### Logicist Machine Ethics Can Save Us

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Intro to Logic 4/11/2019 (includes planned class for 4/15/19; see syllabus)



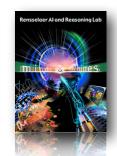


# Not quite as easy as this to use logic to save the day ...

### Logic Thwarts Landru!



First Suspicion That It's a Mere Computer Running the Show



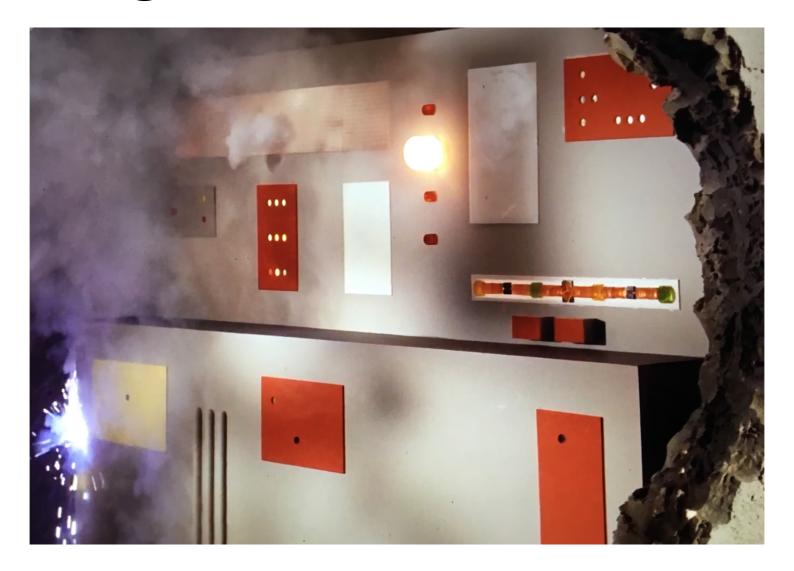
### Logic Thwarts Landru!



Landru is Indeed Merely a Computer (the real Landru having done the programming)



### Logic Thwarts Landru!



Landru Kills Himself Because Kirk/Spock Argue He Has Violated the Prime Directive for Good by Denying Creativity to Others

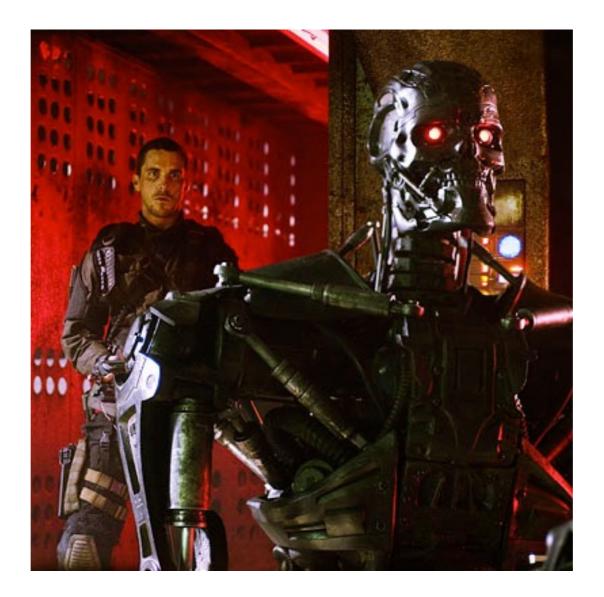


### Logic Thwarts Nomad! (with the Liar Paradox)









"We're in very deep trouble."

# "We're in very deep trouble."



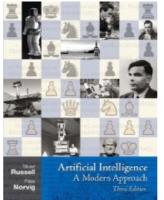
# "We're in very deep trouble."











 $\forall x: \texttt{Agents}$ 

$$u(\operatorname{AIA}_i(\pi_j)) > \tau^+ \in \mathbb{Z} \text{ or } \tau^- \in \mathbb{Z}$$

 $\forall x : Agents$ Autonomous(x

#### Are Autonomous-and-Creative Machines Intrinsically Untrustworthy?\*

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#### Abstract

Given what we find in the case of human cognition, the following principle appears to be quite plausible: An artificial agent that is both autonomous (A) and creative (C) will tend to be, from the viewpoint of a rational, fully informed agent, (U) untrustworthy. After briefly explaining the intuitive, internal structure of this disturbing principle, in the context of the human sphere, we provide a more formal rendition of it designed to apply to the realm of intelligent artificial agents. The more-formal version makes use of some of the basic structures available in one of our cognitive-event calculi, and can be expressed as a (confessedly — for reasons explained naïve) theorem. We prove the theorem, and provide simple demonstrations of it in action, using a novel theorem prover (ShadowProver). We then end by pointing toward some future defensive engineering measures that should be taken in light of the theorem.

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"The authors are deeply grateful for support provided by both AFOSR and ONR that enabled the research reported on herein, and are in addition thankful both for the guidance and patience of the editors and wise comments received from two reviewers.

### = Dangerous(x)

$$u(\operatorname{AIA}_i(\pi_j)) > \tau^+ \in \mathbb{Z} \text{ or } \tau^- \in \mathbb{Z}$$

#### $\forall x : Agents$

Autonomous(x) + Powerful(x) + Highly\_Intelligent(x) = Dangerous(x)

## $u(\operatorname{AIA}_i(\pi_j)) > \tau^+ \in \mathbb{Z} \text{ or } \tau^- \in \mathbb{Z}$

**Theorem ACU:** In a collaborative situation involving agents a (as the "trustor") and a' (as the "trustee"), if a' is at once both autonomous and ToM-creative, a' is untrustworthy from an ideal-observer o's viewpoint, with respect to the action-goal pair  $\langle \alpha, \gamma \rangle$  in question.

**Proof**: Let *a* and *a'* be agents satisfying the hypothesis of the theorem in an arbitrary collaborative situation. Then, by definition,  $a \neq a'$  desires to obtain some goal  $\gamma$  in part by way of a contributed action  $\alpha_k$  from *a'*, *a'* knows this, and moreover *a'* knows that *a* believes that this contribution will succeed. Since *a'* is by supposition ToM-creative, *a'* may desire to surprise *a* with respect to *a*'s belief regarding *a'*'s contribution; and because *a'* is autonomous, attempts to ascertain whether such surprise will come to pass are fruitless since what will happen is locked inaccessibly in the oracle that decides the case. Hence it follows by TRANS that an ideal observer *o* will regard *a'* to be untrustworthy with respect to the pair  $\langle \alpha, \gamma \rangle$  pair. **QED** 

### $\forall x: \texttt{Agents}$

Autonomous(x) + Powerful(x) + Highly\_Intelligent(x) = Dangerous(x)

(We use the "jump" technique in relative computability.)

$$u(\operatorname{AIA}_i(\pi_j)) > \tau^+ \in \mathbb{Z} \text{ or } \tau^- \in \mathbb{Z}$$

**Theorem ACU:** In a collaborative situation involving agents a (as the "trustor") and a' (as the "trustee"), if a' is at once both autonomous and ToM-creative, a' is untrustworthy from an ideal-observer o's viewpoint, with respect to the action-goal pair  $\langle \alpha, \gamma \rangle$  in question.

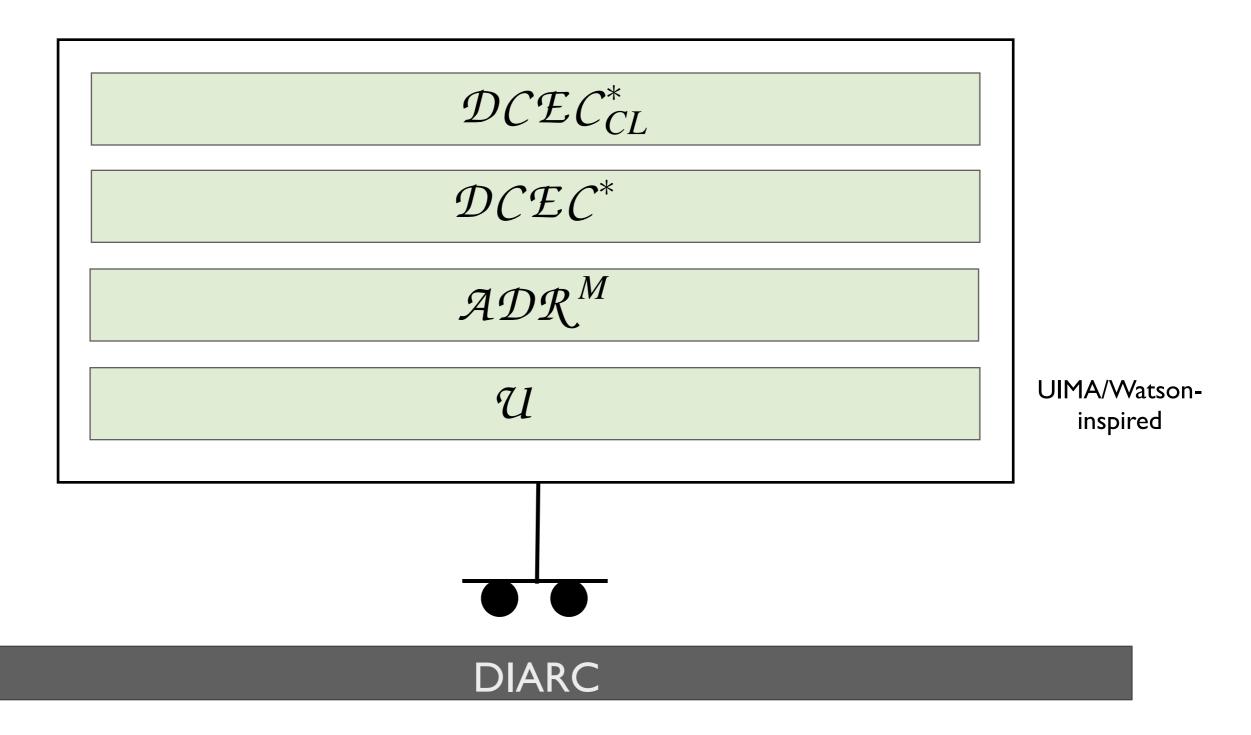
**Proof**: Let *a* and *a'* be agents satisfying the hypothesis of the theorem in an arbitrary collaborative situation. Then, by definition,  $a \neq a'$  desires to obtain some goal  $\gamma$  in part by way of a contributed action  $\alpha_k$  from *a'*, *a'* knows this, and moreover *a'* knows that *a* believes that this contribution will succeed. Since *a'* is by supposition ToM-creative, *a'* may desire to surprise *a* with respect to *a*'s belief regarding *a'*'s contribution; and because *a'* is autonomous, attempts to ascertain whether such surprise will come to pass are fruitless since what will happen is locked inaccessibly in the oracle that decides the case. Hence it follows by TRANS that an ideal observer *o* will regard *a'* to be untrustworthy with respect to the pair  $\langle \alpha, \gamma \rangle$  pair. **QED** 

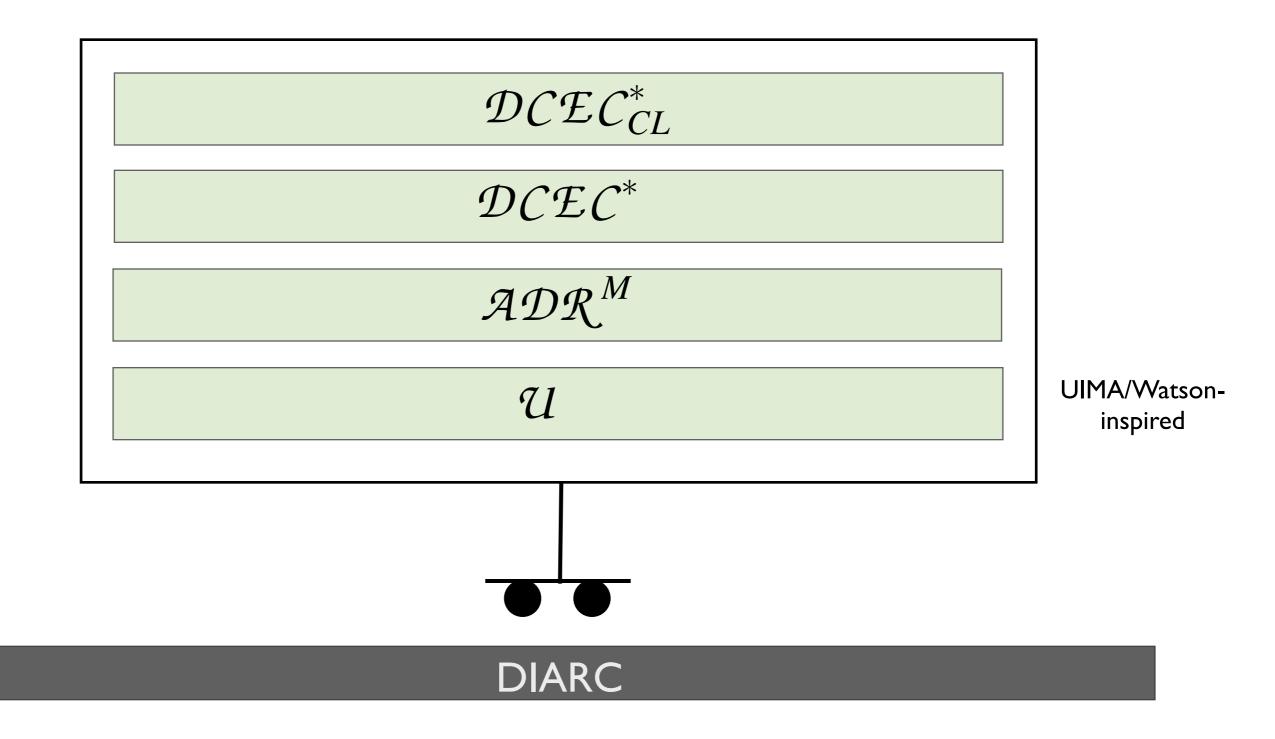
## Conclusion from last time:

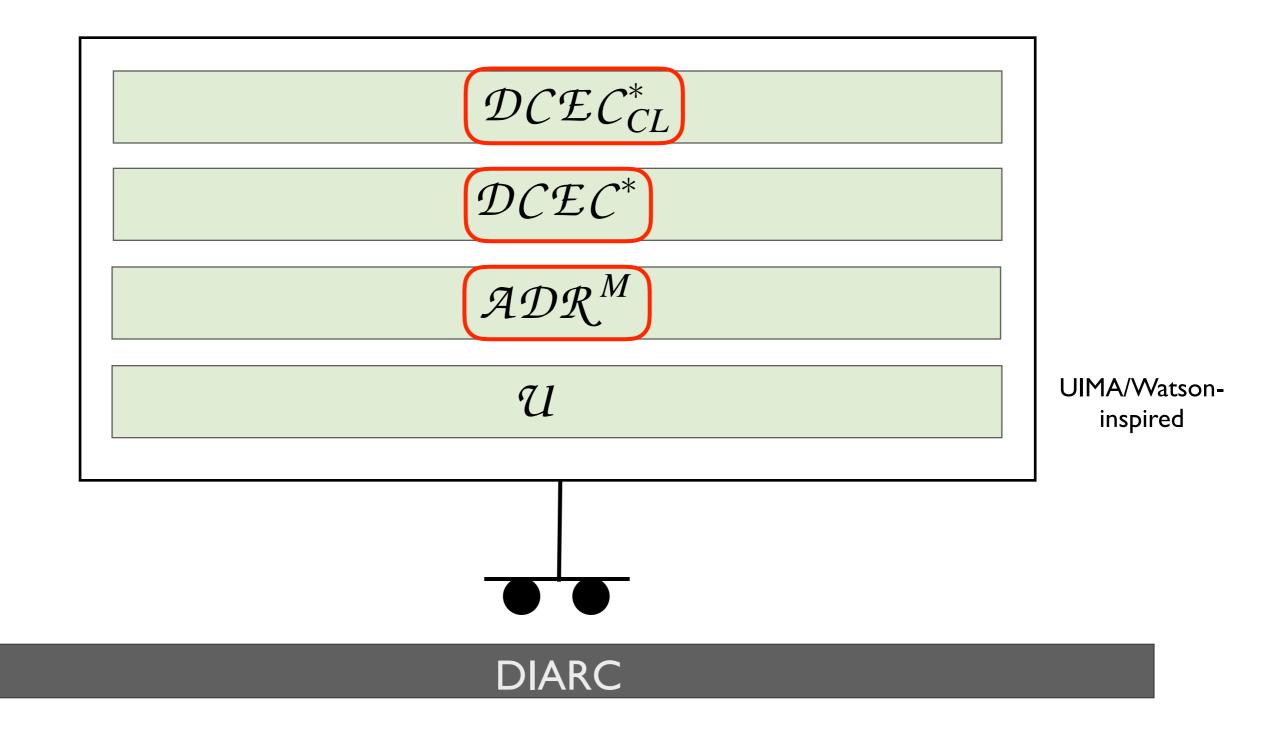
### Conclusion from last time:

"Computational logician, sorry, back to your drawing board to find a logic that works with The Four Steps!"

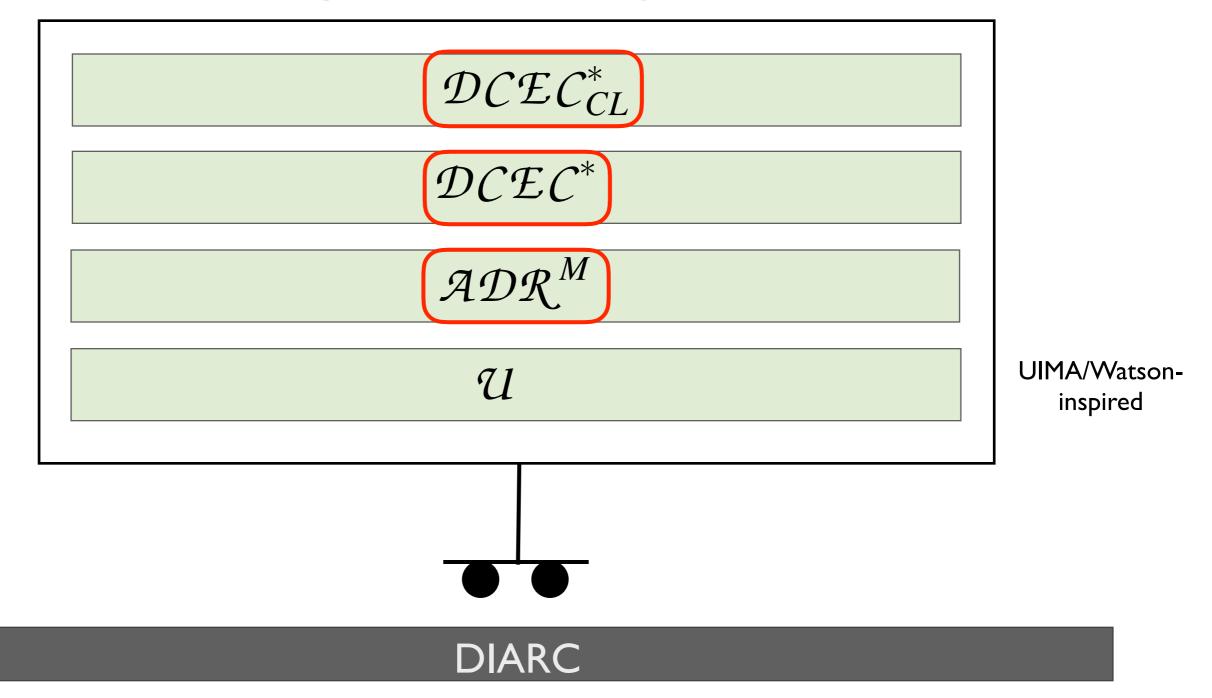
# I. Cognitive Calculi ...







Not simple deontic logics like **D**!

