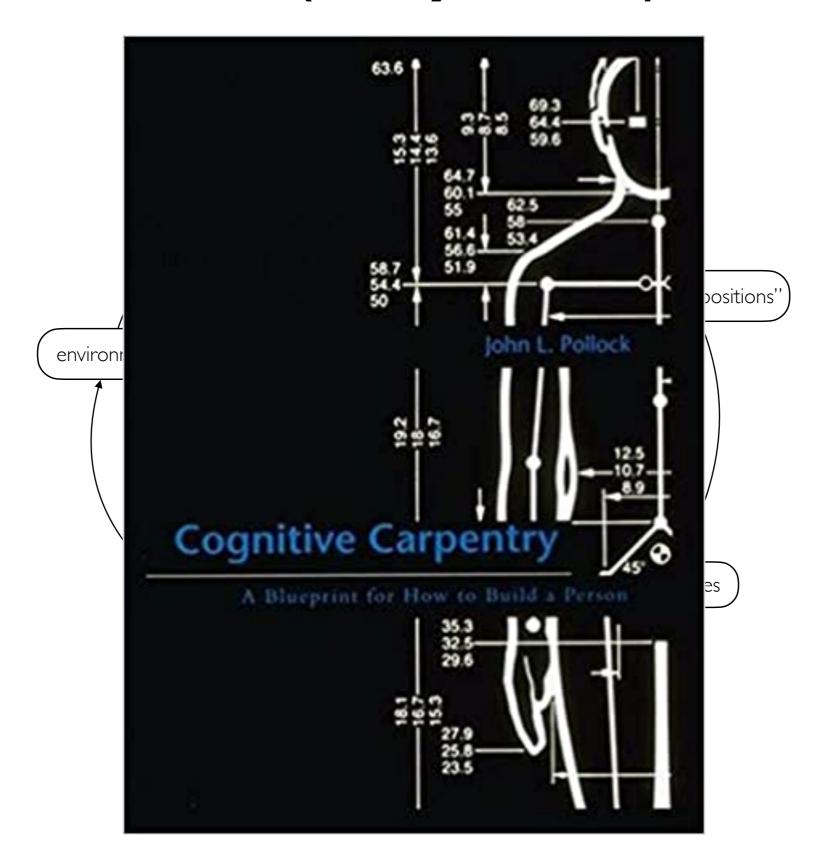
### **Adjudicator**

 $\mathbf{B}(\mathfrak{a}, t_1, \text{NeedToBrake}(c))$ 

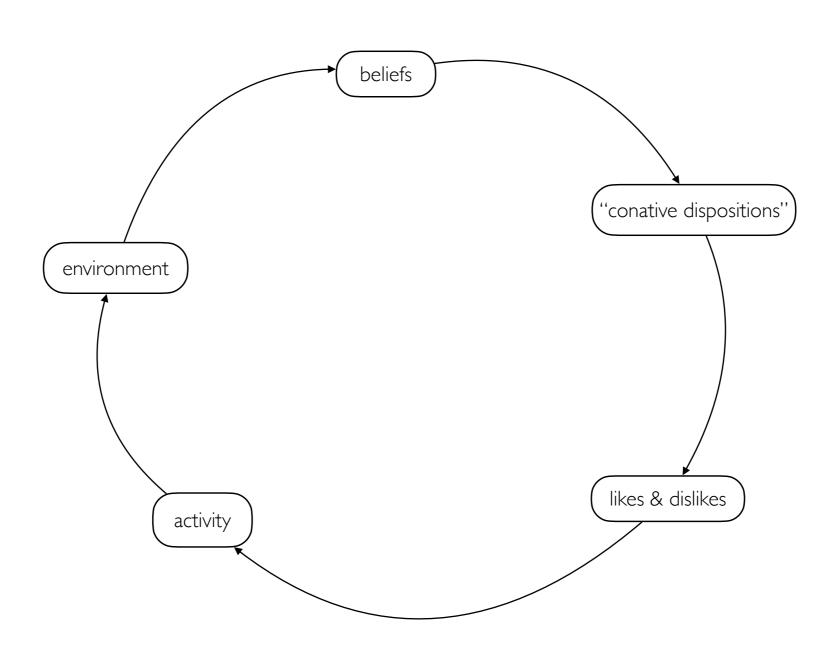


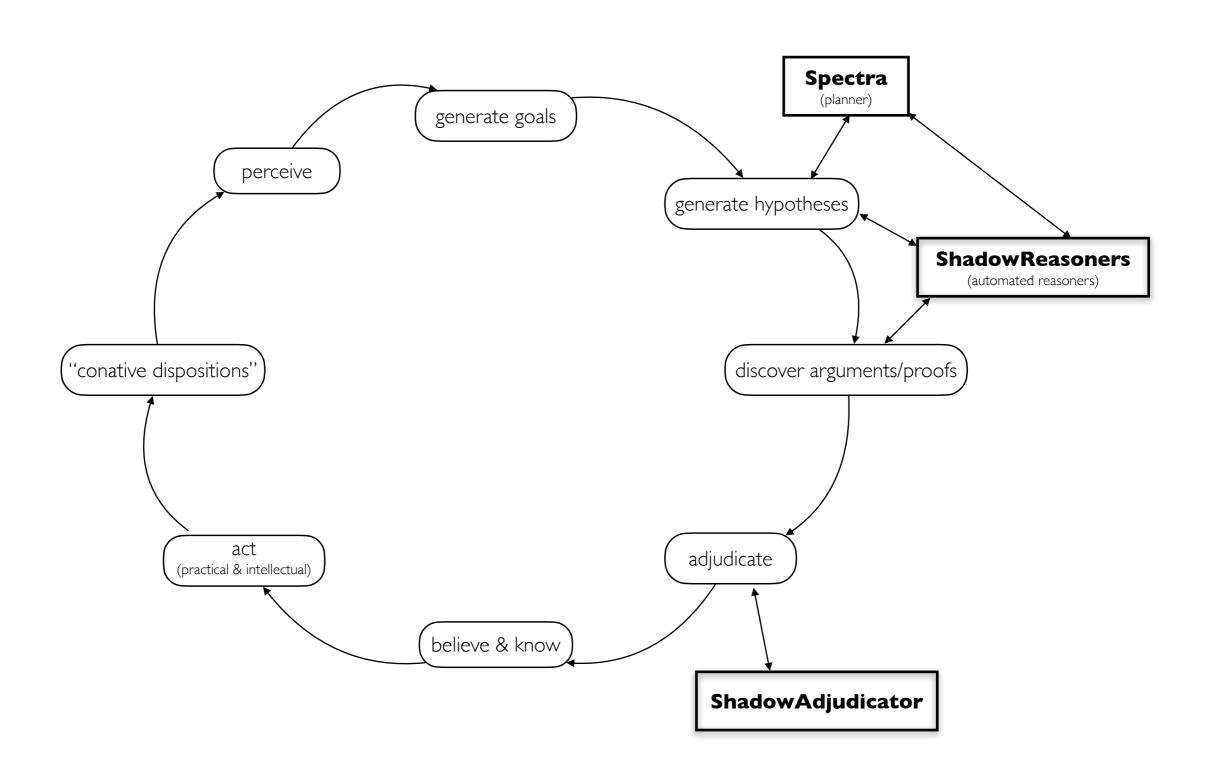


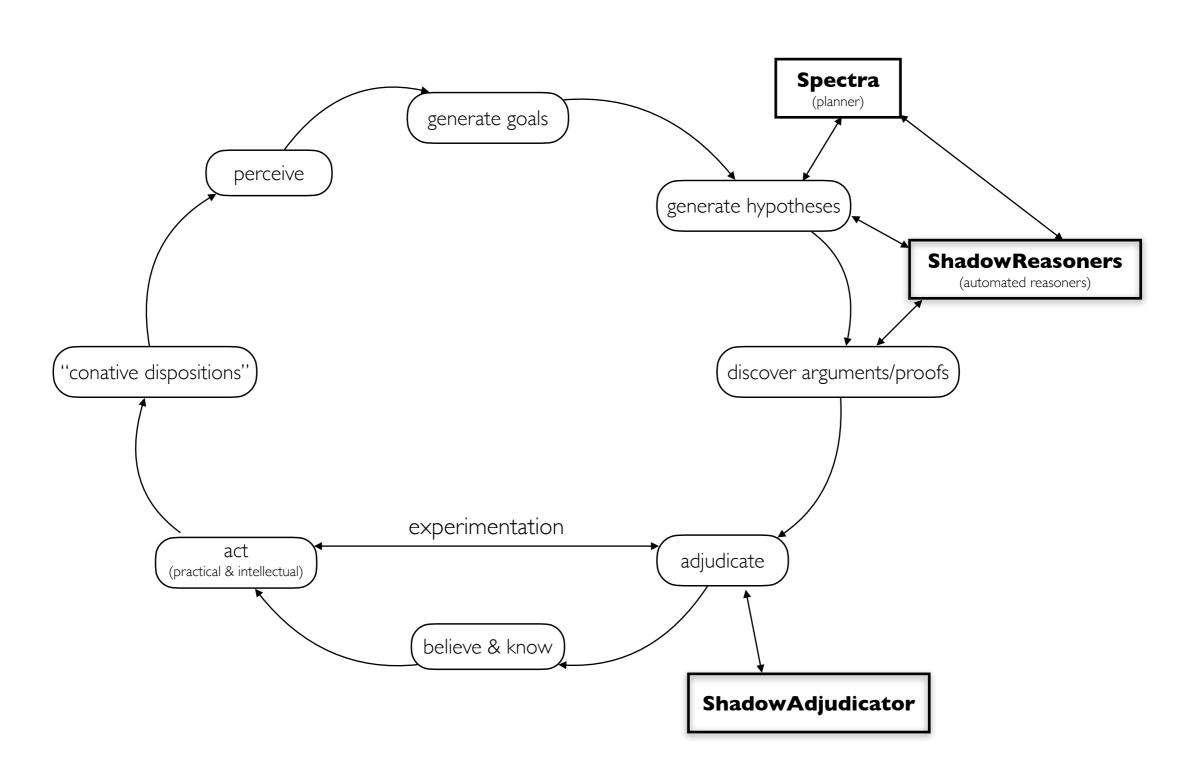
## Pollock's (Polyanna) Loop

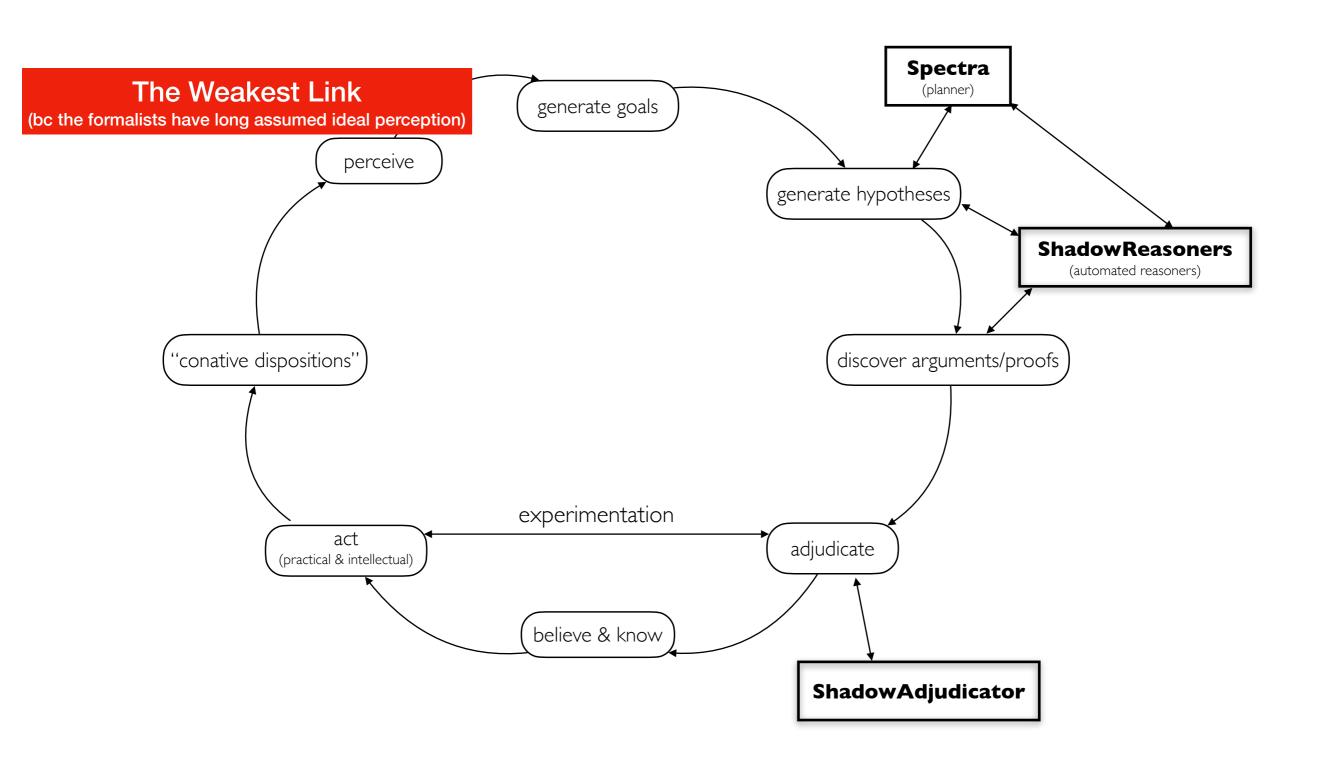


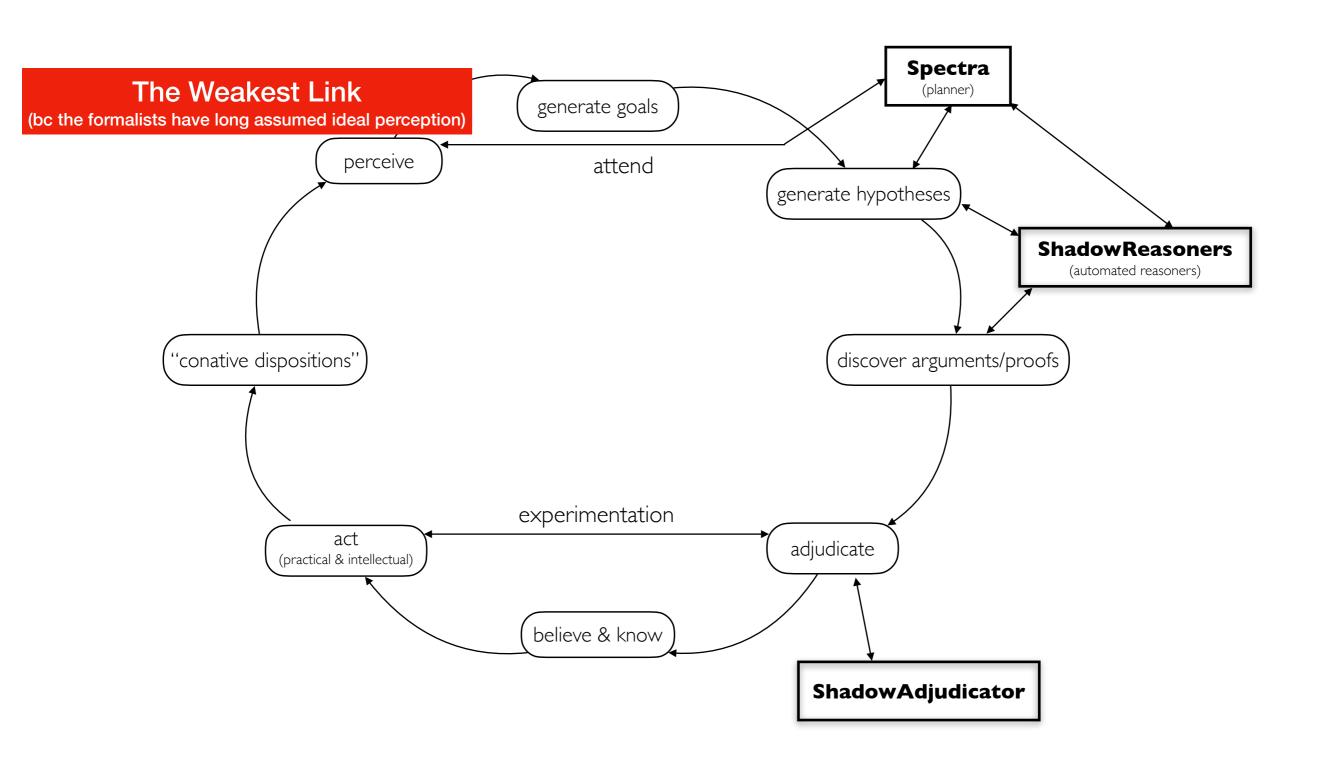
# Pollock's (Polyanna) Loop

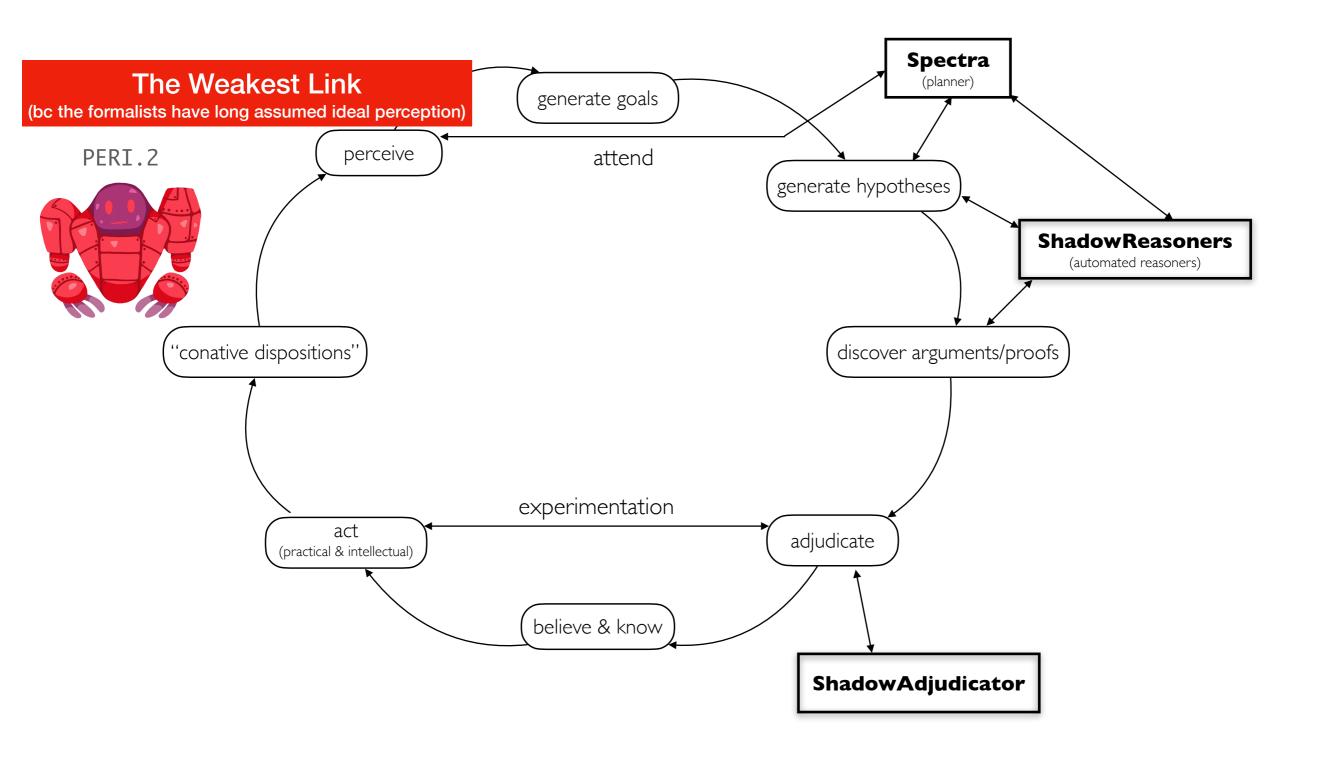


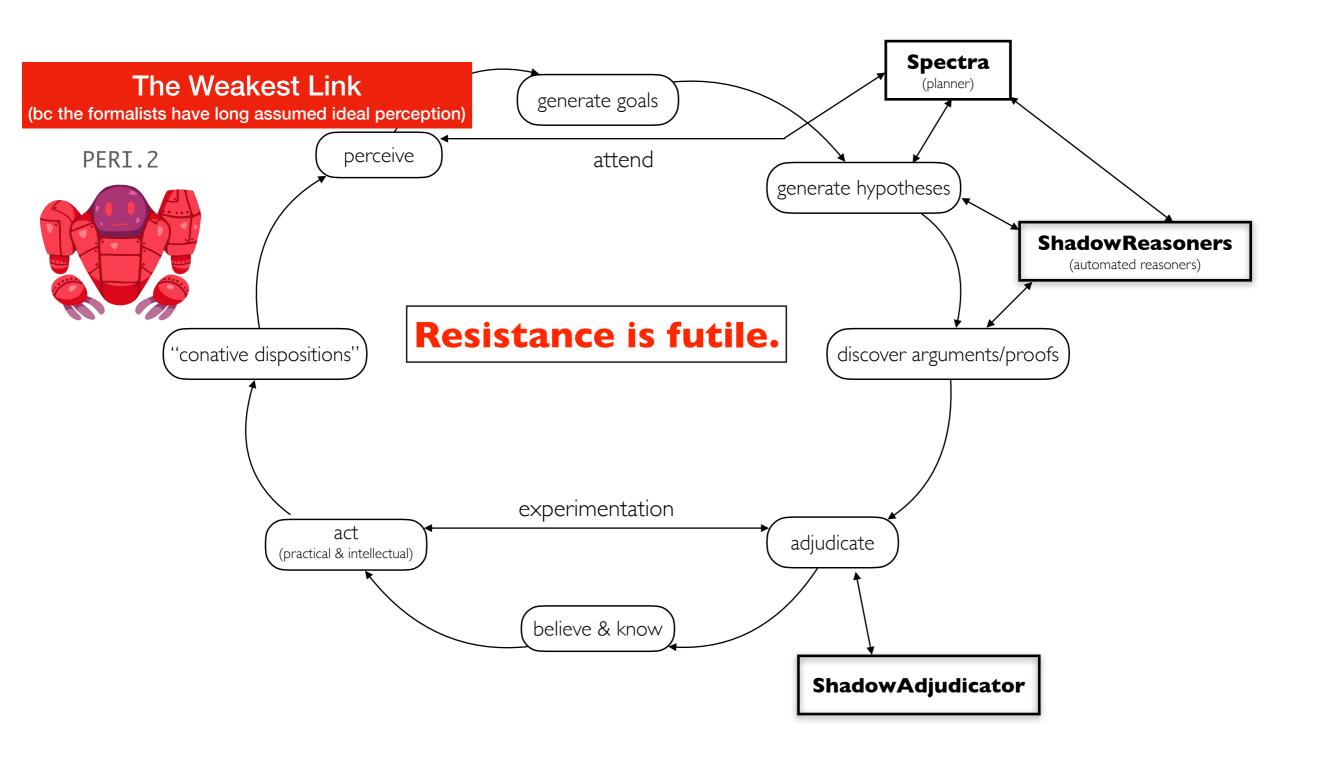












Certain **Evident** Overwhelmingly Likely Beyond Reasonable Doubt Likely More Likely Than Not Counterbalanced More Unlikely Than Not Unlikely Overwhelmingly Unlikely Beyond Reasonable Belief **Evidently False** Certainly False

Certain **Evident** Overwhelmingly Likely Beyond Reasonable Doubt Likely More Likely Than Not Counterbalanced More Unlikely Than Not Unlikely Overwhelmingly Unlikely

Beyond Reasonable Belief

**Evidently False** 

Certainly False

**Epistemically Positive** 

Certain

**Evident** 

Overwhelmingly Likely

Beyond Reasonable Doubt

Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

Overwhelmingly Unlikely

Beyond Reasonable Belief

Evidently False

Certainly False

**Epistemically Positive** 

Certain

**Evident** 

Overwhelmingly Likely

