

Modal Logic: First Steps

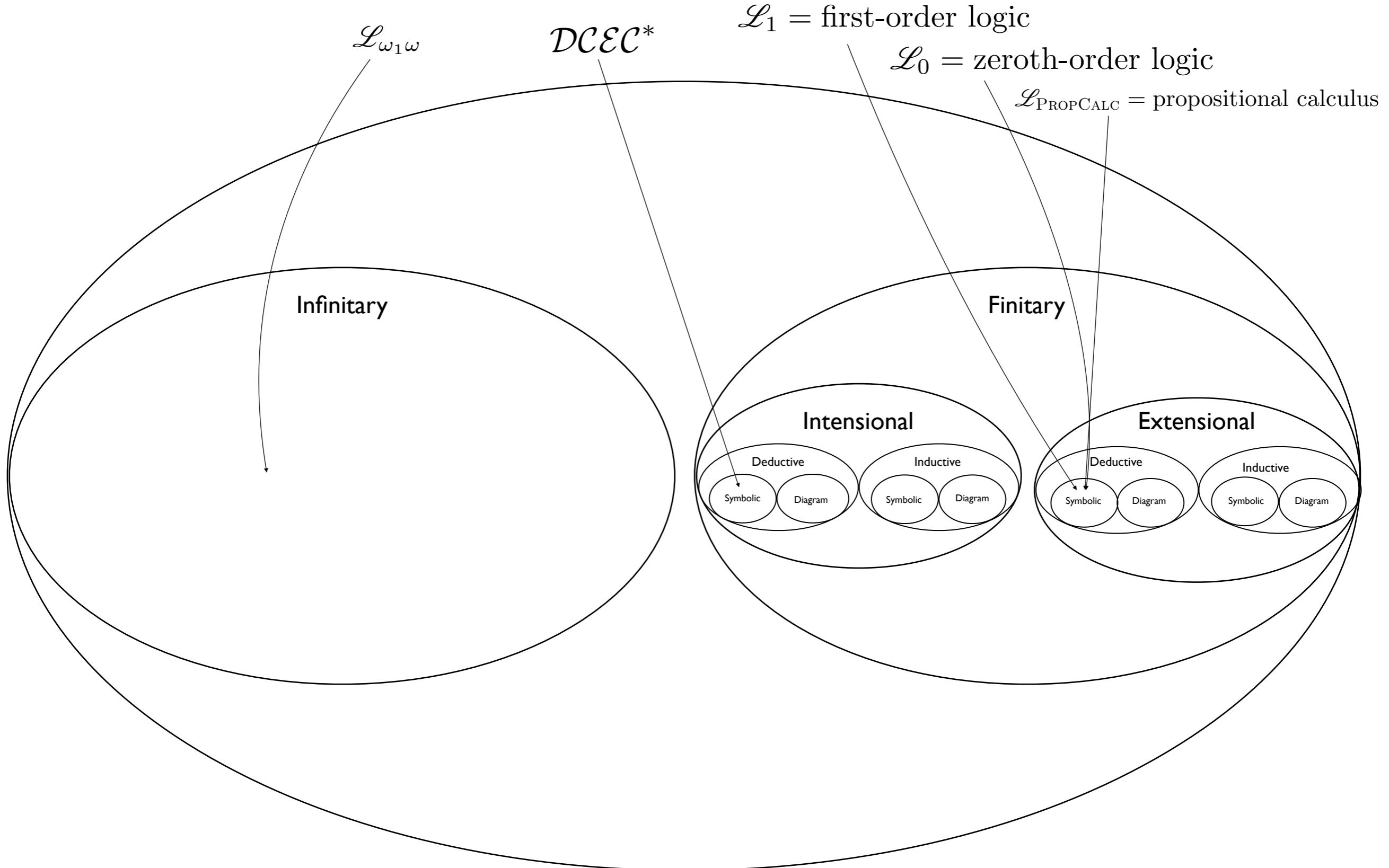
Selmer Bringsjord

Rensselaer AI & Reasoning (RAIR) Lab
Department of Cognitive Science
Department of Computer Science
Lally School of Management & Technology
Rensselaer Polytechnic Institute (RPI)
Troy, New York 12180 USA

Intro to (Formal) Logic (& AI) — (IFLAI)
4/2/2020



The Universe of Logics



The Smallest Infinitary Logic

$\mathcal{L}_{\omega_1\omega}$

The Smallest Infinitary Logic

$$\mathcal{L}_{\omega_1\omega}$$

$\bigvee \Phi$ (where $\Phi := \phi_1, \phi_2, \dots$) is $\phi_1 \vee \phi_2 \dots \infty$

$\bigwedge \Phi$ (where $\Phi := \phi_1, \phi_2, \dots$) is $\phi_1 \wedge \phi_2 \dots \infty$

