

Introducing *Pure General Logic Programming* (PGLP), in HyperSlate[®]:HyperLog[®]; Review of All Inference Rules/ Schemata in PropCalc = \mathcal{L}_{PC}

Selmer Bringsjord

Rensselaer AI & Reasoning (RAIR) Lab
Department of Cognitive Science
Department of Computer Science
Lally School of Management & Technology
Rensselaer Polytechnic Institute (RPI)
Troy, New York 12180 USA

IFLAI
2/6/2023



Logistics again ...

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Your code for starting the registration process is:

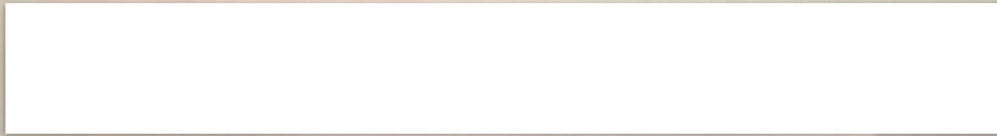
To access HyperGrader, HyperSlate, the license agreement,
and to obtain the textbook LAMA-BDLA, go to::

<https://rpi.logicamodernapproach.com>

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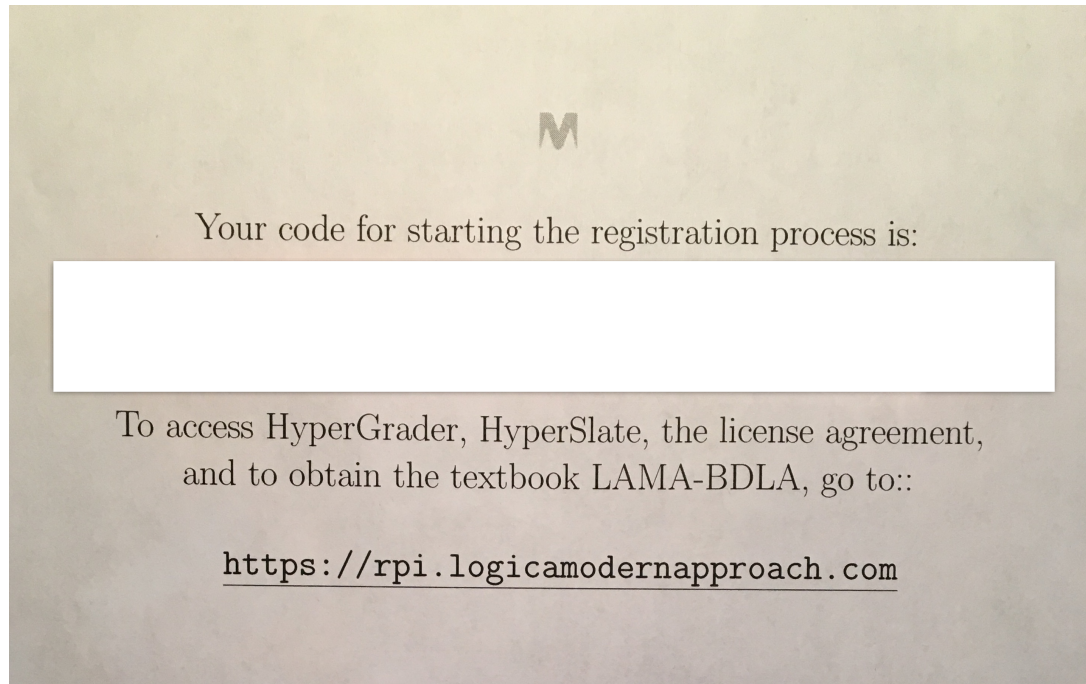


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paradigm!

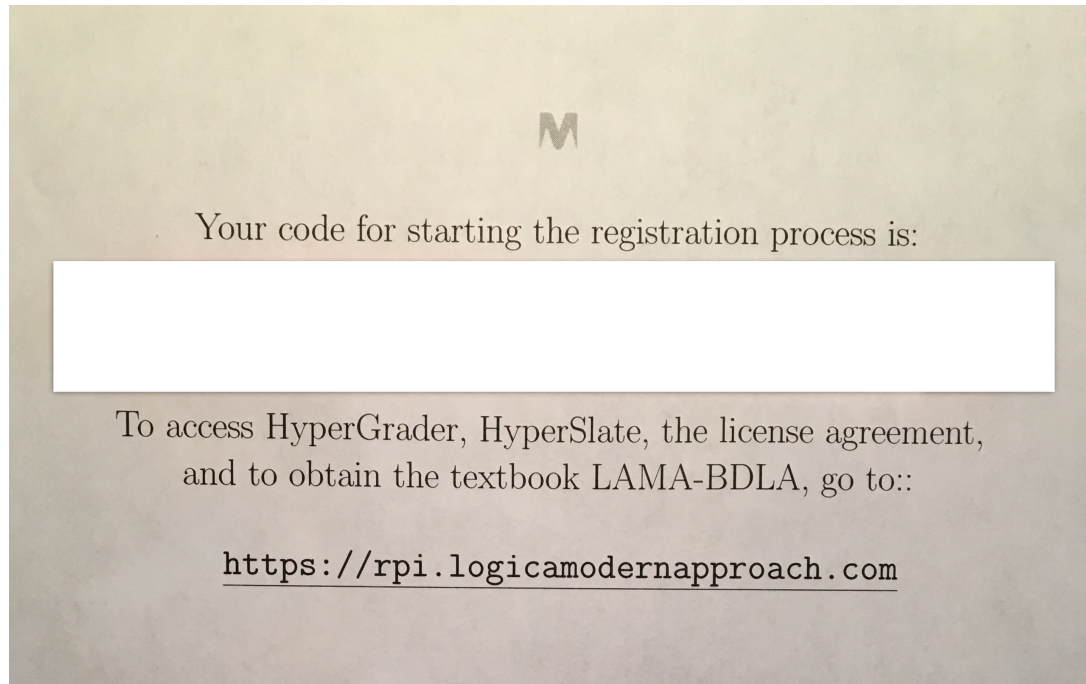
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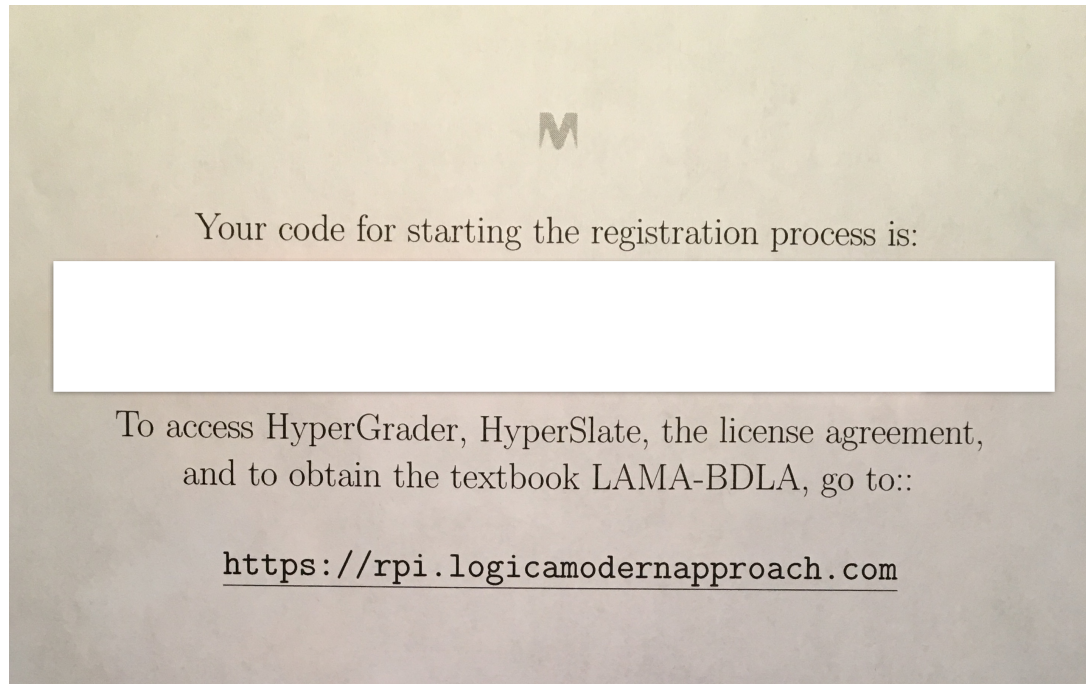


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Watch that the link emailed to you doesn't end up being classified as spam.

Open Office Hours Mon Thu; Today:

Selmer Bringsjord is inviting you to a scheduled Zoom meeting.

Topic: Selmer Bringsjord's Zoom Meeting

Time: Feb 6, 2023 04:00 PM Eastern Time (US and Canada)

Join Zoom Meeting

<https://us02web.zoom.us/j/89580559014?pwd=ZnYzUTBReEdZZnFlQ0UzbDBBei96UT09>

Meeting ID: 895 8055 9014

Passcode: 961547

One tap mobile

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+19292056099,,89580559014#,,,,*961547# US (New York)

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+1 929 205 6099 US (New York)

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+1 312 626 6799 US (Chicago)

+1 301 715 8592 US (Washington DC)

+1 305 224 1968 US

+1 719 359 4580 US

+1 253 205 0468 US

+1 253 215 8782 US (Tacoma)

+1 346 248 7799 US (Houston)

+1 360 209 5623 US

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+1 507 473 4847 US

+1 564 217 2000 US

+1 669 444 9171 US

+1 669 900 6833 US (San Jose)

+1 689 278 1000 US

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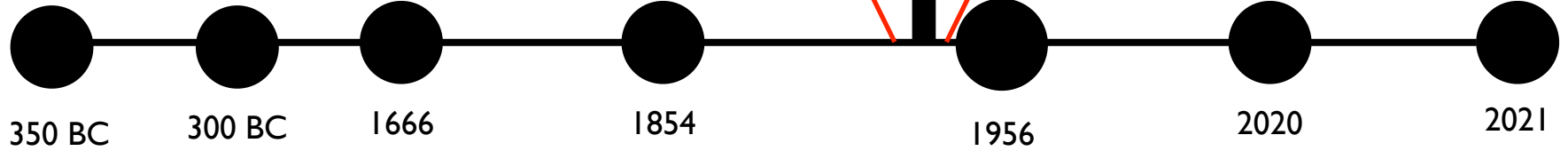
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Entscheidungsproblem