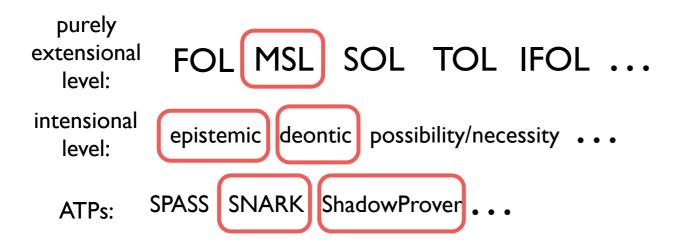


 purely extensional level:
 FOL MSL SOL TOL IFOL ...
 theories: PAZFC axiomatic physics ...

 intensional level:
 epistemic deontic possibility/necessity ...
 model finders: MACE ...

 ATPs:
 SPASS SNARK ShadowProver ...
 nature of representation: symbolic or homomorphic:

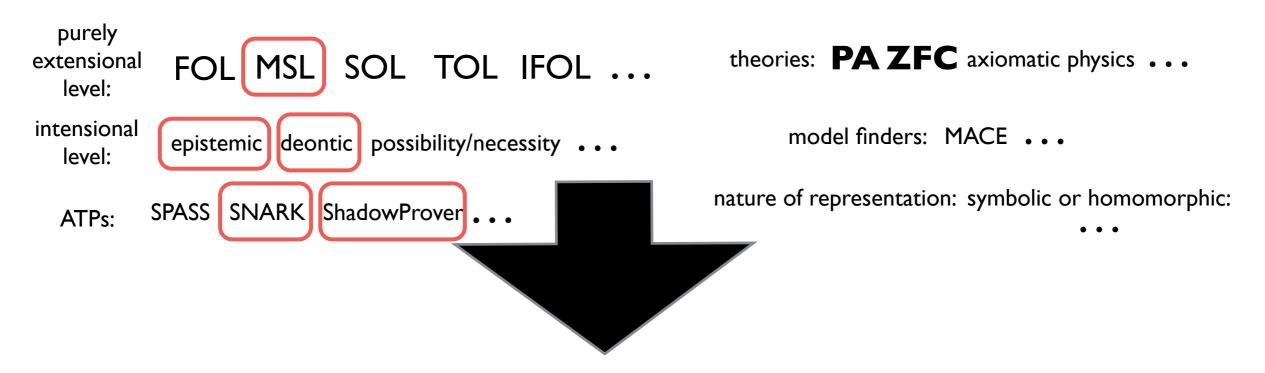


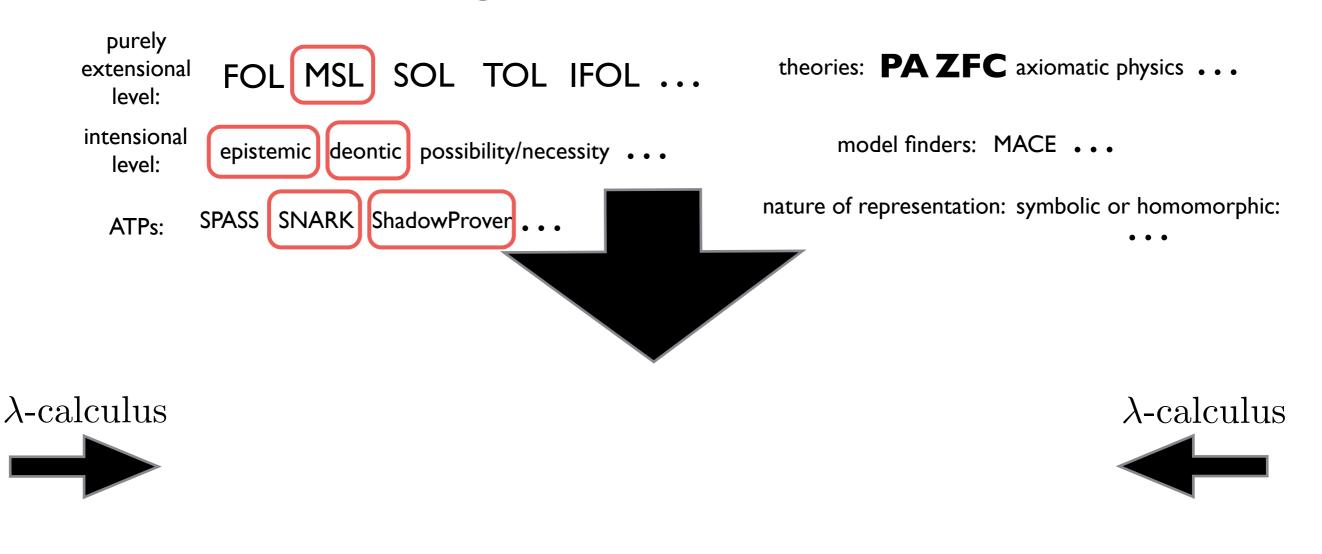
theories: **PAZFC** axiomatic physics ...

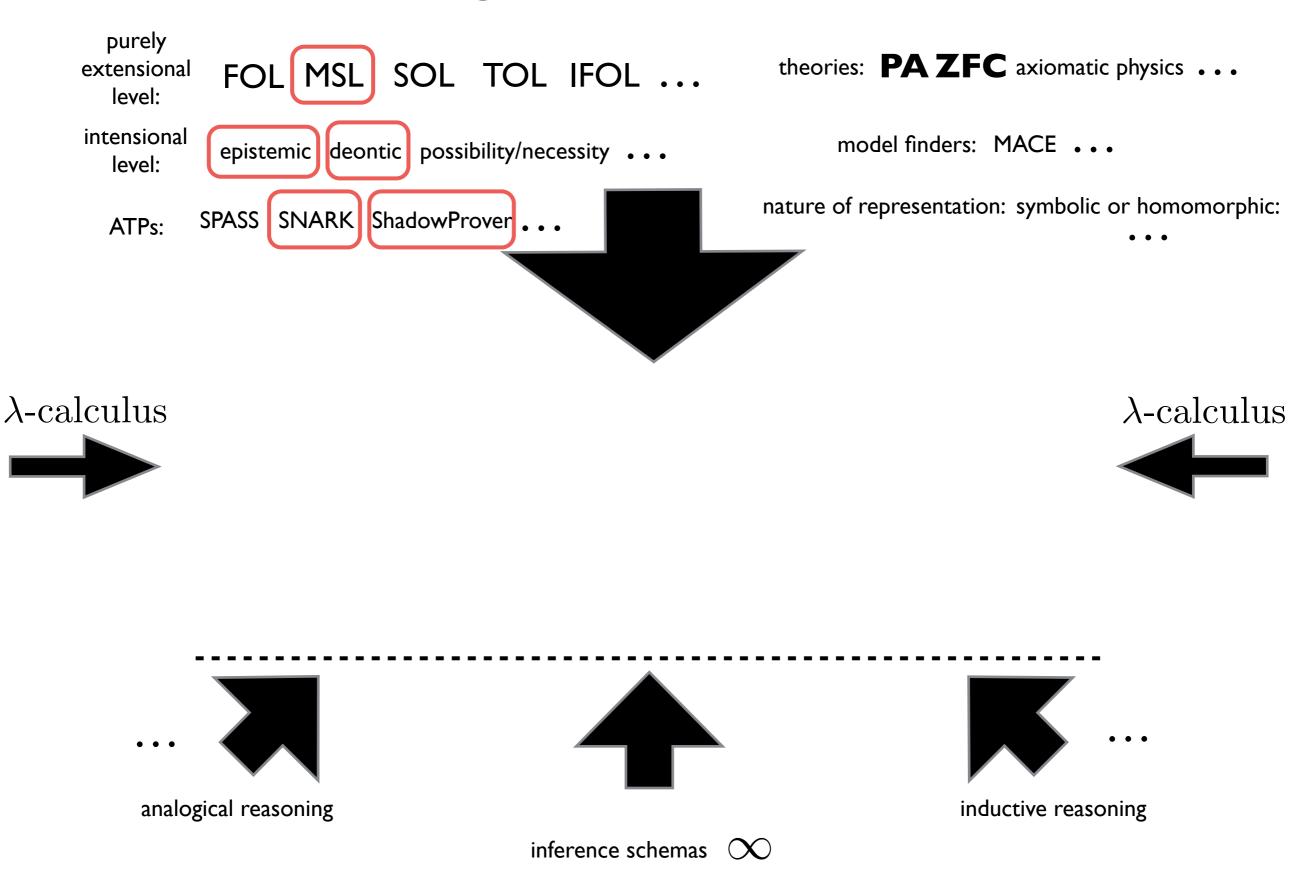
model finders: MACE ...

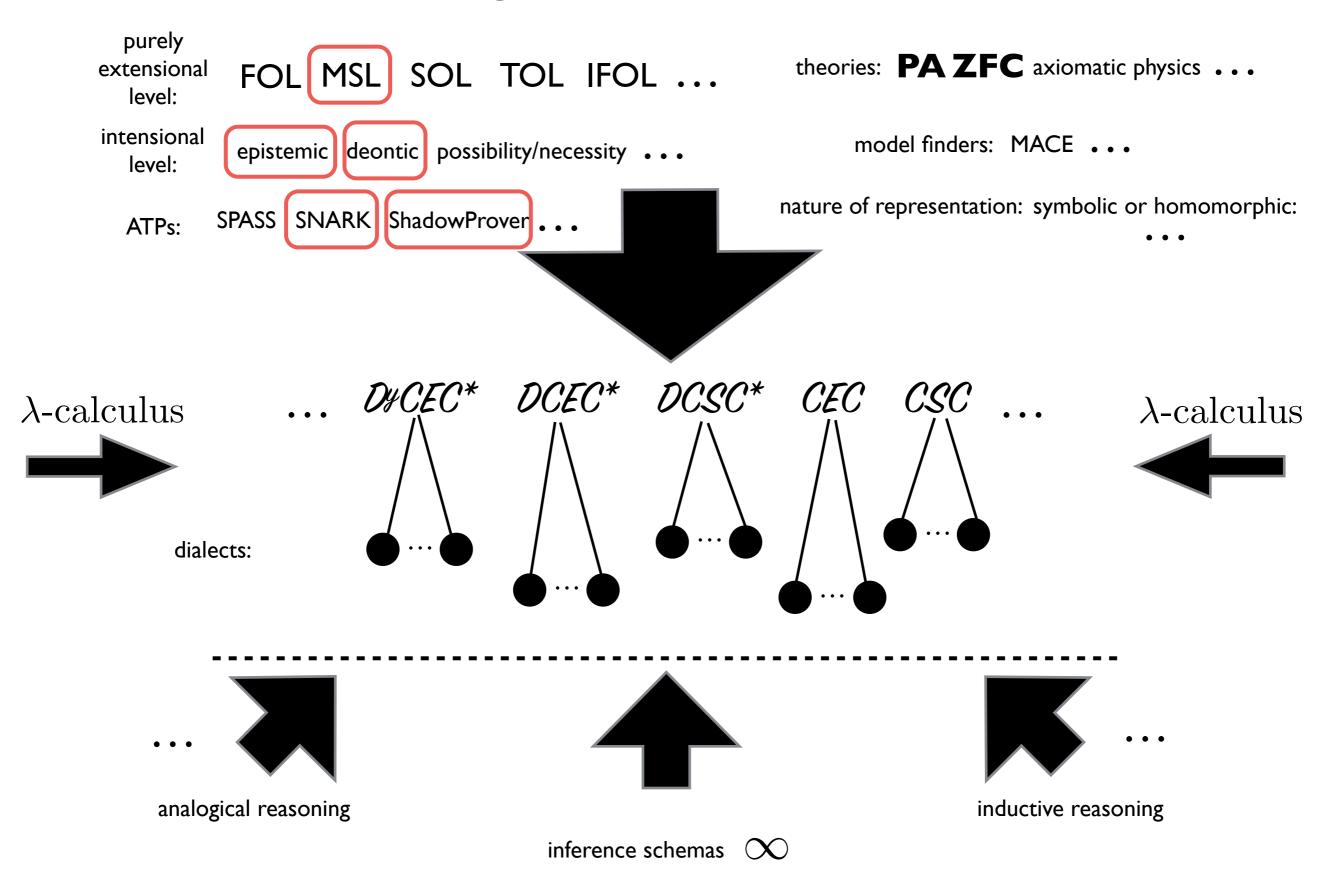
nature of representation: symbolic or homomorphic:

• • •

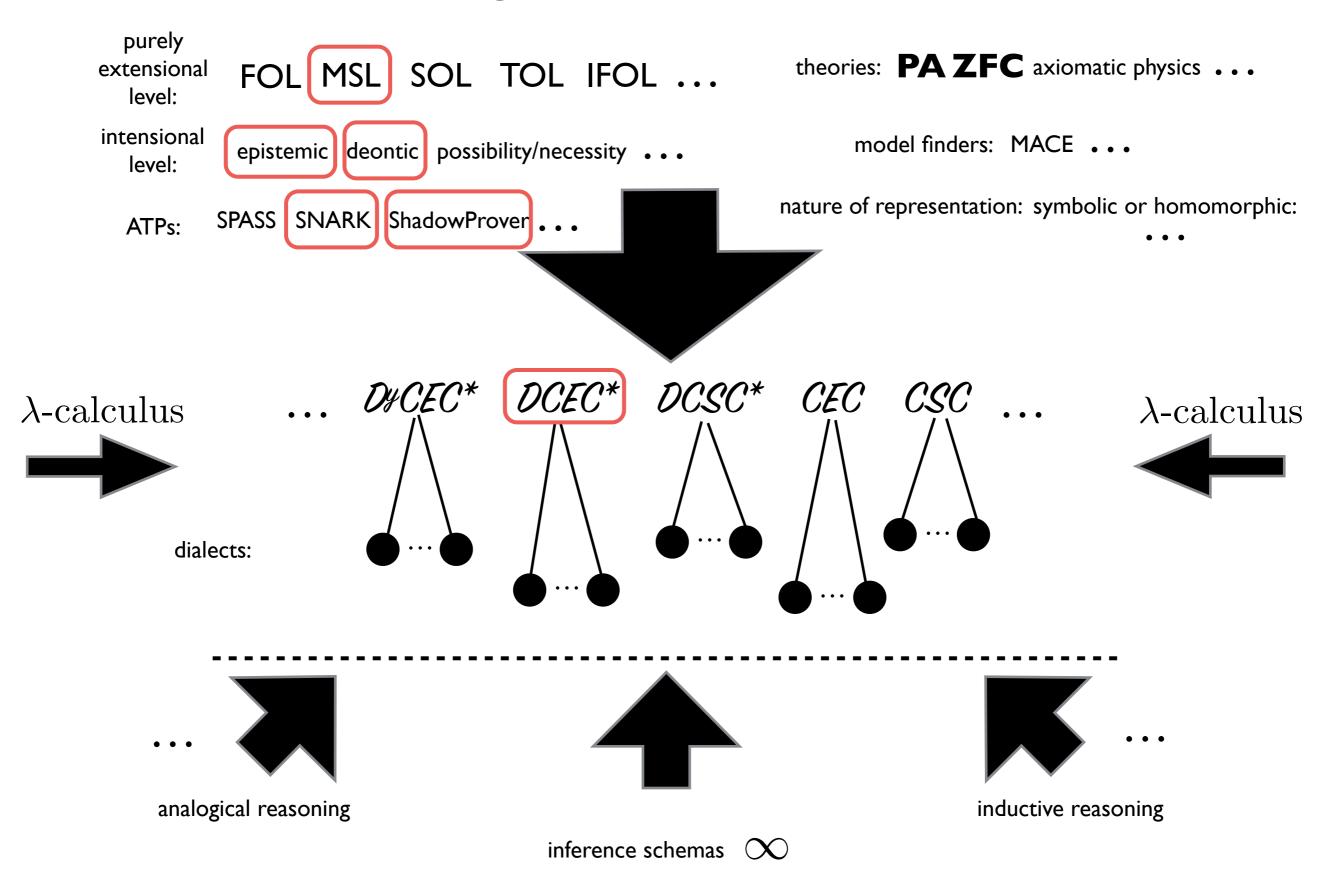




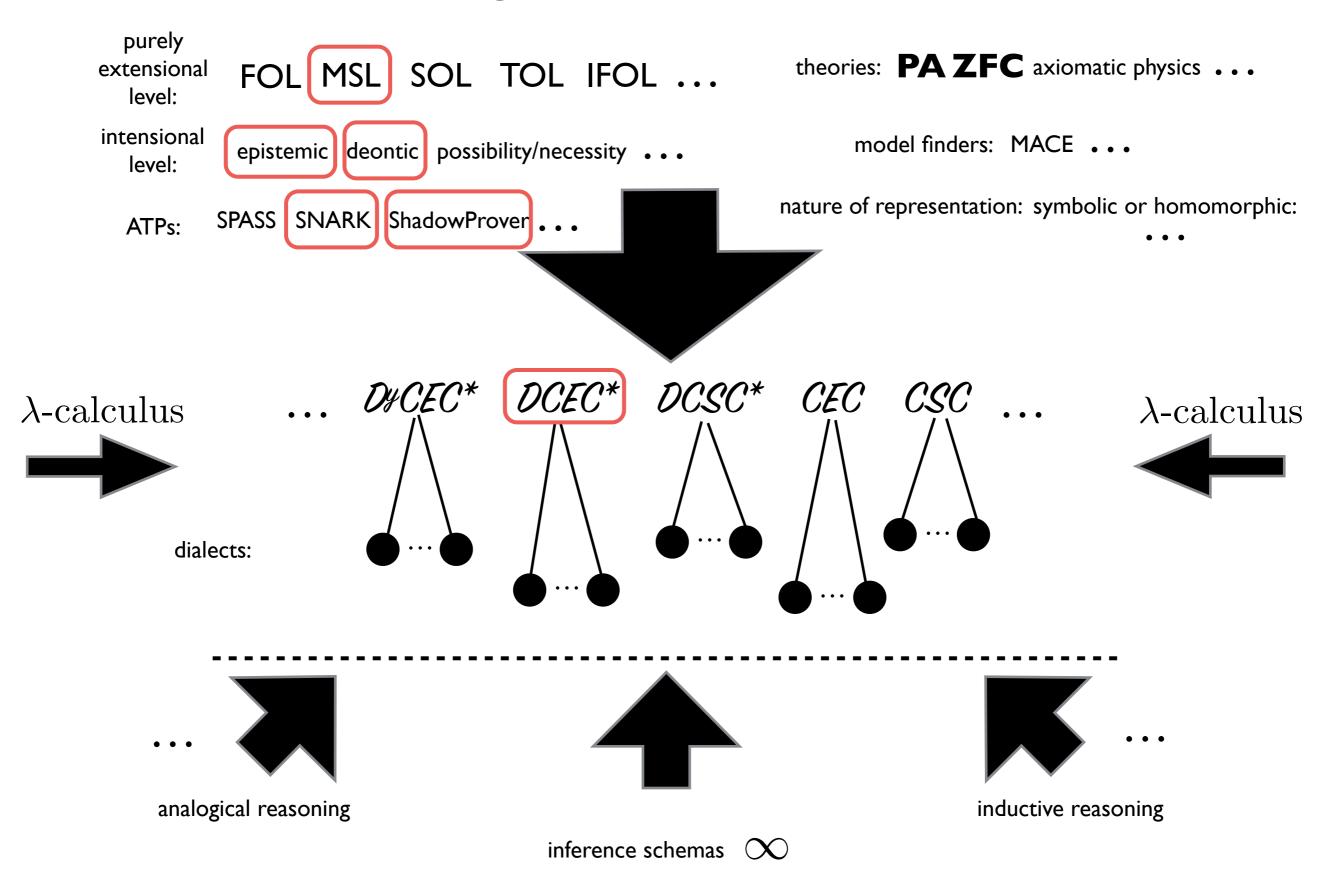


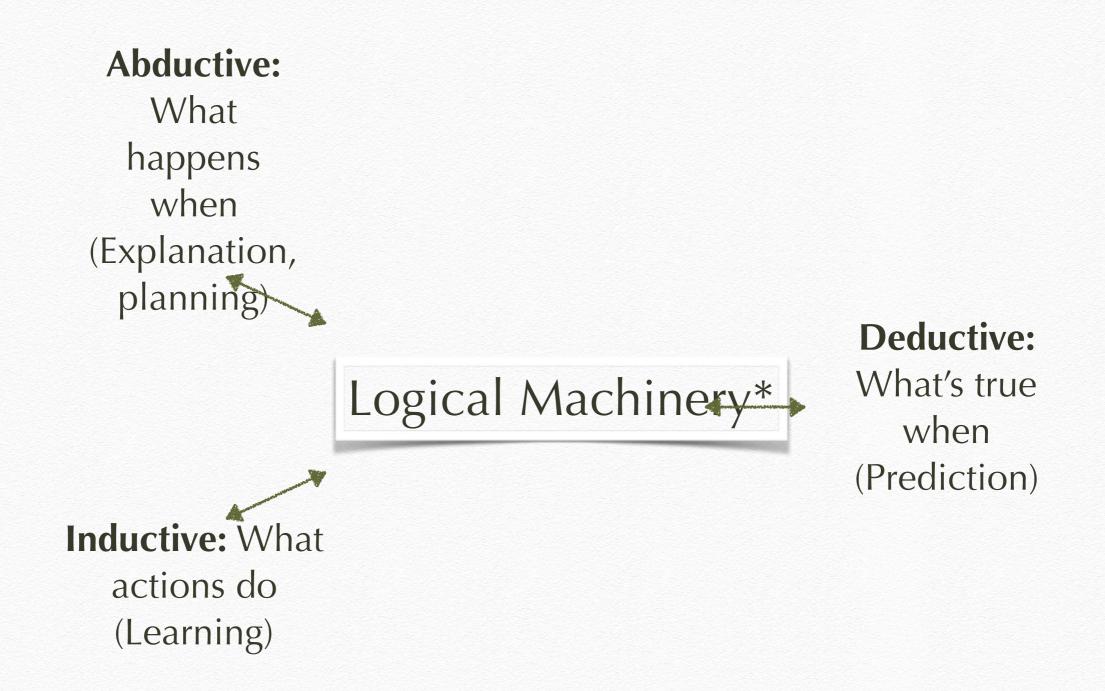


#### Cognitive Calculi 66



#### Cognitive Calculi 66





\*Diagram partly due to Shanahan

# Formal Syntax

## Formal Syntax

 $S ::= \begin{array}{l} \mathsf{Object} \mid \mathsf{Agent} \mid \mathsf{Self} \sqsubseteq \mathsf{Agent} \mid \mathsf{ActionType} \mid \mathsf{Action} \sqsubseteq \mathsf{Event} \mid \\ \mathsf{Moment} \mid \mathsf{Boolean} \mid \mathsf{Fluent} \mid \mathsf{Numeric} \end{array}$ 

action : Agent  $\times$  ActionType  $\rightarrow$  Action

*initially* : Fluent  $\rightarrow$  Boolean

 $\mathit{holds}: \mathsf{Fluent} \times \mathsf{Moment} \rightarrow \mathsf{Boolean}$ 

*happens* : Event  $\times$  Moment  $\rightarrow$  Boolean

clipped: Moment imes Fluent imes Moment  $\rightarrow$  Boolean

 $f ::= initiates : Event \times Fluent \times Moment \rightarrow Boolean$ 

 $\mathit{terminates}: \mathsf{Event} \times \mathsf{Fluent} \times \mathsf{Moment} \to \mathsf{Boolean}$ 

*prior* : Moment imes Moment  $\rightarrow$  Boolean

interval : Moment imes Boolean

 $*: \text{Agent} \to \text{Self}$ 

 $payoff: Agent \times ActionType \times Moment \rightarrow Numeric$ 

 $t ::= x : S | c : S | f(t_1, \dots, t_n)$ 

$$t: \text{Boolean} \mid \neg \phi \mid \phi \land \psi \mid \phi \lor \psi \mid$$
$$\phi ::= \frac{\mathbf{P}(a,t,\phi) \mid \mathbf{K}(a,t,\phi) \mid \mathbf{C}(t,\phi) \mid \mathbf{S}(a,b,t,\phi) \mid \mathbf{S}(a,t,\phi)}{\mathbf{B}(a,t,\phi) \mid \mathbf{D}(a,t,holds(f,t')) \mid \mathbf{I}(a,t,happens(action(a^*,\alpha),t'))}$$
$$\mathbf{O}(a,t,\phi,happens(action(a^*,\alpha),t'))$$

## Inference Schemata

# Inference Schemata

[ <i>R</i> <sub>1</sub> ]	$ [R_{\gamma}]$
$\overline{\mathbf{C}(t,\mathbf{P}(a,t,\phi)\to\mathbf{K}(a,t,\phi))}  [R_1]  \overline{\mathbf{C}(t,\mathbf{K}(a,t,\phi)\to\mathbf{B}(a,t,\phi))}$	$(t, \phi))$
$\frac{\mathbf{C}(t,\phi) \ t \le t_1 \dots t \le t_n}{\mathbf{K}(a_1,t_1,\dots,\mathbf{K}(a_n,t_n,\phi)\dots)}  [R_3]  \frac{\mathbf{K}(a,t,\phi)}{\phi}  [R_4]$	
$\mathbf{K}(a_1,t_1,\ldots\mathbf{K}(a_n,t_n,\phi)\ldots) \qquad $	
$\mathbf{C}(t, \mathbf{K}(a, t_1, \phi_1 \to \phi_2)) \to \mathbf{K}(a, t_2, \phi_1) \to \mathbf{K}(a, t_3, \phi_2) $	25]
${\mathbf{C}(t, \mathbf{B}(a, t_1, \phi_1 \to \phi_2)) \to \mathbf{B}(a, t_2, \phi_1) \to \mathbf{B}(a, t_3, \phi_2)} [R_0]$	5]
$\mathbf{C}(t, \mathbf{C}(t_1, \phi_1 \to \phi_2)) \to \mathbf{C}(t_2, \phi_1) \to \mathbf{C}(t_3, \phi_2) \qquad [R_7]$	
${\mathbf{C}(t,\forall x. \ \phi \to \phi[x \mapsto t])}  [R_8]  {\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \phi_2 \to \neg \phi_2 \to \phi_2 $	$\frac{1}{\phi_1}$ [R <sub>9</sub> ]
	17
$\mathbf{C}(t, [\phi_1 \land \ldots \land \phi_n \to \phi] \to [\phi_1 \to \ldots \to \phi_n \to \psi])  [R_{10}]$	
$\frac{\mathbf{B}(a,t,\phi) \ \phi \to \psi}{\mathbf{B}(a,t,\psi)}  [R_{11a}]  \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\psi)}{\mathbf{B}(a,t,\psi \land \phi)}  [R_{11b}]$	
$\mathbf{B}(a,t,\mathbf{\psi}) \qquad \qquad \mathbf{B}(a,t,\mathbf{\psi}\wedge\phi) \qquad \qquad \mathbf{B}(a,t,$	
$\mathbf{S}(s,h,t,\phi)$ [ <b>P</b> ]	
$\frac{\mathbf{S}(s,h,t,\phi)}{\mathbf{B}(h,t,\mathbf{B}(s,t,\phi))}  [R_{12}]$	
$\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))$ [R12]	
$\frac{\mathbf{P}(a,t,happens(action(a^*,\alpha),t^*))}{\mathbf{P}(a,t,happens(action(a^*,\alpha),t))}  [R_{13}]$	
$\mathbf{B}(a,t,\phi)  \mathbf{B}(a,t,\mathbf{O}(a^*,t,\phi,happens(action(a^*,\alpha),t')))$	
$\mathbf{O}(a,t,\phi,happens(action(a^*,\alpha),t'))$	
$\mathbf{K}(a,t,\mathbf{I}(a^*,t,happens(action(a^*,\alpha),t')))$	$[R_{14}]$
$\phi \leftrightarrow \psi$	
$\overline{\mathbf{O}(a,t,\phi,\gamma)\leftrightarrow\mathbf{O}(a,t,\psi,\gamma)}  [R_{15}]$	

### Event Calculus for Time & Change

[ <i>R</i> <sub>1</sub> ]	$[R_2]$
$\overline{\mathbf{C}(t,\mathbf{P}(a,t,\phi)\to\mathbf{K}(a,t,\phi))}  \begin{bmatrix} R_1 \end{bmatrix}  \overline{\mathbf{C}(t,\mathbf{K}(a,t,\phi)\to\mathbf{B}(a,t,\phi))}$	[**2]
$\frac{\mathbf{C}(t,\phi) \ t \le t_1 \dots t \le t_n}{\mathbf{K}(a_1,t_1,\dots\mathbf{K}(a_n,t_n,\phi)\dots)}  [R_3]  \frac{\mathbf{K}(a,t,\phi)}{\phi}  [R_4]$	
$\overline{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1\to\phi_2))\to\mathbf{K}(a,t_2,\phi_1)\to\mathbf{K}(a,t_3,\phi_2)}  [R_5]$	
$\overline{\mathbf{C}(t, \mathbf{B}(a, t_1, \phi_1 \to \phi_2)) \to \mathbf{B}(a, t_2, \phi_1) \to \mathbf{B}(a, t_3, \phi_2)}  [R_6]$	
$\overline{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1\to\phi_2))\to\mathbf{C}(t_2,\phi_1)\to\mathbf{C}(t_3,\phi_2)}  [R_7]$	
$\overline{\mathbf{C}(t,\forall x.\ \phi \to \phi[x \mapsto t])}  [R_8]  \overline{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \to \neg \phi_2 \to \neg \phi_1)}$	[ <i>R</i> 9]
$\overline{\mathbf{C}(t, [\phi_1 \land \ldots \land \phi_n \to \phi] \to [\phi_1 \to \ldots \to \phi_n \to \psi])}  \begin{bmatrix} R_{10} \end{bmatrix}$	
$\frac{\mathbf{B}(a,t,\phi) \ \phi \to \psi}{\mathbf{B}(a,t,\psi)}  [R_{11a}]  \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\psi)}{\mathbf{B}(a,t,\psi \land \phi)}  [R_{11b}]$	
$\frac{\mathbf{S}(s,h,t,\phi)}{\mathbf{B}(h,t,\mathbf{B}(s,t,\phi))}  [R_{12}]$	
$\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))$	
$\frac{\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))}{\mathbf{P}(a,t,happens(action(a^*,\alpha),t))}  [R_{13}]$	
$\mathbf{B}(a,t,\phi)  \mathbf{B}(a,t,\mathbf{O}(a^*,t,\phi,happens(action(a^*,\alpha),t')))$	
$\mathbf{O}(a,t,\phi,happens(action(a^*,\alpha),t'))$	1
${\mathbf{K}(a,t,\mathbf{I}(a^*,t,happens(action(a^*,\alpha),t')))} [R_1, \mathbf{K}(a,t,\mathbf{I}(a^*,t,happens(action(a^*,\alpha),t')))]$	4]
$\overline{\mathbf{O}(a,t,\phi,\gamma)\leftrightarrow\mathbf{O}(a,t,\psi,\gamma)}  [\kappa_{15}]$	

### Event Calculus for Time & Change

```
 \begin{array}{ll} \hline \mathbf{C}(t,\mathbf{P}(a,t,\phi) \rightarrow \mathbf{K}(a,t,\phi)) & [R_1] & \overline{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \rightarrow \mathbf{B}(a,t,\phi))} & [R_2] \\ \hline \mathbf{C}(t,\phi) \ t \leq t_1 \dots t \leq t_n & [R_3] & \frac{\mathbf{K}(a,t,\phi)}{\phi} & [R_4] \\ \hline \hline \mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1 \rightarrow \phi_2)) \rightarrow \mathbf{K}(a,t_2,\phi_1) \rightarrow \mathbf{K}(a,t_3,\phi_2) & [R_5] \\ \hline \hline \mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1 \rightarrow \phi_2)) \rightarrow \mathbf{B}(a,t_2,\phi_1) \rightarrow \mathbf{B}(a,t_3,\phi_2) & [R_6] \\ \hline \hline \mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \rightarrow \phi_2)) \rightarrow \mathbf{C}(t_2,\phi_1) \rightarrow \mathbf{C}(t_3,\phi_2) & [R_7] \\ \hline \hline \mathbf{C}(t,(\phi_1 \wedge \dots \wedge \phi_n \rightarrow \phi_1) \rightarrow [\phi_1 \rightarrow \dots \rightarrow \phi_n \rightarrow \psi]) & [R_{10}] \\ \hline \hline \mathbf{C}(t,(\phi_1 \wedge \dots \wedge \phi_n \rightarrow \phi_1 \rightarrow [\phi_1 \rightarrow \dots \rightarrow \phi_n \rightarrow \psi]) & [R_{10}] \\ \hline \hline \mathbf{B}(a,t,\psi) & [R_{11a}] & \frac{\mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\psi \wedge \phi)}{\mathbf{B}(a,t,\psi \wedge \phi)} & [R_{11b}] \\ \hline \hline \mathbf{S}(s,h,t,\phi) & [R_{12}] \\ \hline \hline \mathbf{I}(a,t,happens(action(a^*,\alpha),t')) & [R_{13}] \\ \hline \mathbf{B}(a,t,\phi) \ \mathbf{B}(a,t,\mathbf{O}(a^*,t,\phi,happens(action(a^*,\alpha),t')))) & [R_{14}] \\ \hline \hline \mathbf{O}(a,t,\phi,happens(action(a^*,\alpha),t')) & [R_{15}] \end{array}
```

$$\begin{split} & [A_1] \ \mathbf{C}(\forall \ f,t \ . \ initially(f) \land \neg clipped(0,f,t) \Rightarrow holds(f,t)) \\ & [A_2] \ \mathbf{C}(\forall \ e,f,t_1,t_2 \ . \ happens(e,t_1) \land initiates(e,f,t_1) \land t_1 < t_2 \land \neg clipped(t_1,f,t_2) \Rightarrow holds(f,t_2)) \\ & [A_3] \ \mathbf{C}(\forall \ t_1,f,t_2 \ . \ clipped(t_1,f,t_2) \Leftrightarrow [\exists \ e,t \ . \ happens(e,t) \land t_1 < t < t_2 \land terminates(e,f,t)]) \\ & [A_4] \ \mathbf{C}(\forall \ a,d,t \ . \ happens(action(a,d),t) \Rightarrow \mathbf{K}(a,happens(action(a,d),t))) \\ & [A_5] \ \mathbf{C}(\forall \ a,f,t,t' \ . \ \mathbf{B}(a,holds(f,t)) \land \mathbf{B}(a,t < t') \land \neg \mathbf{B}(a,clipped(t,f,t')) \Rightarrow \mathbf{B}(a,holds(f,t'))) \end{split}$$

#### Defs for An Affective Cognitive time&change Calculus

1. <b>Joy</b> : pleased about a desirable event. By 'pleased about a desirable event' the meani will consider is 'pleased about a desirable consequence of the event'.	ing we	9. Satisfaction : (pleased about) the confirmation of the prospect of a desirable even The definition of $holds(AFF(a, satisfaction), t_3)$ is $and(10,11, 7 3)$ .	nt		
$for Some \ c \ B(a,t_3,implies(happens(e,t_1),holds(CON(e,a,c),t_2)))$	(1)	10. Fears-confirmed : (displeased about) the confirmation of the prospect of an un	ndesirable		
$D(a, t_3, holds(CON(e, a, c), t_2))$	(2)	(2) event. (2) The definition of $holds(AFF(a, fears - confirmed), t_3)$ is $and(10,12,9,3)$ .			
$K(a, t_3, happens(e, t_1))$	(3)	11. <b>Relief</b> : (pleased about) the disconfirmation of the prospect of an undesirable even	nt		
The definition of $holds(AFF(a, joy), t_3)$ is therefore and $(1, 2, 3)$ .		$K(a, t_3, not(happens(e, t_1)))$	(13)		
2. <b>Distress</b> : displeased about an undesirable event.		The definition of $holds(AFF(a, relief), t_3)$ is $and(10, 12, 9, 13)$ .			
$not(D(a, t_3, holds(CON(e, a, c), t_3)))$	(4)	<ul> <li>12. Disappointment : (displeased about) the disconfirmation of the prospect of a event</li> <li>The definition of holds(AFF(a, disappointment), t<sub>3</sub>) is and(10, 11, 7, 13).</li> </ul>	desirable		
The definition of $holds(AFF(a, distress), t_3)$ is therefore and $(1,4,3)$ .		13. <b>Pride</b> : (approving of) one's own praiseworthy action			
3. Happy-for: pleased about an event presumed to be desirable for someone else		Here we treat 'approve' as an action event. We also introduce a new predicate $PRAI$ , which will mean that agent a considers x a praiseworthy action by agent b. All the			
for Some $c B(a, t_3, implies(happens(e, t_1), holds(CON(e, a_1, c), t_2)))$	(5)	pretations are shown below.			
$B(a, t_3, D(a_1, t_3, holds(CON(e, a_1, c), t_2)))$	(6)	$happens(action(a, x), t_0)$	(14)		
$D(a, t_3, D(a_1, t_3, holds(CON(e, a_1, c), t_2)))$ $D(a, t_3, holds(CON(e, a_1, c), t_2))$		$for All \ a_x B(a,t_1,implies(happens(action(a_x,x),t_x), PRAISEWORTHY(a,a_x,x))) = (a,t_1,t_2,t_3,t_3,t_3,t_3,t_3,t_3,t_3,t_3,t_3,t_3$	$)), t_x \le t_1 $ (15)		
	(7)	$D(a, t_1, holds(PRAISEWORTHY(a, a, x), t_1))$	(15) $(16)$		
The definition of $holds(AFF(a, happy\_for), t_3)$ is therefore and $(5, 6, 7, 3)$ .					
4. <b>Pity</b> : displeased about an event presumed to be undesirable for someone else. T equivalent to sorry_for in Hobbs-Gordon model.	This is	$happens(action(a, approve(x)), t_1) $ $The definition of holds(AFF(a, pride), t_1) is and(14, B(a, t_1, holds(PRAISEWORTHY(a, a, t_2)))) $ $The definition of holds(AFF(a, pride), t_2) is and(14, B(a, t_1, holds(PRAISEWORTHY(a, a, t_2))))) $ $The definition of holds(AFF(a, pride), t_2) is and(14, B(a, t_1, holds(PRAISEWORTHY(a, a, t_2))))))) $ $The definition of holds(AFF(a, pride), t_2) is and(14, B(a, t_2, holds(PRAISEWORTHY(a, a, t_2))))))))))))))))))))))))))))))))))))$			
		<ul> <li>14. Shame: (disapproving of) one's own blameworthy action</li> </ul>	$I \Pi I (a, a, x), \iota_1)), I ().$		
$B(a, t_3, not(D(a_1, t_3, holds(CON(e, a_1, c), t_2))))$	(8)	This also follows the same explanation as Pride.			
$not(D(a, t_3, holds(CON(e, a_1, c), t_2)))$	(9)	$for All \ a_x B(a,t_1,implies(happens(action(a_x,x),t_x),B(a,t_1,holds(BLAMEWORT))) = 0 \ a_x B(a,t_1,t_1) \ a_x B(a,t_1,t_2) \ a_x B(a,t_1,t_2)$			
The definition of $holds(AFF(a, pity), t_3)$ is therefore and $(5, 8, 9, 3)$ .			(18)		
5. Gloating : pleased about an event presumed to be undesirable for someone else The o	defini-	$not(happens(action(a, approve(x)), t_1))$	(19)		
tion of $holds(AFF(a, gloating), t_3)$ is therefore and $(5, 8, 7, 3)$ .		The definition of $holds(AFF(a, shame), t_1)$ is $and(14, B(a, t_1, holds(BLAMEWOR)))$	$RTHY(a, a, x), t_1)), 19).$		
6. <b>Resentment</b> : displeased about an event presumed to be desirable for someone else definition of holds $(AFE(a, recomment), t_{a})$ is therefore and $(5, 6, 0, 2)$	e The	15. Admiration: (approving of) someone else's praiseworthy action			
definition of $holds(AFF(a, resentment), t_3)$ is therefore and $(5, 6, 9, 3)$ .		$happens(action(a_1, x), t_0)$	(20)		
7. Hope: (pleased about) the prospect of a desirable event		The definition of $holds(AFF(a, admiration), t_1)$ is $and(20, B(a, t_1, holds(PRAISE))) = 0$ .	$WORTHY(a, a_1, x), t_1)), 17).$		
$forSome \ c \ B(a, t_0, implies(happens(e, t_1), \diamond holds(CON(e, a, c), t_2)))$	(10)	16. <b>Reproach</b> : (disapproving of) someone else's blameworthy action The definition of $h$ is $and(20, B(a, t_1, holds(BLAMEWORTHY(a, a_1, x), t_1)), 19)$ .	$olds(AFF(a, reproach), t_1)$		
$D(a, t_0, holds(CON(e, a, c), t_2))$	(11)	17. <b>Gratification</b> : (approving of) one's own praiseworthy action and (being pleased about) the related desirable event. We again interpret 'pleased about the desirable event' as 'pleased			
The definition of $holds(AFF(a, hope), t_0)$ is therefore and (10,11).		about the desirable event. We again interpret pleased about the desirable event as	s pleased		
8. Fear: (displeased about) the prospect of an undesirable event					
$not(D(a,t_0,holds(CON(e,a,c),t_2)))$	(12)	$for Some \ c \ B(a,t_1,implies(happens(action(a,x),t_0),holds(CON(action(a,x),a,t_0),holds(CON(action$	$(c), t_0)))$ (21)		
The definition of $holds(AFF(a, fear), t_0)$ is therefore and (10,12).		$D(a, t_1, holds(CON(action(a, x), a, c), t_0))$	(21) (22)		
The domination of $normalizer (a, j car ), v_0)$ is uncreased and $(10, 12)$ .		The definition of $holds(AFF(a, gratification), t_1)$ is $and(20, B(a, t_1, holds(PRAIS)))$	$EWORTHY(a, a, x), t_1)), 17$		