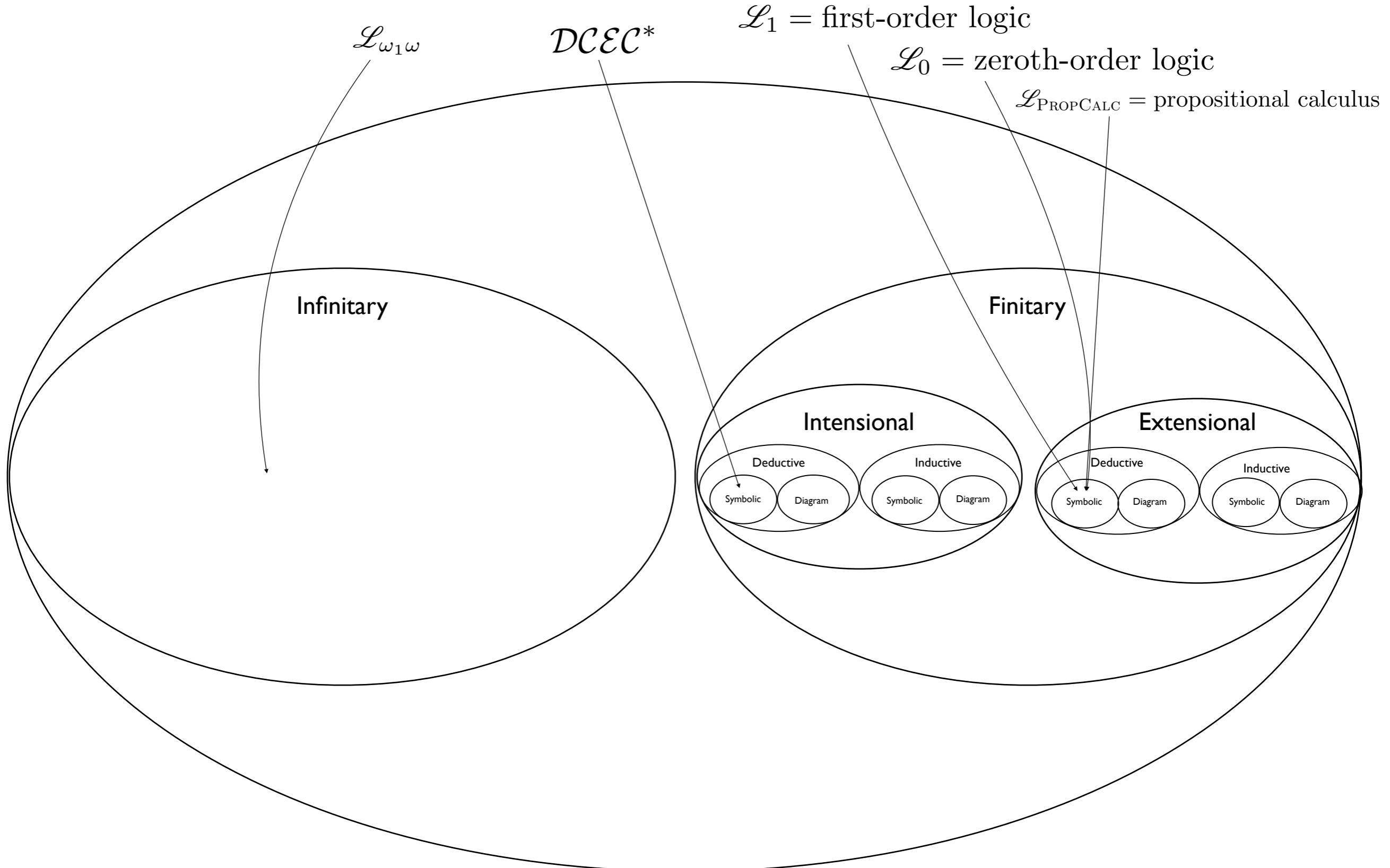
The Smallest Infinitary Logic

$$\mathcal{L}_{\omega_1\omega}$$

$$\begin{aligned} \bigvee \Phi &(\text{where } \Phi := \phi_1, \phi_2, \dots) \text{ is } \phi_1 \vee \phi_2 \dots \infty \\ \bigwedge \Phi &(\text{where } \Phi := \phi_1, \phi_2, \dots) \text{ is } \phi_1 \wedge \phi_2 \dots \infty \end{aligned}$$

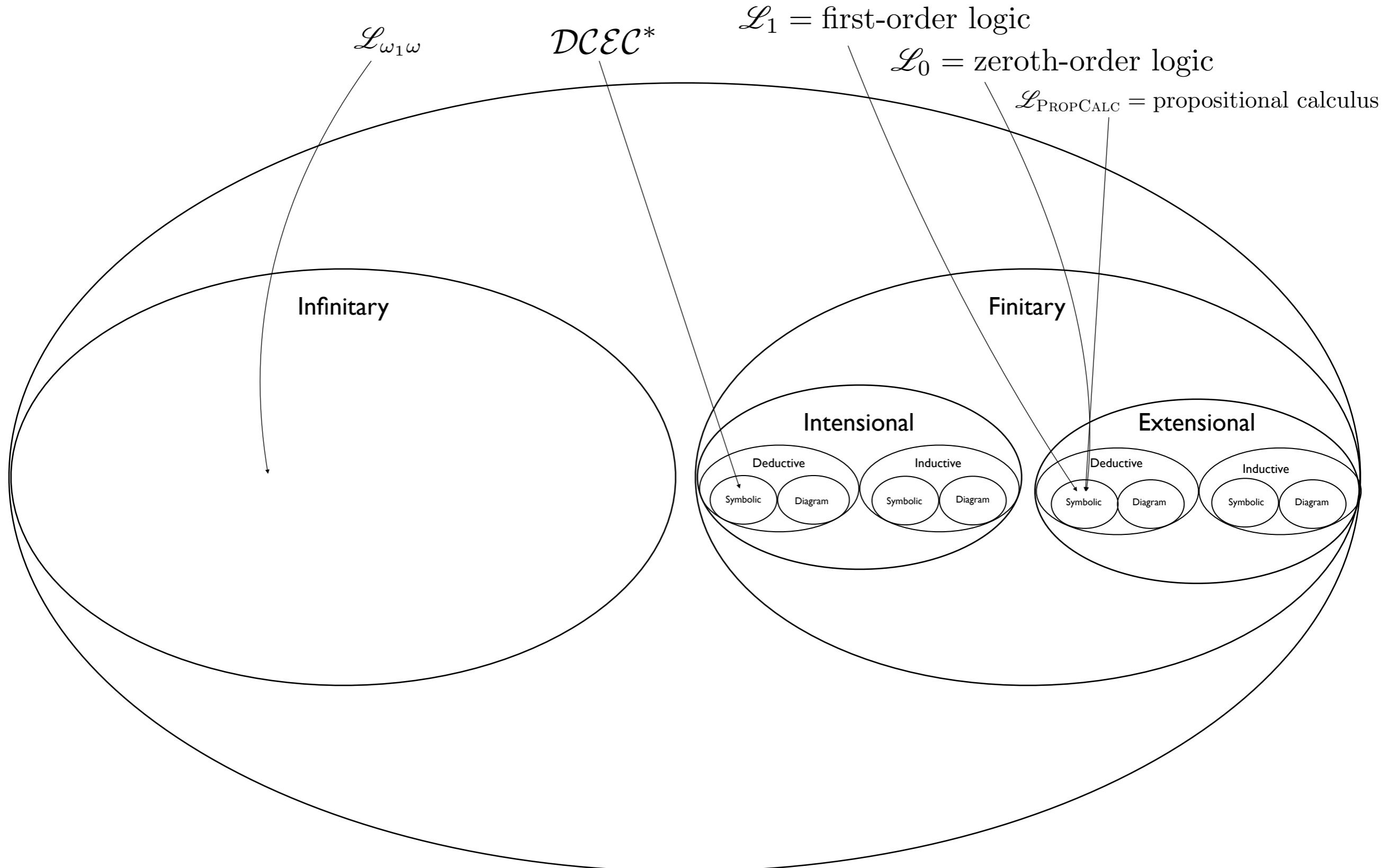
Finitude *can* be captured in *this* logic (but not in \mathcal{L}_1 ; that's why no one won \$1000). How?

