

Propositional Calculus III:

Reductio ad Absurdum

Selmer Bringsjord

Rensselaer AI & Reasoning (RAIR) Lab
Department of Cognitive Science
Department of Computer Science
Lally School of Management & Technology
Rensselaer Polytechnic Institute (RPI)
Troy, New York 12180 USA

Intro to Logic
2/4/2019



Logistics ...

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Any questions?

E-Housekeeping Pts, Again

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- Make sure OS fully up-to-date.

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- Chrome best (but I use Safari).

E-Housekeeping Pts, Again

- Make sure OS fully up-to-date.
- Make sure browser fully up-to-date.
- Chrome best (but I use Safari).
- Always work in the same browser window with multiple tabs; must do this with email and HyperGrader & HyperSlate.

Couple of HyperSlate Pts

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- Consider using right-clicking on nodes.

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- “Where are my files?”

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- We are not covering all inference rules/schemata in class.

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- “Where are my files?”
- We are not covering all inference rules/schemata in class.
- conditional intro, briefly

Schedule ...

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How many have taken
a programming course?

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How many have learned a
logic-programming language?

Schedule ...

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Thursday ...

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Reductio ...

“Reductio ad absurdum, which Euclid loved so much, is one of a mathematician's finest weapons. It is a far finer gambit than any chess gambit: a chess player may offer the sacrifice of a pawn or even a piece, but a mathematician offers the game.”

–G. H. Hardy

A Greek-shocking Example ...

$$\frac{p}{q}$$

What are rational numbers?

$$\frac{p}{q}$$

What are rational numbers?

- Any number that can be expressed in the form of $\frac{p}{q}$ such that we have a numerator p and a non-zero denominator.
- Rational numbers are a subset of real numbers!
- Examples of *irrational* numbers?

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$$\sqrt{2}$$

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$$\sqrt{2} \quad e$$

Prove that:

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$$\sqrt{2}$$

Prove that:

$\sqrt{2}$ is irrational

Suppose $\sqrt{2}$ is **rational**. That means it can be written as the ratio of two integers p and q

$$\sqrt{2} = \frac{p}{q} \quad (1)$$

where we may assume that p and q **have no common factors**. (If there are any common factors we cancel them in the numerator and denominator.) Squaring in (1) on both sides gives

$$2 = \frac{p^2}{q^2} \quad (2)$$

which implies

$$p^2 = 2q^2 \quad (3)$$

Thus p^2 is even. The only way this can be true is that p itself is even. But then p^2 is actually divisible by 4. Hence q^2 and therefore q must be even. So p and q are both even which is a contradiction to our assumption that they have no common factors. The square root of 2 cannot be rational!

And now, what are *prime* numbers?

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- A number that can be divided by only two numbers, one and itself.
- Must be a whole number.
- Example: 2,3,5,7.....

And recall: Euclidean “Magic”

Theorem: There are infinitely many primes.

Proof: We take an indirect route. Let $\Pi = p_1 = 2, p_2 = 3, p_3 = 5, \dots, p_k$ be a finite, exhaustive consecutive sequence of prime numbers. Next, let M_Π be $p_1 \times p_2 \times \dots \times p_k$, and set M'_Π to $M_\Pi + 1$. Either M'_Π is prime, or not; we thus have two (exhaustive) cases to consider.

- C1 Suppose M'_Π is prime. In this case we immediately have a prime number beyond any in Π — contradiction!
- C2 Suppose on the other hand that M'_Π is *not* prime. Then some prime p divides M'_Π . (Why?) Now, p itself is either in Π , or not; we hence have two sub-cases. Supposing that p is in Π entails that p divides M_Π . But we are operating under the supposition that p divides M'_Π as well. This implies that p divides 1, which is absurd (a contradiction). Hence the prime p is outside Π .

Hence for *any* such list Π , there is a prime outside the list. That is, there are infinitely many primes. **QED**

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Notice that this proof uses two applications of indirect proof, and one application of proof by cases.

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Study it word by word until you endorse it with your very soul!

From Algebra 2 in High School ...