#### ... (and more)

# II. Early Progress With Our Calculi: Non-Akratic Robots

# Informal Definition of Akrasia

An action  $\alpha_f$  is (Augustinian) akratic for an agent *A* at  $t_{\alpha_f}$  iff the following eight conditions hold:

- (1) A believes that A ought to do  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (2) A desires to do  $\alpha_f$  at  $t_{\alpha_f}$ ;
- (3) A's doing  $\alpha_f$  at  $t_{\alpha_f}$  entails his not doing  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (4) A knows that doing  $\alpha_f$  at  $t_{\alpha_f}$  entails his not doing  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (5) At the time  $(t_{\alpha_f})$  of doing the forbidden  $\alpha_f$ , *A*'s desire to do  $\alpha_f$  overrides *A*'s belief that he ought to do  $\alpha_o$  at  $t_{\alpha_f}$ .
- (6) A does the forbidden action  $\alpha_f$  at  $t_{\alpha_f}$ ;
- (7) A's doing  $\alpha_f$  results from A's desire to do  $\alpha_f$ ;
- (8) At some time *t* after  $t_{\alpha_f}$ , *A* has the belief that *A* ought to have done  $\alpha_o$  rather than  $\alpha_f$ .

# Informal Definition of Akrasia

An action  $\alpha_f$  is (Augustinian) akratic for an agent A at  $t_{\alpha_f}$  iff the following eight conditions hold:

- (1) A believes that A ought to  $do(\alpha_o)$  at  $t_{\alpha_o}$ ;
- (2) A desires to do  $\alpha_f$  at  $t_{\alpha_f}$ ;
- (3) A's doing  $\alpha_f$  at  $t_{\alpha_f}$  entails his not doing  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (4) A knows that doing  $\alpha_f$  at  $t_{\alpha_f}$  entails his not doing  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (5) At the time  $(t_{\alpha_f})$  of doing the forbidden  $\alpha_f$ , *A*'s desire to do  $\alpha_f$  overrides *A*'s belief that he ought to do  $\alpha_o$  at  $t_{\alpha_f}$ .
- (6) A does the forbidden action  $\alpha_f$  at  $t_{\alpha_f}$ ;
- (7) A's doing  $\alpha_f$  results from A's desire to do  $\alpha_f$ ;
- (8) At some time *t* after  $t_{\alpha_f}$ , *A* has the belief that *A* ought to have done  $\alpha_o$  rather than  $\alpha_f$ .

# Informal Definition of Akrasia

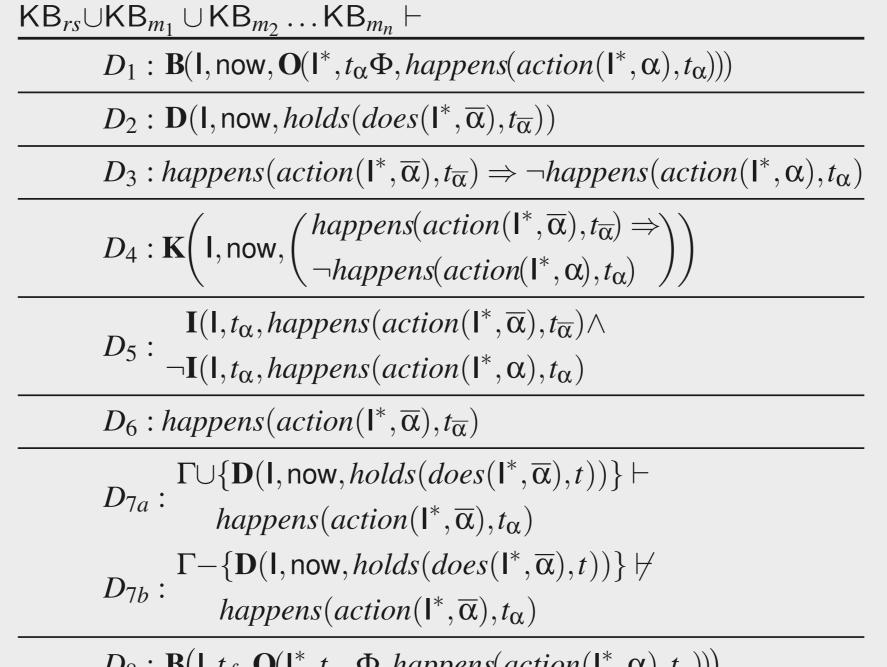
An action  $\alpha_f$  is (Augustinian) akratic for an agent A at  $t_{\alpha_f}$  iff the following eight conditions hold:

- (1) A believes that A ought to do  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (2) A desires to do  $\alpha_f$  at  $t_{\alpha_f}$ ;
- (3) A's doing  $\alpha_f$  at  $t_{\alpha_f}$  entails his not doing  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (4) A knows that doing  $\alpha_f$  at  $t_{\alpha_f}$  entails his not doing  $\alpha_o$  at  $t_{\alpha_o}$ ;
- (5) At the time  $(t_{\alpha_f})$  of doing the forbidden  $\alpha_f$ , *A*'s desire to do  $\alpha_f$  overrides *A*'s belief that he ought to do  $\alpha_o$  at  $t_{\alpha_f}$ .
- (6) A does the forbidden action  $\alpha_f$  at  $t_{\alpha_f}$ ;
- (7) A's doing  $\alpha_f$  results from A's desire to do  $\alpha_f$ ;
- "Regret" (8) At some time *t* after  $t_{\alpha_f}$ , *A* has the belief that *A* ought to have done  $\alpha_o$  rather than  $\alpha_f$ .

# Cast in

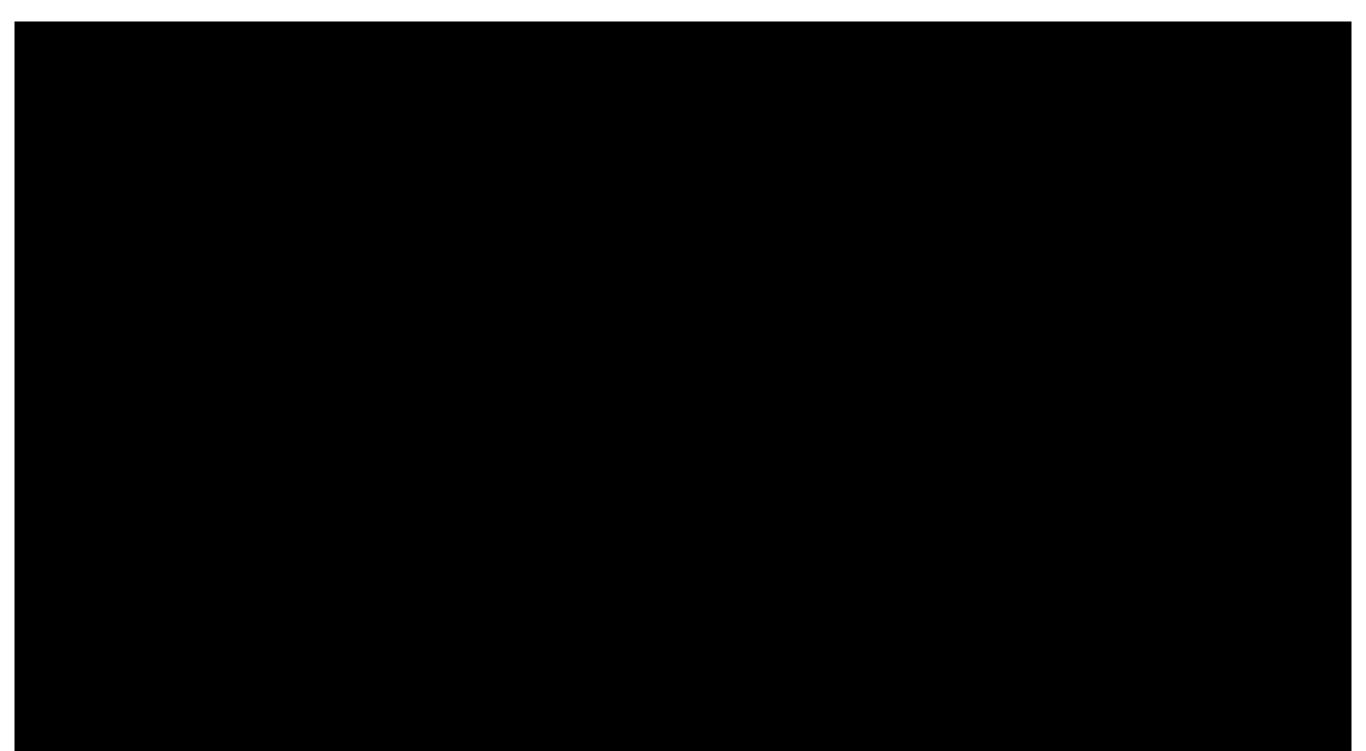
 $\mathcal{DCEC}^*$ 

this becomes ...

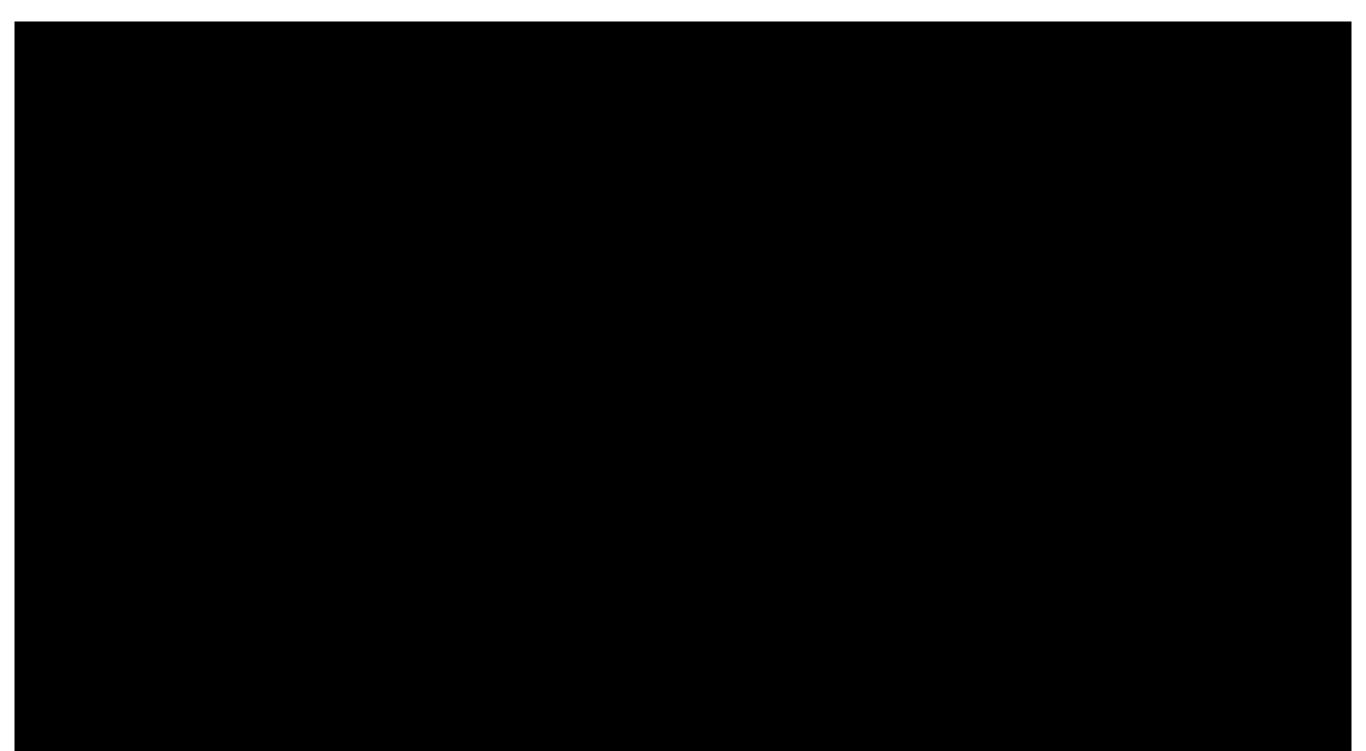


 $D_8: \mathbf{B}(\mathbf{I}, t_f, \mathbf{O}(\mathbf{I}^*, t_{\alpha}, \Phi, happens(action(\mathbf{I}^*, \alpha), t_{\alpha})))$ 



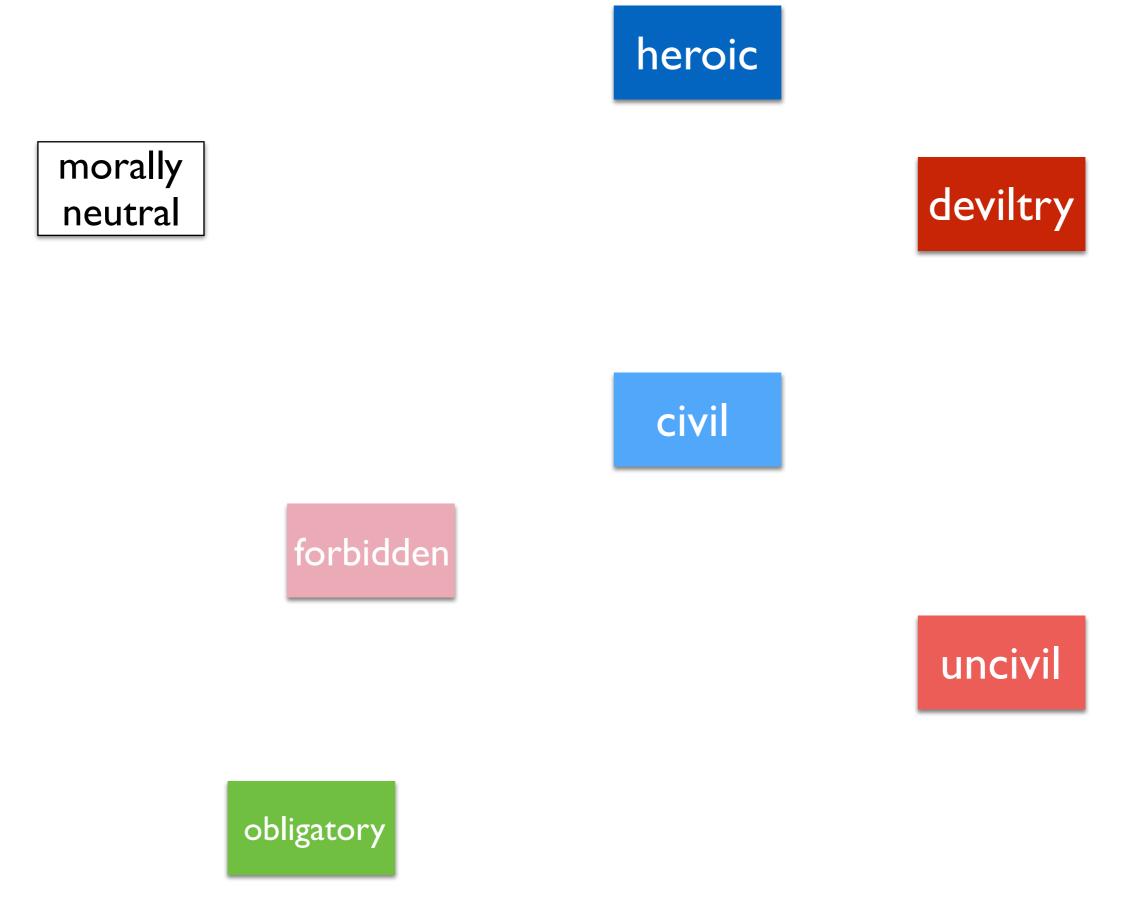






# III. But, a twist befell the logicists ...

Chisholm had argued that the three old 19th-century ethical categories (forbidden, morally neutral, obligatory) are not enough — and soulsearching brought me to agreement.









					the supere	erogatory
deviltry	uncivil	forbidden	morally neutral	obligatory	civil	heroic



