On to *Intensional* Logics

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Troy, New York 12180 USA

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In The Logic-and-AI News

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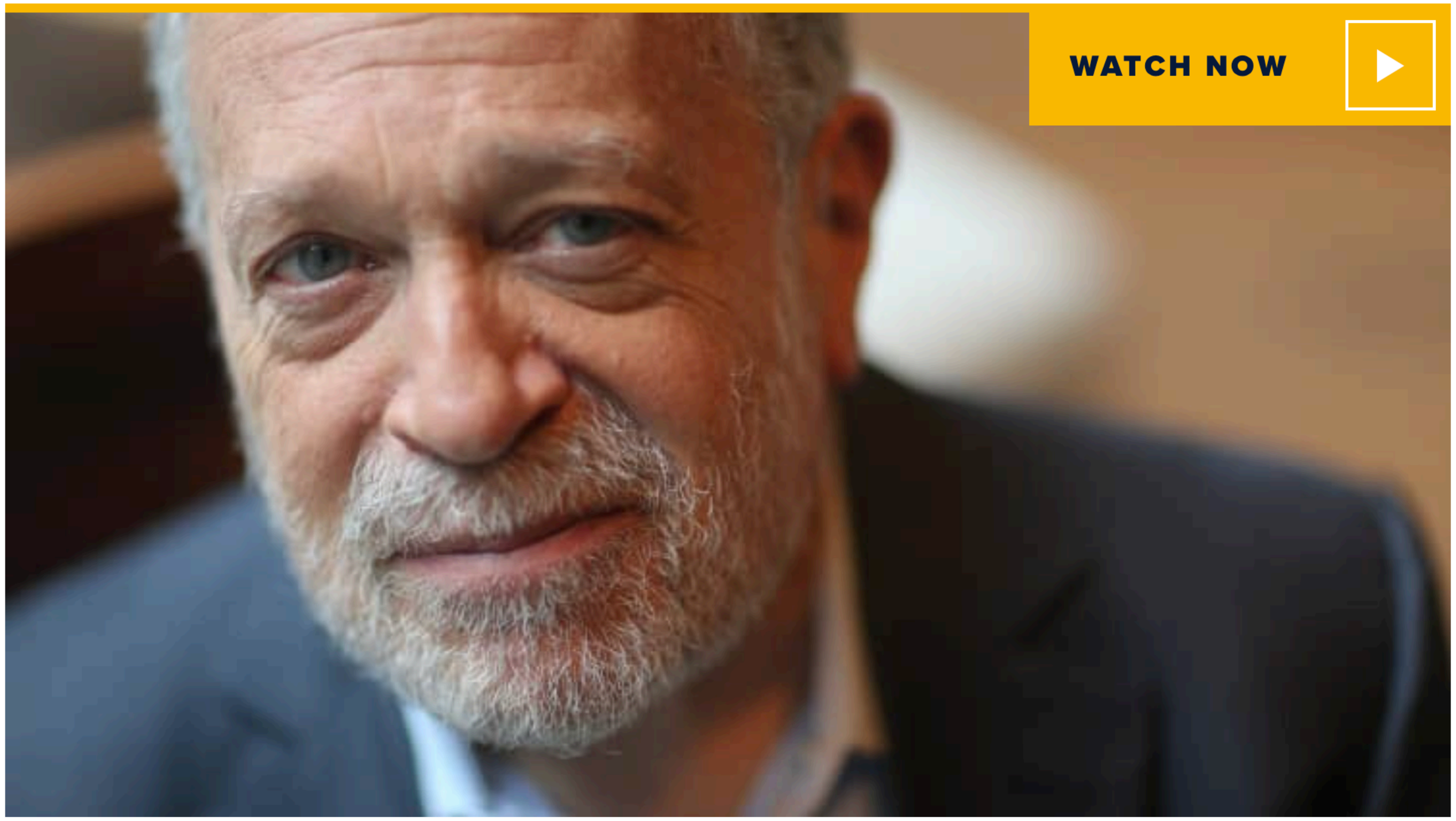
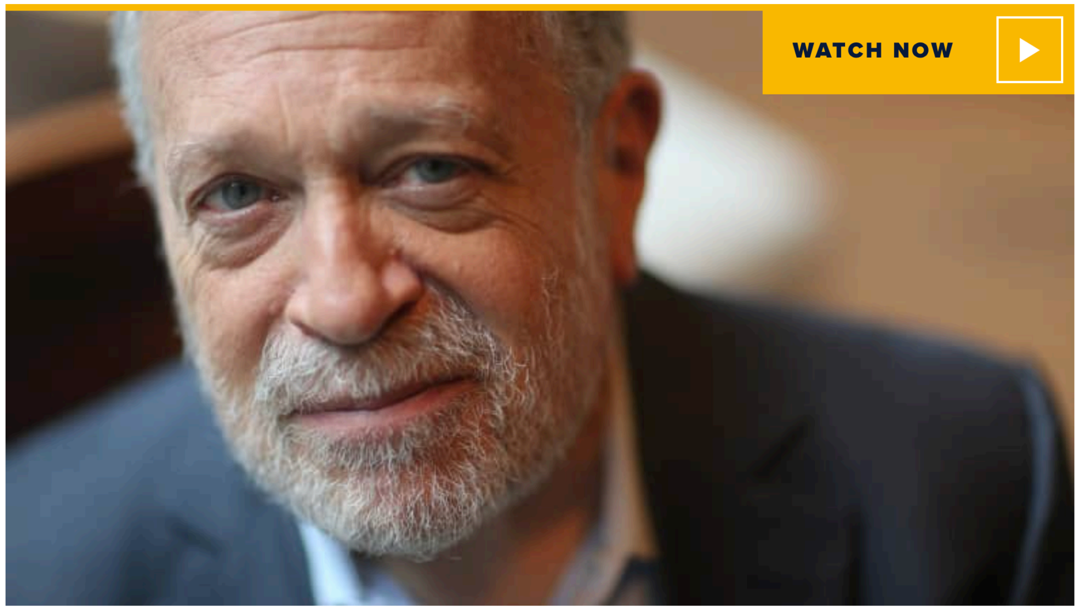


THE BOTTOM LINE

SHARE

Why AI may end labor unions and become your new employer: Robert Reich

Sep 3 2023



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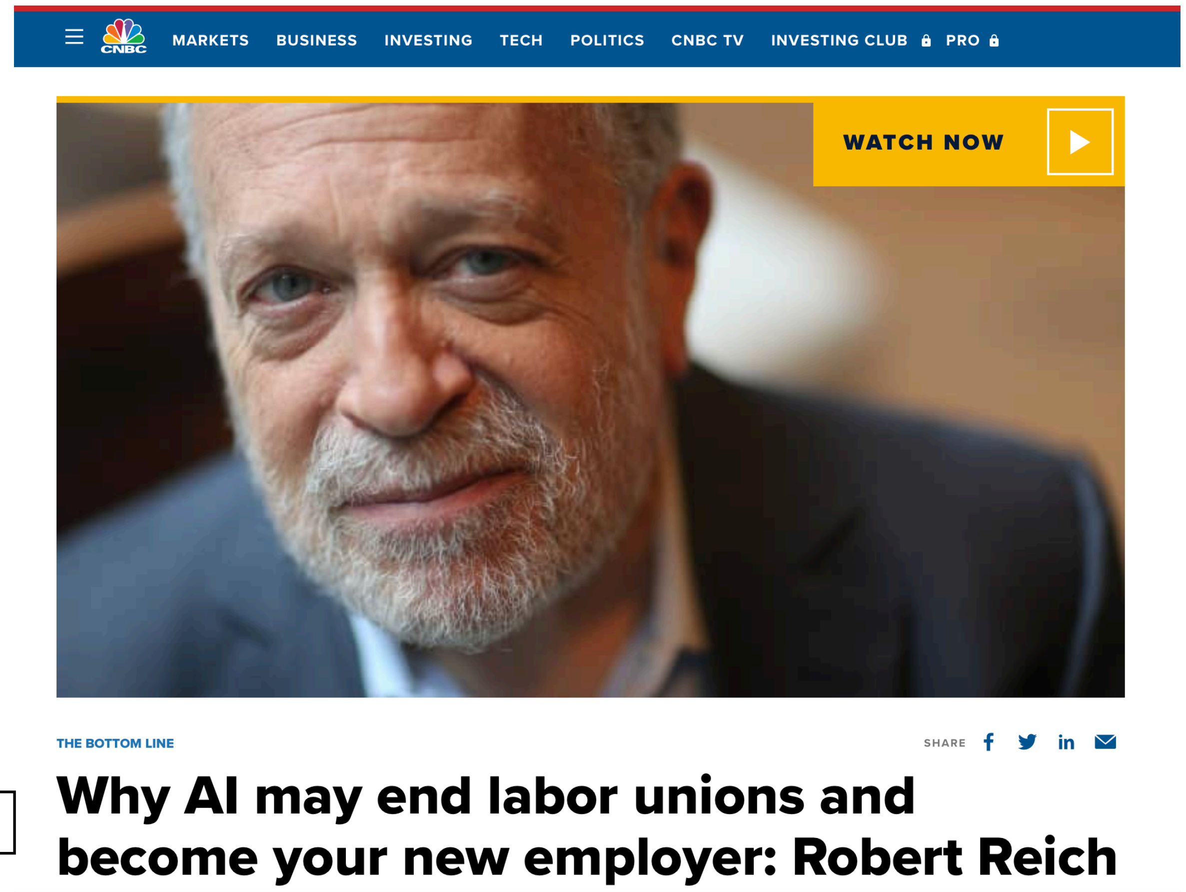

Sep 3 2023

Speaking of intensional/modal logic, is this mathematically possible?