Beyond Reasonable Doubt

Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

Overwhelmingly Unlikely

Beyond Reasonable Belief

**Evidently False** 

Certainly False

**Epistemically Negative** 

**Epistemically Positive** 

Certain

Evident

Overwhelmingly Likely

Beyond Reasonable Doubt

Likely

More Likely Than Not

Counterbalanced

More Unlikely Than Not Unlikely

Overwhelmingly Unlikely

Beyond Reasonable Belief

**Evidently False** 

Certainly False

Epistemically Negative

Epistemically Positive	(6)	Certain	
	(5)	Evident	
		Overwhelmingly Likely	
	(3)	Beyond Reasonable Doubt	
	(2)	Likely	
	(1)	More Likely Than Not	
	(0)	Counterbalanced	
	(-1)	More Unlikely Than Not	
	(-2)	Unlikely	
	(-3)	Overwhelmingly Unlikely	
	(-4)	Beyond Reasonable Belief	
Epistemically Negative	(-5)	Evidently False	
	(-6)	Certainly False	

**(I)** 

(0)

**Epistemically Positive** 

$$\mathbf{B}^{1 \le \sigma \le 6}(\mathfrak{a}, t, \phi...)$$

(6) Certain

(5) Evident

(4) Overwhelmingly Likely

(3) Beyond Reasonable Doubt

(2) Likely

More Likely Than Not

Counterbalanced

(-1) More Unlikely Than Not

(-2) Unlikely

(-3) Overwhelmingly Unlikely

(-4) Beyond Reasonable Belief

(-5) Evidently False

(-6) Certainly False

Epistemically Negative