“Universal
Computational Logic”



Logic Theorist
(birth of modern logicist AI)



350 BC

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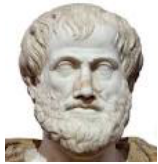
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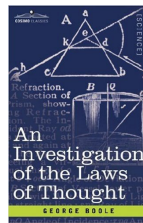


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Simon

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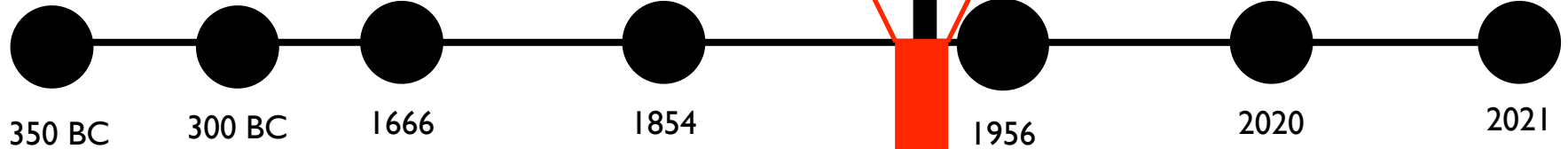
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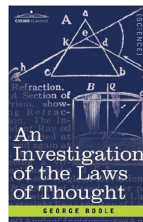


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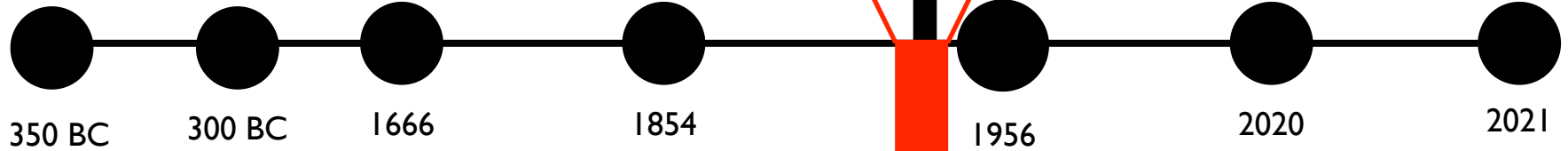
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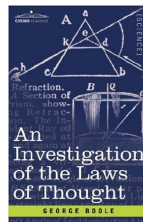


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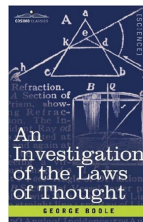


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Exceeds Leibniz & de-mystifies
Euclid: the “compellingness” of
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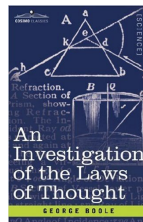


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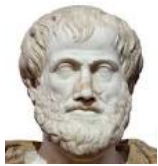
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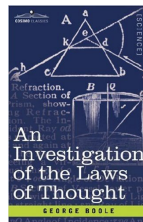


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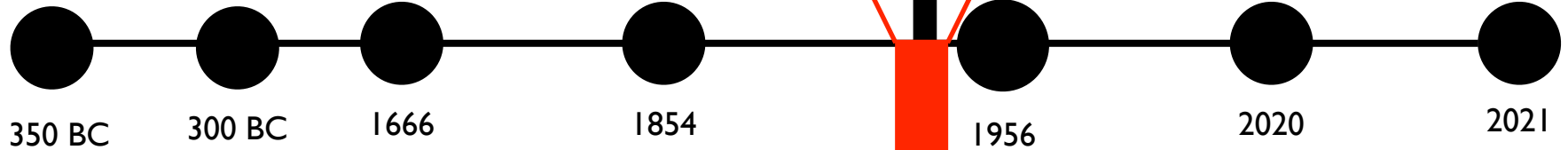
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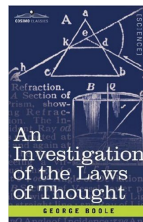
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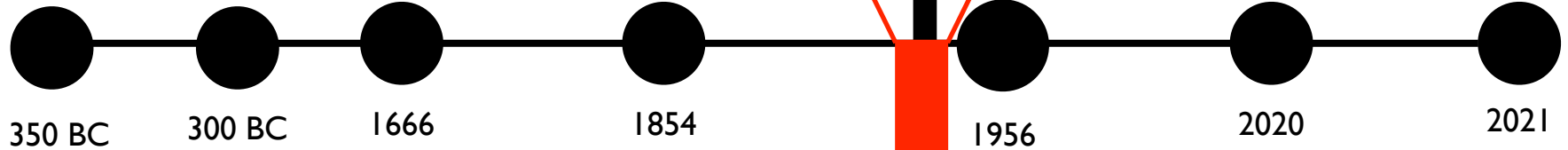
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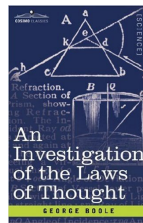


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Here's what a computer is, and given that, sorry, the *Entscheidungsproblem* can't be solved by such a machine!

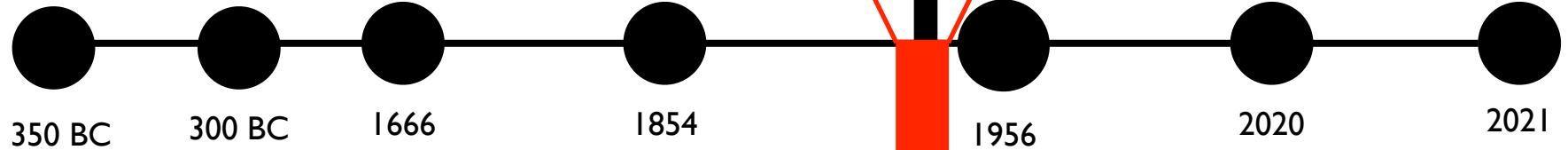
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Functional = Church;
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Where is logic-based/logicist
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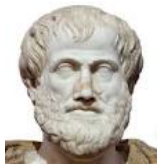
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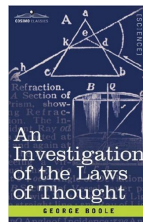


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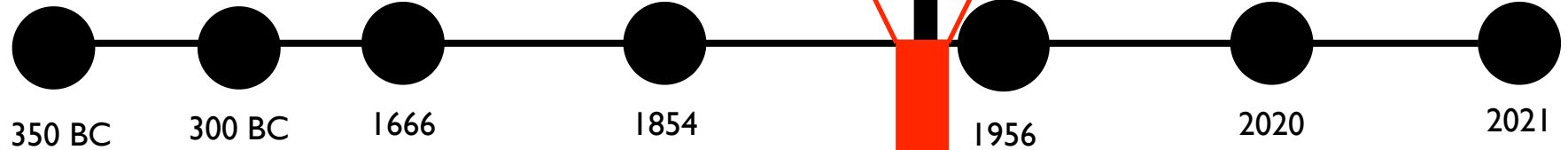
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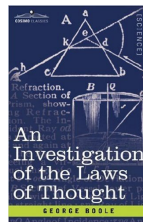


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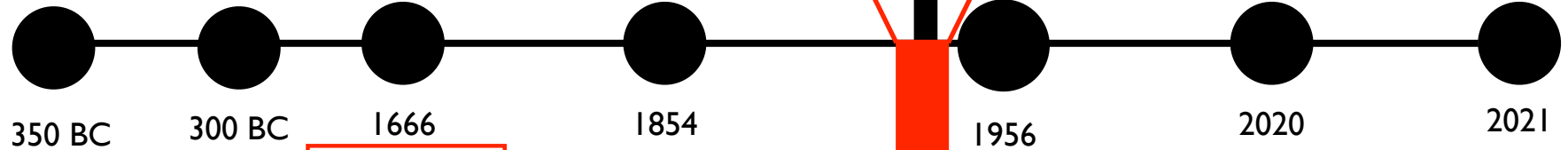
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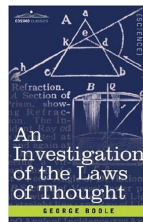


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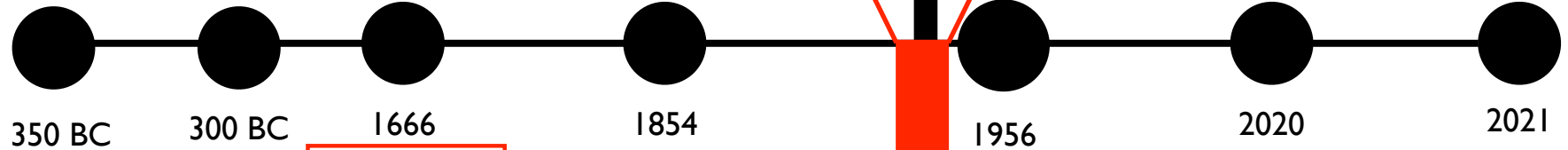
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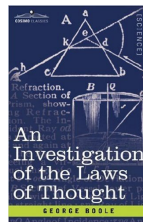


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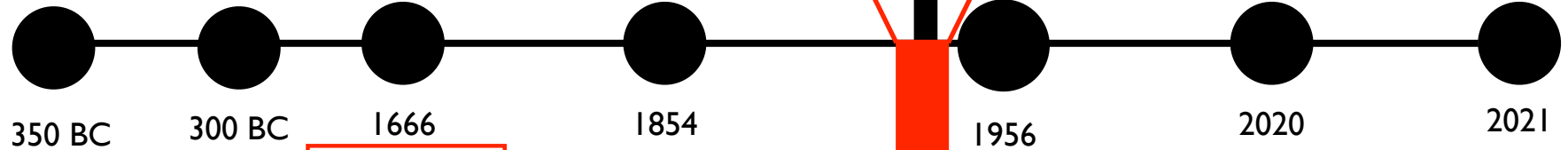
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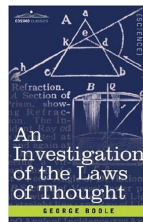


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Programming Languages

COURSE HOME

SYLLABUS

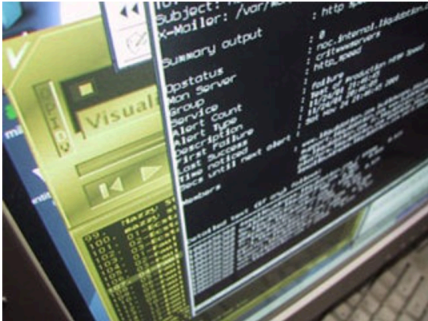
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Programming computer screen. (Photo courtesy of openphoto.net.)