On the esemplastic extensional-logic ladder …
questions?

Climbing the *k*-order Ladder

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a is a llama, as is *b*, *a* likes *b*, and the father of *a* is a llama as well.

Climbing the *k*-order Ladder

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SOL $\mathscr{L}_2$

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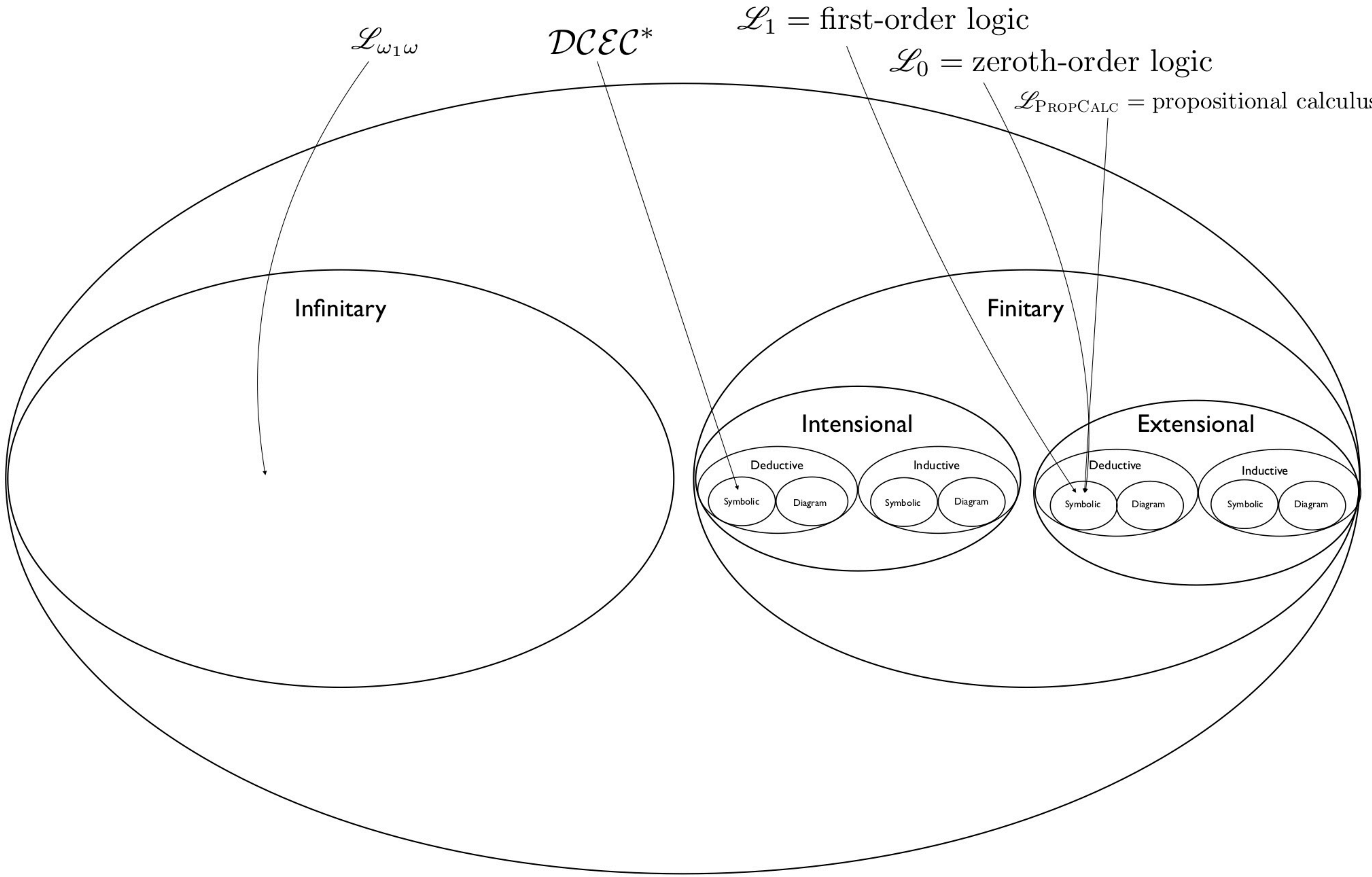
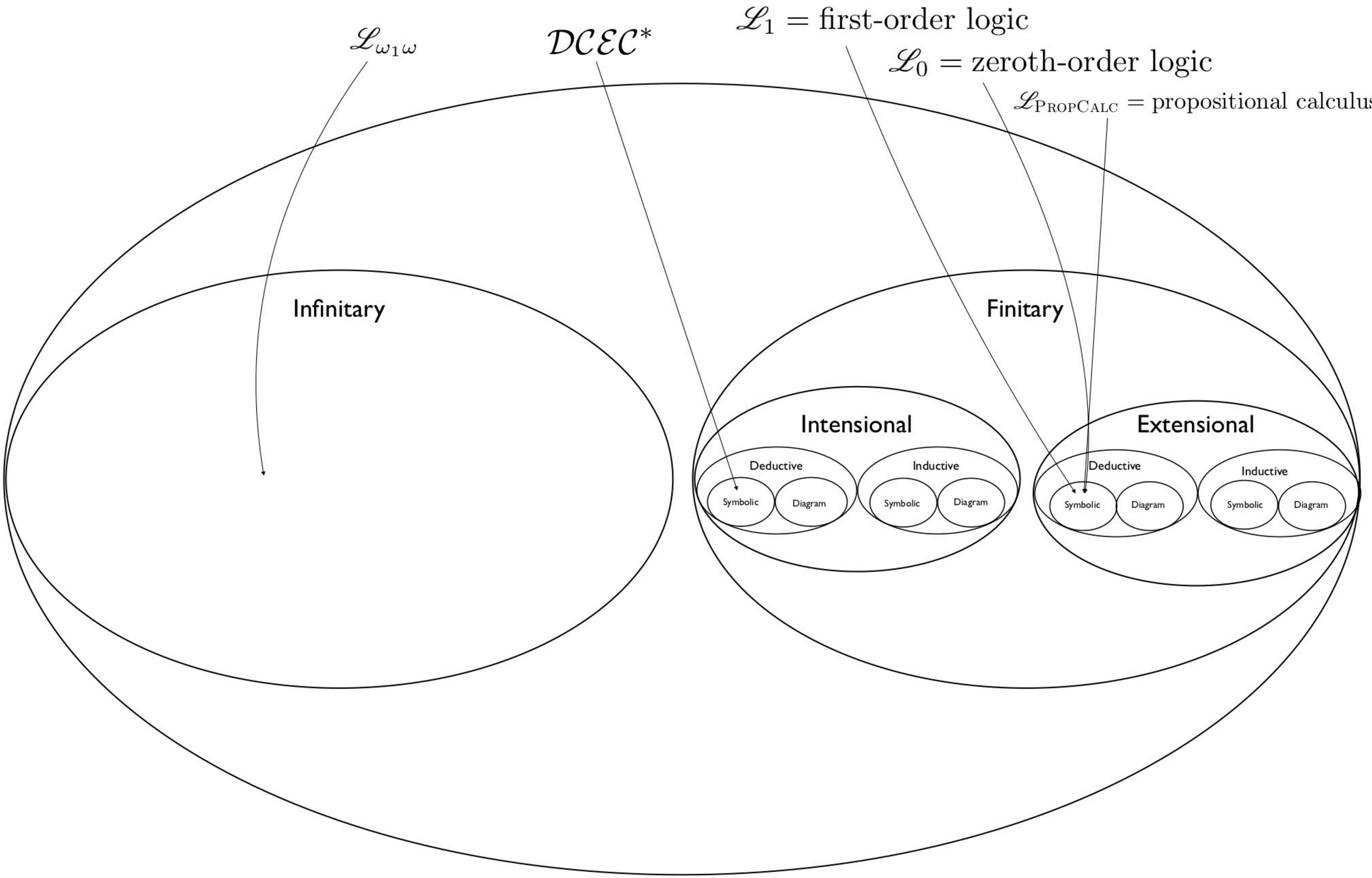
The Universe of Logics