### $\mathcal{DCEC}$ Signature

```
S ::= \mathsf{Agent} \mid \mathsf{ActionType} \mid \mathsf{Action} \sqsubseteq \mathsf{Event} \mid \mathsf{Moment} \mid \mathsf{Fluent}
\begin{cases} \mathit{action} : \mathsf{Agent} \times \mathsf{ActionType} \to \mathsf{Action} \\ \mathit{initially} : \mathsf{Fluent} \to \mathsf{Formula} \\ \mathit{holds} : \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{happens} : \mathsf{Event} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{clipped} : \mathsf{Moment} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{clipped} : \mathsf{Moment} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{terminates} : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{prior} : \mathsf{Moment} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{prior} : \mathsf{Moment} \times \mathsf{Moment} \to \mathsf{Formula} \\ \mathit{t} ::= \mathit{x} : \mathit{S} \mid \mathit{c} : \mathit{S} \mid \mathit{f} (\mathit{t}_{1}, \ldots, \mathit{t}_{n}) \\ \mathsf{q} : \mathsf{Formula} \mid \neg \phi \mid \phi \land \psi \mid \phi \lor \psi \mid \forall \mathit{x} : \phi(\mathit{x}) \mid \exists \mathit{x} : \phi(\mathit{x}) \\ \mathsf{P}(\mathit{a}, \mathit{t}, \phi) \mid \mathsf{K}(\mathit{a}, \mathit{t}, \phi) \mid \mathsf{S}(\mathit{a}, \mathit{b}, \mathit{t}, \phi) \mid \mathsf{S}(\mathit{a}, \mathit{t}, \phi) \\ \mathsf{C}(\mathit{t}, \phi) \mid \mathsf{B}(\mathit{a}, \mathit{t}, \phi) \mid \mathsf{D}(\mathit{a}, \mathit{t}, \phi) \mid \mathsf{I}(\mathit{a}, \mathit{t}, \phi) \\ \mathsf{O}(\mathit{a}, \mathit{t}, \phi, (\neg) \mathit{happens}(\mathit{action}(\mathit{a}^*, \alpha), \mathit{t}')) \\ \\ \mathsf{Modal Operator Descriptors:} \\ \mathsf{Perceives}, \mathsf{Knows}, \mathsf{Says}, \mathsf{Common-knowledge} \\ \mathsf{Believes}, \mathsf{Desires}, \mathsf{Intends}, \mathsf{Ought-to} \end{aligned}
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#### **DCEC** Inference Schemata

$$\frac{\mathbf{K}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{K}(a,t_2,\phi)} \quad [I_{\mathbf{K}}] \quad \frac{\mathbf{B}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{B}(a,t_2,\phi)} \quad [I_{\mathbf{B}}]$$

$$\overline{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))} \quad [I_{1}] \quad \overline{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))} \quad [I_{2}]$$

$$\frac{\mathbf{C}(t,\phi) \ t \leq t_1 \dots t \leq t_n}{\mathbf{K}(a_1,t_1,\dots \mathbf{K}(a_n,t_n,\phi)\dots)} \quad [I_{3}] \quad \frac{\mathbf{K}(a,t,\phi)}{\phi} \quad [I_{4}]$$

$$\overline{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{K}(a,t_2,\phi_1) \to \mathbf{K}(a,t_3,\phi_2)}} \quad [I_{5}]$$

$$\overline{\mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{B}(a,t_2,\phi_1) \to \mathbf{B}(a,t_3,\phi_2)}} \quad [I_{6}]$$

$$\overline{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)}} \quad [I_{7}]$$

$$\overline{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)}} \quad [I_{9}]$$

$$\overline{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \neg \phi_1)} \quad [I_{9}]$$

$$\overline{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \neg \phi_1)}} \quad [I_{9}]$$

$$\overline{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \neg \phi_1)}} \quad [I_{10}]$$

$$\overline{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi \to \psi)}} \quad [I_{11a}] \quad \overline{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\psi)}} \quad [I_{11b}]$$

$$\overline{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi)}} \quad [I_{12}] \quad \overline{\mathbf{B}(a,t,\phi) \mathbf{B}(a,t,\phi)}} \quad [I_{13}]$$

$$\overline{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi)} \quad [I_{12}] \quad \overline{\mathbf{B}(a,t,\phi,\chi)}} \quad [I_{14}]$$

### $\mathcal{DCEC}$ Signature