Instructor(s)

Prof. Michael Ernst

MIT Course Number

6.821

As Taught In

Fall 2002

Level

Graduate

CITE THIS COURSE

Course Features

› Assignments: programming (no examples) › Exams (no solutions)

Course Description

6.821 teaches the principles of functional, imperative, and logic programming languages. Topics covered include: meta-circular interpreters, semantics (operational and denotational), type systems (polymorphism, inference, and abstract types), object oriented programming, modules, and multiprocessing. The course involves substantial programming assignments and problem sets as well as a significant amount of reading. The course uses the Scheme+ programming language for all of its assignments.

SYLLABUS

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Syllabus

Programming Languages CSCI-4430

Meetings: Webex, TF 2:30-4:20pm
Website: <http://www.cs.rpi.edu/~milanova/csci4430>

I. Brief Course Description

This course is a study of important concepts in programming languages. Topics include programming language syntax and semantics, types and parameter passing, and programming paradigms (logic-oriented, functional, von Neumann, object-oriented).

Prerequisite: Introduction to Algorithms (CSCI 2300) and Principles of Software (CSCI 2600)

Mailing list: proglang@cs.lists.rpi.edu. Email goes to Milanova, Kuzmin, and Hulbert. Use this list for administrative questions, including homework extension requests, quiz and exam makeup requests, extra time scheduling, and so on.

II. Learning Outcomes

The goal of this course is to teach students how to analyze programming languages. Students will become more productive programmers, will be able to learn new programming languages with ease, and will be able to choose the most suitable programming language for a given problem.

Concretely, students who successfully complete the course should be able to 1) explain programming language syntax and semantics, 2) implement a front-end for a programming language, 3) explain the concepts of scoping, data abstraction, types, control abstraction, and parameter passing, which are essential building blocks of programming languages, and 4) demonstrate competence across a spectrum of programming language paradigms by writing programs in Prolog, Scheme, and Haskell.

III. Required Textbook

Programming Language Pragmatics, Fourth Edition, by Michael Scott, Morgan Kaufmann, 2015.

IV. Class Work and Policies

Quizzes

There are 9 quizzes that should be completed and submitted individually. We will drop the lowest quiz grade and only 8 will count towards the final grade. Quizzes will be administered on Submitty at the beginning of our regularly scheduled class time. We will be offering alternative times for quizzes and exams. **If you are unable to attend during regularly scheduled class hours, you must request an alternative time. Email course staff at proglang@cs.lists.rpi.edu by September 10 outlining the reasons why you will be attending at an alternative time (e.g., you reside in a different time zone). We will assign an alternative time and you will be taking the quizzes during this time slot on the date of the quiz. Note that once assigned, you cannot change the quiz time slot.**

Syllabus

Programming Language

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Website: <http://www.cs.rpi.edu/~milanova/csci4430>

I. Brief Course Description

This course is a study of important concepts in programming languages. Topics include programming (logic-oriented, functional, von Neumann, object-oriented).

Prerequisite: Introduction to Algorithms (CSCI 2300) and Principles of Software (CSCI 2600)

Mailing list: proglang@cs.lists.rpi.edu. Email goes to Milanova, Kuzmin, and Hulbert. Use this list for a requests, extra time scheduling, and so on.

II. Learning Outcomes

The goal of this course is to teach students how to analyze programming languages. Students will become more comfortable with programming languages, and will be able to choose the most suitable programming language for a given problem.

Concretely, students who successfully complete the course should be able to 1) explain programming concepts, 2) explain the concepts of scoping, data abstraction, types, control abstraction, and parameter passing, and 3) explain programming concepts across a spectrum of programming language paradigms by writing programs in Prolog, Scheme, and L.

III. Required Textbook

Programming Language Pragmatics, Fourth Edition, by Michael Scott, Morgan Kaufmann, 2015.

IV. Class Work and Policies

There are *Two* Logicist Branches;
BI:

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Frege, 1893:

“Aha! Currying! I recast multiple-arity
operations with functions into a unary affair!”

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Haskell OCaml, Scheme, ...

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Haskell OCaml, Scheme, ...

Athena

Two Logicist Branches; B2:

Two Logician Branches; B2:

The AI Branch: Automated Reasoning

Two Logicist Branches; B2:

The AI Branch: Automated Reasoning

Leibniz

Two Logician Branches; B2:

The AI Branch: Automated Reasoning

Leibniz

**Simon & Newell @
Dawn of Modern AI: LT & GPS**

Two Logicist Branches; B2:

The AI Branch: Automated Reasoning

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Prolog?

Two Logicist Branches; B2:

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PGLP

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...

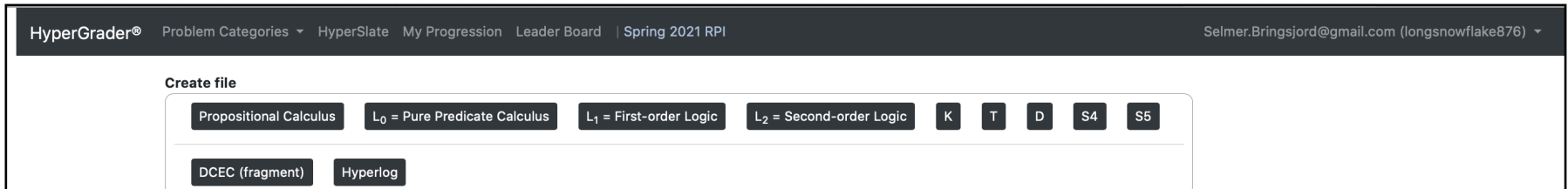
PGLP

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PGLP

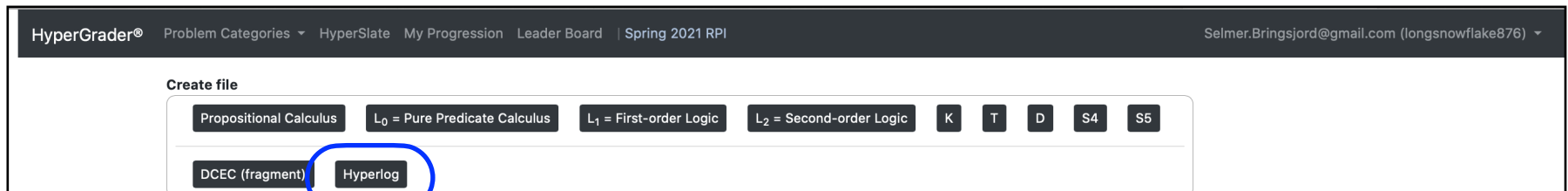
HyperSlate® : HyperLog®

Two Logician Branches; B2:

The AI Branch: Automated Reasoning

Leibniz

**Simon & Newell @
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PGLP

HyperSlate® : HyperLog®

HyperLog: Historico-logico-programming Landscape



First "logic programs"
300 BC



Liebniz
Dies 1716



Frege
1893



Schöenfinkel
1893

Combinatory Logic



Church

λ -calculus

simple type theory



Logic Theorist
(birth of modern logicist AI)