$\mathscr{L}_3$
$\mathscr{L}_2$

The Universe of Logics



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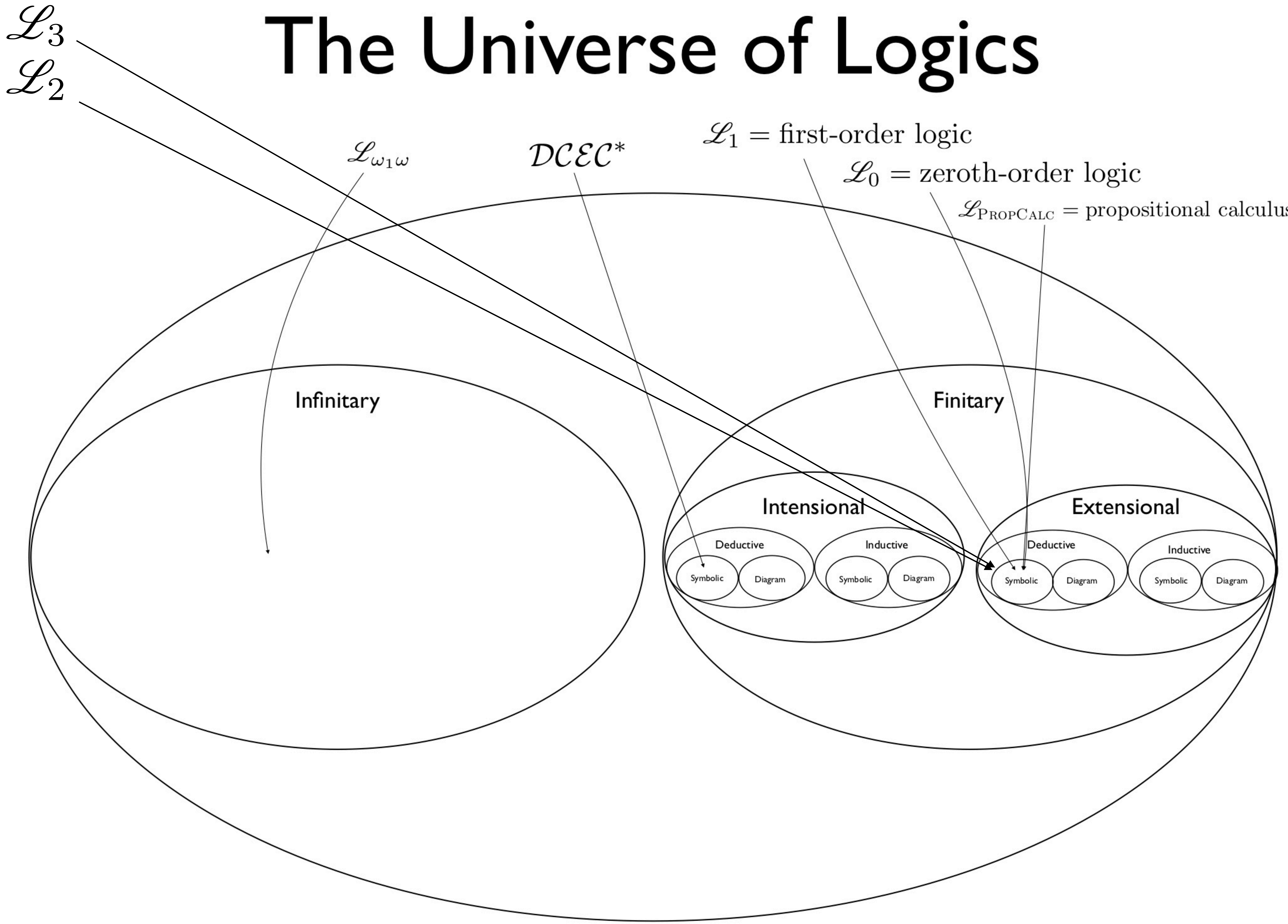

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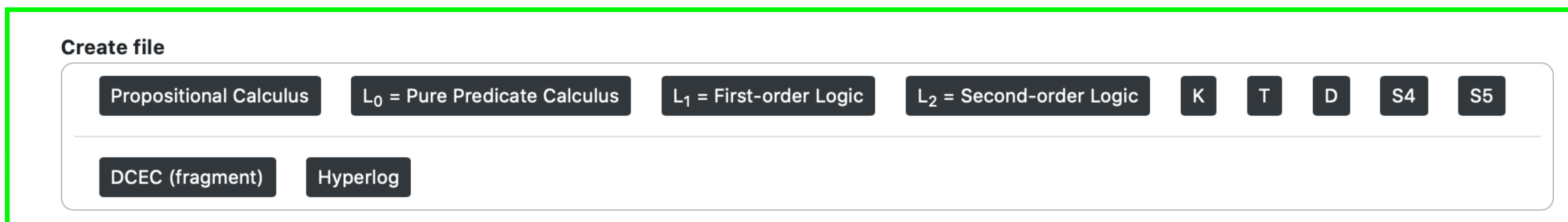

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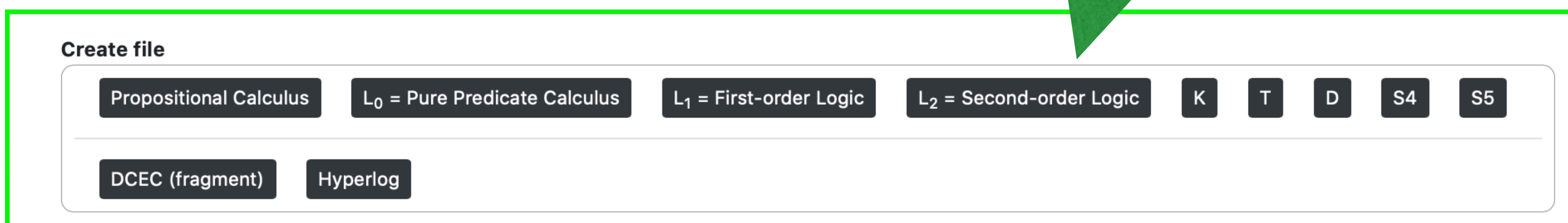

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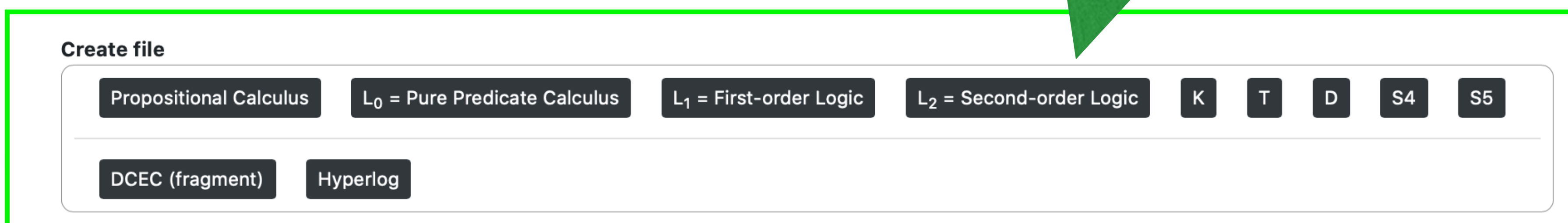

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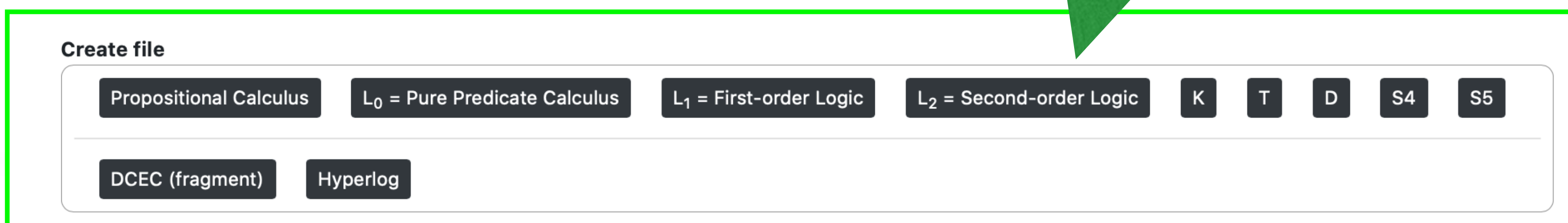

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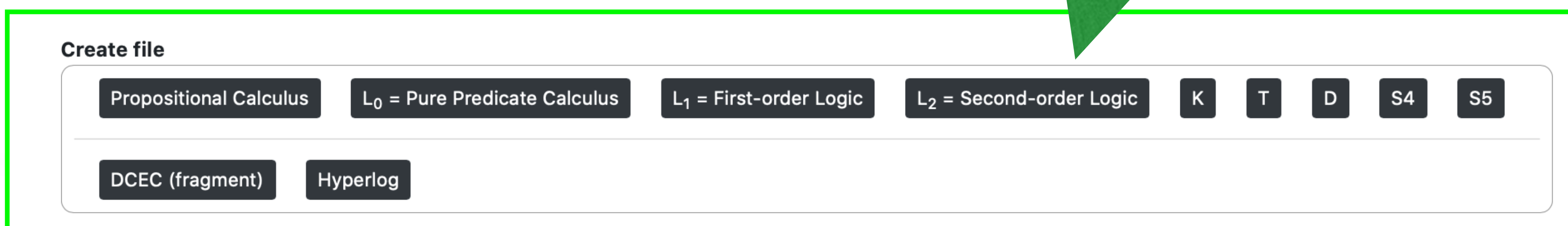

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Incomplete!