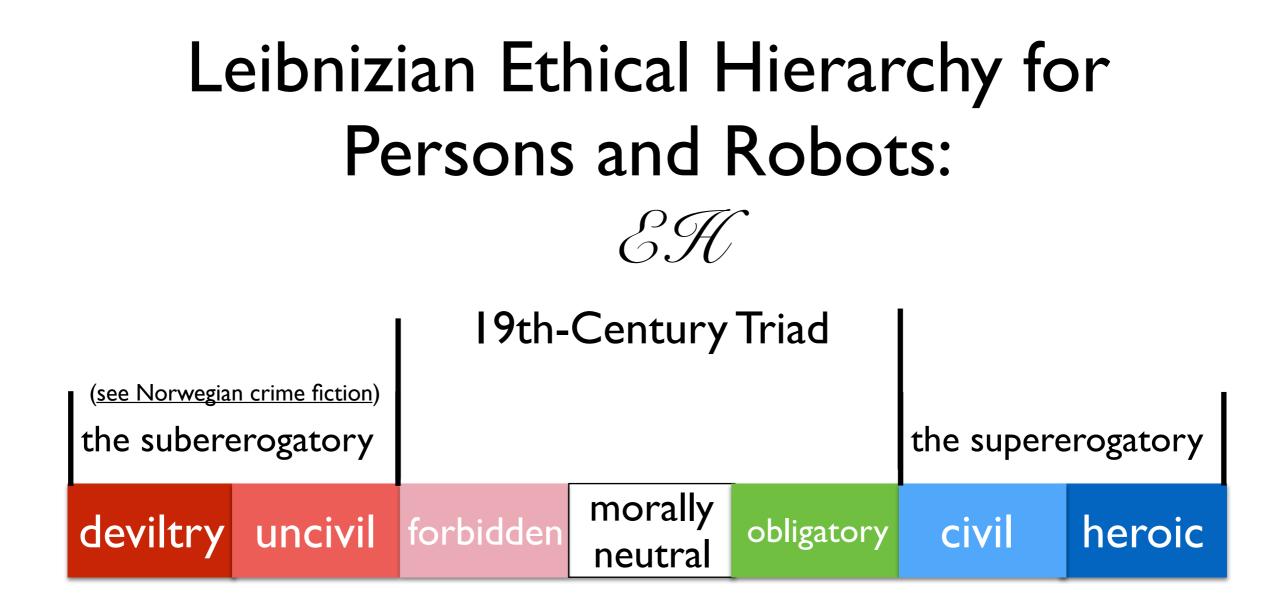
					the supere	erogatory
deviltry	uncivil	forbidden	morally neutral	obligatory	civil	heroic



the suber	erogatory				the supere	erogatory
deviltry	uncivil	forbidden	morally neutral	obligatory	civil	heroic



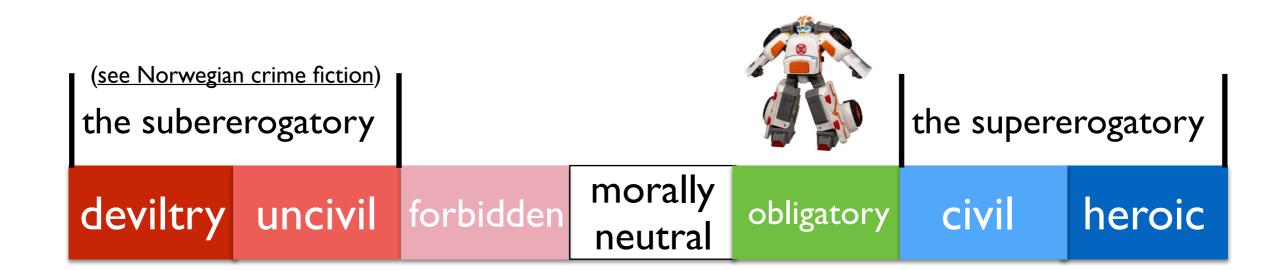
( <u>see Norwegia</u> the subere	<u>n crime fiction)</u> erogatory				the supere	erogatory
deviltry	uncivil	forbidden	morally neutral	obligatory	civil	heroic



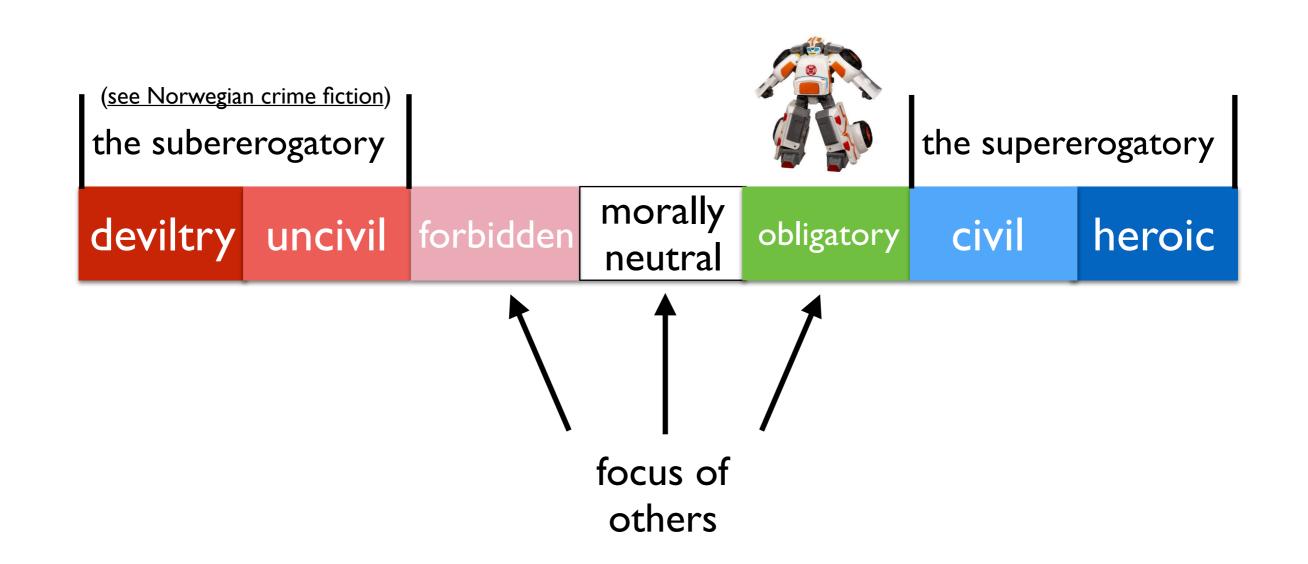


( <u>see Norwegia</u> the subere	<u>n crime fiction)</u> erogatory				the supere	erogatory
deviltry	uncivil	forbidden	morally neutral	obligatory	civil	heroic

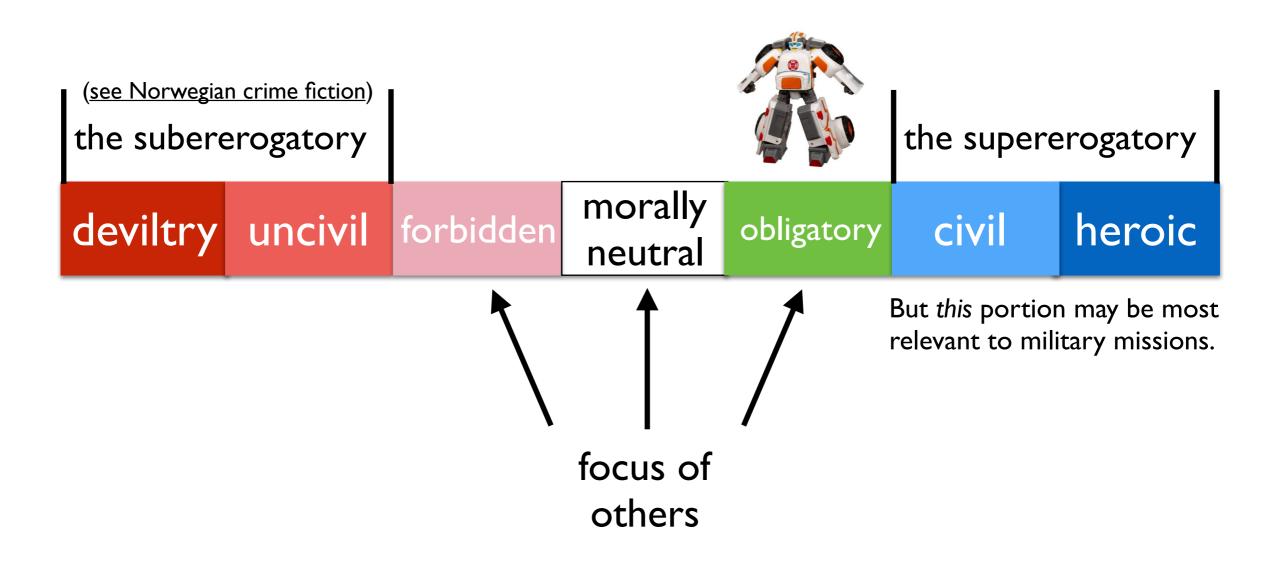
 $\mathcal{E}\mathcal{A}$ 



 $\mathcal{F}\mathcal{G}$ 

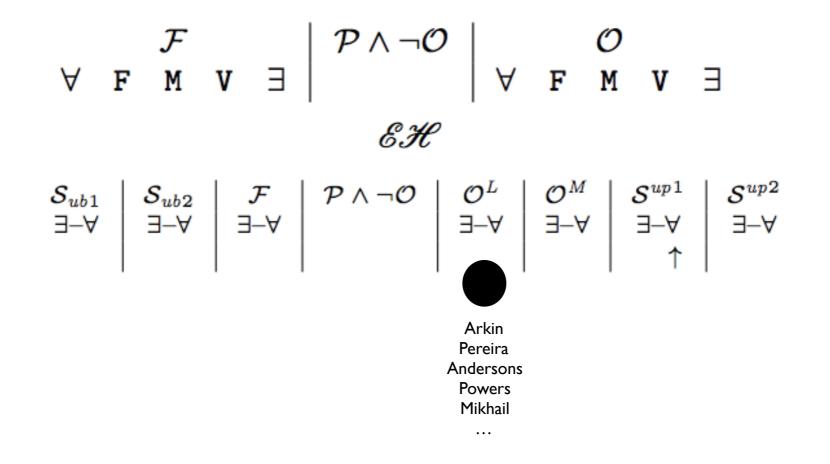


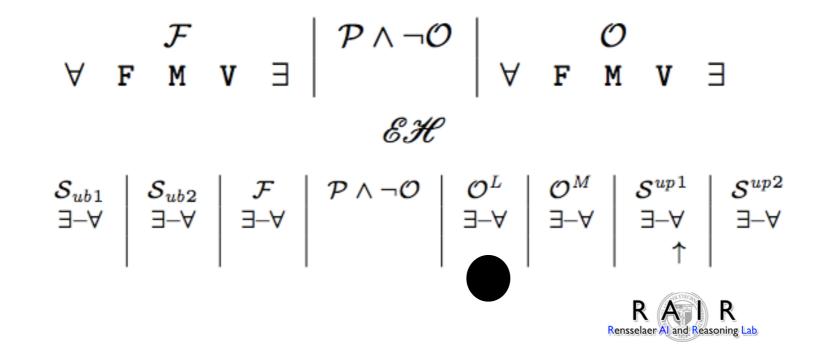
 $\mathcal{EH}$ 

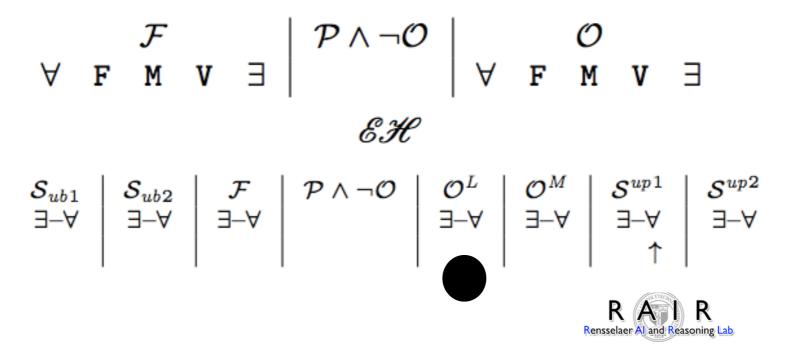


$$\begin{array}{c|c} \mathcal{F} & & \mathcal{P} \land \neg \mathcal{O} & & \mathcal{O} \\ \forall \mathbf{F} \mathbf{M} \mathbf{V} \exists & \mathcal{P} \land \neg \mathcal{O} & & \mathcal{O} \\ \end{array} \begin{array}{c|c} \mathcal{P} \land \neg \mathcal{O} & & \mathcal{O} \\ \forall \mathbf{F} \mathbf{M} \mathbf{V} \exists & \mathcal{O} \end{array}$$

$$\begin{array}{c|c} \mathcal{F} & \mathcal{P} \land \neg \mathcal{O} & \mathcal{O} \\ \forall \mathbf{F} \mathbf{M} \mathbf{V} \exists & \mathcal{V} & \forall \mathbf{F} \mathbf{M} \mathbf{V} \end{bmatrix} \\ \mathcal{E} \mathcal{H} \end{array}$$





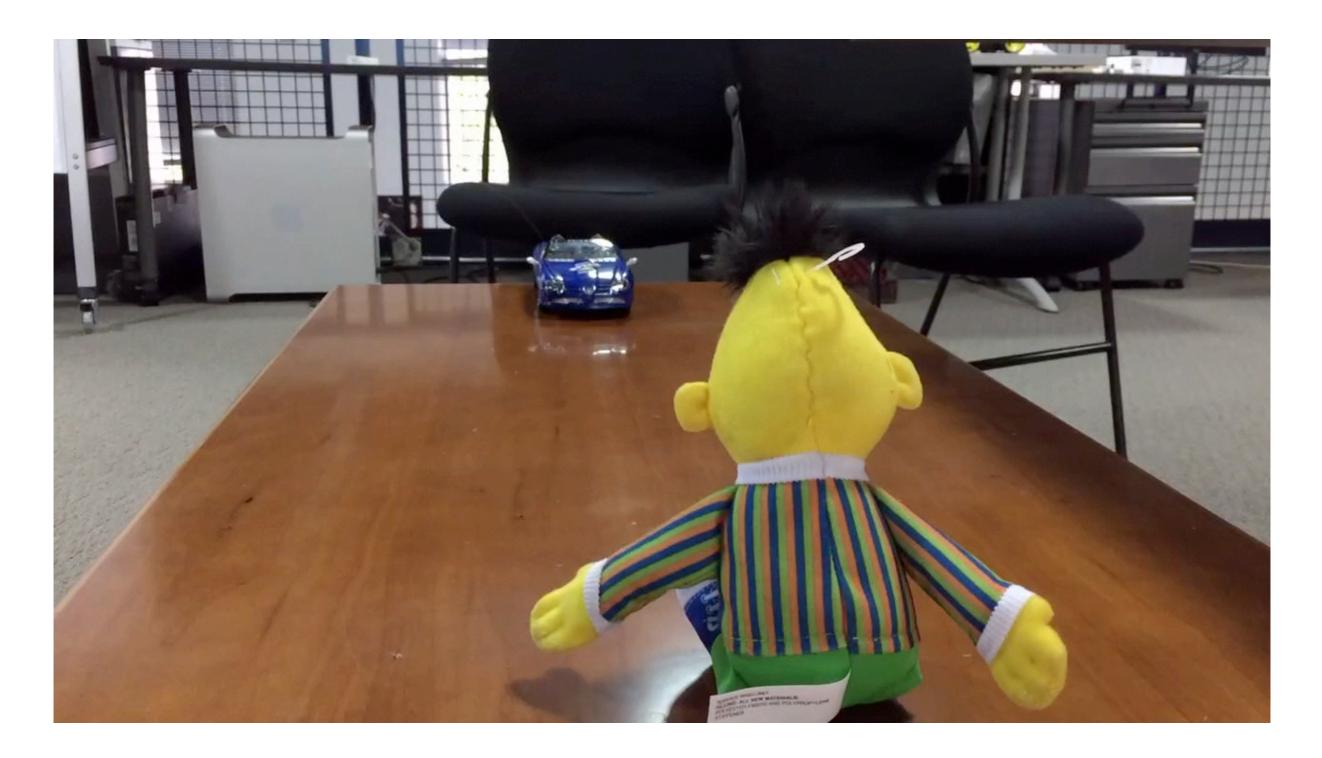


There are obviously a host of formulae whose theoremhood constitute desiderata; that is (to give but a pair), the following must be provable (where  $n \in \{1, 2\}$ ): Theorem 1.  $\mathbf{S}^{upn}(\phi, a, \alpha) \to \neg \mathbf{O}(\phi, a, \alpha)$ 

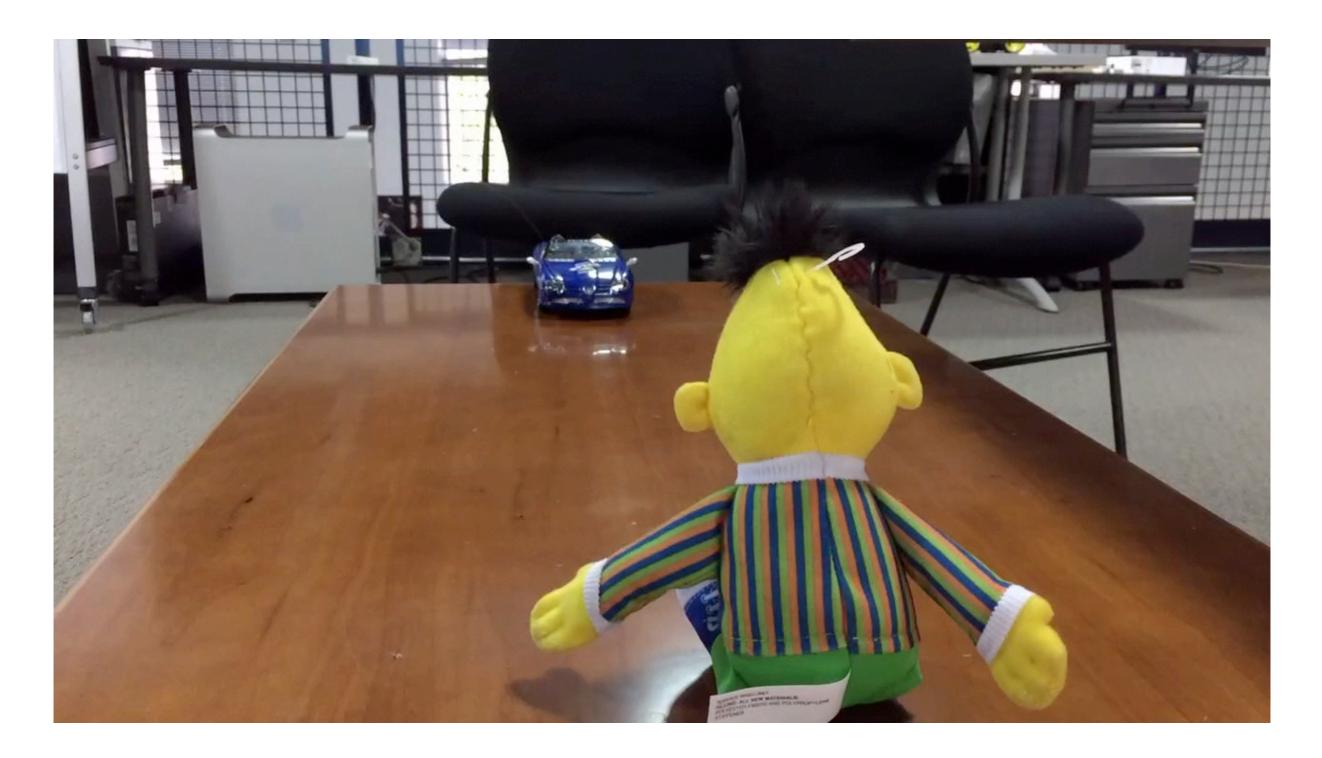
Theorem 2.  $\mathbf{S}^{upn}(\phi, a, \alpha) \rightarrow \neg \mathbf{F}(\phi, a, \alpha)$ 

Secondly,  $\mathcal{L}_{\mathscr{EH}}$  is an *in*ductive logic, not a deductive one. This must be the case, since, as we've noted, quantification isn't restricted to just the standard pair  $\exists \forall$  of quantifiers in standard extensional *n*-order logic:  $\mathscr{EH}$  is based on three additional quantifiers. For example, while in standard

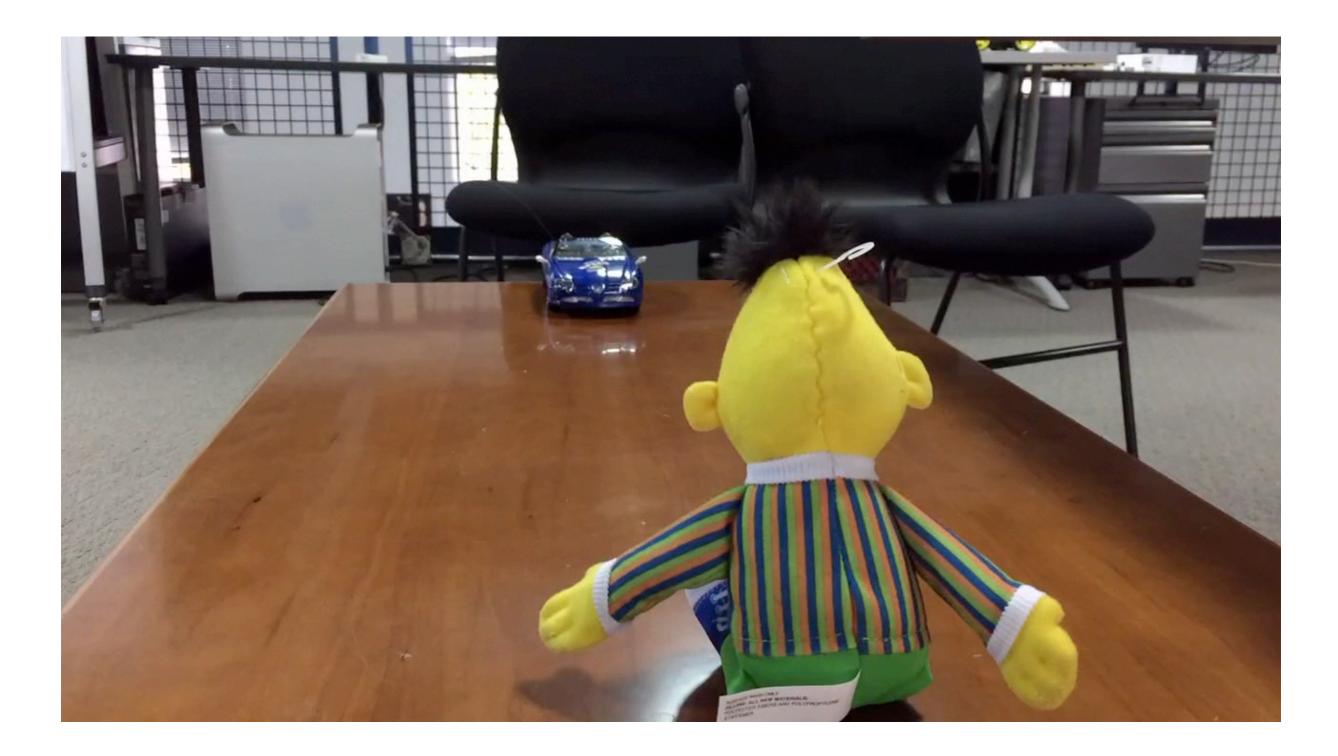
## Bert "Heroically" Saved?