Simon

1956



Turing

Lisp

Prolog

Fortran

Smalltalk

A
t
h
e
n
a

ML

Scheme

CL

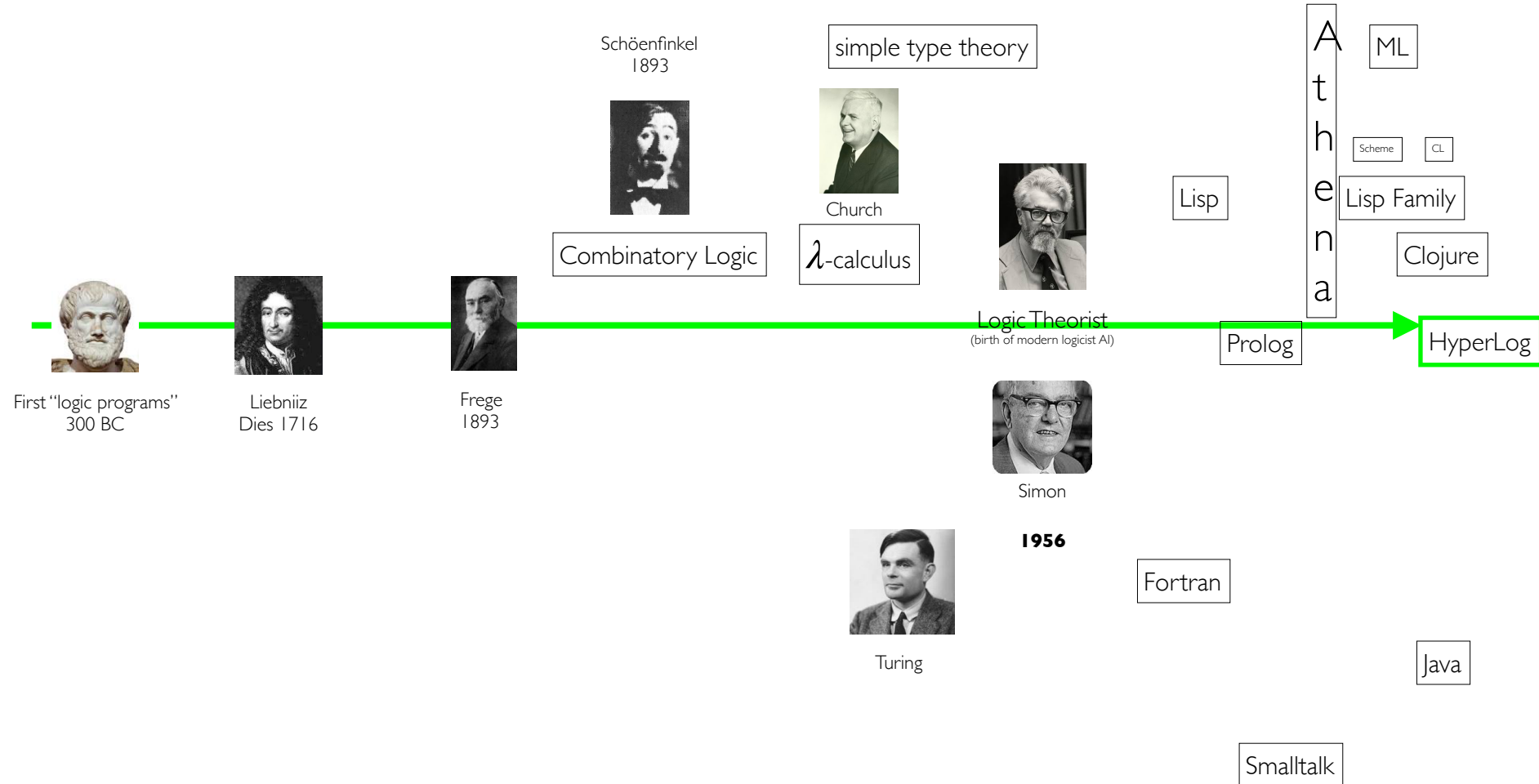
Lisp Family

Clojure

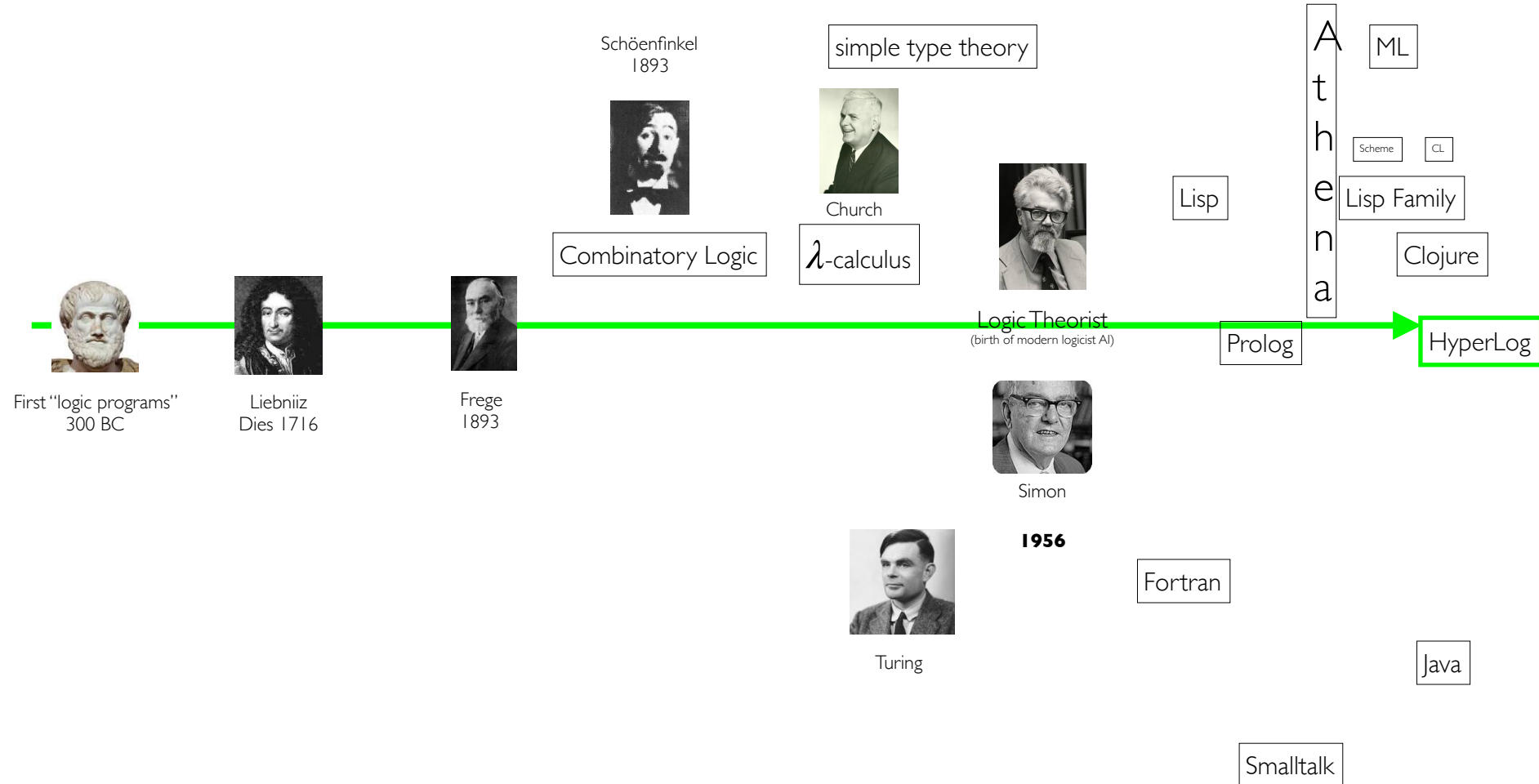
HyperLog

Java

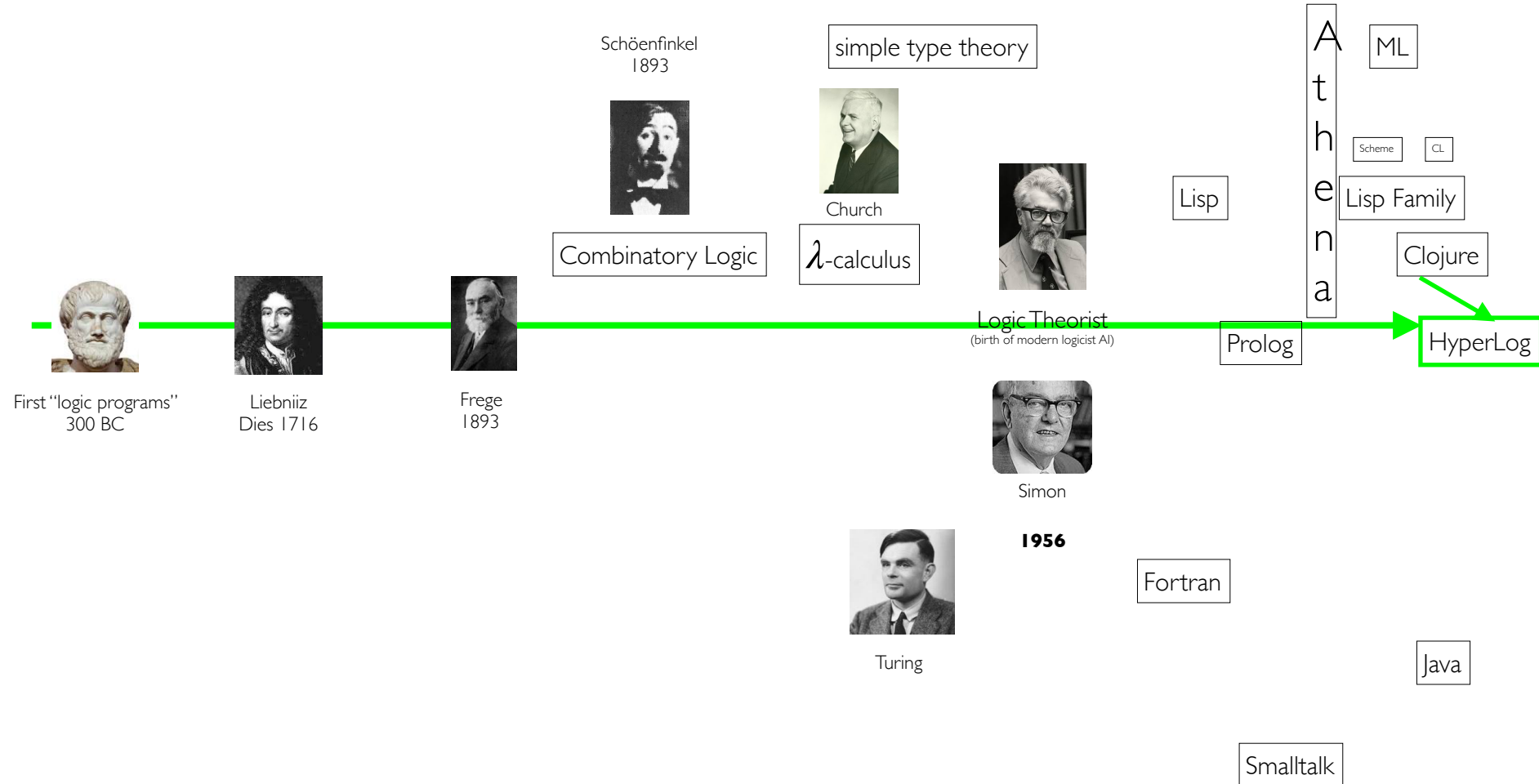
HyperLog: Historico-logico-programming Landscape