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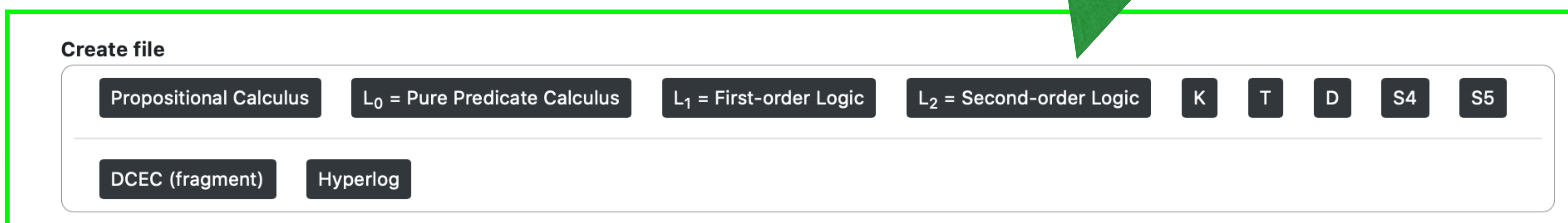

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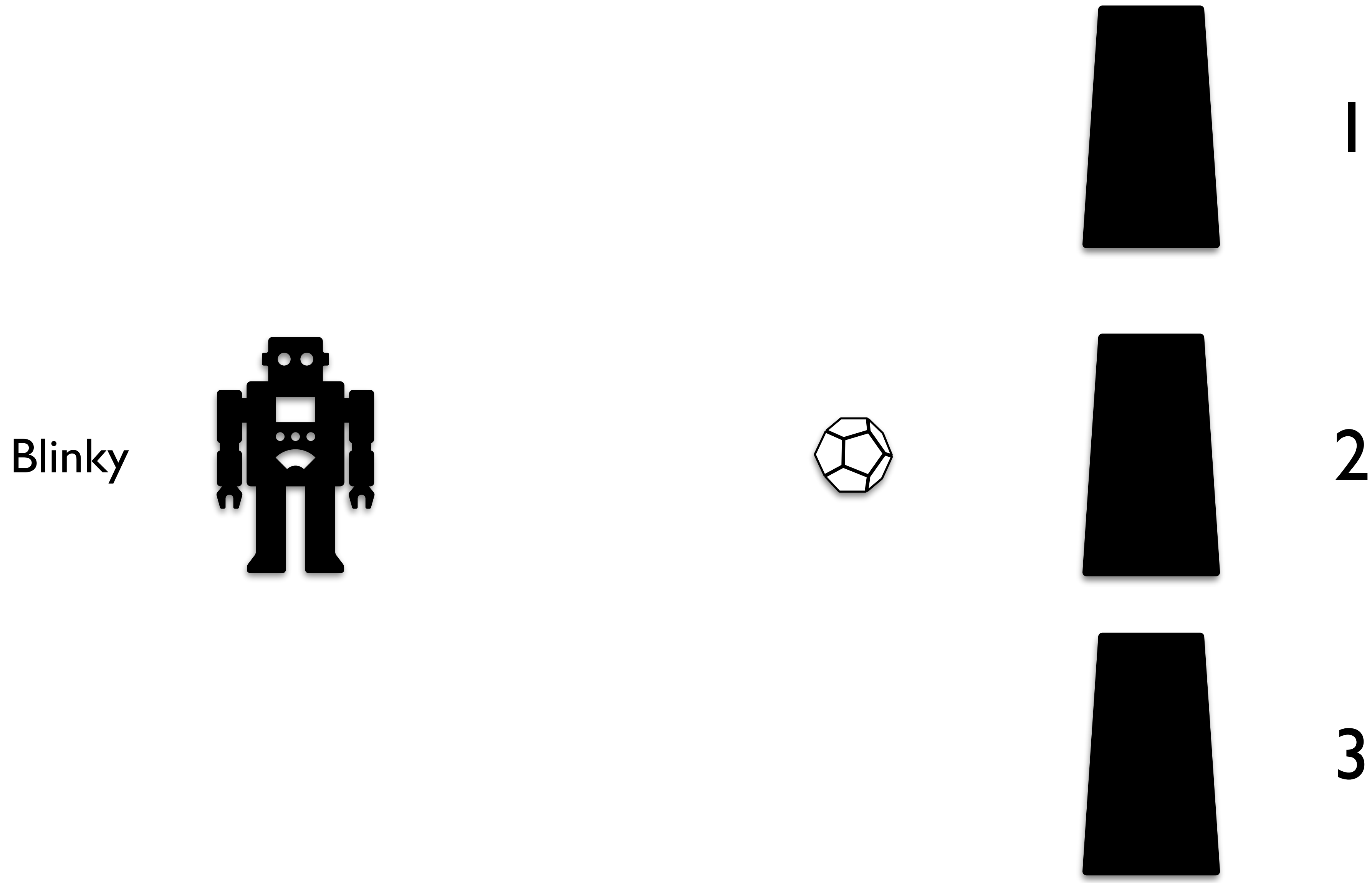
Blinky as portal to intensional logics …