The Universe of Logics



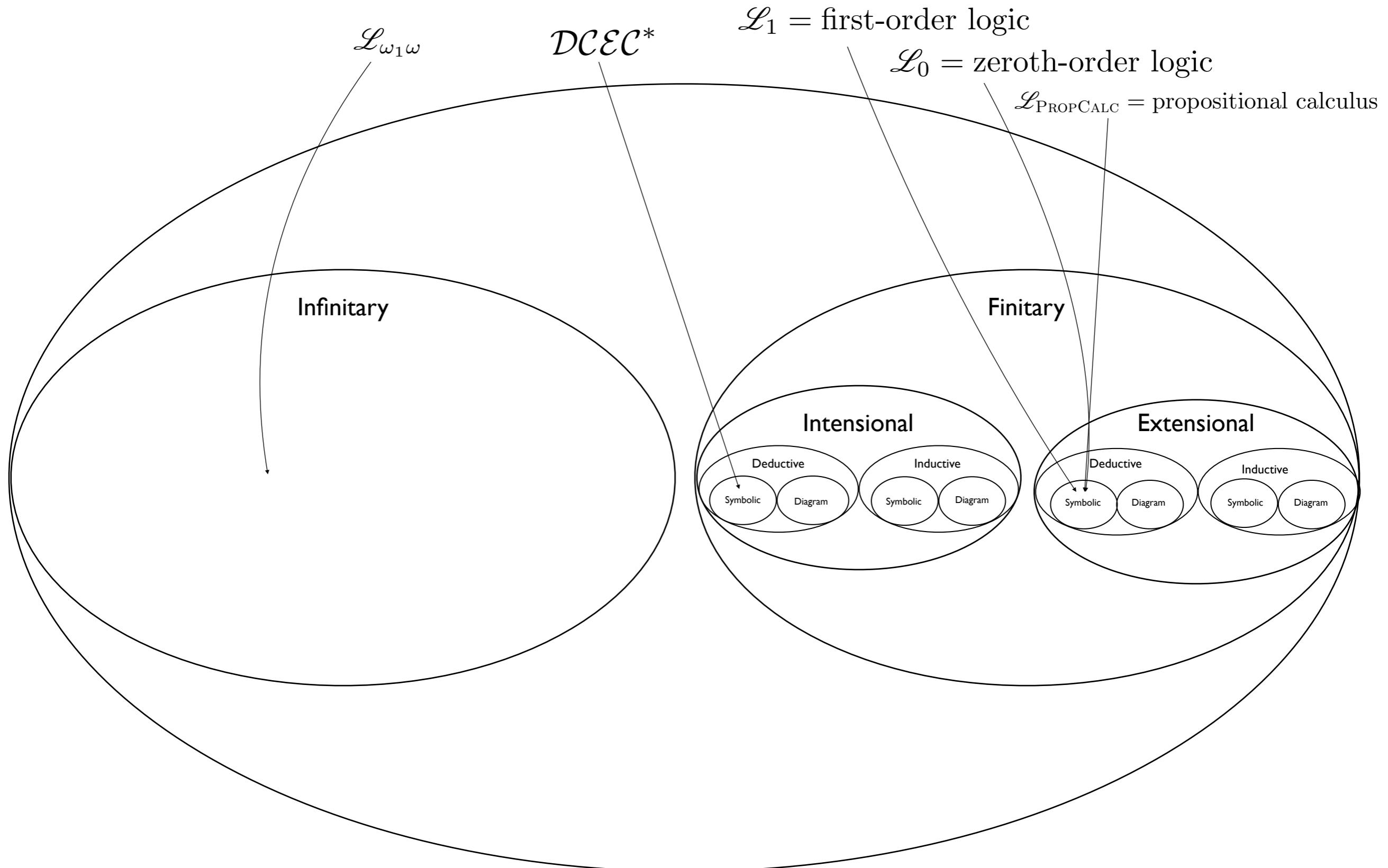
\mathcal{L}_3 = third-order logic

The Universe of Logics



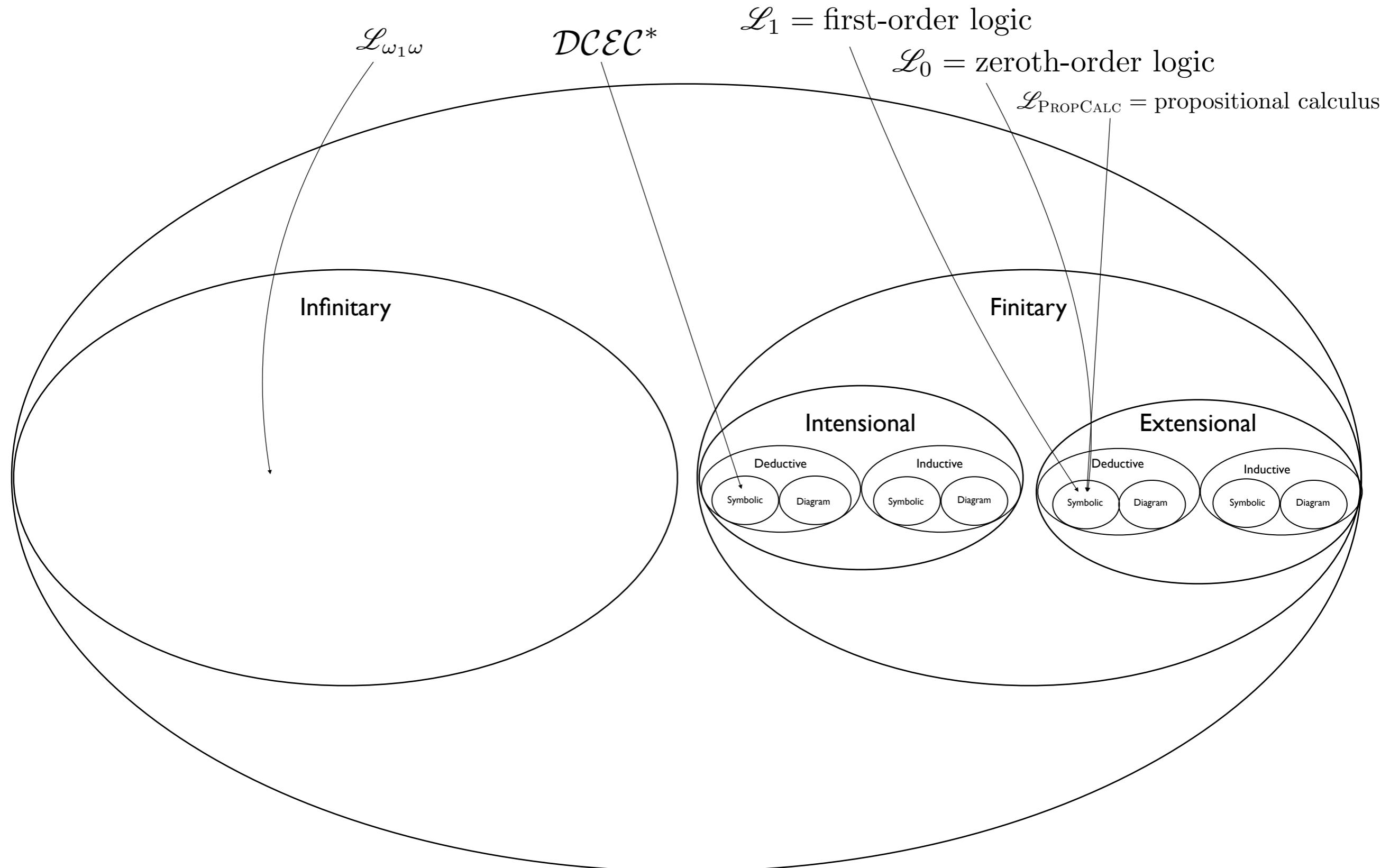
\mathcal{L}_2 = second-order logic \mathcal{L}_3 = third-order logic