(Pearson Common-Core Compliant Textbook)

A proof involving indirect reasoning is an **indirect proof**. Often in an indirect proof, a statement and its negation are the only possibilities. When you see that one of these possibilities leads to a conclusion that contradicts a fact you know to be true, you can eliminate that possibility. For this reason, indirect proof is sometimes called *proof by contradiction*.

TAKE NOTE Key Concept

Writing an Indirect Proof

- Step 1** State as a temporary assumption the opposite (negation) of what you want to prove.
- Step 2** Show that this temporary assumption leads to a contradiction.
- Step 3** Conclude that the temporary assumption must be false and that what you want to prove must be true.

Problem 3 Writing an Indirect Proof

Proof

Given: $\triangle ABC$ is scalene.

Prove: $\angle A$, $\angle B$, and $\angle C$ all have different measures.

THINK

Assume temporarily the opposite of what you want to prove.

Show that this assumption leads to a contradiction.

Conclude that the temporary assumption must be false and that what you want to prove must be true.

WRITE

Assume temporarily that two angles of $\triangle ABC$ have the same measure. Assume that $m\angle A = m\angle B$.

By the Converse of the Isosceles Triangle Theorem, the sides opposite $\angle A$ and $\angle B$ are congruent. This contradicts the given information that $\triangle ABC$ is scalene.

The assumption that two angles of $\triangle ABC$ have the same measure must be false. Therefore, $\angle A$, $\angle B$, and $\angle C$ all have different measures.

▶ SHOW SOLUTION

▶ GOT IT?