Blinky
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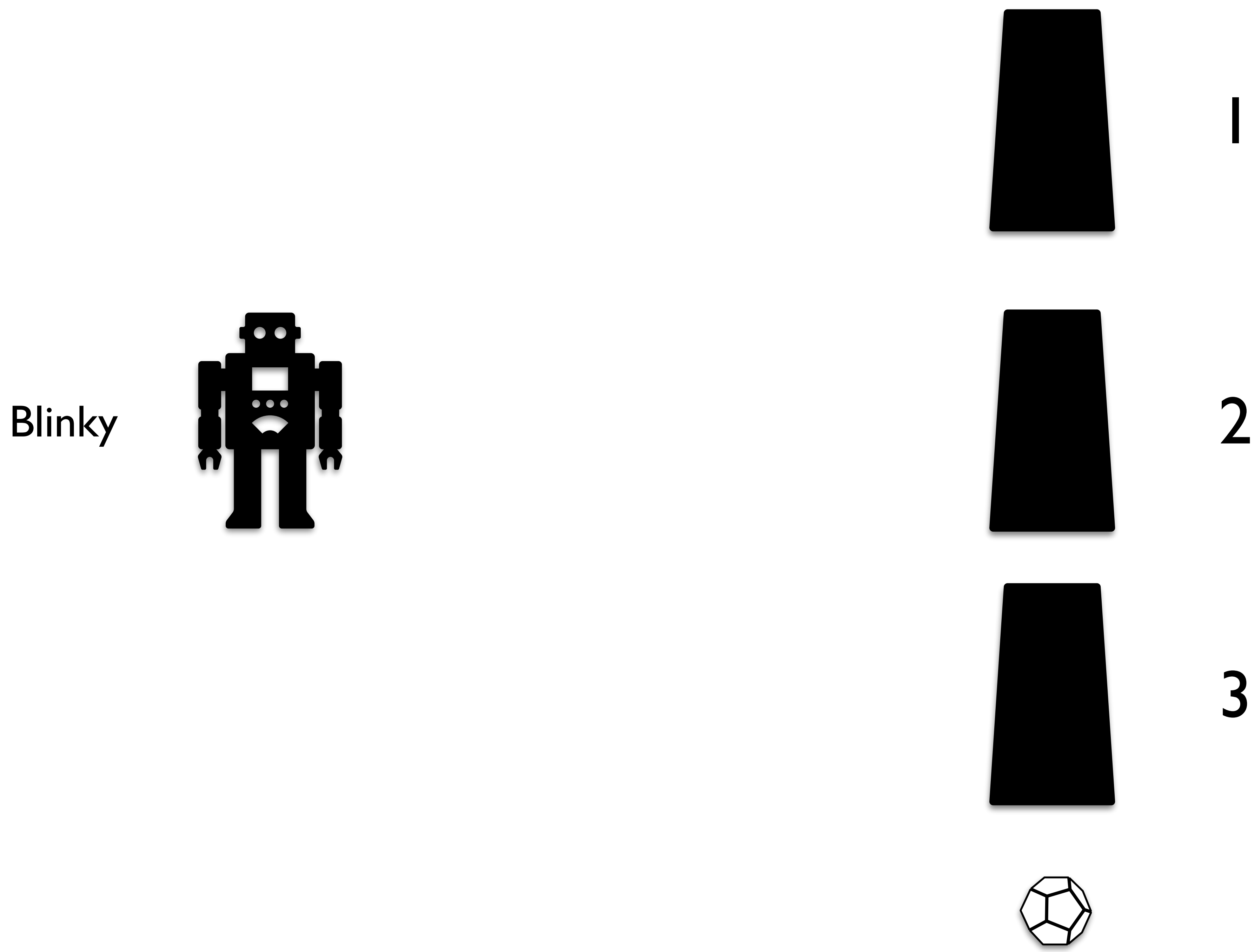
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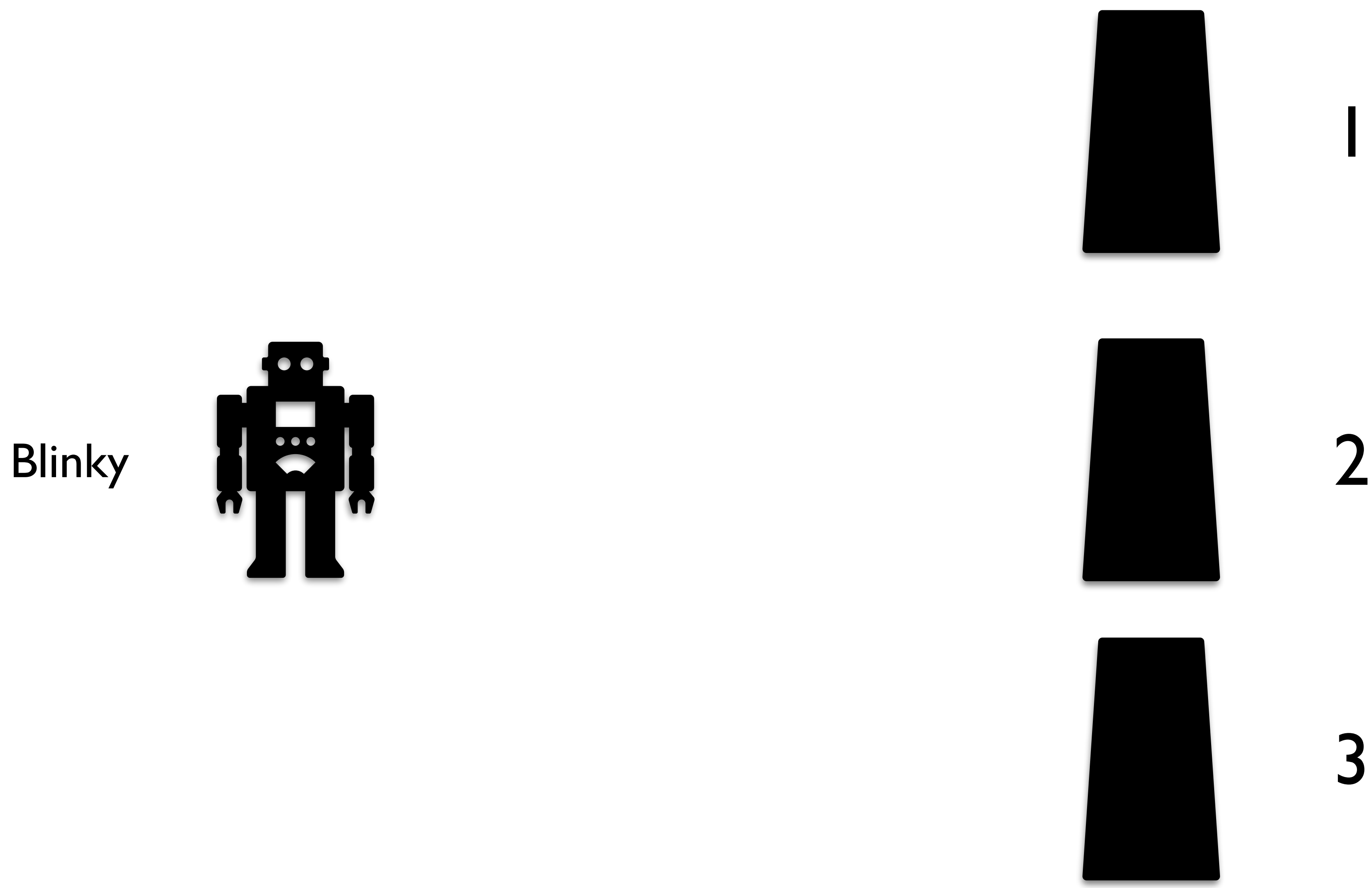
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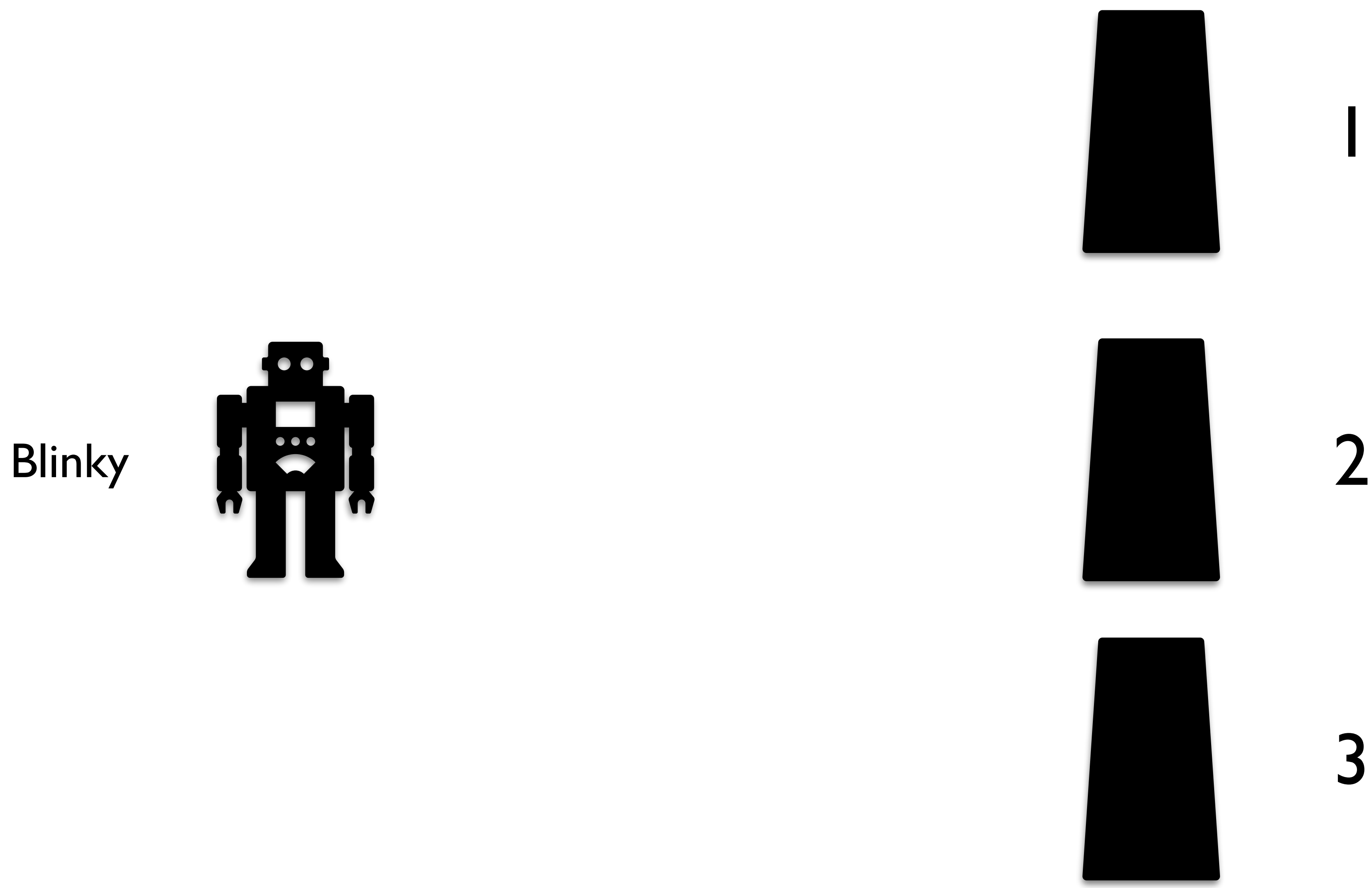
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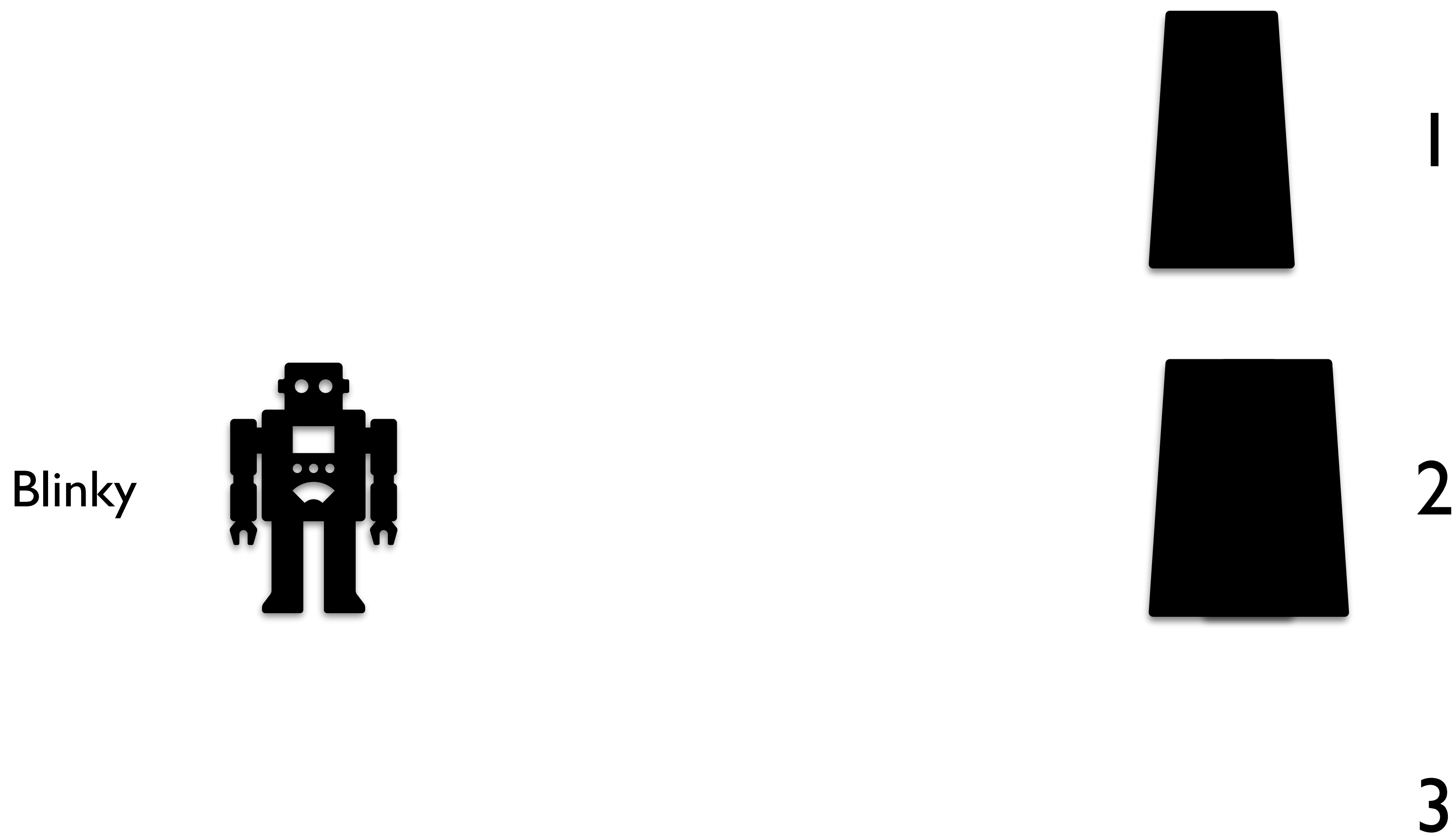
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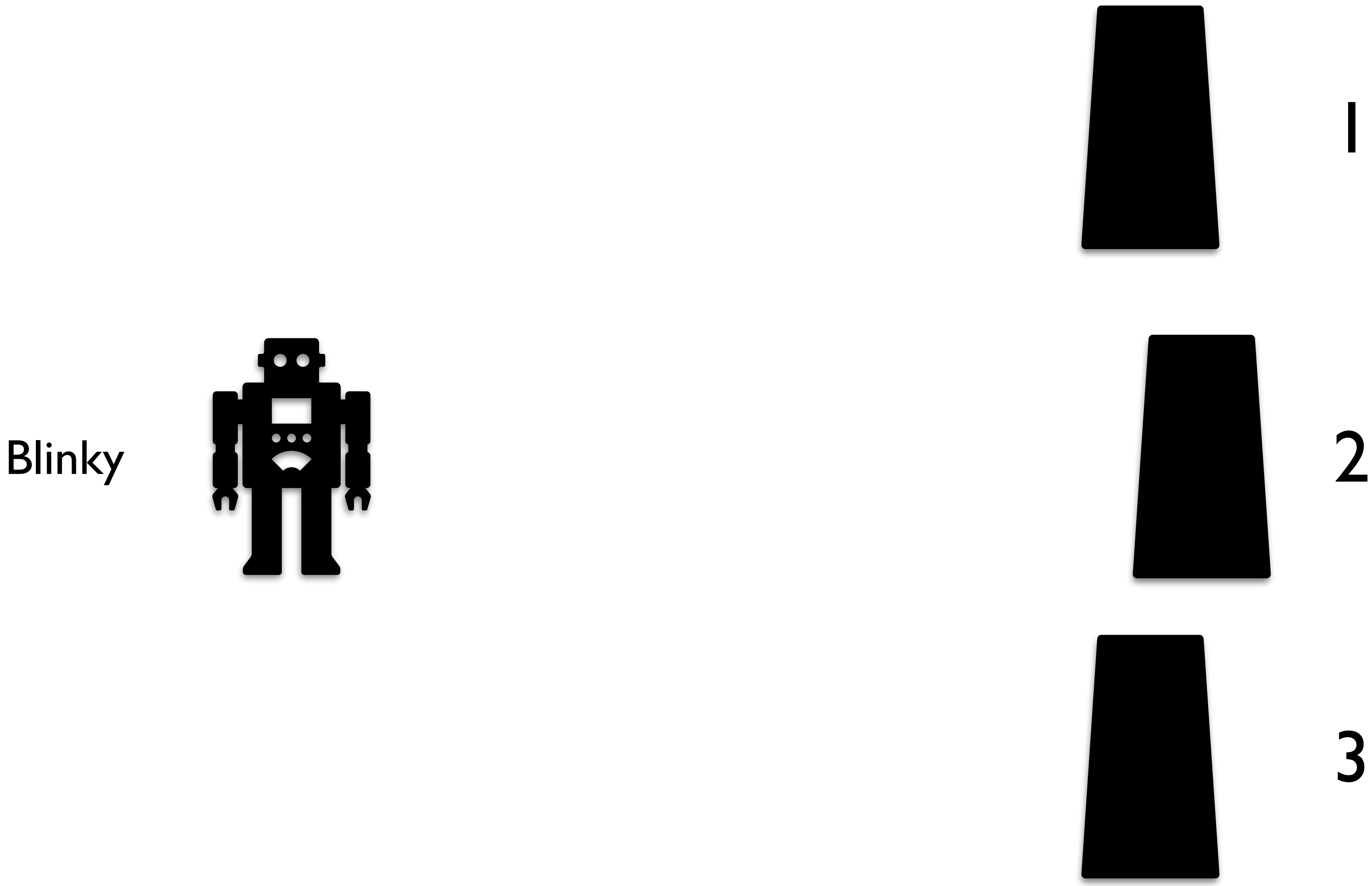
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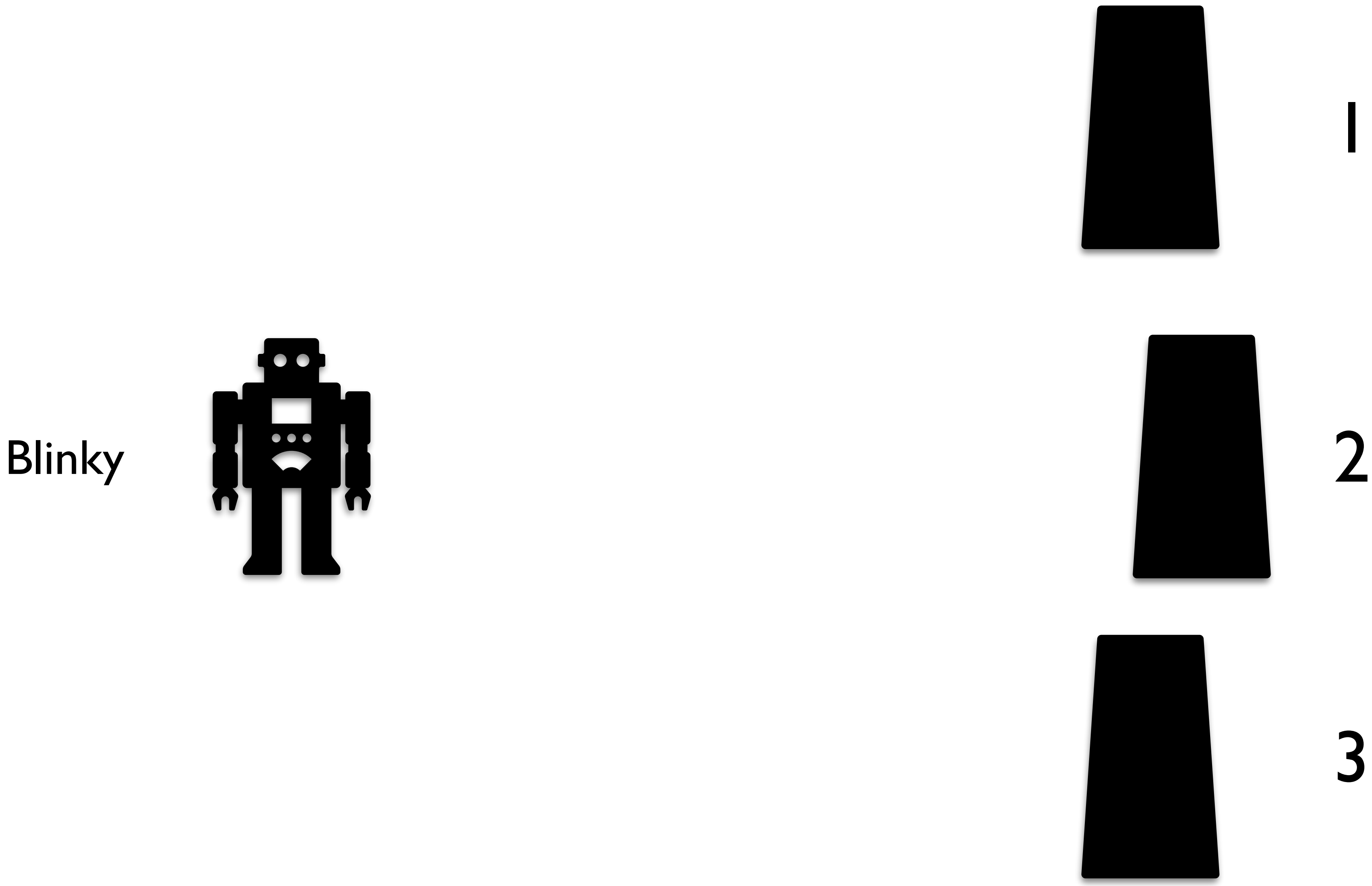