The Universe of Logics



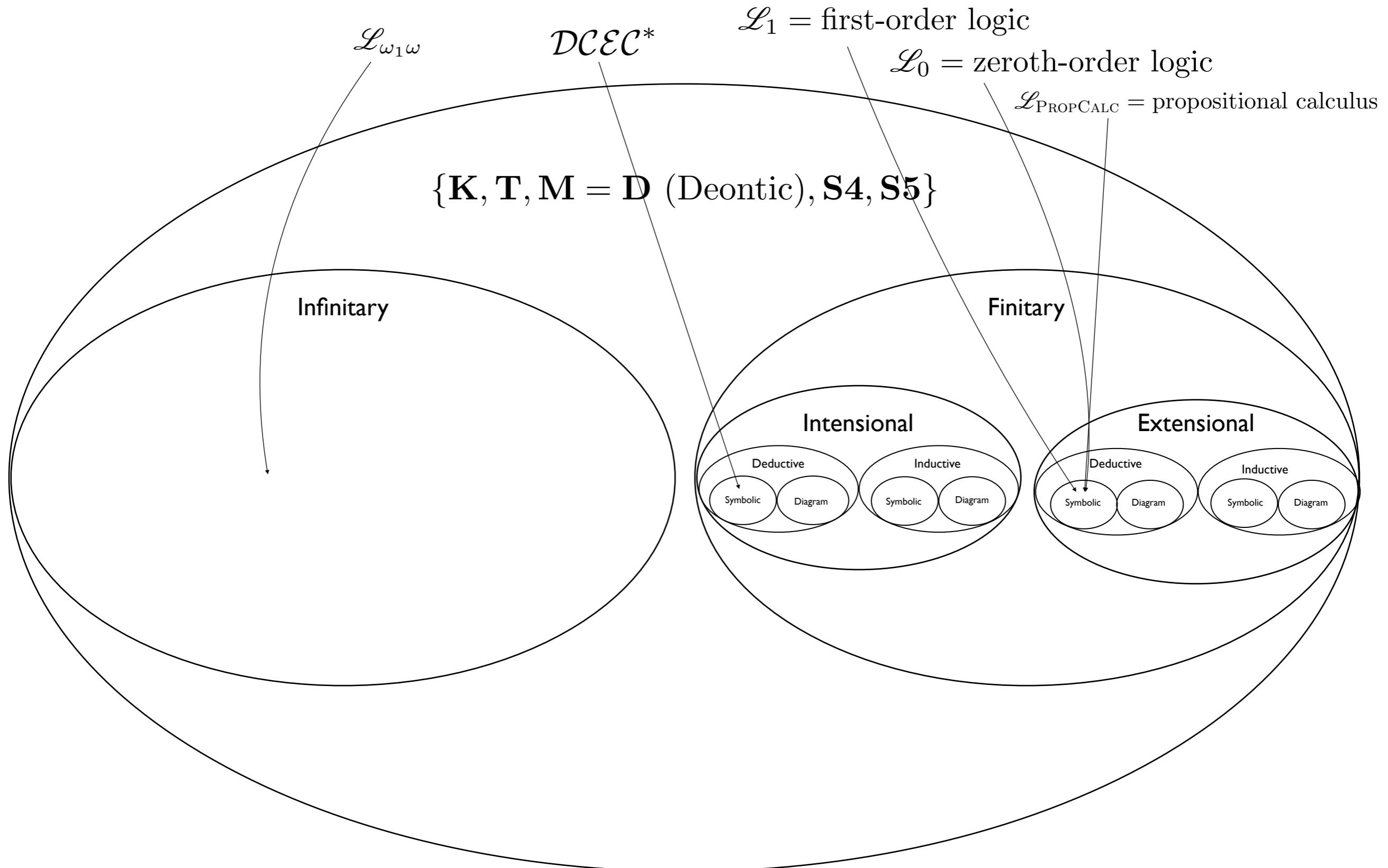
$\{\mathbf{K}, \mathbf{T}, \mathbf{M} = \mathbf{D}$ (Deontic), $\mathbf{S4}, \mathbf{S5}\}$ \mathcal{L}_2 = second-order logic \mathcal{L}_3 = third-order logic