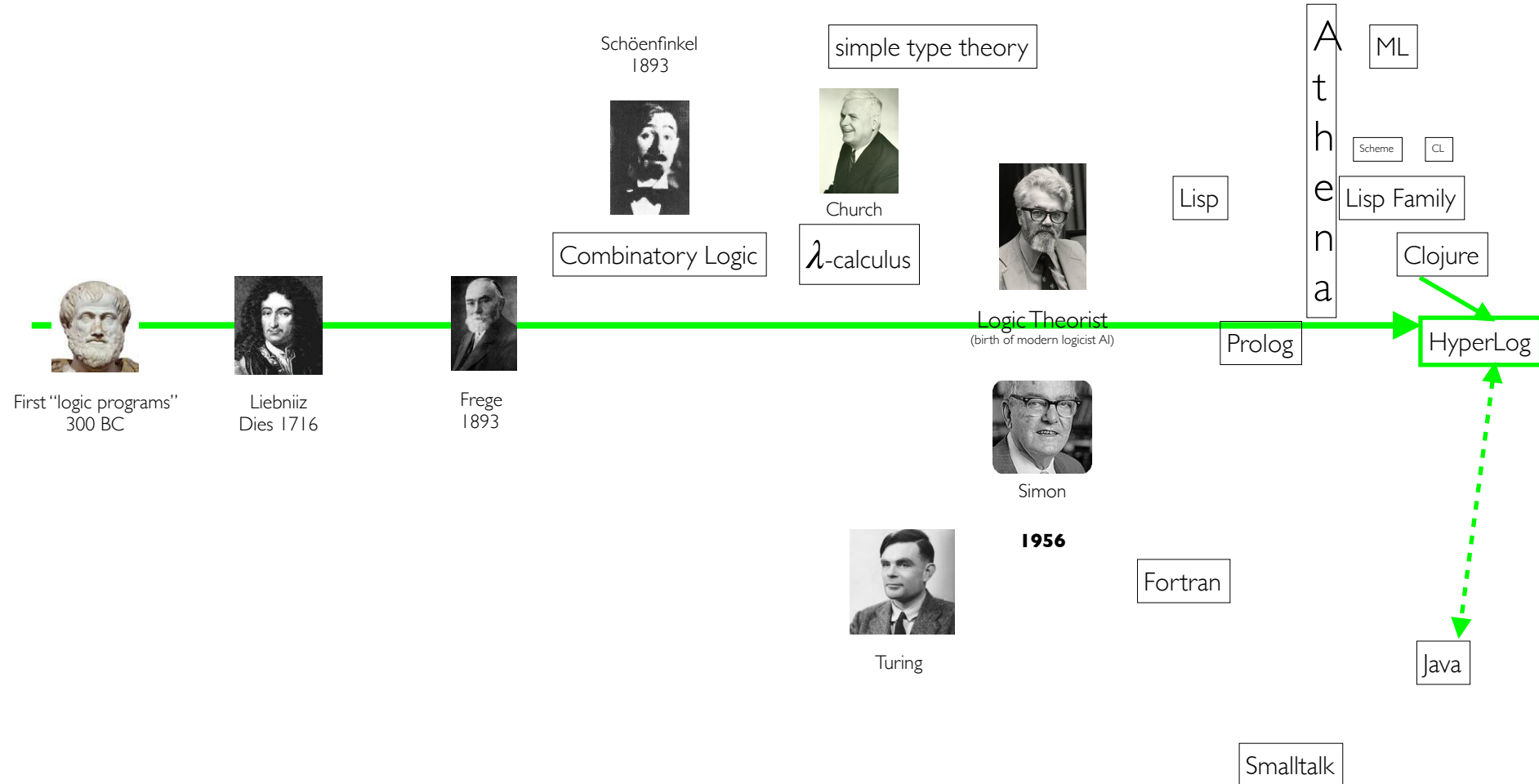
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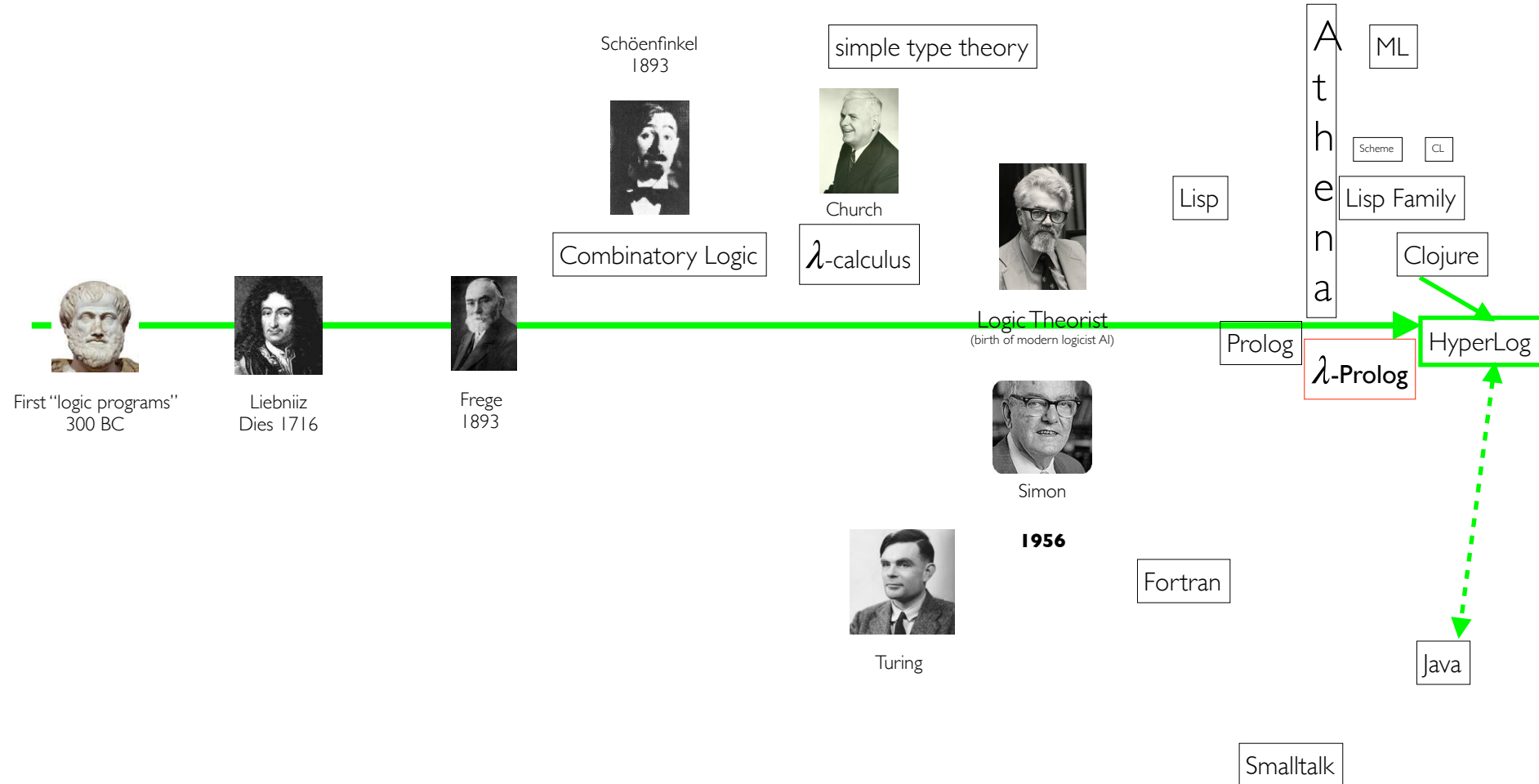
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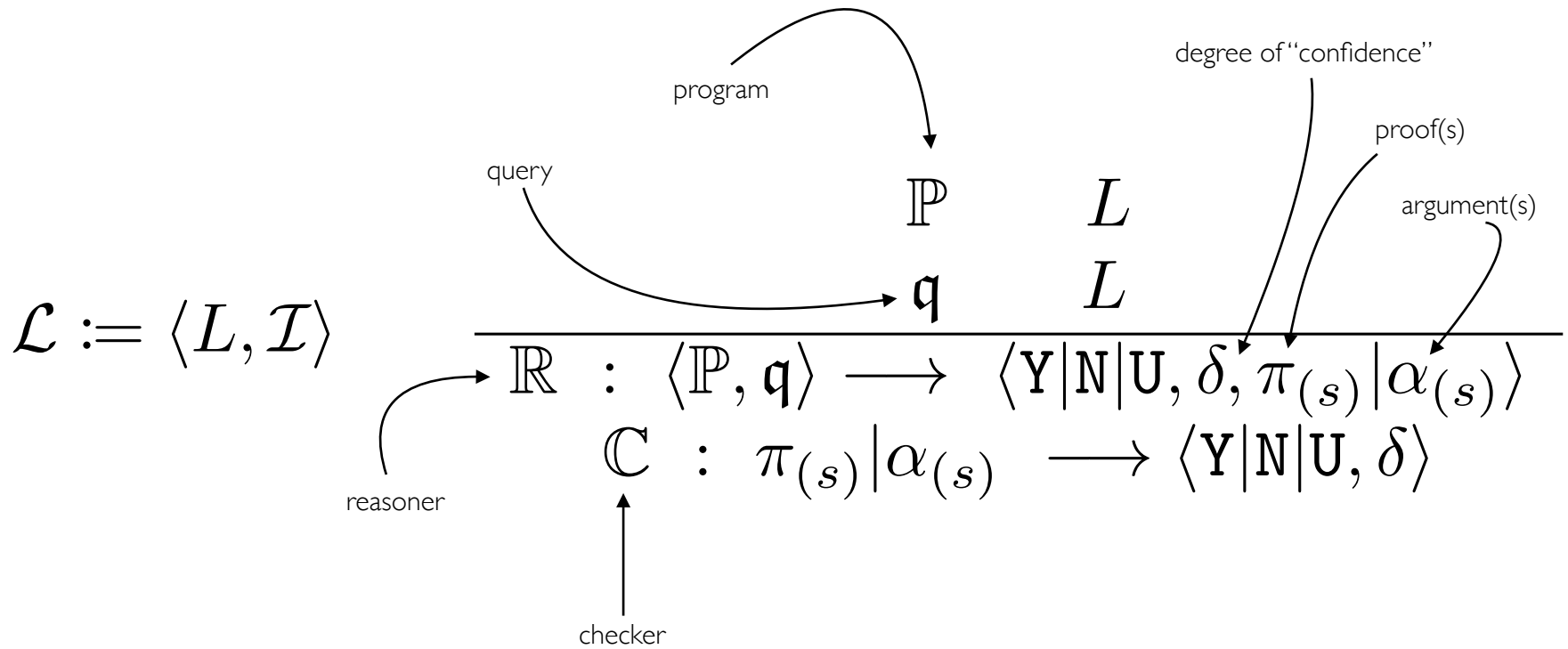
HyperLog: Historico-logico-programming Landscape



Single-Slide Encapsulation ...

$$\mathcal{L} := \langle L, \mathcal{I} \rangle$$

$$\begin{array}{c} \mathbb{P} \quad L \\ \mathfrak{q} \quad L \\ \hline \mathbb{R} : \langle \mathbb{P}, \mathfrak{q} \rangle \longrightarrow \langle \mathbf{Y} | \mathbf{N} | \mathbf{U}, \delta, \pi_{(s)} | \alpha_{(s)} \rangle \\ \mathbb{C} : \pi_{(s)} | \alpha_{(s)} \longrightarrow \langle \mathbf{Y} | \mathbf{N} | \mathbf{U}, \delta \rangle \end{array}$$



A Hard Question ...

Easy Question

Easy Question

What is pure procedural programming?

Another Easy Question

Another Easy Question

What is pure functional programming?

A Hard Question

A Hard Question

What is pure *logic* programming?

A Hard Question

What is pure *logic* programming?

A Hard Question

What is pure *logic* programming?

A: ...



B: ...

C: ...

...

Naveen: “Using automated theorem provers; in fact, you can just use HyperSlate.®”



“Direct” Programming in HyperSlate®

$$\forall m \forall i \forall o \exists \Phi \exists \phi_o [m : i \longrightarrow o \Leftrightarrow \Phi \vdash ? \phi_o]$$

Collection of nodes in HyperSlate®

Single node in HyperSlate®.

And just use the oracles to collaborate with you!

Ingredients for Making a PGLP Program ...