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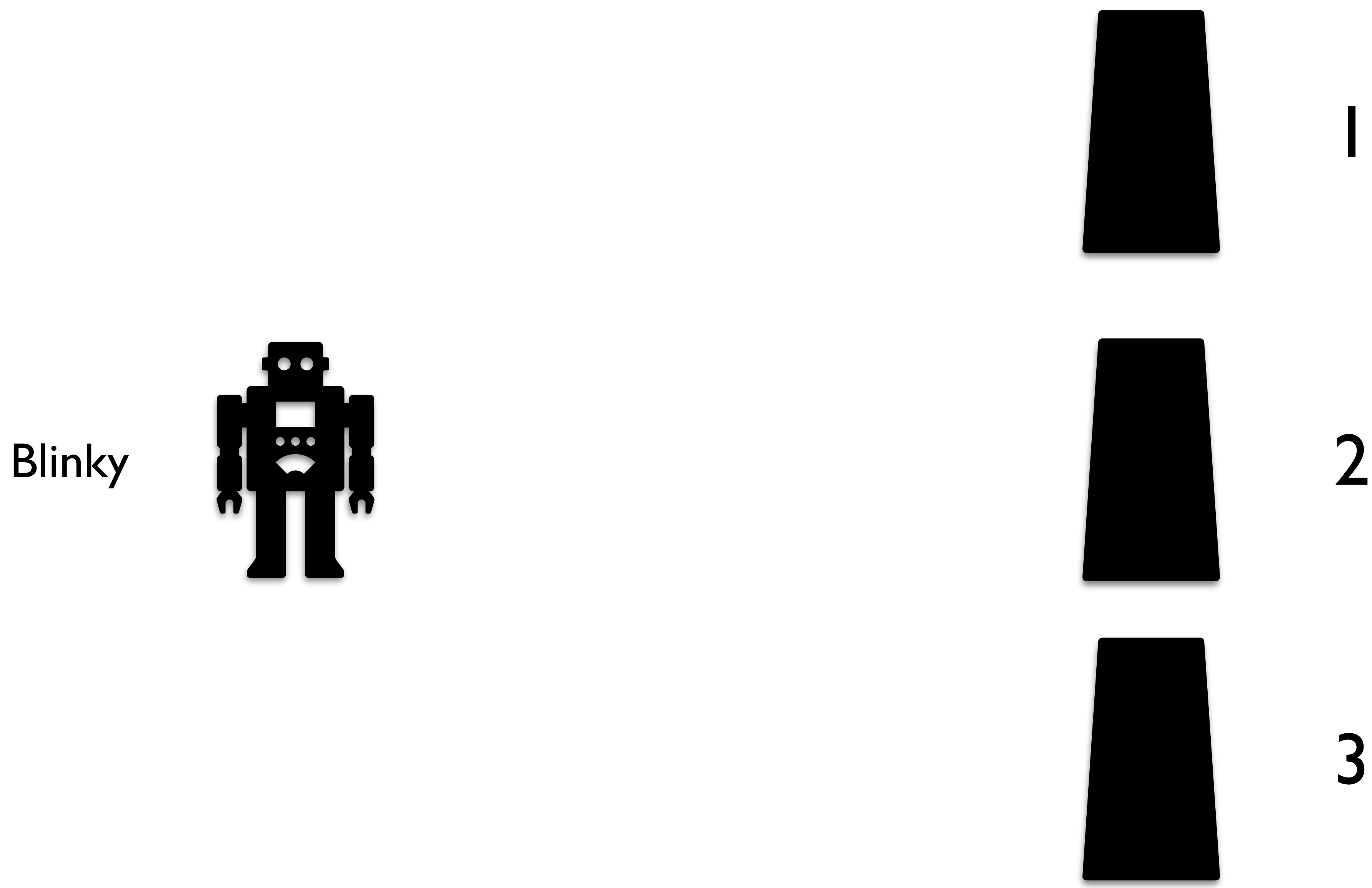
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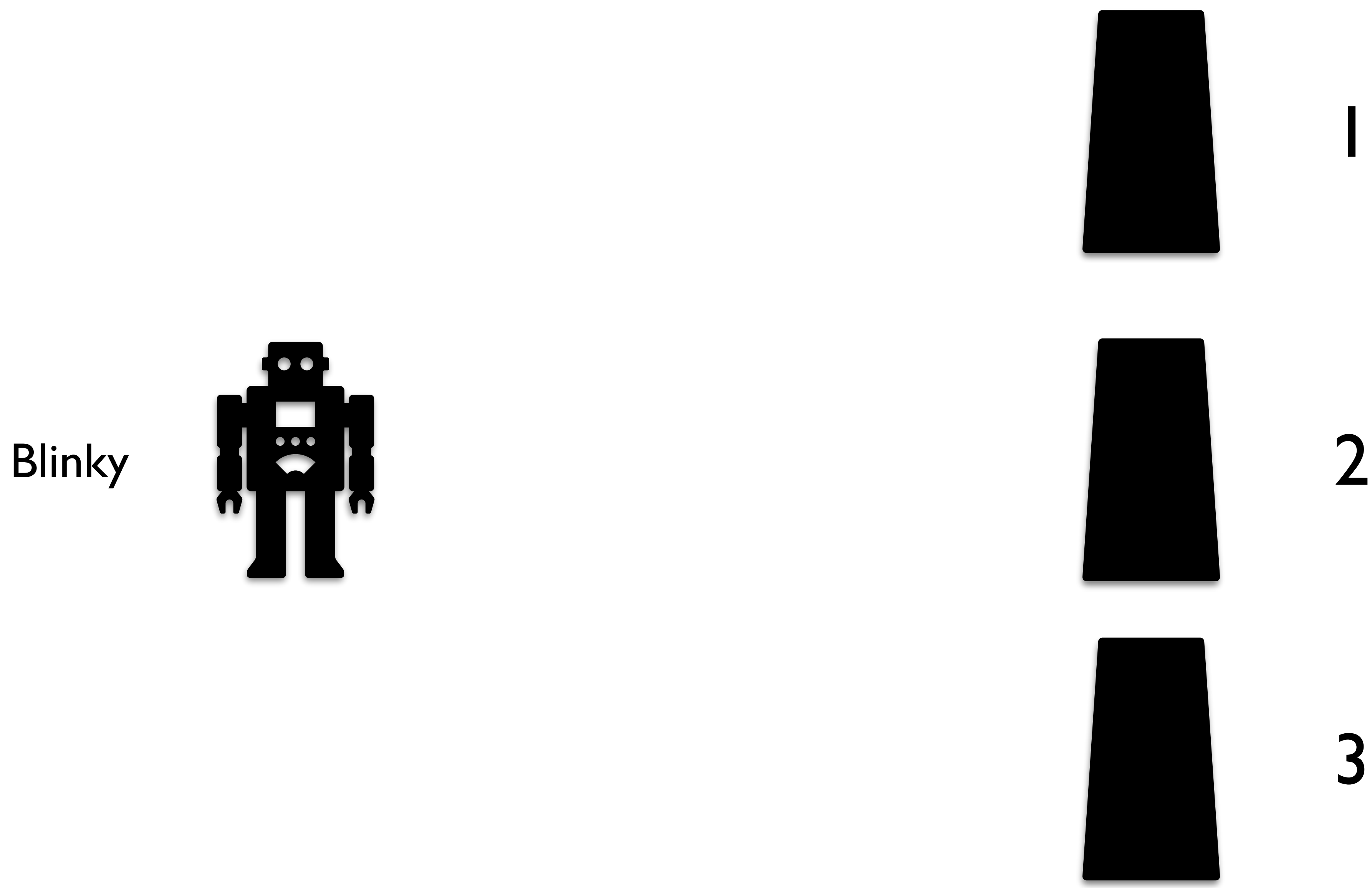
Blinky believes that the ball is in the cup at location #1.