The Universe of Logics



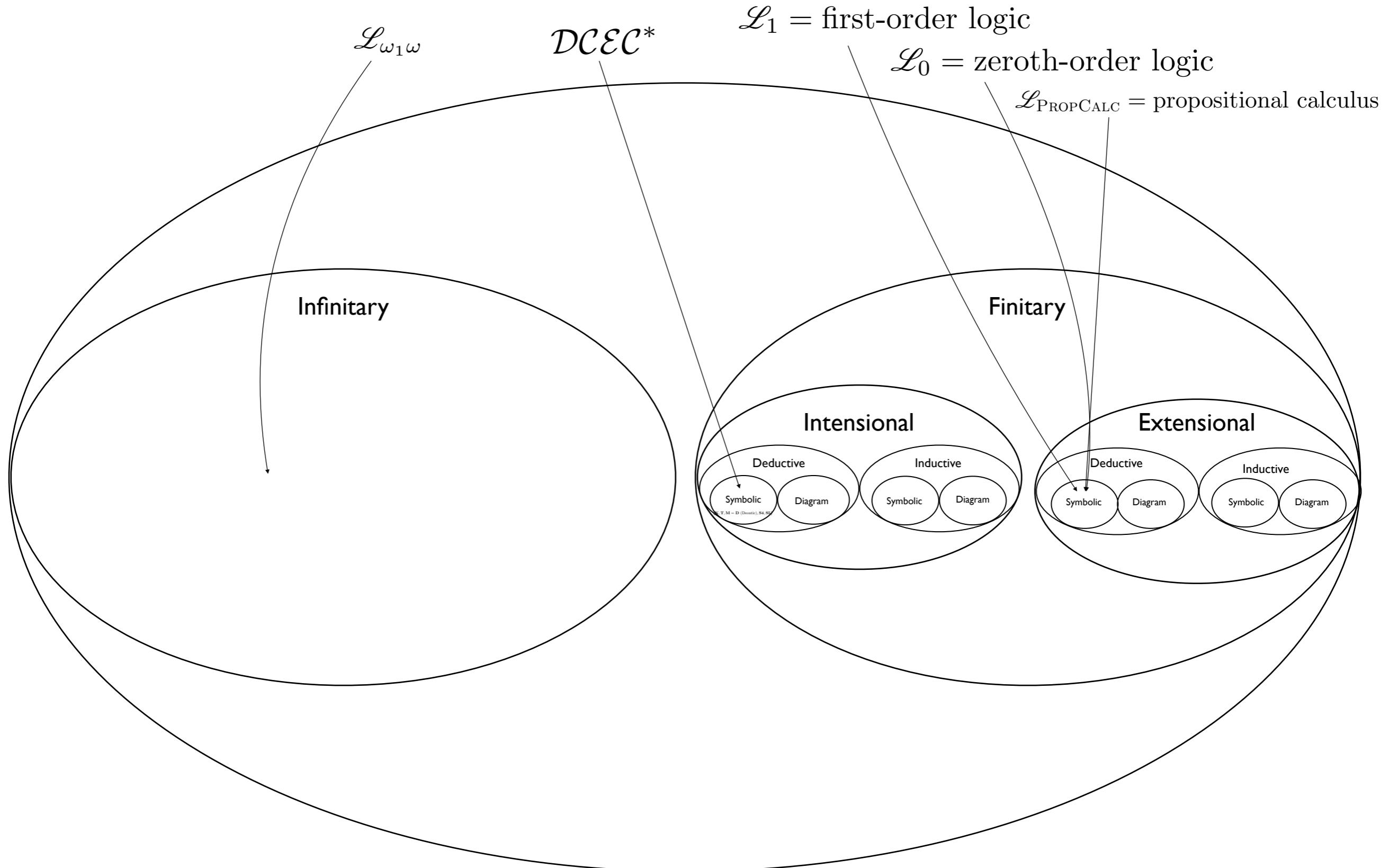
\mathcal{L}_2 = second-order logic \mathcal{L}_3 = third-order logic