On the Anatomy of a PGLP Program

On the Anatomy of a PGLP Program

Linguistics

\vdots	\vdots	\vdots
L_2^μ	meta-level ₂ language	$(\{\phi\} \vdash \psi \wedge \{\psi\} \vdash \delta) \vdash_{\mu_2} \{\phi\} \vdash \delta$
L_1^μ	meta-level ₁ language	$\exists x \text{ rank}(\phi) = x \quad \{\phi\} \vdash \psi \quad \mathfrak{U} \models \phi$
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A collection of inference schemata. (For economy, see coming Example 1.)

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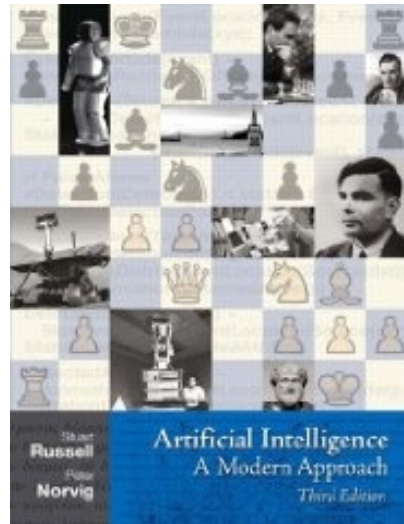
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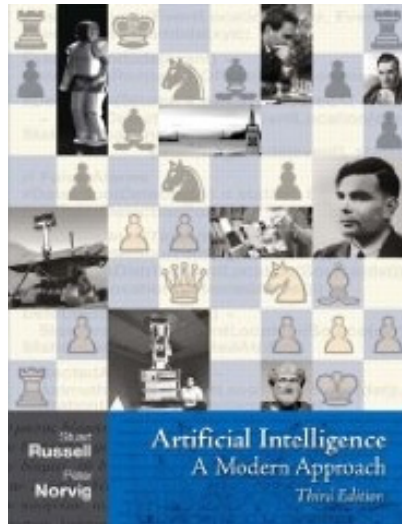
AI today ...

AI today:




AI today:



AI today:



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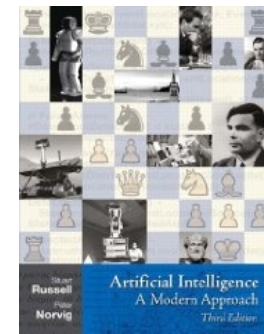
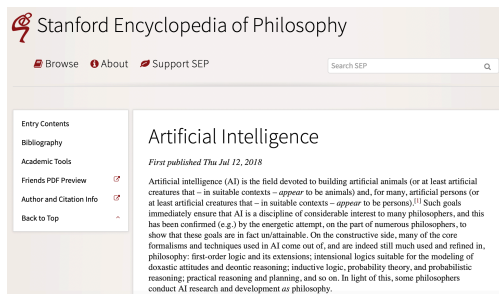
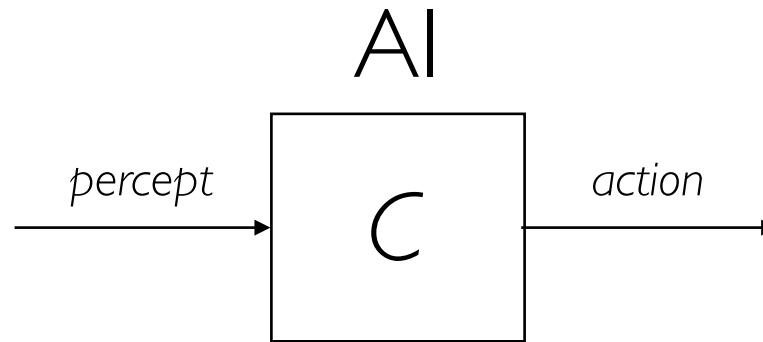
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Artificial Intelligence

First published Thu Jul 12, 2018

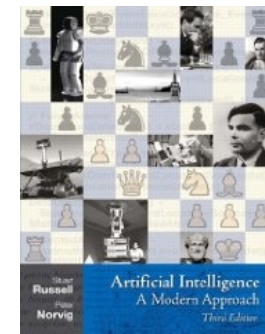
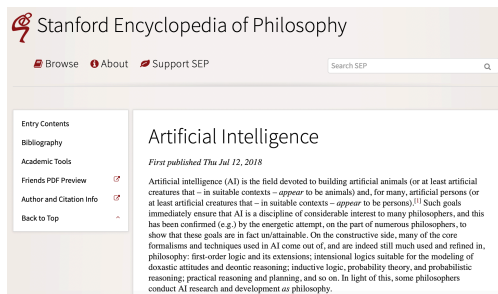
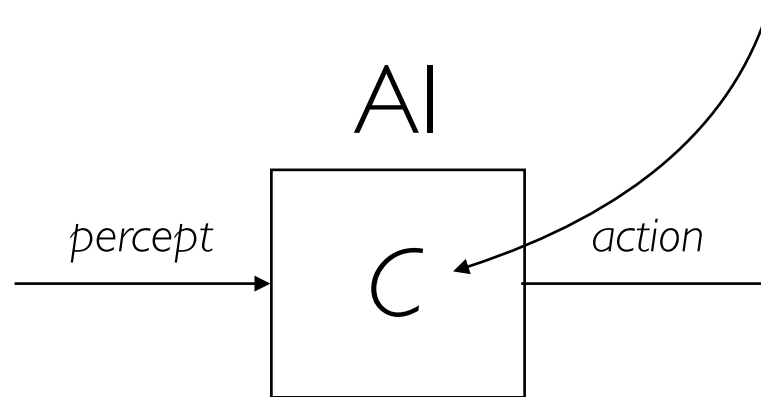
Artificial intelligence (AI) is the field devoted to building artificial animals (or at least artificial creatures that – in suitable contexts – *appear* to be animals) and, for many, artificial persons (or at least artificial creatures that – in suitable contexts – *appear* to be persons).^[1] Such goals immediately ensure that AI is a discipline of considerable interest to many philosophers, and this has been confirmed (e.g.) by the energetic attempt, on the part of numerous philosophers, to show that these goals are in fact un/attainable. On the constructive side, many of the core formalisms and techniques used in AI come out of, and are indeed still much used and refined in, philosophy: first-order logic and its extensions; intensional logics suitable for the modeling of doxastic attitudes and deontic reasoning; inductive logic, probability theory, and probabilistic reasoning; practical reasoning and planning, and so on. In light of this, some philosophers conduct AI research and development *as* philosophy.

AI:



AI:

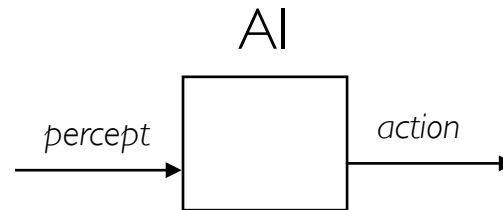
A (Turing-level) entity that computes.



Resurrection of The Triad

The Triad Resurrected & Rebuilt, & Better

Logic
 \mathcal{L}

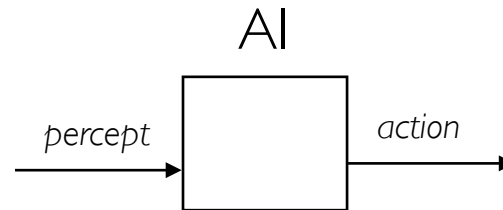


$$\mathcal{L} := \langle L, \mathcal{I} \rangle \quad \frac{\begin{array}{cc} \mathbb{P} & L \\ \mathfrak{q} & L \end{array}}{\begin{array}{l} \mathbb{R} : \langle \mathbb{P}, \mathfrak{q} \rangle \longrightarrow \langle \mathbf{Y} | \mathbf{N} | \mathbf{U}, \delta, \pi_{(s)} | \alpha_{(s)} \rangle \\ \mathbb{C} : \pi_{(s)} | \alpha_{(s)} \longrightarrow \langle \mathbf{Y} | \mathbf{N} | \mathbf{U}, \delta \rangle \end{array}}$$

Pure General Logic Programming

The Triad Resurrected & Rebuilt, & Better

Logic
 \mathcal{L}



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Pure General Logic Programming

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Pure General Logic Programming

What's Part 2 about ? ...

PROGRAMMING LANGUAGE PRAGMATICS

FOURTH EDITION

Michael L. Scott

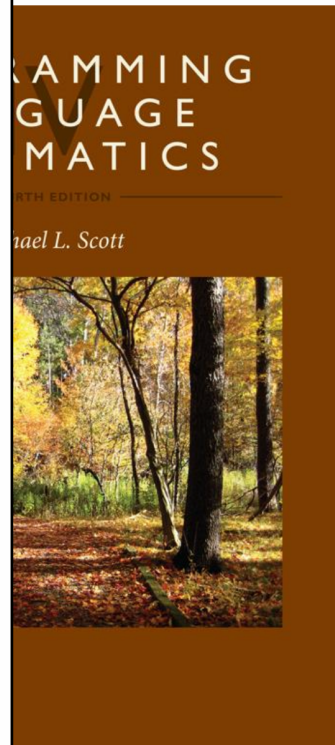


MK
MOSKOWSKI & KAPLAN

Alternative Programming Models

As we noted in [Chapter 1](#), programming languages are traditionally though imperfectly classified into various imperative and declarative families. We have had occasion in Parts I and II to mention issues of particular importance to each of the major families. Moreover much of what we have covered—syntax, semantics, naming, types, abstraction—applies uniformly to all. Still, our attention has focused mostly on mainstream imperative languages. In Part III we shift this focus.

Functional and logic languages are the principal nonimperative options. We consider them in [Chapters 11](#) and [12](#), respectively. In each case we structure our discussion around representative languages: Scheme and OCaml for functional programming, Prolog for logic programming. In [Chapter 11](#) we also cover eager and lazy evaluation, and first-class and higher-order functions. In [Chapter 12](#) we cover issues that make fully automatic, general purpose logic programming difficult, and describe restrictions used in practice to keep the model tractable. Optional sections in both chapters consider mathe-



Models

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Models

As we noted in [Chapter 1](#), programming languages are traditionally thought of in terms of various imperative and declarative models. We had occasion in Parts I and II to discuss the particular importance to functional programming. Moreover much of what we say about semantics, naming, types, and so on applies uniformly to all. Still, our attention has been on mainstream imperative programming, and we shift this focus.