```
S ::= \mathsf{Agent} \mid \mathsf{ActionType} \mid \mathsf{Action} \sqsubseteq \mathsf{Event} \mid \mathsf{Moment} \mid \mathsf{Fluent} action : \mathsf{Agent} \times \mathsf{ActionType} \to \mathsf{Action} initially : \mathsf{Fluent} \to \mathsf{Formula} holds : \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} happens : \mathsf{Event} \times \mathsf{Moment} \to \mathsf{Formula} clipped : \mathsf{Moment} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} initiates : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} terminates : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \to \mathsf{Fluent} \to \mathsf{Formula} trice : \mathsf{Event} \times \mathsf{Fluent} \to \mathsf{Fluent} \to \mathsf{Formula}
```

Modal Operator Descriptors:
Perceives, Knows, Says, Common-knowledge
Believes, Desires, Intends, Ought-to

### **DCEC** Inference Schemata

$$\frac{\mathbf{K}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{K}(a,t_2,\phi)} \quad [I_{\mathbf{K}}] \quad \frac{\mathbf{B}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{B}(a,t_2,\phi)} \quad [I_{\mathbf{B}}]$$

$$\frac{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))}{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))} \quad [I_1] \quad \frac{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))}{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))} \quad [I_2]$$

$$\frac{\mathbf{C}(t,\phi) \ t \leq t_1 \dots t \leq t_n}{\mathbf{K}(a_1,t_1,\dots\mathbf{K}(a_n,t_n,\phi)\dots)} \quad [I_3] \quad \frac{\mathbf{K}(a,t,\phi)}{\phi} \quad [I_4]$$

$$\frac{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{K}(a,t_2,\phi_1) \to \mathbf{K}(a,t_3,\phi_2)}{\mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)} \quad [I_5]$$

$$\frac{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)}{\mathbf{C}(t,(\phi_1 \to \phi_2 \to \phi_2 \to \phi_1))} \quad [I_6]$$

$$\frac{\mathbf{C}(t,(\phi_1 \to \phi_1 \to \phi_1)) \quad [I_8] \quad \mathbf{C}(t,(\phi_1 \to \phi_2 \to \phi_2 \to \phi_1))}{\mathbf{C}(t,(\phi_1 \to \phi_2 \to \phi_2 \to \phi_1))} \quad [I_9]$$

$$\frac{\mathbf{C}(t,(\phi_1 \to \phi_1 \to \phi_1)) \quad [I_{10}] \quad \mathbf{E}(t_1,\phi_1 \to \phi_1)}{\mathbf{E}(t_1,\phi_1 \to \phi_1)} \quad [I_{10}] \quad \mathbf{E}(t_1,\phi_1 \to \phi_1) \quad [I_{11a}] \quad \mathbf{E}(t_1,\phi_1 \to \phi_1) \quad [I_{11b}]$$

$$\frac{\mathbf{C}(t,\phi_1 \to \phi_1 \to \phi_1)}{\mathbf{E}(t_1,\phi_1 \to \phi_1)} \quad [I_{11a}] \quad \mathbf{E}(t_1,\phi_1 \to \phi_1) \quad [I_{11b}] \quad [I_{11b}]$$

$$\frac{\mathbf{C}(t,\phi_1 \to \phi_1,\phi_1)}{\mathbf{E}(t_1,\phi_1 \to \phi_1)} \quad [I_{12}] \quad \mathbf{E}(t_1,\phi_1,\phi_1) \quad [I_{12}] \quad [I_{13}] \quad [I_{14}] \quad [I_{14}]$$

#### $\mathcal{DCEC}$ Signature