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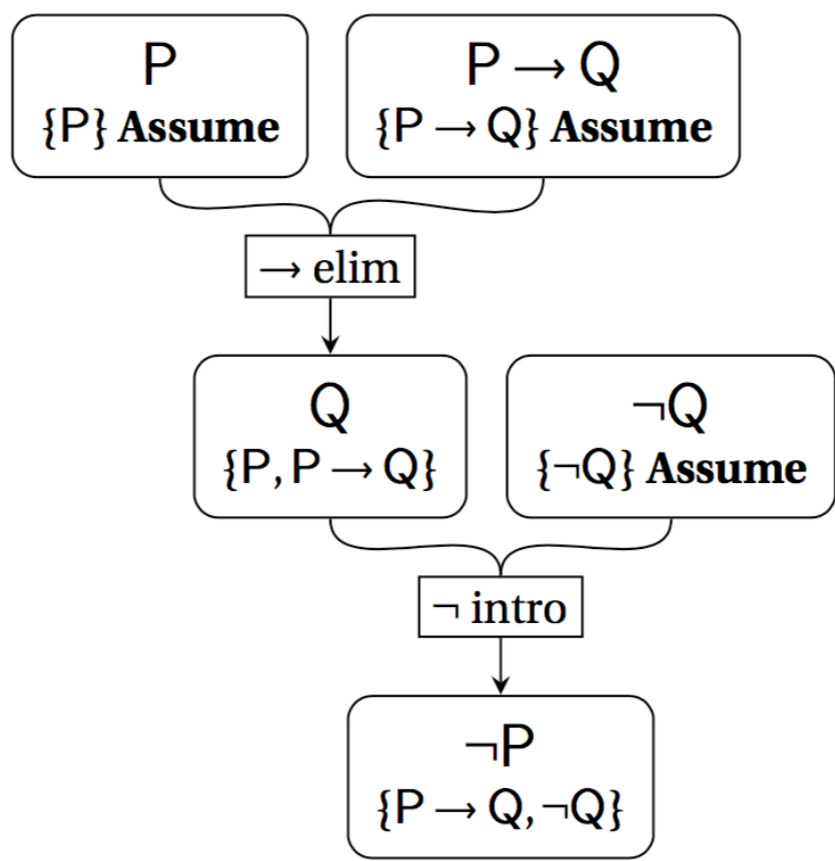
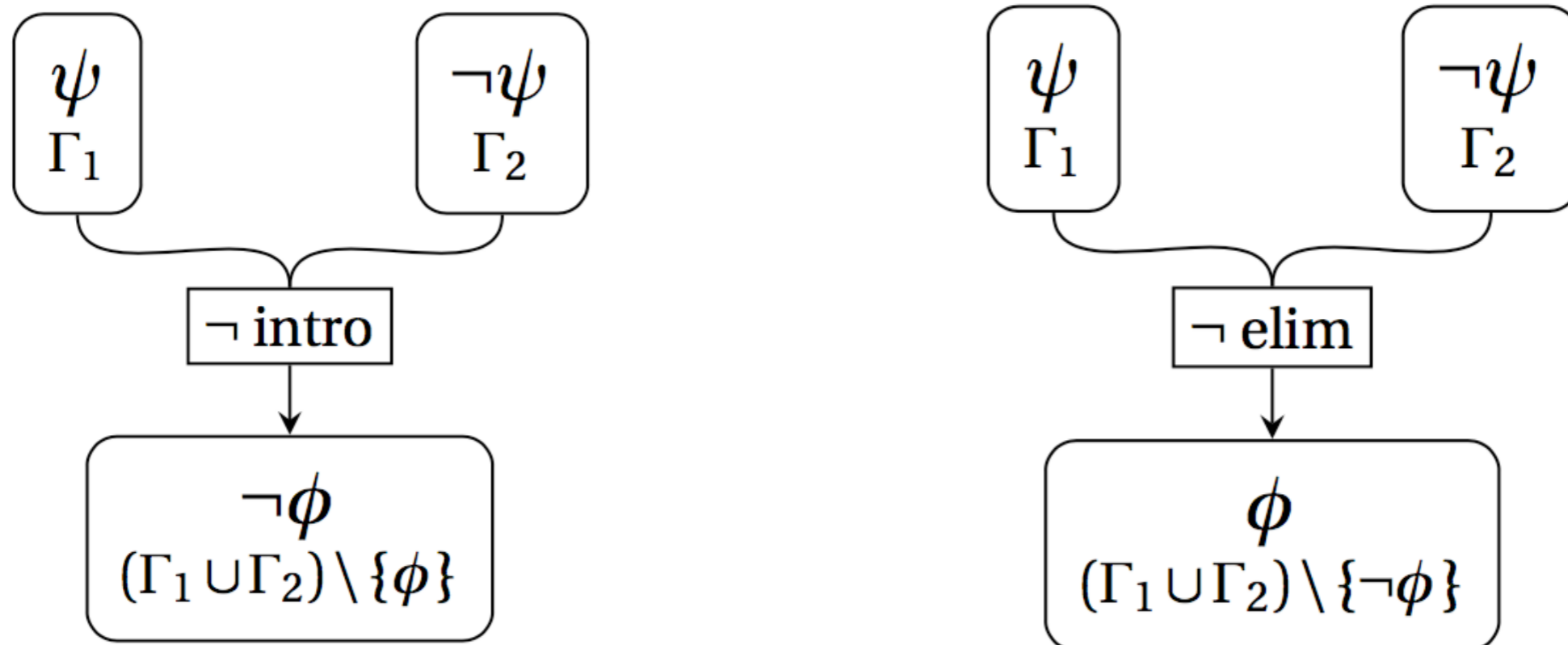
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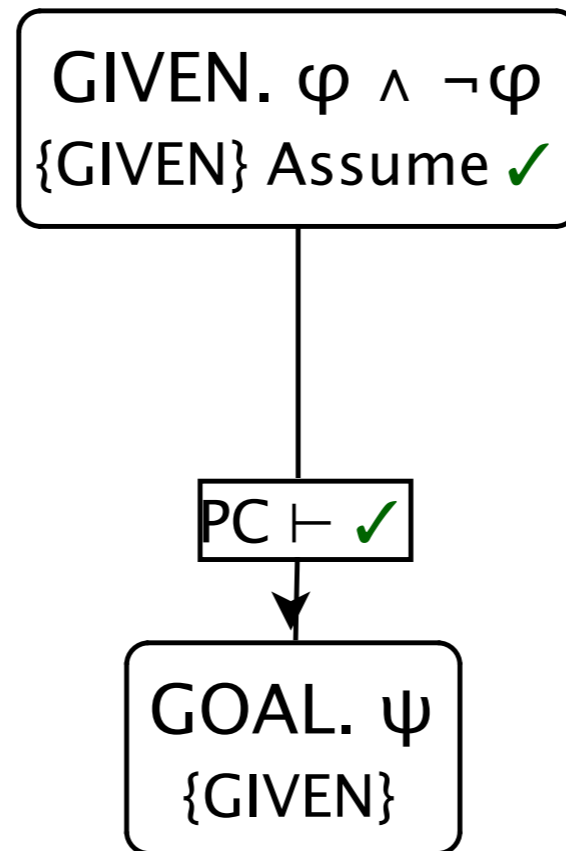
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  Journal = {Educational Studies in Mathematics},  
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  Title = {Making the Transition to Formal Proof},  
  Volume = {27.3},  
  Year = 1994}
```

▶ SHOW SOLUTION

▶ GOT IT?



Explosion



Explosion: Partial Proof Plan