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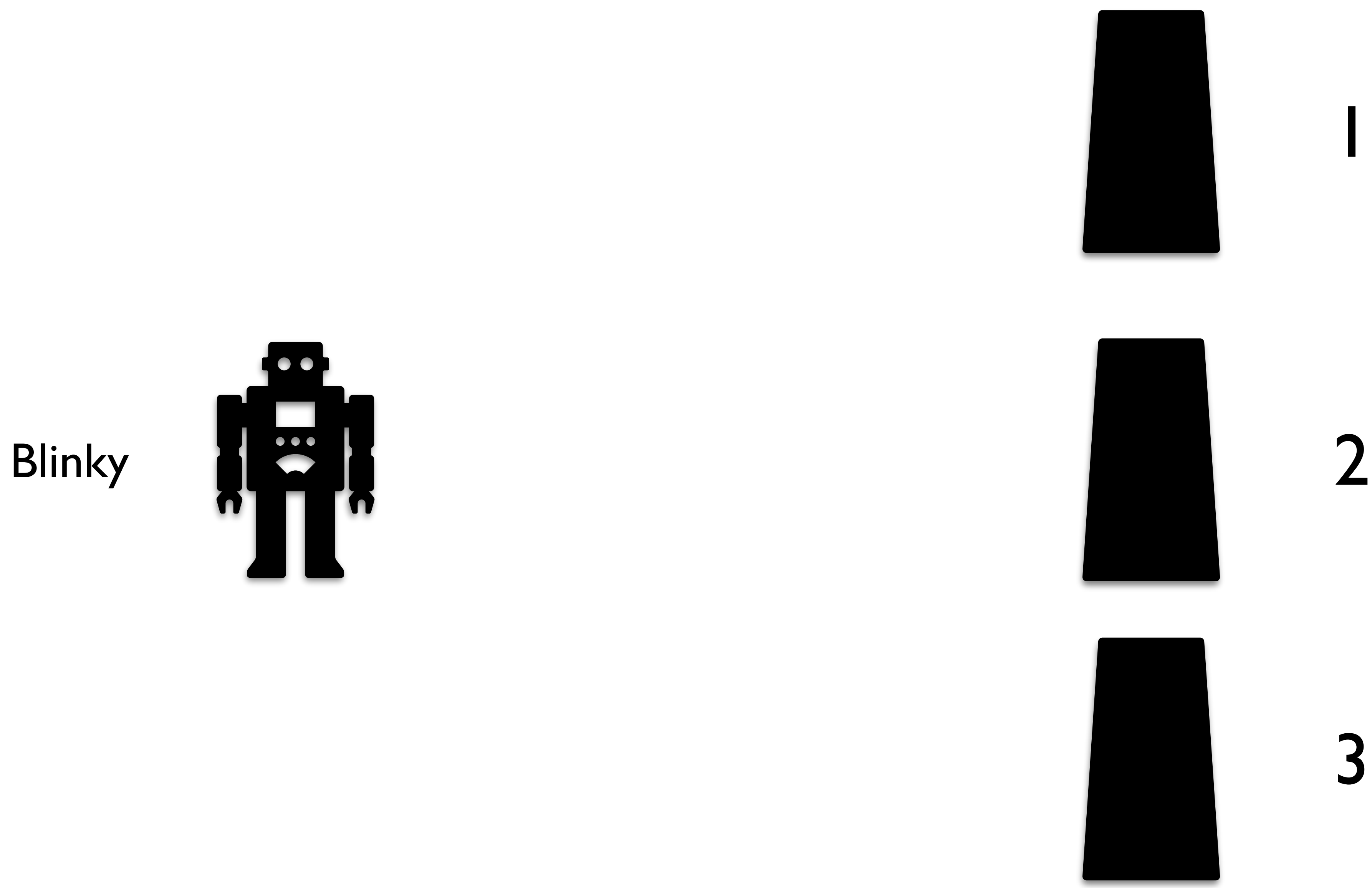
B(blinky, loc-ball-1)