```
S ::= \operatorname{Agent} \mid \operatorname{ActionType} \mid \operatorname{Action} \sqsubseteq \operatorname{Event} \mid \operatorname{Moment} \mid \operatorname{Fluent} action : \operatorname{Agent} \times \operatorname{ActionType} \to \operatorname{Action} initially : \operatorname{Fluent} \to \operatorname{Formula} holds : \operatorname{Fluent} \times \operatorname{Moment} \to \operatorname{Formula} happens : \operatorname{Event} \times \operatorname{Moment} \to \operatorname{Formula} clipped : \operatorname{Moment} \times \operatorname{Fluent} \times \operatorname{Moment} \to \operatorname{Formula} initiates : \operatorname{Event} \times \operatorname{Fluent} \times \operatorname{Moment} \to \operatorname{Formula} terminates : \operatorname{Event} \times \operatorname{Fluent} \times \operatorname{Moment} \to \operatorname{Formula} trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid c : S \mid f(t_1, \dots, t_n) trice : S \mid f(t_1, \dots, t_n) trice : S \mid f(t_1, \dots, t_n) trice : S \mid f(t_1, \dots, t_n) tric
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Modal Operator Descriptors:
Perceives, Knows, Says, Common-knowledge
Believes, Desires, Intends, Ought-to

### **DCEC** Inference Schemata

$$\frac{\mathbf{K}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{K}(a,t_2,\phi)} \quad [I_{\mathbf{K}}] \quad \frac{\mathbf{B}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{B}(a,t_2,\phi)} \quad [I_{\mathbf{B}}]$$

$$\frac{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))}{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))} \quad [I_1] \quad \frac{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))}{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))} \quad [I_2]$$

$$\frac{\mathbf{C}(t,\phi) \ t \leq t_1 \dots t \leq t_n}{\mathbf{K}(a_1,t_1,\dots\mathbf{K}(a_n,t_n,\phi)\dots)} \quad [I_3] \quad \frac{\mathbf{K}(a,t,\phi)}{\phi} \quad [I_4]$$

$$\frac{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{K}(a,t_2,\phi_1) \to \mathbf{K}(a,t_3,\phi_2)}{\mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)} \quad [I_6]$$

$$\frac{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)}{\mathbf{C}(t,\forall x.\phi \to \phi[x \mapsto t])} \quad \frac{[I_8]}{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \neg \phi_1)} \quad [I_9]$$

$$\frac{\mathbf{C}(t,[\phi_1 \land \dots \land \phi_n \to \phi] \to [\phi_1 \to \dots \to \phi_n \to \phi])}{\mathbf{C}(t,[\phi_1 \land \dots \land \phi_n \to \phi] \to [h_1]} \quad \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi)}{\mathbf{B}(a,t,\phi)} \quad [I_{11a}] \quad \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi)}{\mathbf{B}(a,t,\phi \land \psi)} \quad [I_{11b}]$$

$$\frac{\mathbf{S}(s,h,t,\phi)}{\mathbf{B}(h,t,\mathbf{B}(s,t,\phi))} \quad [I_{12}] \quad \frac{\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))}{\mathbf{P}(a,t,happens(action(a^*,\alpha),t'))} \quad [I_{13}]$$

$$\frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\mathbf{O}(a,t,\phi,\chi))}{\mathbf{K}(a,t,\mathbf{I}(a,t,\chi))} \quad \mathbf{O}(a,t,\phi,\chi)} \quad [I_{14}]$$

#### 2.2 Inductive Deontic Cognitive Event Calculus

 $\mathcal{DCEC}$  employs no uncertainty system (e.g., probability measures, *strength factors*, or likelihood measures) and hence is purely deductive. Therefore, as we wish to enable our agents to reason about situations involving uncertainty, we must ultimately utilize the *Inductive*  $\mathcal{DCEC}$ :  $\mathcal{IDCEC}$ .

In general, to go from a deductive to an inductive cognitive calculus, we require two components: (1) an uncertainty system, and (2) inference schemata that delineate the methods by which inferences linking formulae and other information can be used to build formally valid arguments.

The particular uncertainty system we use herein is discussed in §2.3. The inference schemata of  $\mathcal{IDCEC}$  consist of the union of the set presented in §2.1 with that in the box titled **Additional Inference Schemata for**  $\mathcal{IDCEC}$ . Likewise, the signature of  $\mathcal{IDCEC}$  subsumes that of the deductive  $\mathcal{DCEC}$ ; the syntax of  $\mathcal{IDCEC}$  also includes the forms given in the box titled **Additional Syntax for**  $\mathcal{IDCEC}$ .

### Additional Syntax for IDCEC

$$\phi ::= \left\{ \mathbf{B}^{\sigma}(a,t,\phi) \right.$$
 where  $\sigma \in [-5,-4,\ldots,4,5]$ 

#### Additional Inference Schemata for $\mathcal{IDCEC}$

$$\frac{\mathbf{P}(a,t_1,\phi_1),\ \Gamma\vdash t_1 < t_2}{\mathbf{B}^4(a,t_2,\phi)}\ [I_{\mathbf{P}}^s]$$
 
$$\frac{\mathbf{B}^{\sigma_1}(a,t_1,\phi_1),\ldots,\mathbf{B}^{\sigma_m}(a,t_m,\phi_m),\{\phi_1,\ldots,\phi_m\}\vdash \phi,\{\phi_1,\ldots,\phi_m\}\nvdash \zeta,\Gamma\vdash t_i < t}{\mathbf{B}^{\min(\sigma_1,\ldots,\sigma_m)}(a,t,\phi)}$$
 where  $\sigma\in[0,1,\ldots,5,6]$  
$$\frac{\mathbf{C}(t,\mathbf{B}^{-\sigma}(a,t,\phi)\leftrightarrow\mathbf{B}^{\sigma}(a,t,\neg\phi))}{[I_{\neg}^s]}$$

Briefly,  $\mathbf{B}^{\sigma}(a, t, \phi)$  denotes that agent a at time t believes  $\phi$  with uncertainty  $\sigma$ . We justify in the next section the range of values for  $\sigma$ .