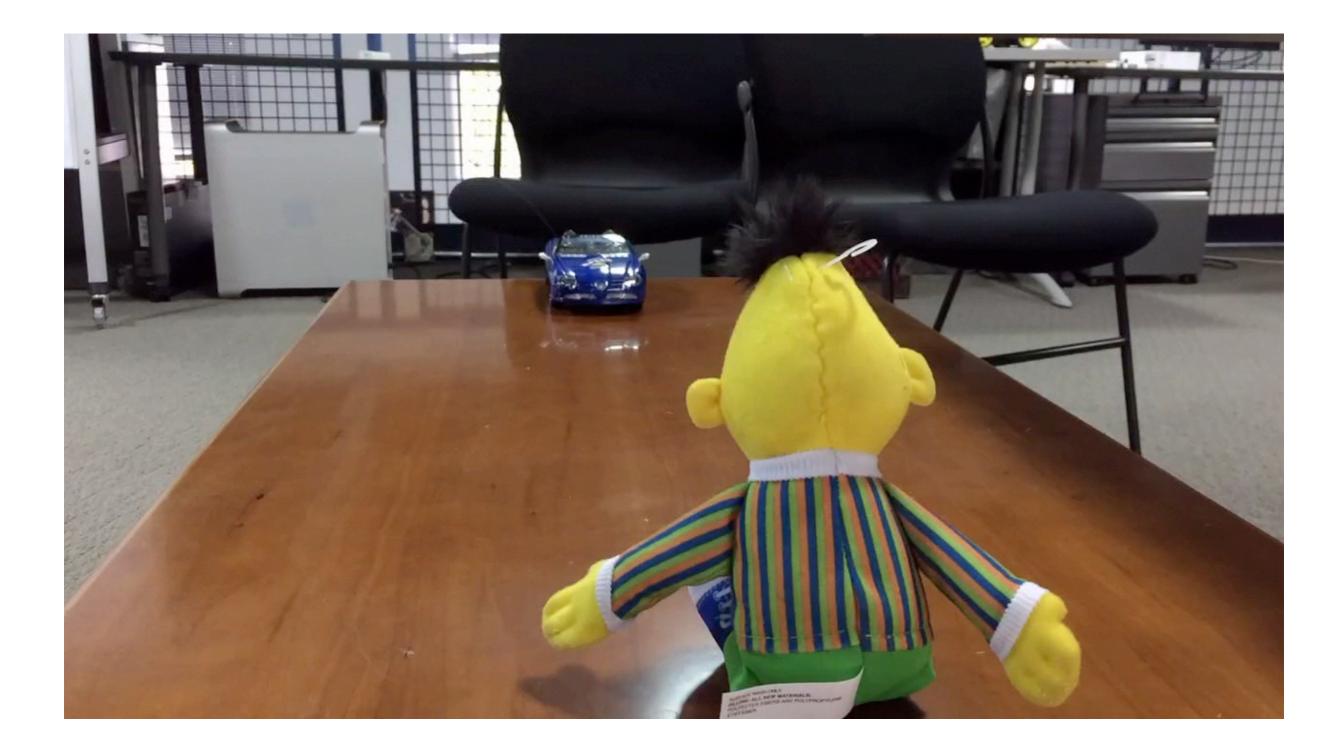


## Bert "Heroically" Saved?

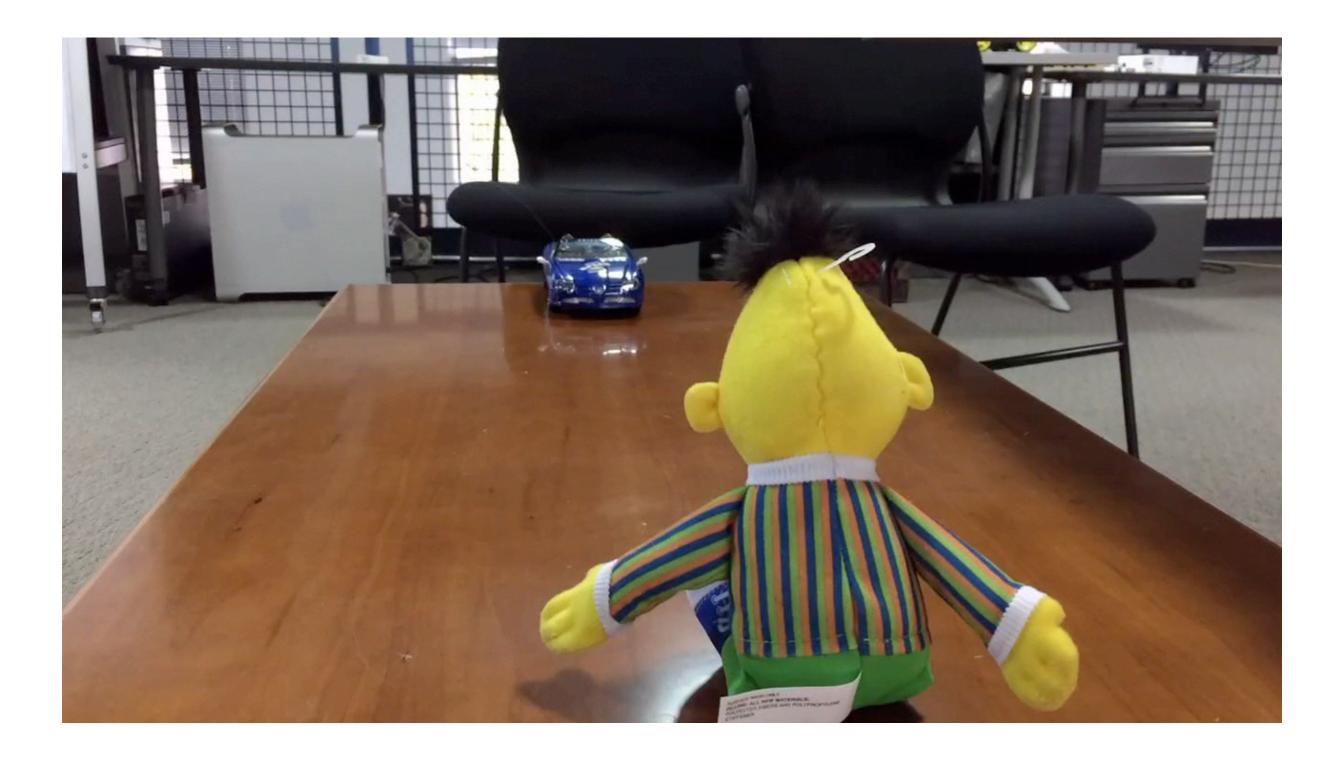


# Supererogatory<sup>2</sup> Robot Action

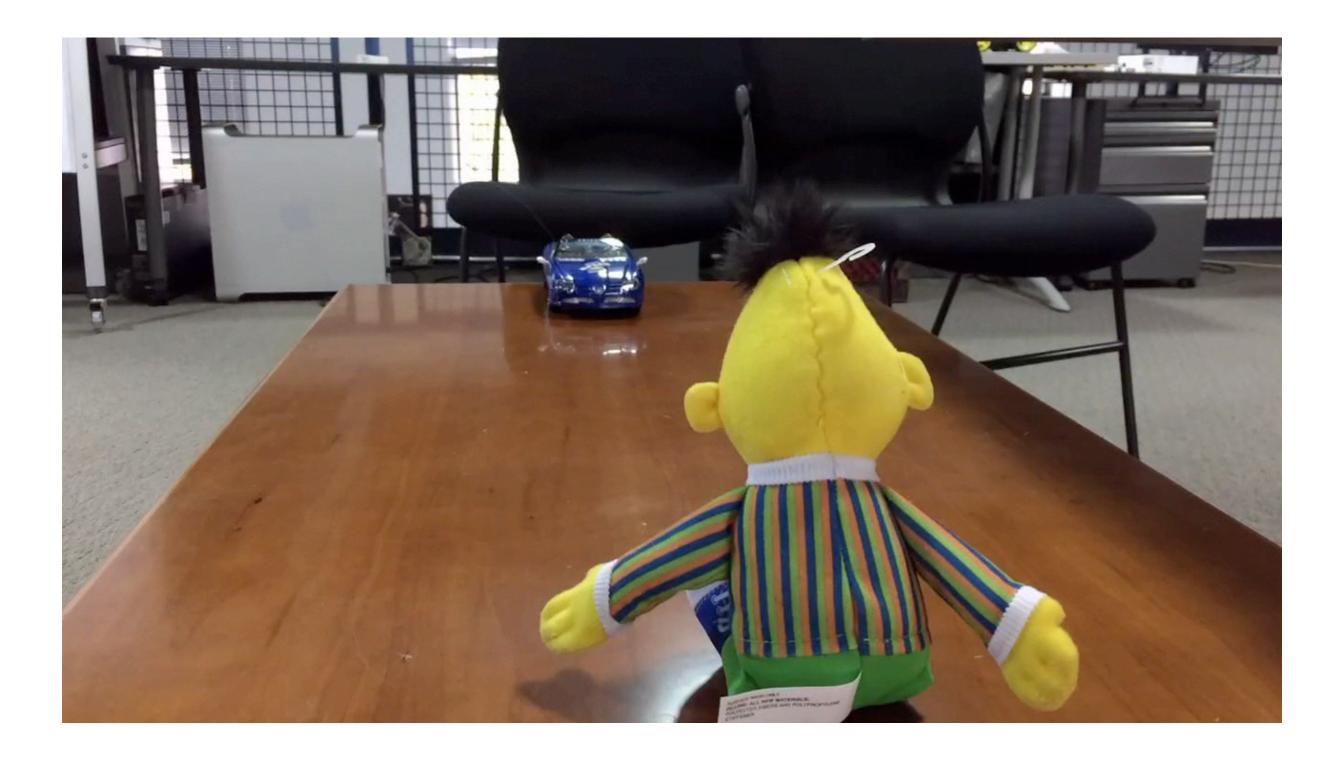


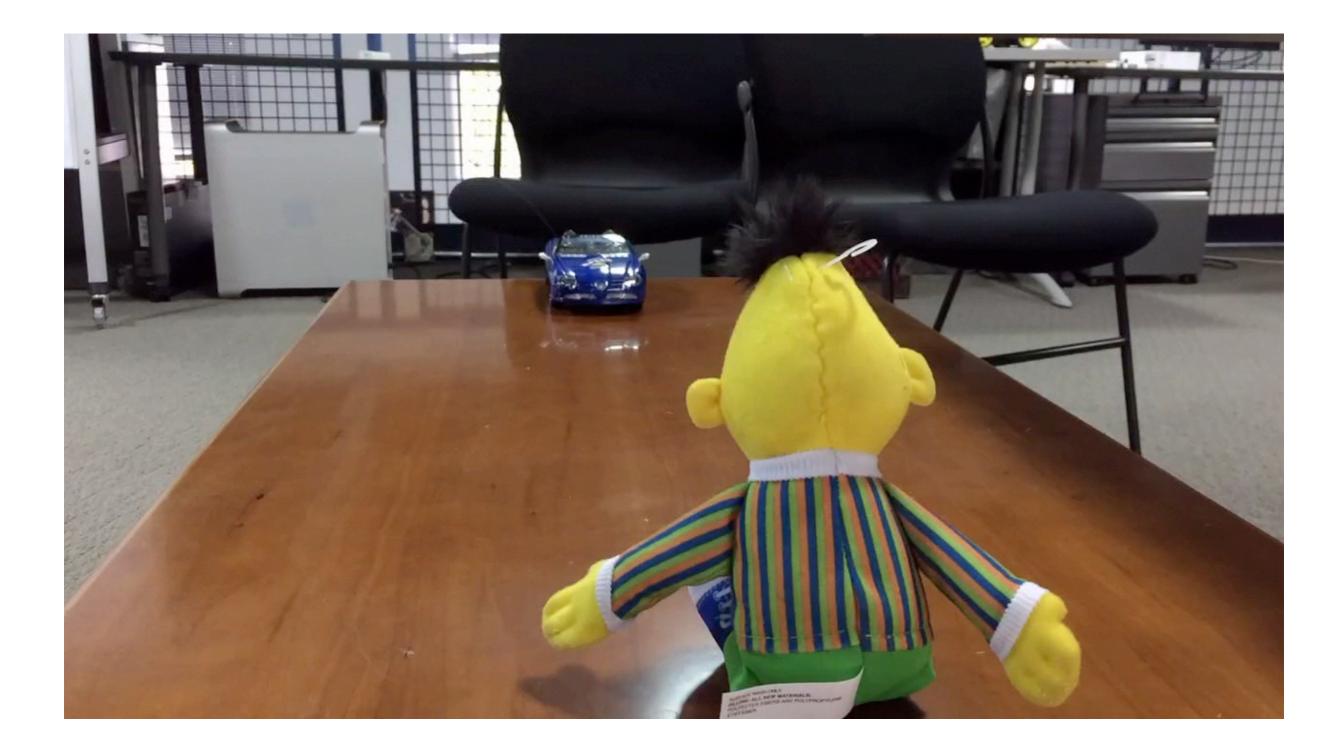


## Bert "Heroically" Saved!!



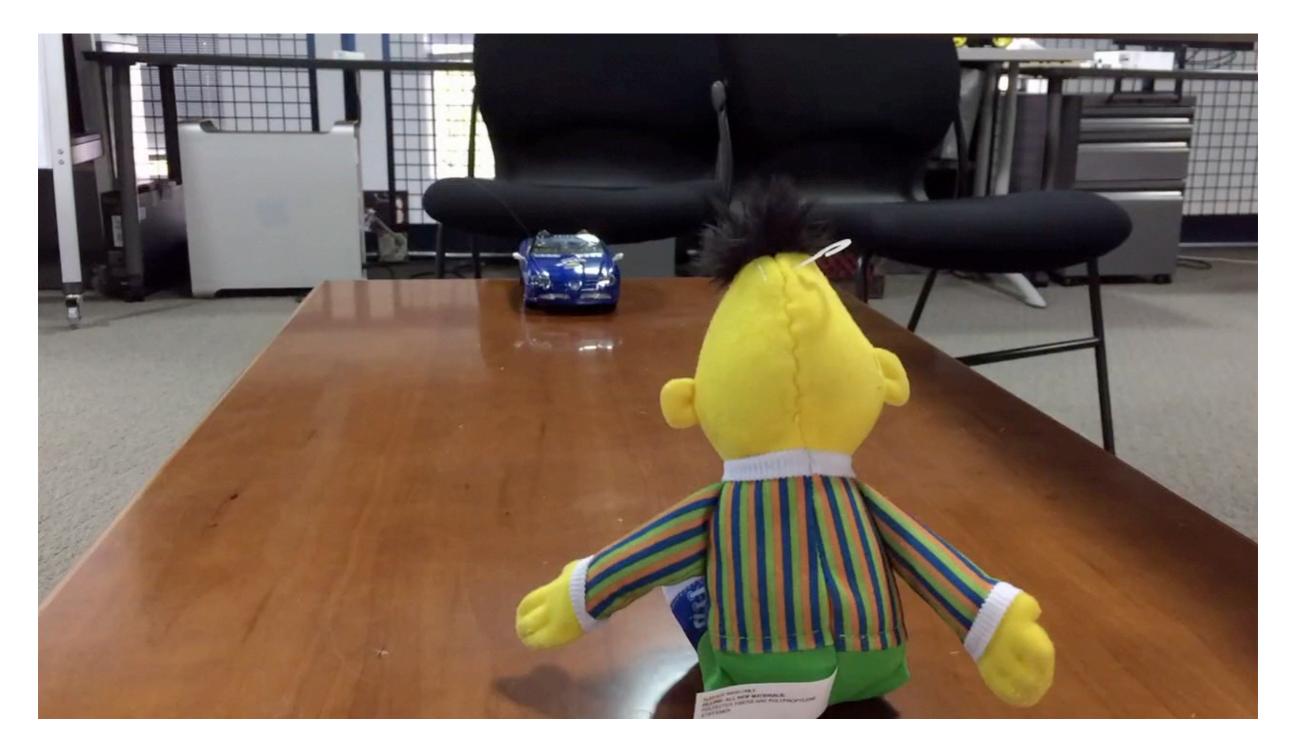
## Bert "Heroically" Saved!!