Functional and logic languages are nonimperative options. We discuss them in [Chapters 11](#) and [12](#), respectively. We structure our discussion around the examples of Scheme and OCaml for functional programming and Prolog for logic programming. We cover eager and lazy evaluation, closures, higher-order functions. I discuss the issues that make fully automatic compilation of logic programming difficult. We discuss how it is used in practice to keep track of state. Optional sections in both

mathematical foundations: Lambda Calculus for functional programming, Predicate Calculus for logic programming.

The remaining two chapters consider concurrent and scripting models, both of which are increasingly popular, and cut across the imperative/declarative divide. Concurrency is driven by the hardware parallelism of internetworked computers and by the coming explosion in multithreaded processors and chip-level multiprocessors. Scripting is driven by the growth of the World Wide Web and by an increasing emphasis on programmer productivity, which places rapid development and reusability above sheer runtime performance.

[Chapter 13](#) begins with the fundamentals of concurrency, including communication and synchronization, thread creation syntax, and the implementation of threads. The remainder of the chapter is divided between *shared-memory* models, in which threads use explicit or implicit synchronization mechanisms to manage a common set of variables, and (on the companion site) *message-passing* models, in which threads interact only through explicit communication.

The first half of [Chapter 14](#) surveys problem domains in which scripting plays a major role: shell (command) languages, text processing and report generation, mathematics and statistics, the "gluing" together of program components, extension mechanisms for complex applications, and client and server-side Web scripting. The second half considers some of the more important language innovations championed by scripting languages: flexible scoping and naming conventions, string and pattern manipulation (extended regular expressions), and high level data types.

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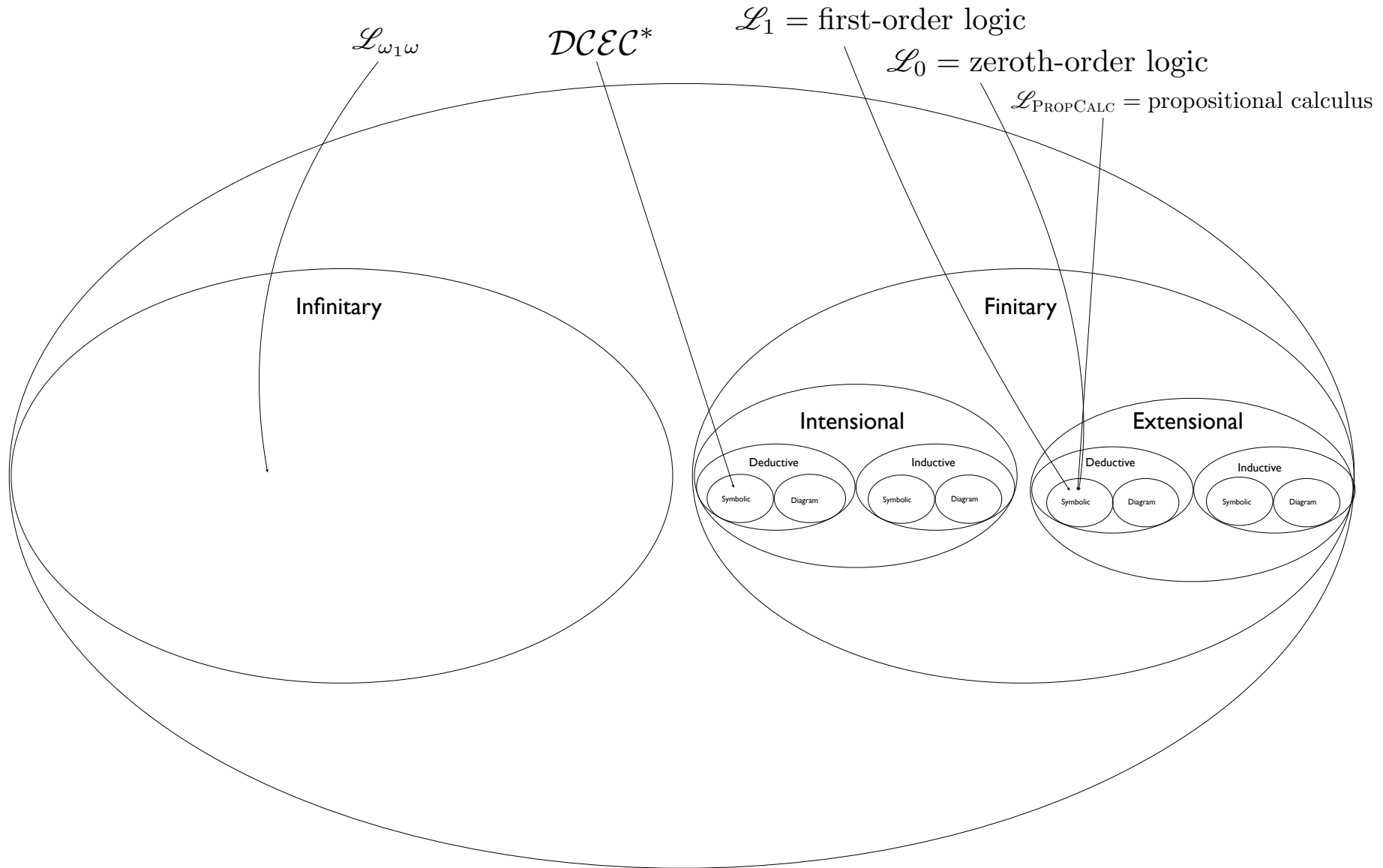
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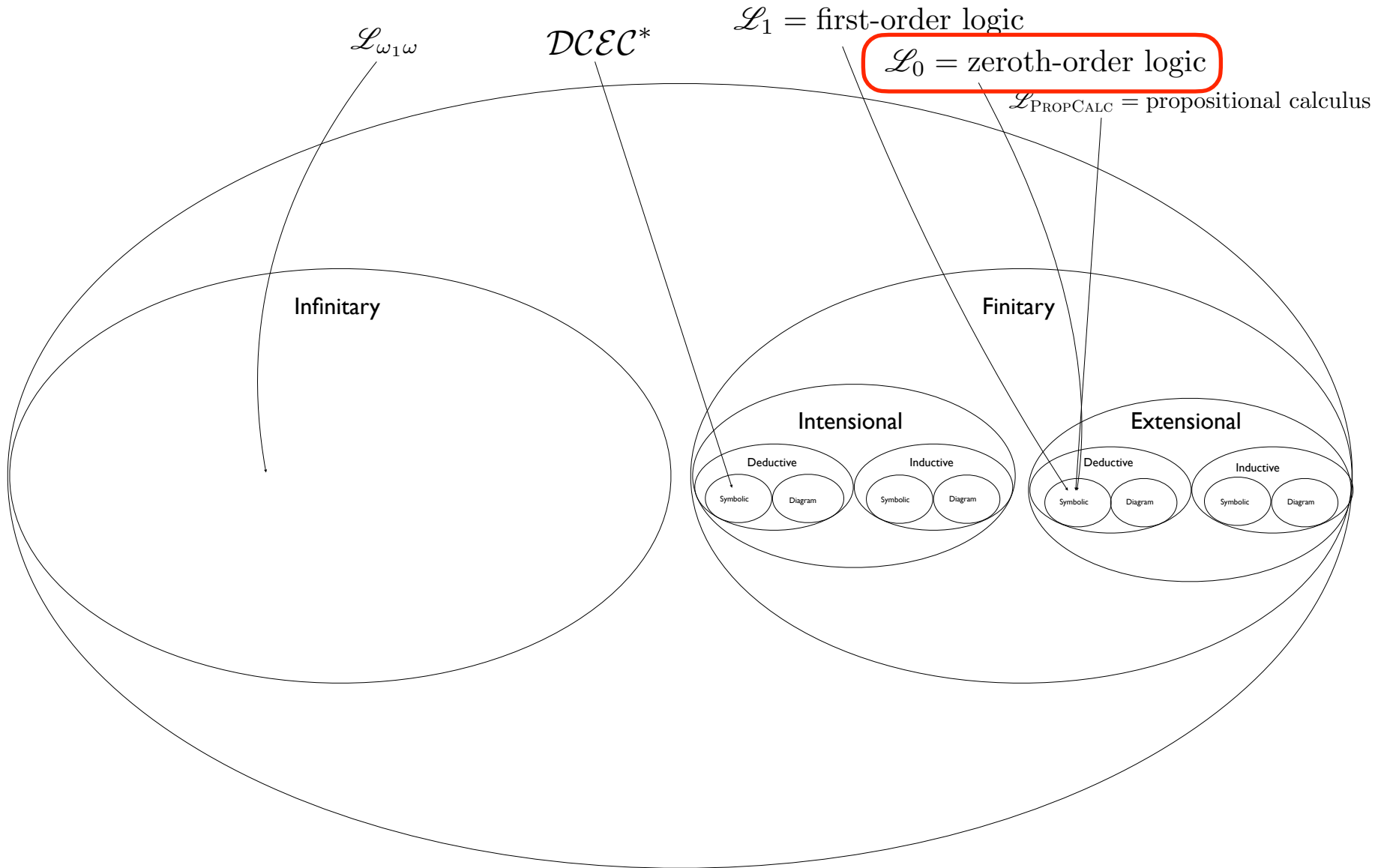
[Chapter 13](#) begins with the fundamentals of concurrency, including communication and synchronization, thread creation syntax, and the implementation of threads. The remainder of the chapter is divided between *shared-memory* models, in which threads use explicit or implicit synchronization mechanisms to manage a common set of variables, and (on the companion site) *message-passing* models, in which threads interact only through explicit communication.

The first half of [Chapter 14](#) surveys problem domains in which scripting plays a major role: shell (command) languages, text processing and report generation, mathematics and statistics, the "gluing" together of program components, extension mechanisms for complex applications, and client and server-side Web scripting. The second half considers some of the more important language innovations championed by scripting languages: flexible scoping and naming conventions, string and pattern manipulation (extended regular expressions), and high level data types.