The Universe of Logics

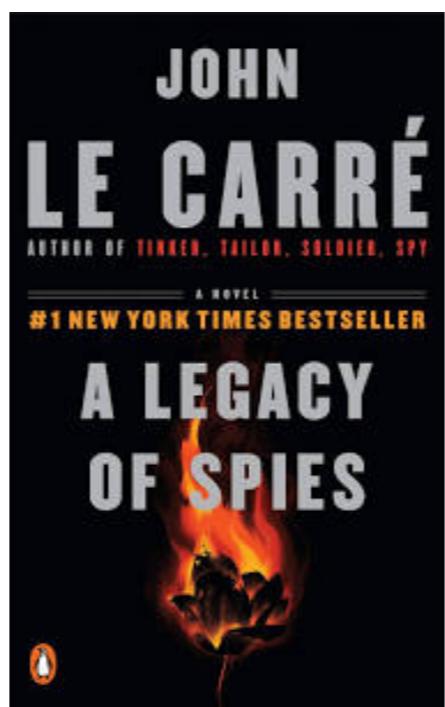


\mathcal{L}_2 = second-order logic \mathcal{L}_3 = third-order logic

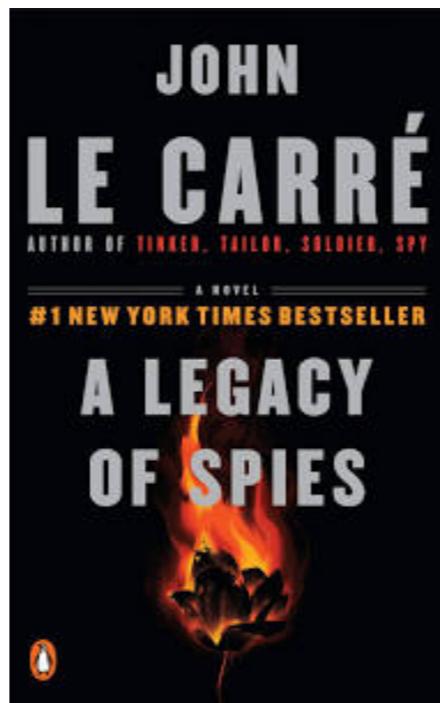
The Universe of Logics



Uh oh

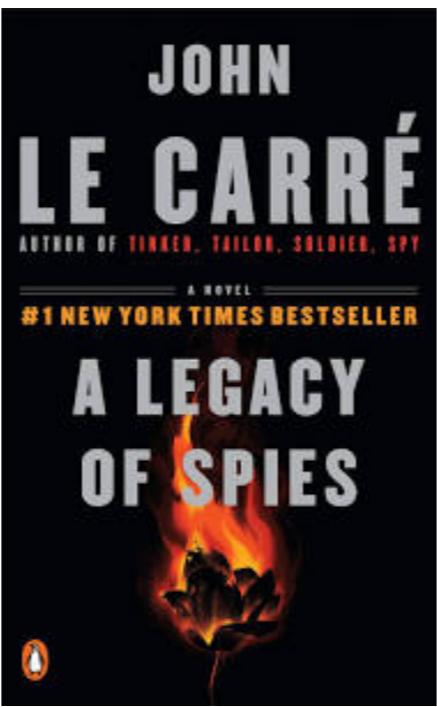


Uh oh



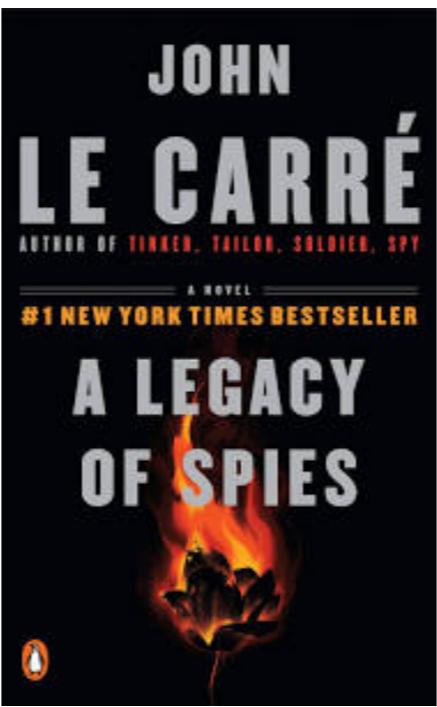
- | | | |
|-------|--|------|
| (1) | $lecarre = cornwall$ | fact |
| (2) | $authOf(legacyofspies) = lecarre$ | fact |
| (3) | $Knows(alvin, authOf(legacyofspies) = lecarre)$ | fact |
| ∴ (4) | $Knows(alvin, authOf(legacyofspies) = cornwall)$ | fact |

Uh oh



- | | | |
|--------------|--|------|
| (1) | $lecarre = cornwall$ | fact |
| (2) | $authOf(legacyofspies) = lecarre$ | fact |
| (3) | $Knows(alvin, authOf(legacyofspies)) = lecarre$ | fact |
| \therefore | $Knows(alvin, authOf(legacyofspies)) = cornwall$ | fact |

Uh oh



- | | | |
|--------------|--|------|
| (1) | $lecarre = cornwall$ | fact |
| (2) | $authOf(legacyofspies) = lecarre$ | fact |
| (3) | $Knows(alvin, authOf(legacyofspies)) = lecarre$ | fact |
| \therefore | $Knows(alvin, authOf(legacyofspies)) = cornwall$ | fact |
- ??

Uh oh



- | | | |
|-------|---|------|
| (1) | $lecarre = \text{cornwall}$ | fact |
| (2) | $\text{authOf}(\text{legacyofspies}) = lecarre$ | fact |
| (3) | $\text{Knows}(\text{alvin}, \text{authOf}(\text{legacyofspies})) = lecarre$ | fact |
| ∴ (4) | $\text{Knows}(\text{alvin}, \text{authOf}(\text{legacyofspies})) = \text{cornwall}$ | fact |

This is only the *first* wave of some *really* bad news that befalls those in AI (& CogSci) who might think that extensional logic is going to work in the realm of human-level cognition ...

Given the Web, What Is Intelligence, Really?

Article (PDF Available) in [Metaphilosophy](#) 43(4) · July 2012 with 23 Reads ⓘ

DOI: 10.1111/j.1467-9973.2012.01760.x

 [Cite this publication](#)



Selmer Bringsjord



Naveen Sundar Govindarajulu

18.4

Abstract

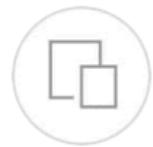
This article argues that existing systems on the Web cannot approach human-level intelligence, as envisioned by Descartes, without being able to achieve genuine problem solving on unseen problems. The article argues that this entails committing to a strong intensional logic. In addition to revising extant arguments in favor of intensional systems, it presents a novel mathematical argument to show why extensional systems can never hope to capture the inherent complexity of natural language. The argument makes its case by focusing on representing, with increasing degrees of complexity, knowledge in a first-order language. Nevertheless, the attempts at representation fail to achieve consistency, making the case for an intensional representation system for natural language clear.

Given the Web, What Is Intelligence, Really?

Article (PDF Available) in [Metaphilosophy](#) 43(4) · July 2012 with 23 Reads ⓘ

DOI: 10.1111/j.1467-9973.2012.01760.x

 [Cite this publication](#)



Selmer Bringsjord



Naveen Sundar Govindarajulu

18.4

Abstract

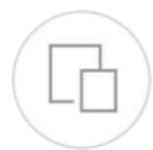
This article argues that existing systems on the Web cannot approach human-level intelligence, as envisioned by Descartes, without being able to achieve genuine problem solving on unseen problems. The article argues that this entails committing to a strong intensional logic. In addition to revising extant arguments in favor of intensional systems, it presents a novel mathematical argument to show why extensional systems can never hope to capture the inherent complexity of natural language. The argument makes its case by focusing on representing, with increasing degrees of complexity, knowledge in a first-order language. Nevertheless, the attempts at representation fail to achieve consistency, making the case for an intensional representation system for natural language clear.

Given the Web, What Is Intelligence, Really?

Article (PDF Available) in [Metaphilosophy](#) 43(4) · July 2012 with 23 Reads ⓘ

DOI: 10.1111/j.1467-9973.2012.01760.x

[↓ Cite this publication](#)



Selmer Bringsjord



Naveen Sundar Govindarajulu

18.4

Abstract

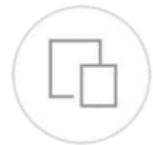
This article argues that existing systems on the Web cannot approach human-level intelligence, as envisioned by Descartes, without being able to achieve genuine problem solving on unseen problems. The article argues that this entails committing to a strong intensional logic. In addition to revising extant arguments in favor of intensional systems, it presents a novel mathematical argument to show why extensional systems can never hope to capture the inherent complexity of natural language. The argument makes its case by focusing on representing, with increasing degrees of complexity, knowledge in a first-order language. Nevertheless, the attempts at representation fail to achieve consistency, making the case for an intensional representation system for natural language clear.

Given the Web, What Is Intelligence, Really?