Slate - explosion.slt

GIVEN. $\varphi \wedge \neg\varphi$
{GIVEN} Assume ✓

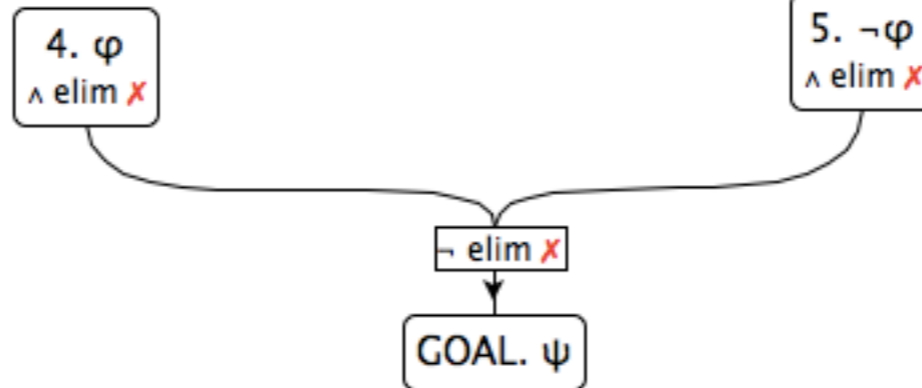
3. $\neg\psi$
{3} Assume ✓

4. φ
 \wedge elim ✗

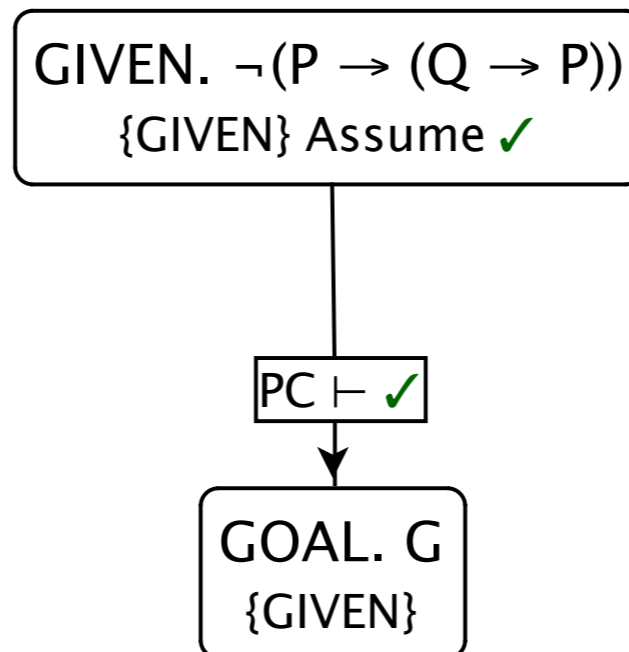
5. $\neg\varphi$
 \wedge elim ✗

\neg elim ✗

GOAL. ψ



GreenCheeseMoon I



GreenCheeseMoon I: Partial Proof Plan

Slate - GreenCheeseMoon1.slt

GIVEN. $\neg(P \rightarrow (Q \rightarrow P))$
{GIVEN} Assume ✓

Sub-Proof Here

5. P
{5} Assume ✓

3. $P \rightarrow (Q \rightarrow P)$
PC = ✓

4. $\neg G$
{4} Assume ✓

GOAL. G
 \neg elim ✗