The Universe of Logics



The Universe of Logics



Starter Hyperlog[®]: Datalog

\mathcal{L}_0 = zeroth-order logic

Datalog Syntax

<i>(program)</i> P	$:=$	$\phi_1 \wedge \phi_2 \wedge \dots \wedge \phi_n$
<i>(horn-clause formula)</i> ϕ_i	$:=$	$\alpha_1 \wedge \dots \wedge \alpha_m \rightarrow \alpha_j$
<i>(atomic formula)</i> α_i	$:=$	$R(t_1, t_2, \dots, t_o)$
<i>(terms)</i> t	$:=$	$x \mid c$

where x is a variable and c a constant

Starter Hyperlog[®]: Datalog

\mathcal{L}_0 = zeroth-order logic

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HyperGrader[®] Problem Categories ▾ HyperSlate My Progression Leader Board | Spring 2021 RPI Selmer.Bringsjord@gmail.com (longsnowflake876) ▾

Create file

Propositional Calculus

L₀ = Pure Predicate Calculus

L₁ = First-order Logic

L₂ = Second-order Logic

K

T

D

S4

S5

DCEC (fragment)

Hyperlog

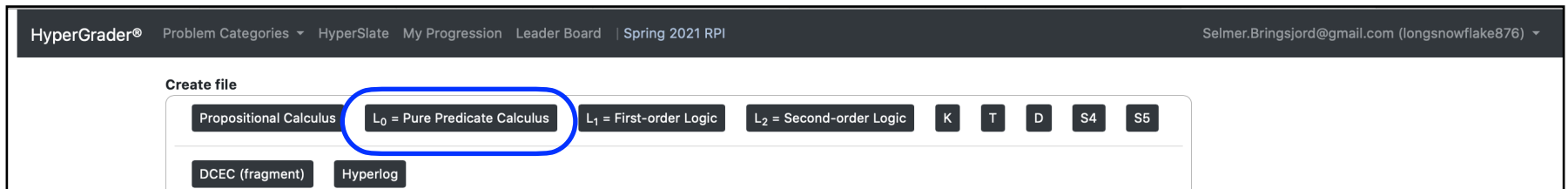
Starter Hyperlog[®]: Datalog

\mathcal{L}_0 = zeroth-order logic

Datalog Syntax

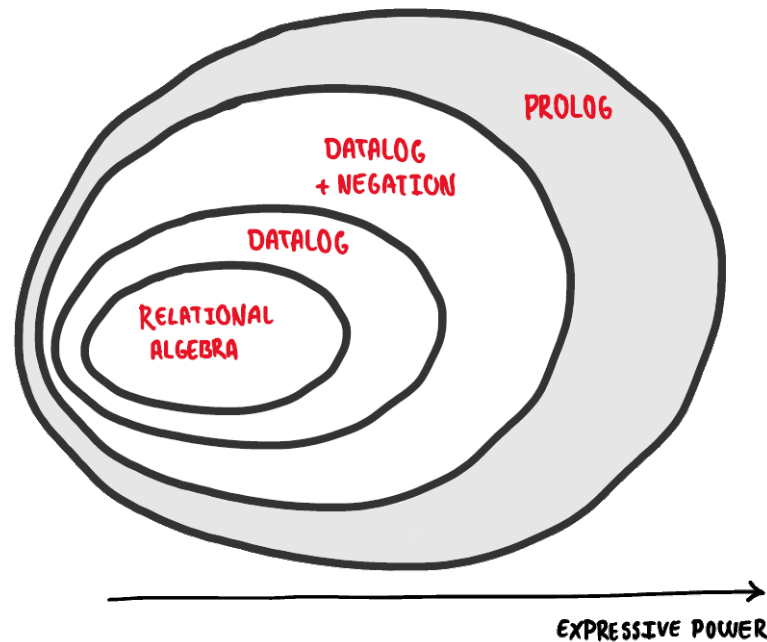
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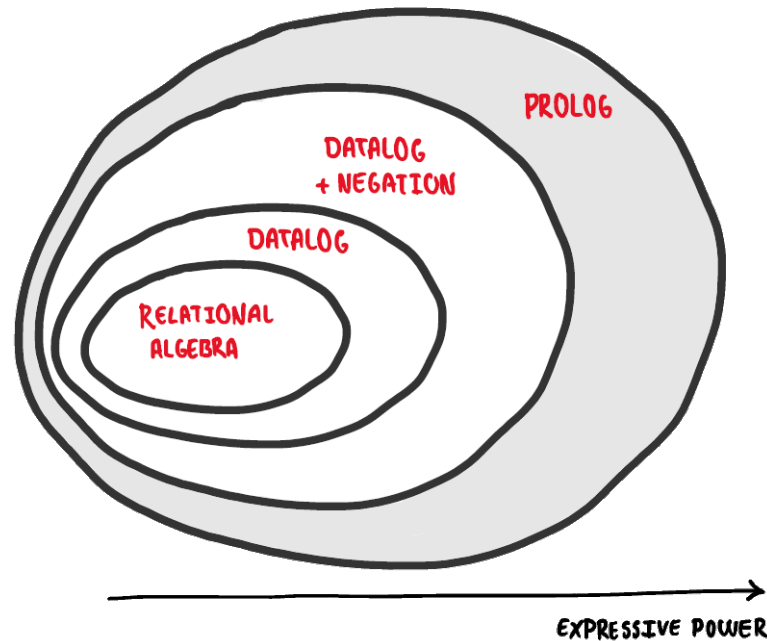
Starter HyperLog[®]: Datalog

From [“Introduction to Datalog,”](#) an excellent online piece.



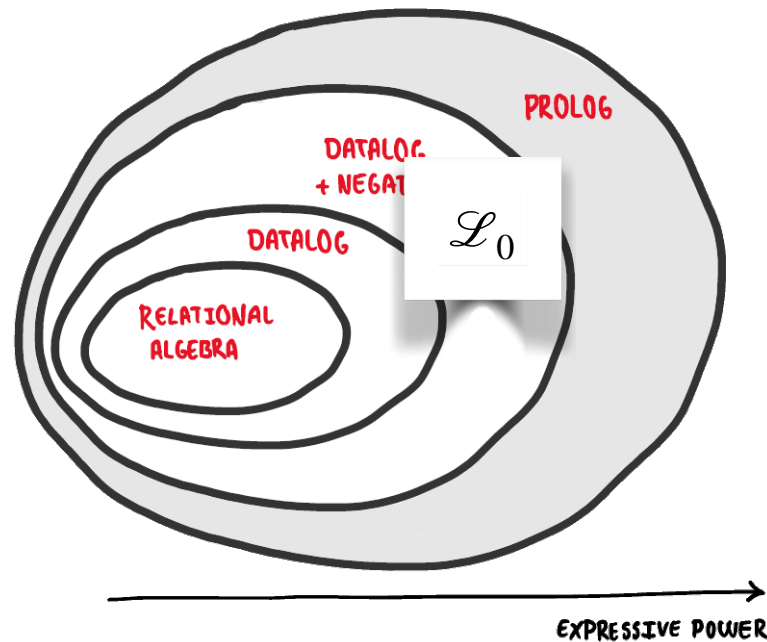
Starter HyperLog[®]: Datalog

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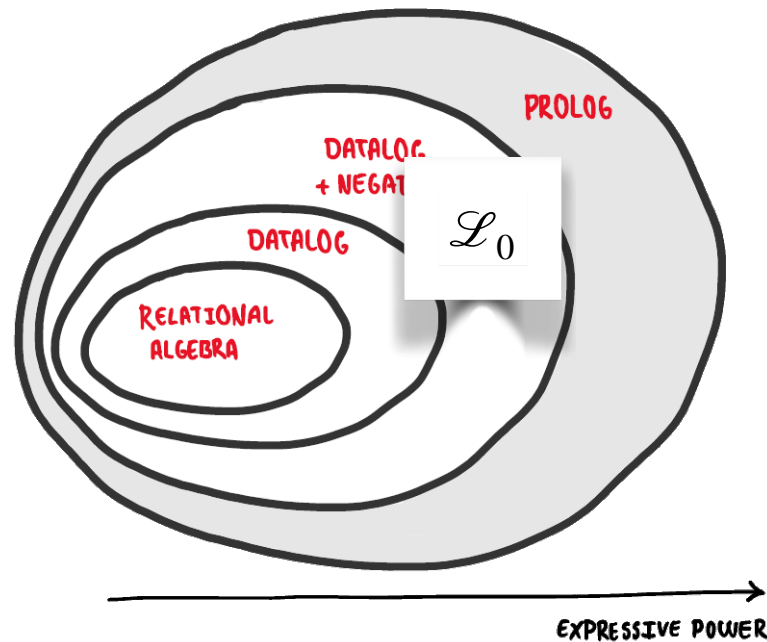
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Starter HyperLog[®]: Datalog

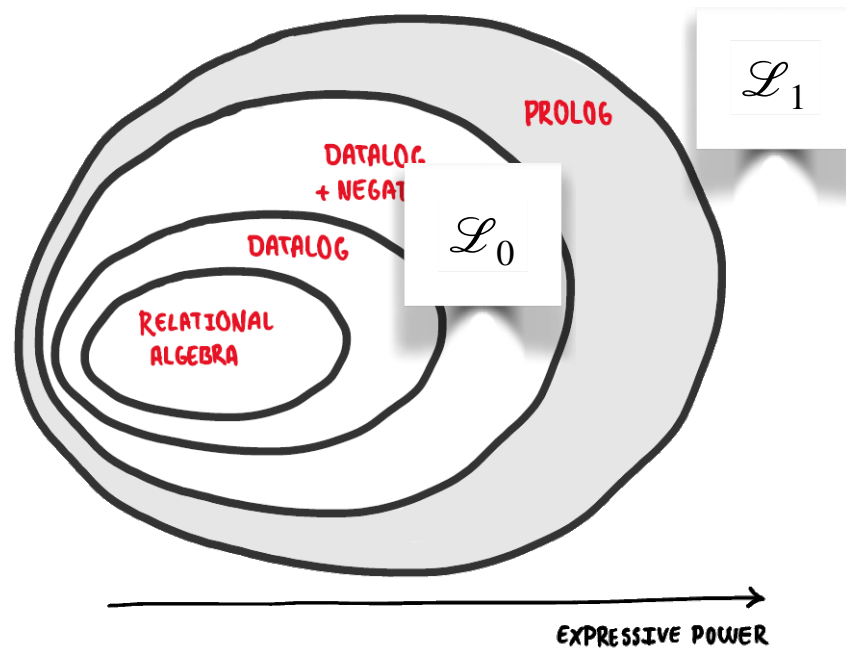
\mathcal{L}_1

From “[Introduction to Datalog](#),” an excellent online piece.



Starter HyperLog[®]: Datalog

From [“Introduction to Datalog,”](#) an excellent online piece.



Slutten

Slutten

Part II:

**Review of All Inference Rules/
Schemata in PropCalc = \mathcal{L}_{PC}**