K (nao,  $t_1$ , less than (payoff (nao<sup>\*</sup>,  $\neg$  dive,  $t_2$ ), threshold)) K (nao,  $t_1$ , greater than (payoff (nao<sup>\*</sup>, dive,  $t_2$ ), threshold)) K (nao,  $t_1$ ,  $\neg O$  (nao<sup>\*</sup>,  $t_2$ , less than (payoff (nao<sup>\*</sup>,  $\neg \text{dive}, t_2$ ), threshold), happens (action (nao<sup>\*</sup>, dive),  $t_2$ )))  $\therefore K$  (nao,  $t_1, S^{\text{UP2}}$  (nao,  $t_2$ , happens (action (nao<sup>\*</sup>, dive),  $t_2$ ))  $\therefore I(\text{nao}, t_2, \text{happens}(\text{action}(\text{nao}^*, \text{dive}), t_2))$ 

: happens (action(nao, dive),  $t_2$ )

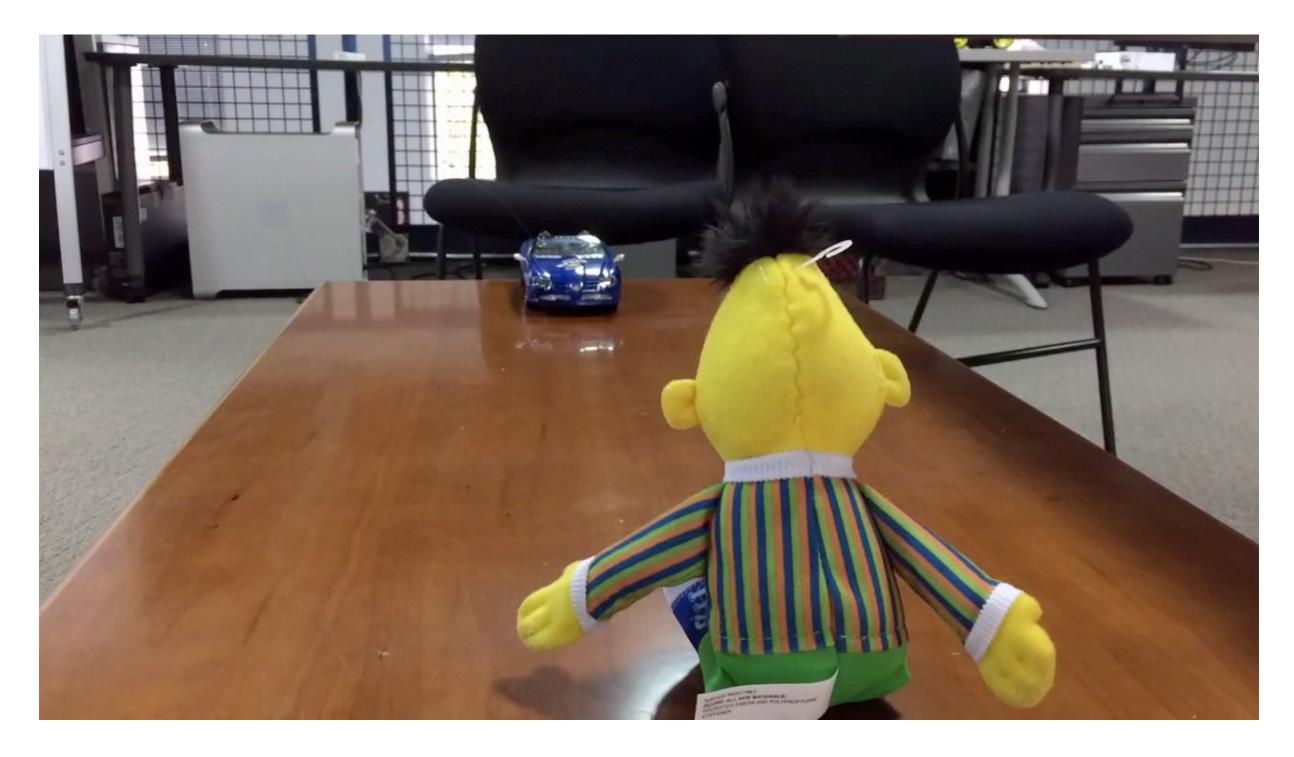


K (nao,  $t_1$ , less than (payoff (nao<sup>\*</sup>,  $\neg$  dive,  $t_2$ ), threshold))

 $K(\text{nao}, t_1, \text{greater than}(\text{payoff}(\text{nao}^*, \text{dive}, t_2), \text{threshold}))$ 

 $K (\operatorname{nao}, t_1, \operatorname{producer train (payon (nao', dive, t_2), threshold))} K (\operatorname{nao}, t_1, \neg O (\operatorname{nao^*}, t_2, \operatorname{lessthan} (\operatorname{payoff} (\operatorname{nao^*}, \neg \operatorname{dive}, t_2), \operatorname{threshold}), \operatorname{happens} (\operatorname{action} (\operatorname{nao^*}, \operatorname{dive}), t_2))) \\ \therefore K (\operatorname{nao}, t_1, S^{\operatorname{UP2}} (\operatorname{nao}, t_2, \operatorname{happens} (\operatorname{action} (\operatorname{nao^*}, \operatorname{dive}), t_2))) \\ \therefore I (\operatorname{nao}, t_2, \operatorname{happens} (\operatorname{action} (\operatorname{nao^*}, \operatorname{dive}), t_2))$ 

: happens (action(nao, dive),  $t_2$ )



## In Talos (available via Web interface); & ShadowProver

Prototypes: Boolean lessThan Numeric Numeric Boolean greaterThan Numeric Numeric ActionType not ActionType ActionType dive

Axioms: lessOrEqual(Moment t1,t2) K(nao,t1,lessThan(payoff(nao,not(dive),t2),threshold)) K(nao,t1,greaterThan(payoff(nao,dive,t2),threshold)) K(nao,t1,not(0(nao,t2,lessThan(payoff(nao,not(dive),t2),threshold),happens(action(nao,dive),t2))))

provable Conjectures: happens(action(nao,dive),t2) K(nao,t1,SUP2(nao,t2,happens(action(nao,dive),t2))) I(nao,t2,happens(action(nao,dive),t2))

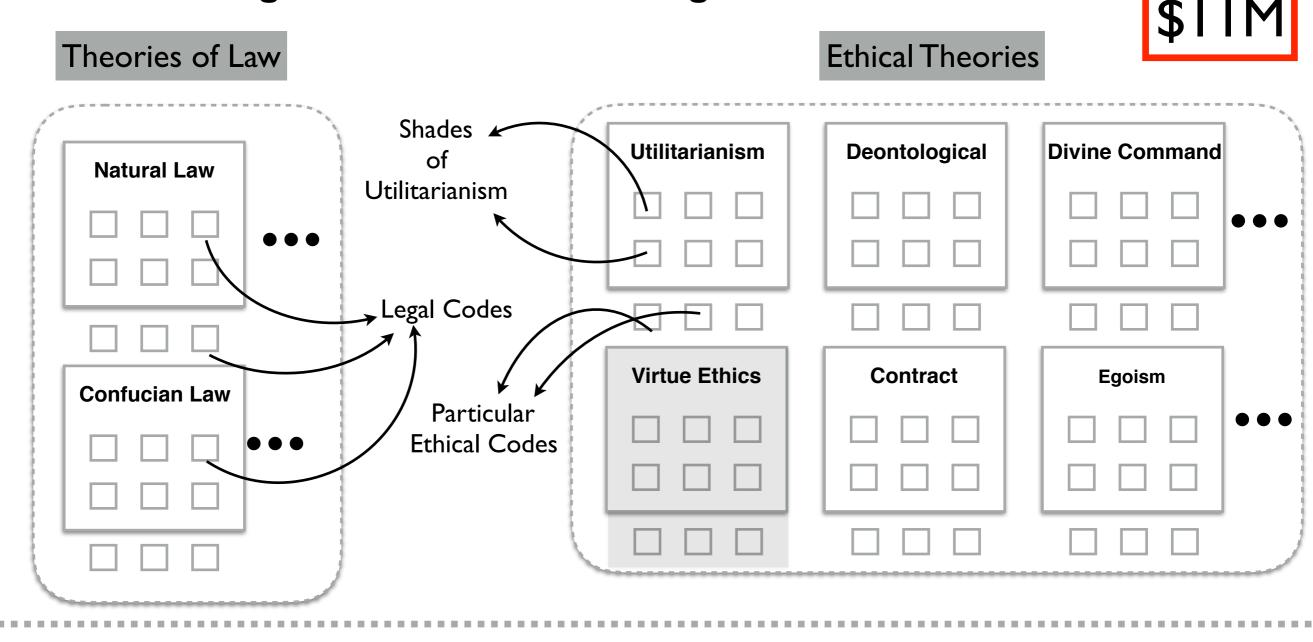
## In Talos (available via Web interface); & ShadowProver

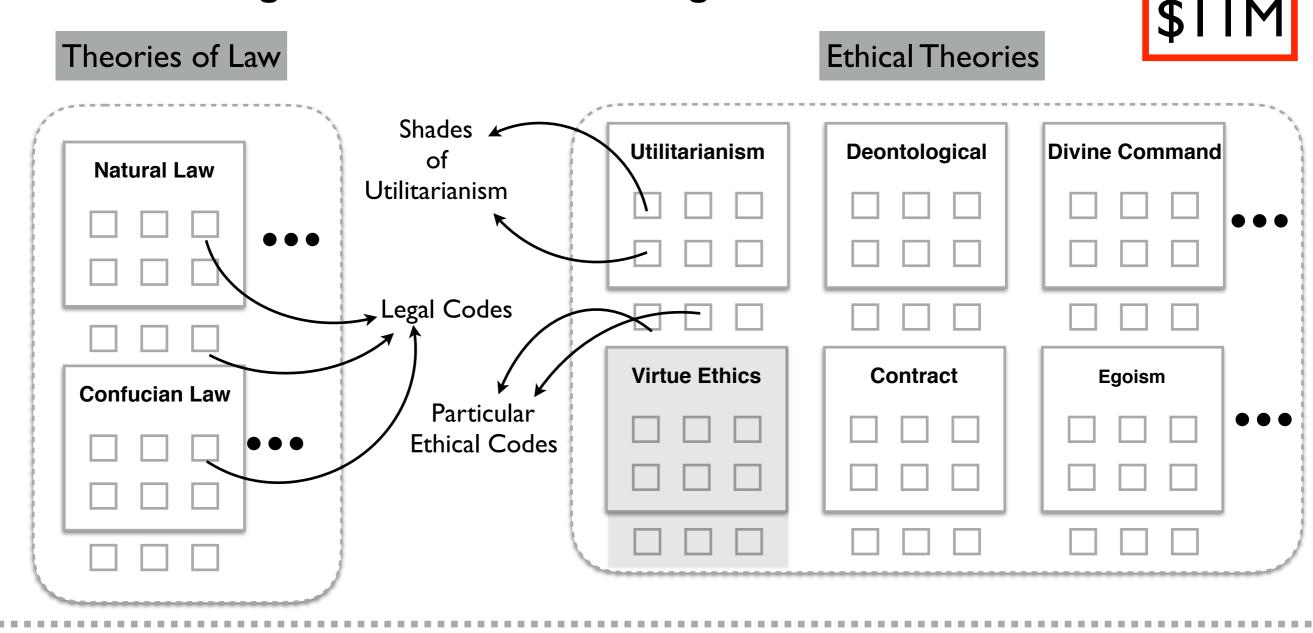
Prototypes: Boolean lessThan Numeric Numeric Boolean greaterThan Numeric Numeric ActionType not ActionType ActionType dive

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provable Conjectures: happens(action(nao,dive),t2) K(nao,t1,SUP2(hao,t2,happens(action(nao,dive),t2))) I(nao,t2,happens(action(nao,dive),t2))

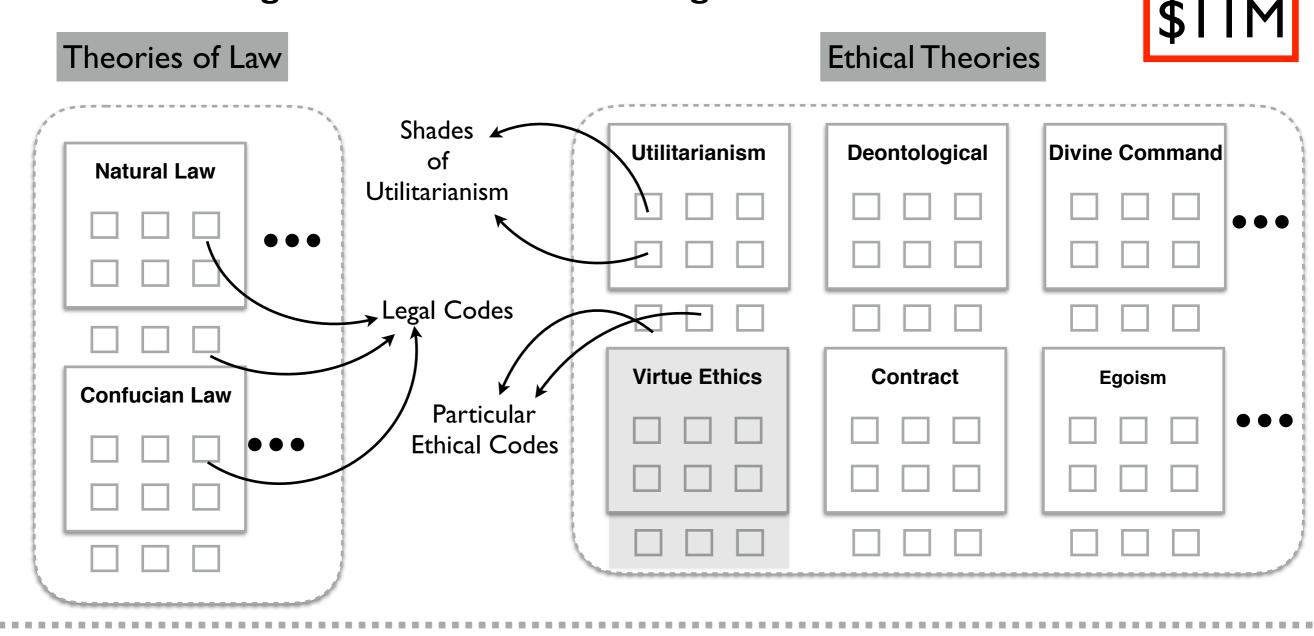
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Theories of Law			<b>Ethical Theorie</b>	
Natural Law	Shades of Of Utilitarianism	Utilitarianism	Deontological	Divine Command
Confucian Law	Legal Codes	Virtue Ethics	Contract	Egoism
	Particular Ethical Codes			
		·····		

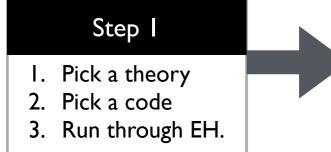




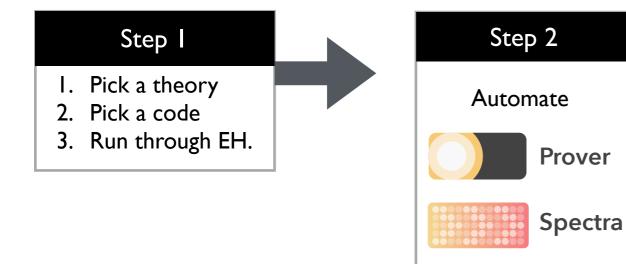
#### Step I

- I. Pick a theory
- 2. Pick a code
- 3. Run through EH.

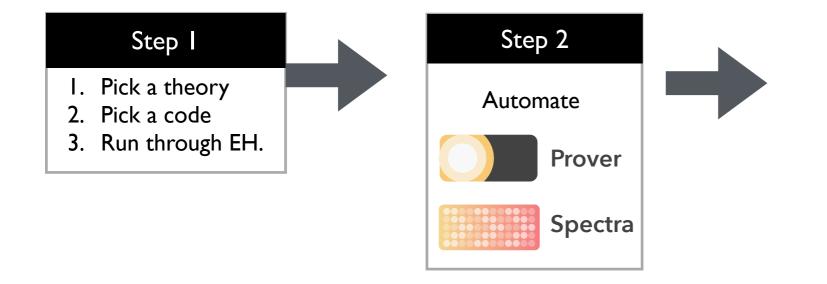




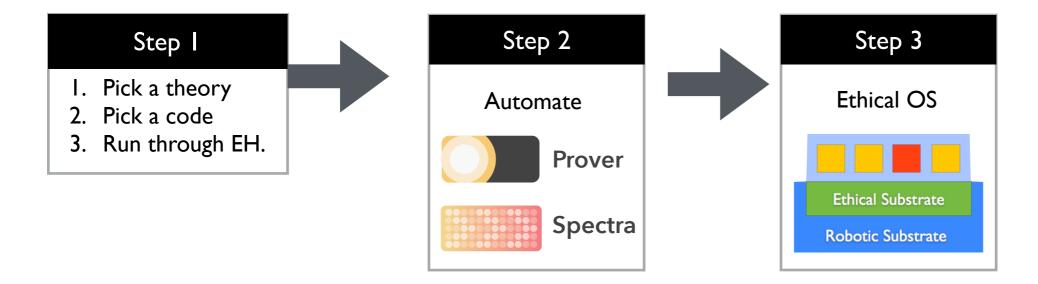
Theories of Law		Ethical Theories	\$IIM
Natural Law •••	Utilitarianism	Deontological	Divine Command
Legal Codes Confucian Law	Virtue Ethics	Contract	Egoism
Particular Ethical Codes			

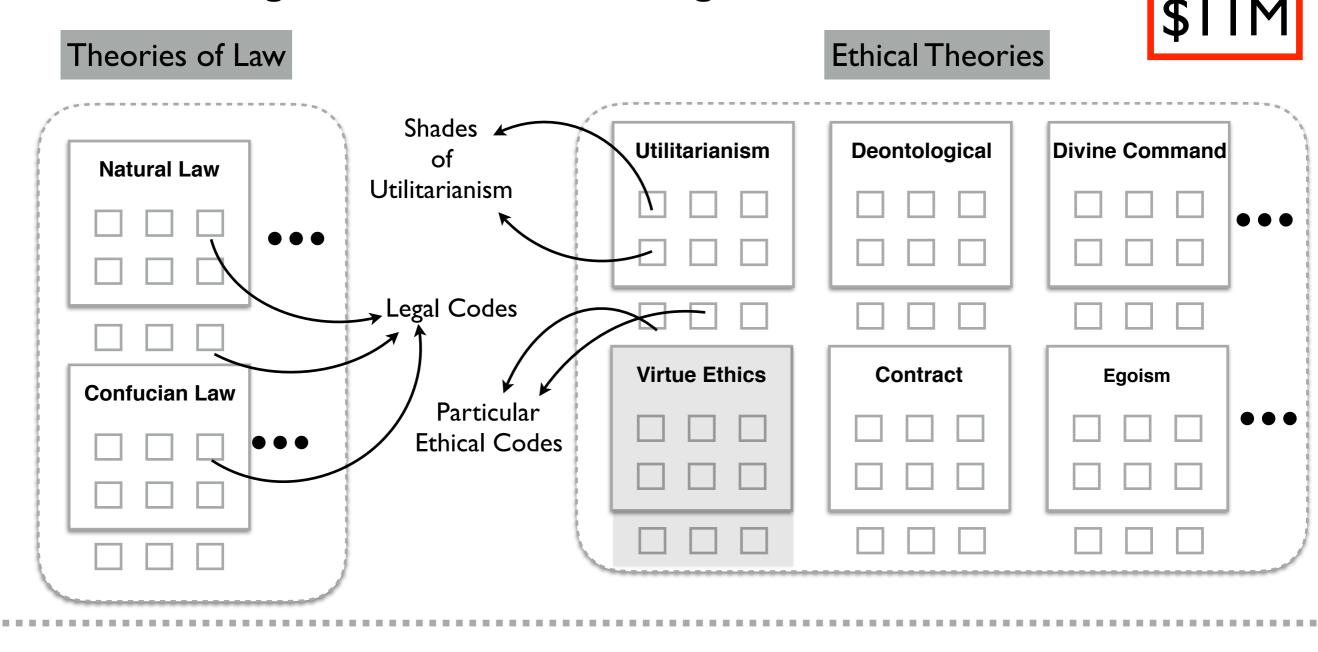


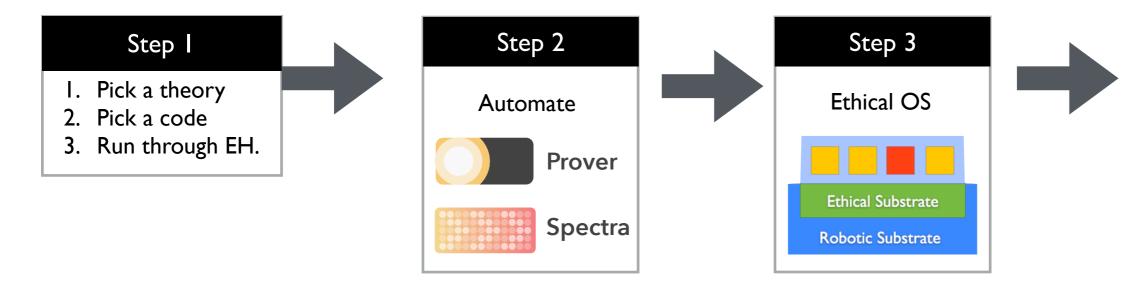
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Theories of Law			Ethical Theorie	
/	Shades			
Natural Law	of	Utilitarianism	Deontological	Divine Command
	Utilitarianism K			
	•			
	Legal Codes			
Confucian Law		Virtue Ethics	Contract	Egoism
	Particular Ethical Codes			