Article (PDF Available) in [Metaphilosophy](#) 43(4) · July 2012 with 23 Reads ⓘ

DOI: 10.1111/j.1467-9973.2012.01760.x

[↓ Cite this publication](#)



Selmer Bringsjord



Naveen Sundar Govindarajulu

18.4

Abstract

This article argues that existing systems on the Web cannot approach human-level intelligence, as envisioned by Descartes, without being able to achieve genuine problem solving on unseen problems. The article argues that this entails committing to a strong intensional logic. In addition to revising extant arguments in favor of intensional systems, it presents a novel mathematical argument to show why extensional systems can never hope to capture the inherent complexity of natural language. The argument makes its case by focusing on representing, with increasing degrees of complexity, knowledge in a first-order language. Nevertheless, the attempts at representation fail to achieve consistency, making the case for an intensional representation system for natural language clear.

Let's investigate the first theorem in this paper, in HyperSlate™ ...

Modal Logic to Save Lives

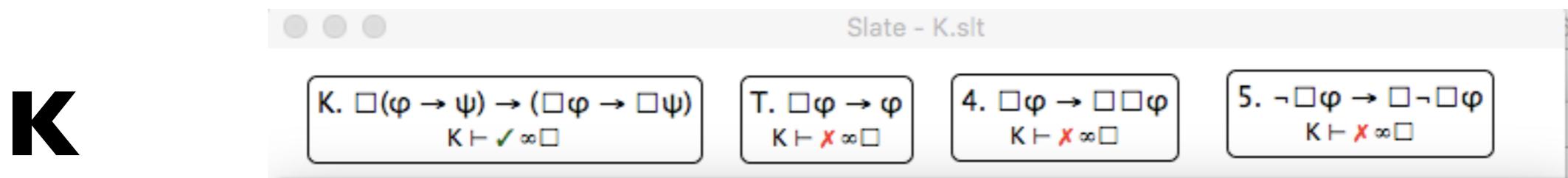
https://en.wikipedia.org/wiki/Air_France_Flight_447

Modal Logic to Save Lives

https://en.wikipedia.org/wiki/Air_France_Flight_447

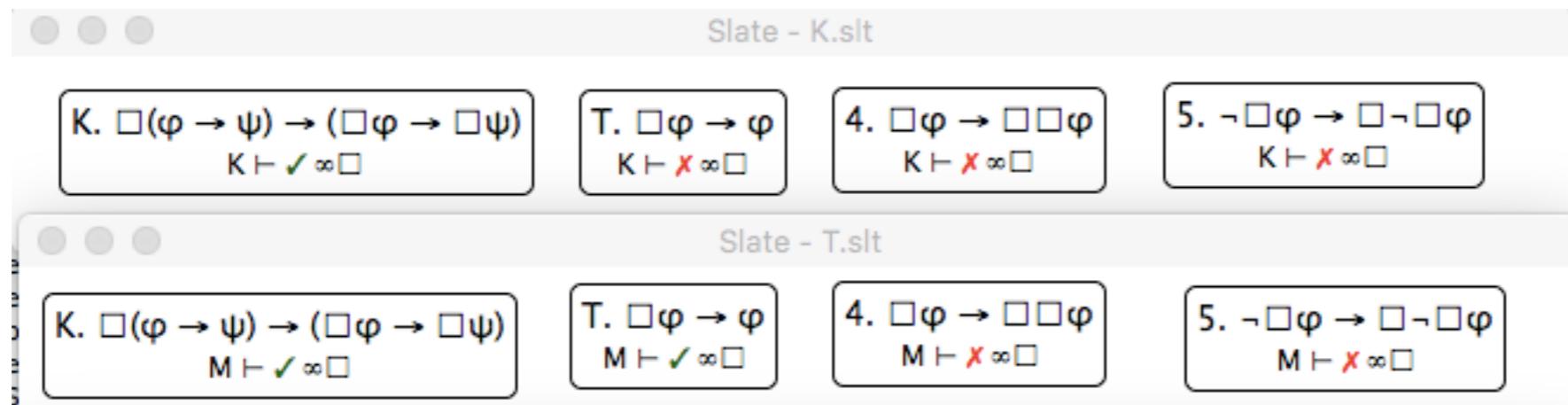
..	(1)	Knows (<i>ai</i> , Knows (<i>pilot2</i> , <i>nosepulledup</i>))	fact
..	(2)	Knows (<i>pilot2</i> , <i>nosepulledup</i>)	from (1)
..	(3)	<i>nosepulledup</i>	from (2) in S4
	(4)	<i>nosepulledup</i> → <i>aiooverride!</i>	premise
	(5)	<i>aiooverride!</i>	from (3), (4) \mathcal{L}_{PC}

Encapsulation to Memorize



Encapsulation to Memorize

K
T



Encapsulation to Memorize

K

T

D

The screenshot shows the Slate application interface with three tabs: **Slate - K.slt**, **Slate - T.slt**, and **Slate - D.slt**. Each tab contains five boxes, each representing a logical axiom:

- Slate - K.slt:**
 - K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
K $\vdash \checkmark \infty \Box$
 - T. $\Box\varphi \rightarrow \varphi$
K $\vdash \textcolor{red}{X} \infty \Box$
 - 4. $\Box\varphi \rightarrow \Box\Box\varphi$
K $\vdash \textcolor{red}{X} \infty \Box$
 - 5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
K $\vdash \textcolor{red}{X} \infty \Box$
- Slate - T.slt:**
 - K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
M $\vdash \checkmark \infty \Box$
 - T. $\Box\varphi \rightarrow \varphi$
M $\vdash \checkmark \infty \Box$
 - 4. $\Box\varphi \rightarrow \Box\Box\varphi$
M $\vdash \textcolor{red}{X} \infty \Box$
 - 5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
M $\vdash \textcolor{red}{X} \infty \Box$
- Slate - D.slt:**
 - K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
D $\vdash \checkmark \infty \Box$
 - T. $\Box\varphi \rightarrow \varphi$
D $\vdash \textcolor{red}{X} \infty \Box$
 - D. $\Box\varphi \rightarrow \Diamond\varphi$
D $\vdash \checkmark \infty \Box$
 - 4. $\Box\varphi \rightarrow \Box\Box\varphi$
D $\vdash \textcolor{red}{X} \infty \Box$
 - 5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
D $\vdash \textcolor{red}{X} \infty \Box$
 - INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
D $\vdash \checkmark \infty \Box$

Encapsulation to Memorize

K

T

D

4 = S4

Slate - K.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
K $\vdash \checkmark \infty \Box$