The first inference schema allows agents to infer evident beliefs ( $\sigma = 4$ , as defined in the next section) from what they perceive [5] The second schema allows agents to infer a belief that is provable from the beliefs they currently assert, so long as the belief set is not inconsistent. In practice, we usually check that the belief set is consistent by attempting to prove a reserved propositional atom  $\zeta$  which does not

#### $\mathcal{DCEC}$ Signature

```
S ::= \mathsf{Agent} \mid \mathsf{ActionType} \mid \mathsf{Action} \sqsubseteq \mathsf{Event} \mid \mathsf{Moment} \mid \mathsf{Fluent}
action : \mathsf{Agent} \times \mathsf{ActionType} \to \mathsf{Action}
initially : \mathsf{Fluent} \to \mathsf{Formula}
holds : \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
happens : \mathsf{Event} \times \mathsf{Moment} \to \mathsf{Formula}
clipped : \mathsf{Moment} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
initiates : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
terminates : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Moment} \to \mathsf{Formula}
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trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
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trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
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trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
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trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
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trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula}
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trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Fluent} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Fluent} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Fluent} \to \mathsf{Fluent} \to \mathsf{Fluent} \to \mathsf{Formula}
trice : \mathsf{Event} \times \mathsf{Fluent} \to \mathsf{Fluent}
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Modal Operator Descriptors:
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### **DCEC** Inference Schemata

$$\frac{\mathbf{K}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{K}(a,t_2,\phi)} \quad [I_{\mathbf{K}}] \quad \frac{\mathbf{B}(a,t_1,\Gamma), \ \Gamma \vdash \phi, \ t_1 \leq t_2}{\mathbf{B}(a,t_2,\phi)} \quad [I_{\mathbf{B}}]$$

$$\frac{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))}{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \to \mathbf{K}(a,t,\phi))} \quad [I_1] \quad \frac{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))}{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \to \mathbf{B}(a,t,\phi))} \quad [I_2]$$

$$\frac{\mathbf{C}(t,\phi) \ t \leq t_1 \dots t \leq t_n}{\mathbf{K}(a_1,t_1,\dots\mathbf{K}(a_n,t_n,\phi)\dots)} \quad [I_3] \quad \frac{\mathbf{K}(a,t,\phi)}{\phi} \quad [I_4]$$

$$\frac{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{K}(a,t_2,\phi_1) \to \mathbf{K}(a,t_3,\phi_2)}{\mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)} \quad [I_6]$$

$$\frac{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \to \phi_2)) \to \mathbf{C}(t_2,\phi_1) \to \mathbf{C}(t_3,\phi_2)}{\mathbf{C}(t,\forall x.\phi \to \phi[x \mapsto t])} \quad \frac{[I_8]}{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \neg \phi_1)} \quad [I_9]$$

$$\frac{\mathbf{C}(t,[\phi_1 \land \dots \land \phi_n \to \phi] \to [\phi_1 \to \dots \to \phi_n \to \phi])}{\mathbf{C}(t,[\phi_1 \land \dots \land \phi_n \to \phi] \to [h_1]} \quad \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi)}{\mathbf{B}(a,t,\phi)} \quad [I_{11a}] \quad \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\phi)}{\mathbf{B}(a,t,\phi \land \psi)} \quad [I_{11b}]$$

$$\frac{\mathbf{S}(s,h,t,\phi)}{\mathbf{B}(h,t,\mathbf{B}(s,t,\phi))} \quad [I_{12}] \quad \frac{\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))}{\mathbf{P}(a,t,happens(action(a^*,\alpha),t'))} \quad [I_{13}]$$

$$\frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\mathbf{O}(a,t,\phi,\chi))}{\mathbf{K}(a,t,\mathbf{I}(a,t,\chi))} \quad \mathbf{O}(a,t,\phi,\chi)} \quad [I_{14}]$$

### 2 Inductive Deontic Cognitive Event Calculus

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#### Additional Inference Schemata for $\mathcal{IDCEC}$

$$\frac{\mathbf{P}(a,t_1,\phi_1),\ \Gamma\vdash t_1 < t_2}{\mathbf{B}^4(a,t_2,\phi)} \ [I_{\mathbf{P}}^s]$$
 
$$\frac{\mathbf{B}^{\sigma_1}(a,t_1,\phi_1),\ldots,\mathbf{B}^{\sigma_m}(a,t_m,\phi_m),\{\phi_1,\ldots,\phi_m\}\vdash\phi,\{\phi_1,\ldots,\phi_m\}\vdash\zeta,\Gamma\vdash t_i < t}{\mathbf{B}^{\min(\sigma_1,\ldots,\sigma_m)}(a,t,\phi)} \ \text{where } \sigma\in[0,1,\ldots,5,6]$$
 
$$\frac{\mathbf{C}(t,\mathbf{B}^{-\sigma}(a,t,\phi)\leftrightarrow\mathbf{B}^{\sigma}(a,t,\neg\phi))}{\mathbf{C}(t,\mathbf{B}^{-\sigma}(a,t,\phi)\leftrightarrow\mathbf{B}^{\sigma}(a,t,\neg\phi))} \ [I_{\gamma}^s]$$

Briefly,  $\mathbf{B}^{\sigma}(a, t, \phi)$  denotes that agent a at time t believes  $\phi$  with uncertainty  $\sigma$ . We justify in the next section the range of values for  $\sigma$ .

The first inference schema allows agents to infer evident beliefs ( $\sigma = 4$ , as defined in the next section) from what they perceive [5] The second schema allows agents to infer a belief that is provable from the beliefs they currently assert, so long as the belief set is not inconsistent. In practice, we usually check that the belief set is consistent by attempting to prove a reserved propositional atom  $\zeta$  which does not

Loggik kan hjelpe deg å leve for alltid.