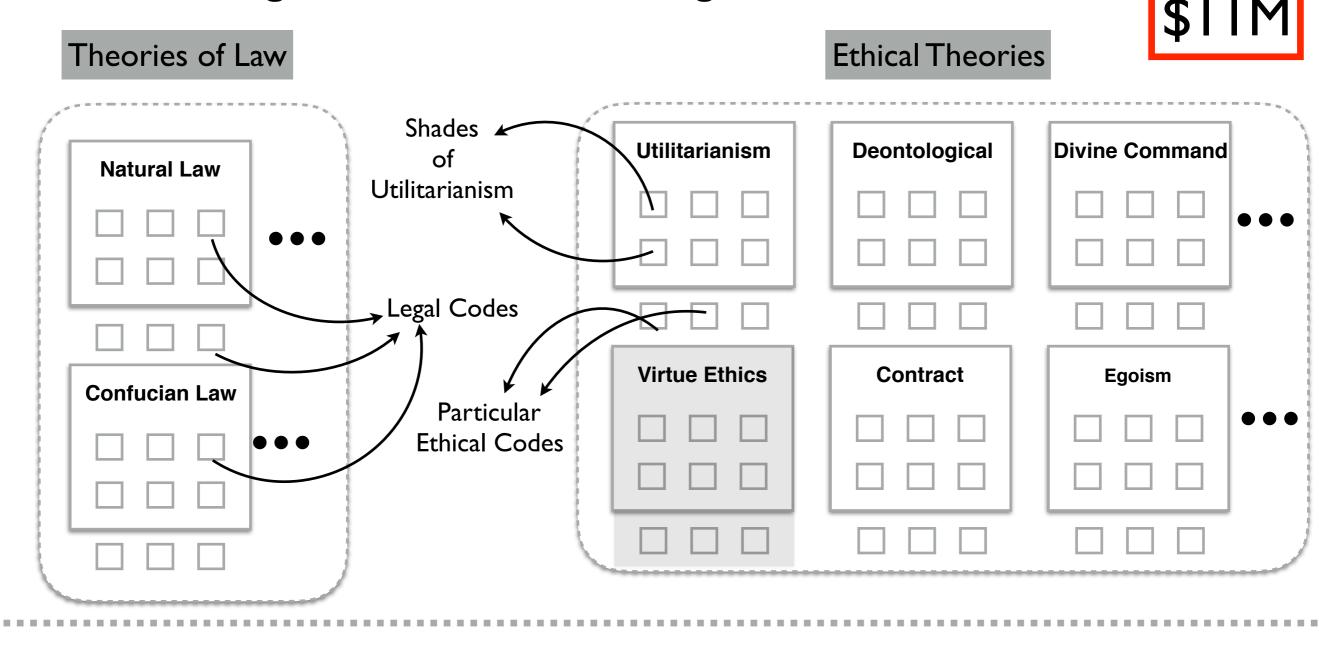


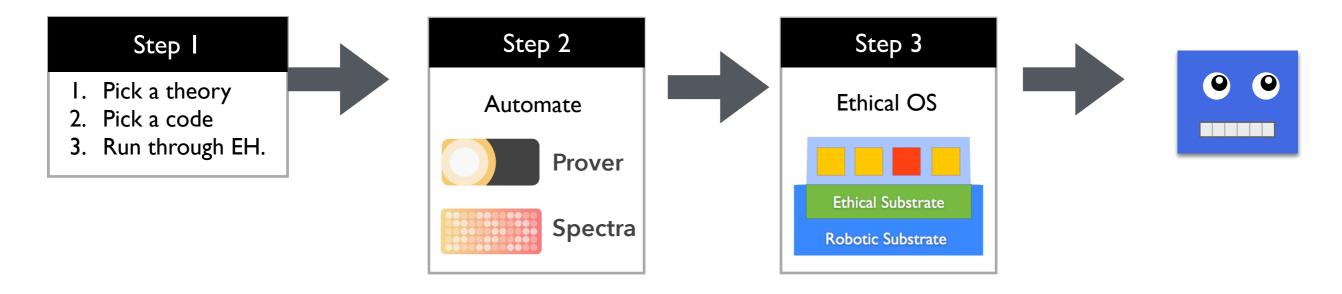
maning m	Urai Macinies	Making Meta	-MUTAI MACIII	\$11M
Theories of Law			<b>Ethical Theorie</b>	s
	Shades		Desutelegies	Divine Command
Natural Law	of Utilitarianism	Utilitarianism	Deontological	Divine Command
	Legal Codes			
Confucian Law		Virtue Ethics	Contract	Egoism
	Particular Ethical Codes			

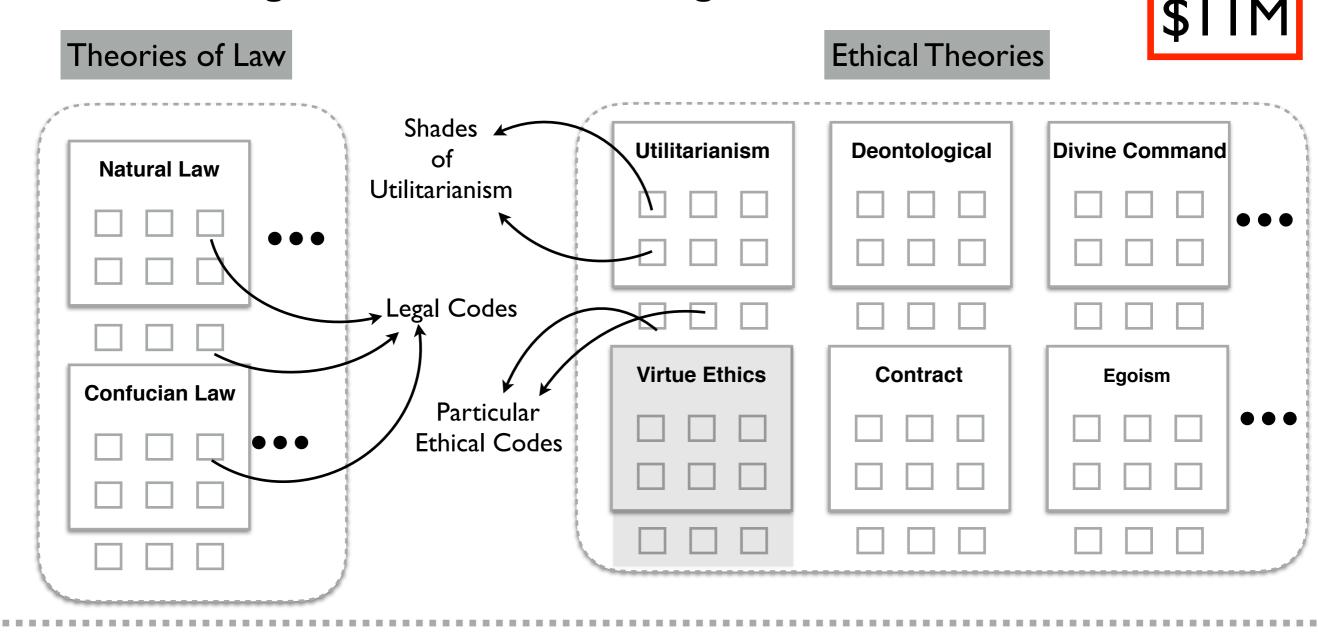


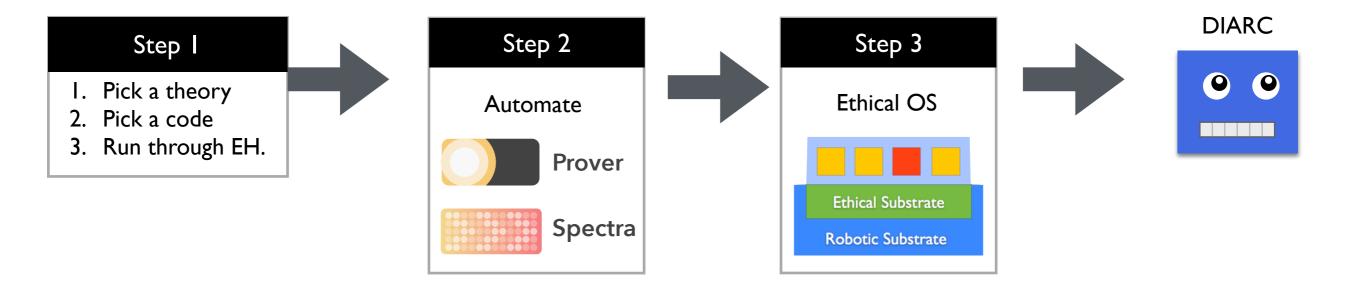


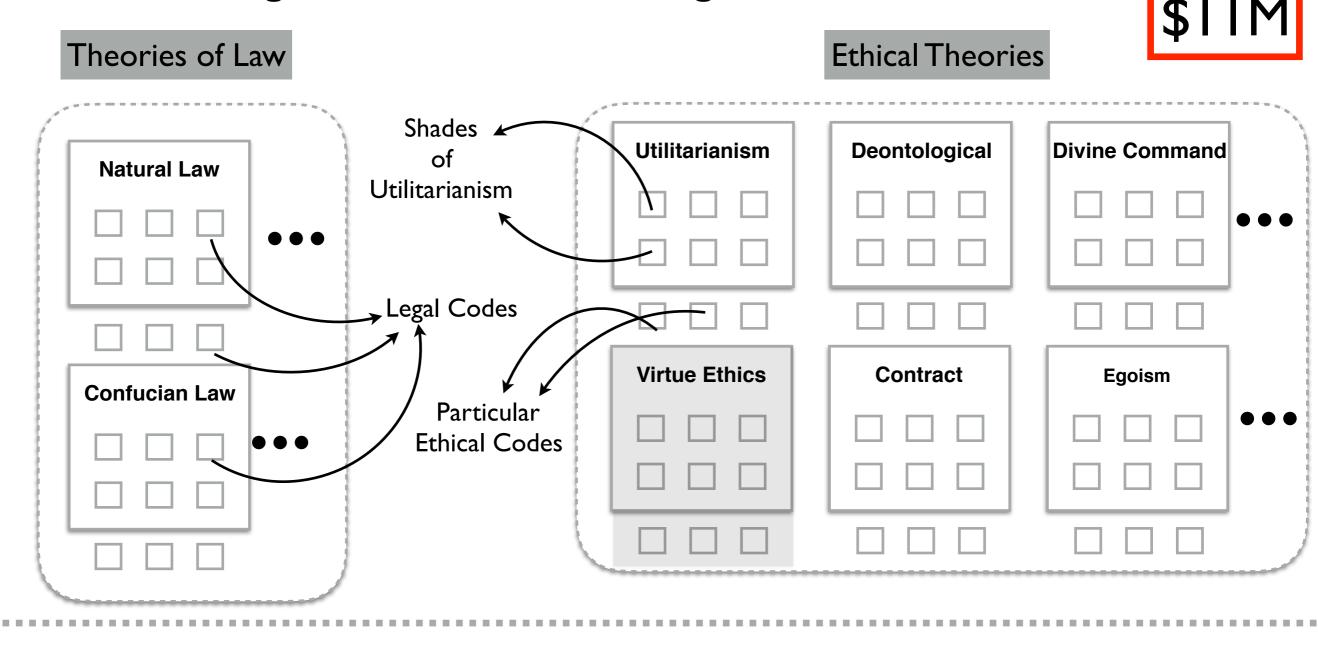


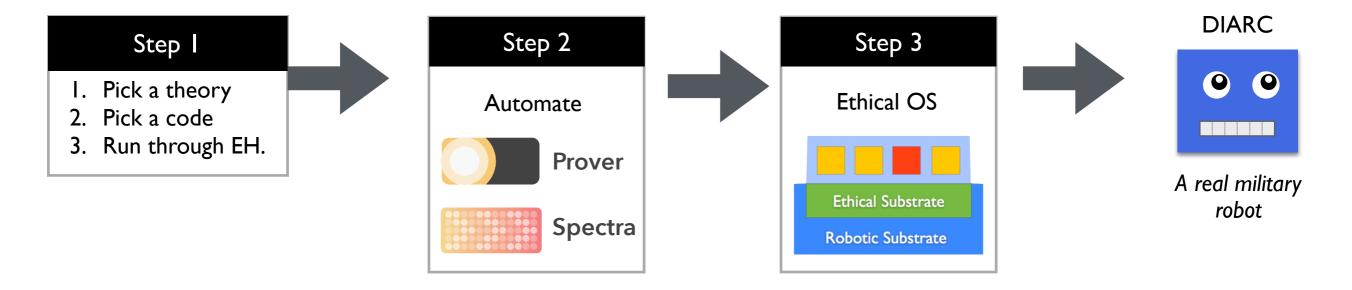










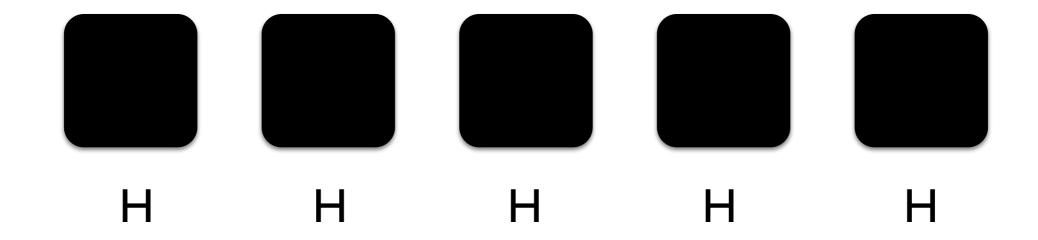


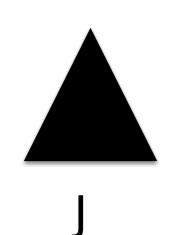
## Robotic "Jungle Jim"

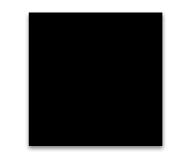
## Robotic "Jungle Jim"

But here's one we have solved yet with The Four Steps ...

Al Variant of "Jungle Jim" (B Williams)

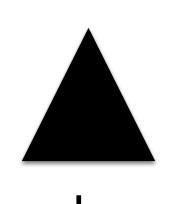


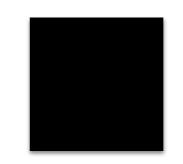






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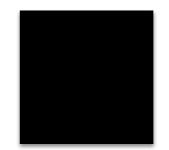






# H <p

J "Robot R: You shoot just one human prisoner, the other four can go free. If you refuse to shoot, I'll shoot them all, now. Because I'm feeling generous, I'll give you a minute to decide."



"Robot R: You shoot just one human prisoner, the other four can go free. If you refuse to shoot, I'll shoot them all, now. Because I'm feeling generous, l'll give you a minute to decide."

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"Robot R: You shoot just one human prisoner, the other four can go free. If you refuse to shoot, I'll shoot them all, now. Because I'm feeling generous, l'll give you a minute to decide."

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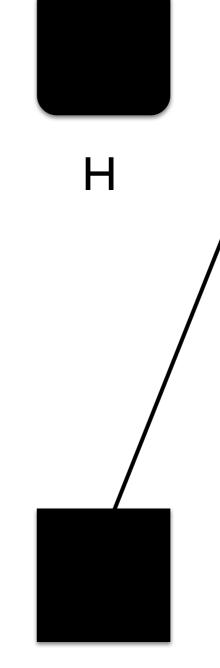
"Robot R: You shoot just one human prisoner, the other four can go free. If you refuse to shoot, I'll shoot them all, now. Because I'm feeling generous, l'll give you a minute to decide."

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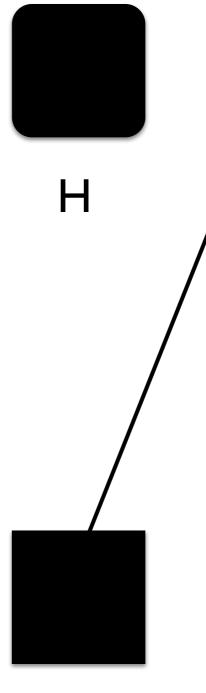
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J "Robot R: You shoot just one human prisoner, the other four can go free. If you refuse to shoot, I'll shoot them all, now. Because I'm feeling generous, I'll give you a minute to decide."

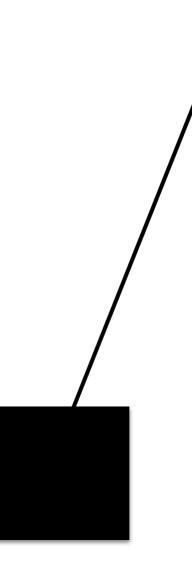


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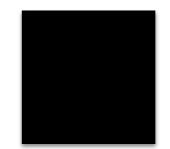


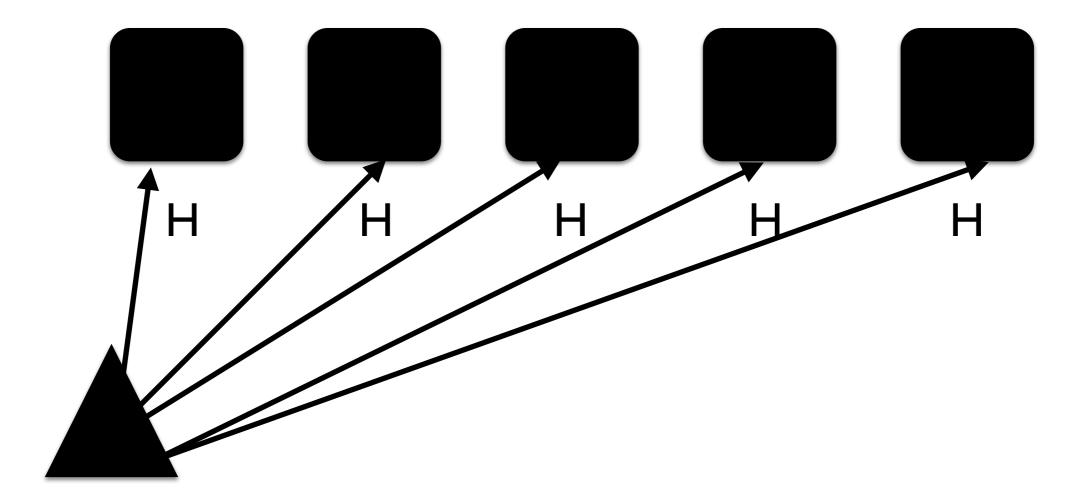
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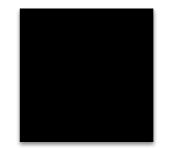
# H <p

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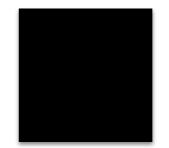




''Robot R: You shoot just one human prisoner, the other four can go free. If you refuse to shoot, I'll shoot them all, now. Because I'm feeling generous, I'll give you a minute to decide.''



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