T. $\Box\varphi \rightarrow \varphi$
K $\vdash \text{X} \infty \Box$

4. $\Box\varphi \rightarrow \Box\Box\varphi$
K $\vdash \text{X} \infty \Box$

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
K $\vdash \text{X} \infty \Box$

Slate - T.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
M $\vdash \checkmark \infty \Box$

T. $\Box\varphi \rightarrow \varphi$
M $\vdash \checkmark \infty \Box$

4. $\Box\varphi \rightarrow \Box\Box\varphi$
M $\vdash \text{X} \infty \Box$

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
M $\vdash \text{X} \infty \Box$

Slate - D.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
D $\vdash \checkmark \infty \Box$

T. $\Box\varphi \rightarrow \varphi$
D $\vdash \text{X} \infty \Box$

D. $\Box\varphi \rightarrow \Diamond\varphi$
D $\vdash \checkmark \infty \Box$

4. $\Box\varphi \rightarrow \Box\Box\varphi$
D $\vdash \text{X} \infty \Box$

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
D $\vdash \text{X} \infty \Box$

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
D $\vdash \checkmark \infty \Box$

Slate - S4.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
S4 $\vdash \checkmark \infty \Box$

T. $\Box\varphi \rightarrow \varphi$
S4 $\vdash \checkmark \infty \Box$

D. $\Box\varphi \rightarrow \Diamond\varphi$
S4 $\vdash \checkmark \infty \Box$

4. $\Box\varphi \rightarrow \Box\Box\varphi$
S4 $\vdash \checkmark \infty \Box$

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
S4 $\vdash \text{X} \infty \Box$

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
{INTER} Assume \checkmark

Encapsulation to Memorize

K

T

D

4 = S4

5 = S5

Slate - K.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
K ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
K ⊢ X ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
K ⊢ X ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
K ⊢ X ∞□

Slate - T.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
M ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
M ⊢ ✓ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
M ⊢ X ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
M ⊢ X ∞□

Slate - D.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
D ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
D ⊢ X ∞□

D. $\Box\varphi \rightarrow \Diamond\varphi$
D ⊢ ✓ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
D ⊢ X ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
D ⊢ X ∞□

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
D ⊢ ✓ ∞□

Slate - S4.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
S4 ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
S4 ⊢ ✓ ∞□

D. $\Box\varphi \rightarrow \Diamond\varphi$
S4 ⊢ ✓ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
S4 ⊢ ✓ ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
S4 ⊢ X ∞□

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
{INTER} Assume ✓

Slate - S5.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
S5 ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
S5 ⊢ ✓ ∞□

D. $\Box\varphi \rightarrow \Diamond\varphi$
{D} Assume ✓

4. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
{4} Assume ✓

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
S5 ⊢ ✓ ∞□

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
{INTER} Assume ✓

Encapsulation to Memorize

K

T

D

4 = S4

5 = S5

Slate - K.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
K ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
K ⊢ ✗ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
K ⊢ ✗ ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
K ⊢ ✗ ∞□

Slate - T.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
M ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
M ⊢ ✓ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
M ⊢ ✗ ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
M ⊢ ✗ ∞□

Slate - D.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
D ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
D ⊢ ✗ ∞□

D. $\Box\varphi \rightarrow \Diamond\varphi$
D ⊢ ✓ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
D ⊢ ✗ ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
D ⊢ ✗ ∞□

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
D ⊢ ✓ ∞□

Slate - S4.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
S4 ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
S4 ⊢ ✓ ∞□

D. $\Box\varphi \rightarrow \Diamond\varphi$
S4 ⊢ ✓ ∞□

4. $\Box\varphi \rightarrow \Box\Box\varphi$
S4 ⊢ ✓ ∞□

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
S4 ⊢ ✗ ∞□

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
{INTER} Assume ✓

Slate - S5.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
S5 ⊢ ✓ ∞□

T. $\Box\varphi \rightarrow \varphi$
S5 ⊢ ✓ ∞□

D. $\Box\varphi \rightarrow \Diamond\varphi$
{D} Assume ✓

4. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$
{4} Assume ✓

5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$
S5 ⊢ ✓ ∞□

INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$
{INTER} Assume ✓

Encapsulation to Memorize

K**T****D****4 = S4****5 = S5**

Slate - K.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$ K $\vdash \checkmark \infty \Box$	T. $\Box\varphi \rightarrow \varphi$ K $\vdash \times \infty \Box$	4. $\Box\varphi \rightarrow \Box\Box\varphi$ K $\vdash \times \infty \Box$	5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$ K $\vdash \times \infty \Box$
---	---	---	---

Slate - T.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$ M $\vdash \checkmark \infty \Box$	T. $\Box\varphi \rightarrow \varphi$ M $\vdash \checkmark \infty \Box$	4. $\Box\varphi \rightarrow \Box\Box\varphi$ M $\vdash \times \infty \Box$	5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$ M $\vdash \times \infty \Box$
---	---	---	---

Slate - D.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$ D $\vdash \checkmark \infty \Box$	T. $\Box\varphi \rightarrow \varphi$ D $\vdash \times \infty \Box$	D. $\Box\varphi \rightarrow \Diamond\varphi$ D $\vdash \checkmark \infty \Box$	4. $\Box\varphi \rightarrow \Box\Box\varphi$ D $\vdash \times \infty \Box$
5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$ D $\vdash \times \infty \Box$	INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$ D $\vdash \checkmark \infty \Box$		

Slate - S4.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$ S4 $\vdash \checkmark \infty \Box$	T. $\Box\varphi \rightarrow \varphi$ S4 $\vdash \checkmark \infty \Box$	D. $\Box\varphi \rightarrow \Diamond\varphi$ S4 $\vdash \checkmark \infty \Box$	4. $\Box\varphi \rightarrow \Box\Box\varphi$ S4 $\vdash \checkmark \infty \Box$
5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$ S4 $\vdash \times \infty \Box$	INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$ {INTER} Assume \checkmark		

Slate - S5.slt

K. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$ S5 $\vdash \checkmark \infty \Box$	T. $\Box\varphi \rightarrow \varphi$ S5 $\vdash \checkmark \infty \Box$	D. $\Box\varphi \rightarrow \Diamond\varphi$ {D} Assume \checkmark	4. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$ {4} Assume \checkmark
5. $\neg\Box\varphi \rightarrow \Box\neg\Box\varphi$ S5 $\vdash \checkmark \infty \Box$	INTER. $\Box\varphi \leftrightarrow \neg\Diamond\neg\varphi$ {INTER} Assume \checkmark		