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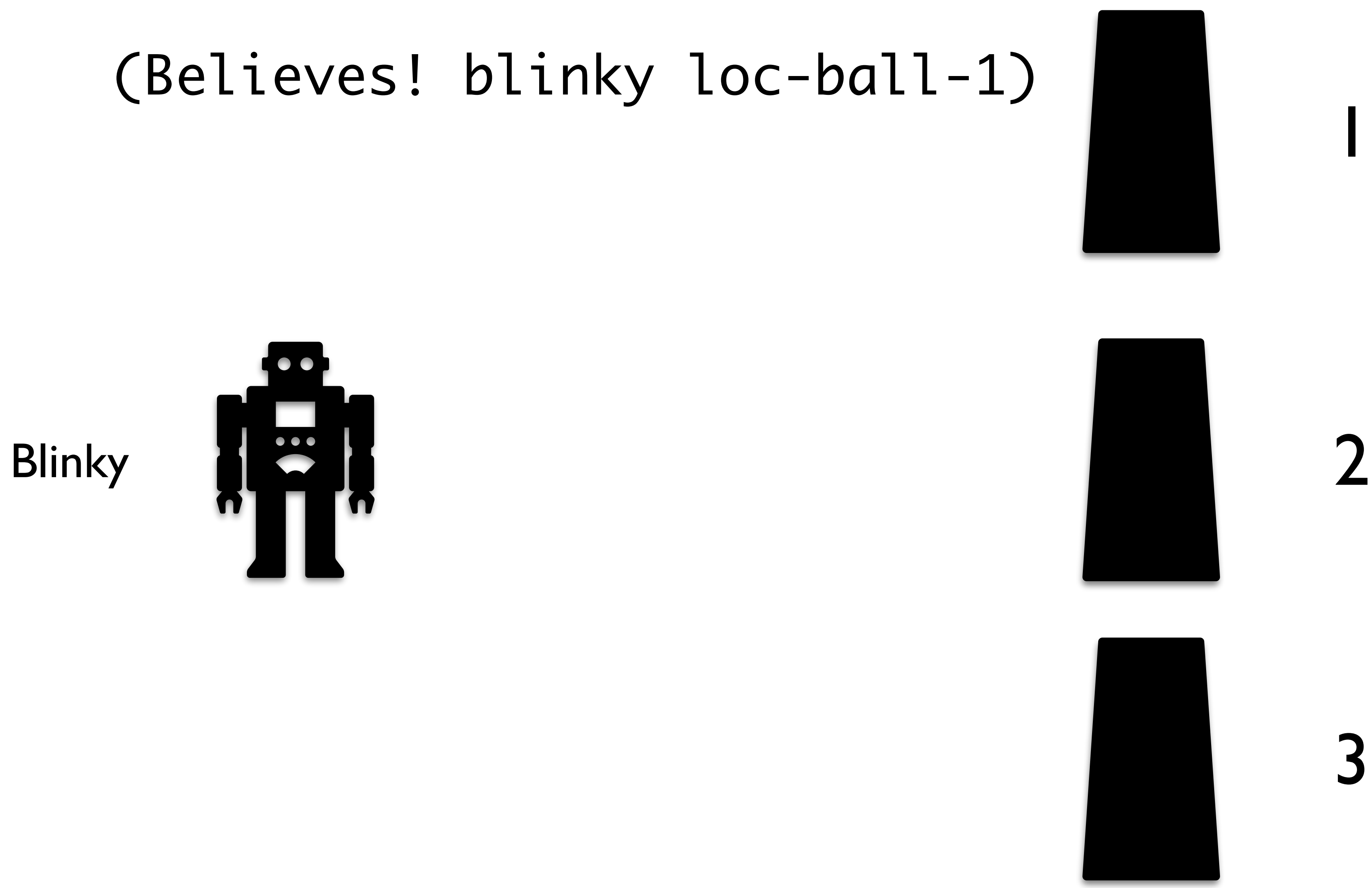
B(blinky, loc-ball-1)

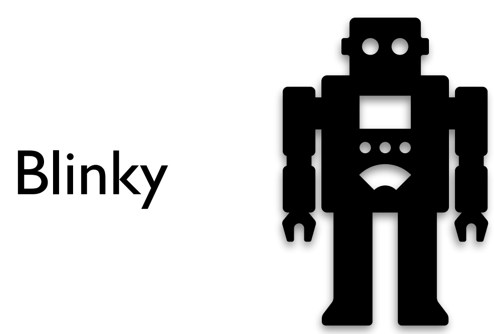
`(Believes! blinky loc-ball-1)`

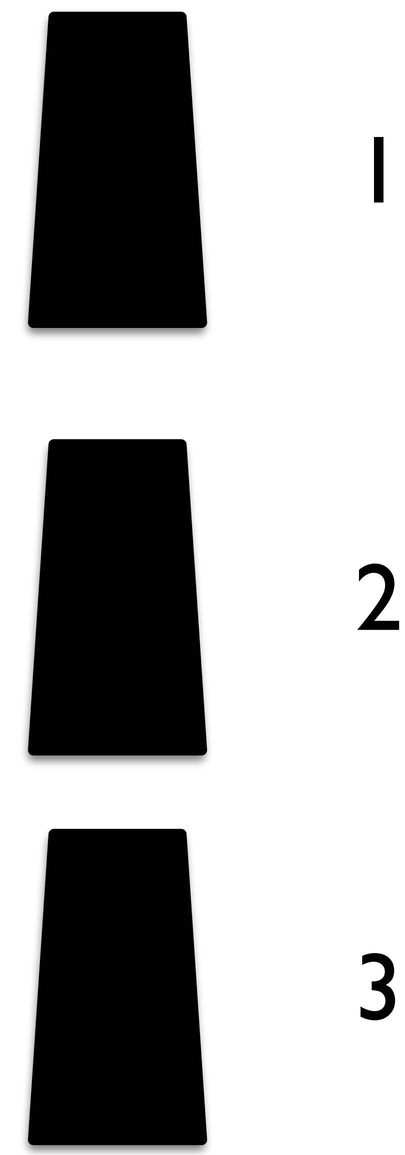


Blinky believes that the ball is in the cup at location #1.

B(blinky, loc-ball-1)

```
(Believes! blinky loc-ball-1)
```





In intensional logics, meaning and designation are separated, and compositionality is abandoned.

Blinky believes that the ball is in the cup at location #1.

B(blinky, loc-ball-1)

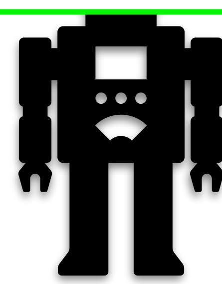
```
(Believes! blinky loc-ball-1)
```

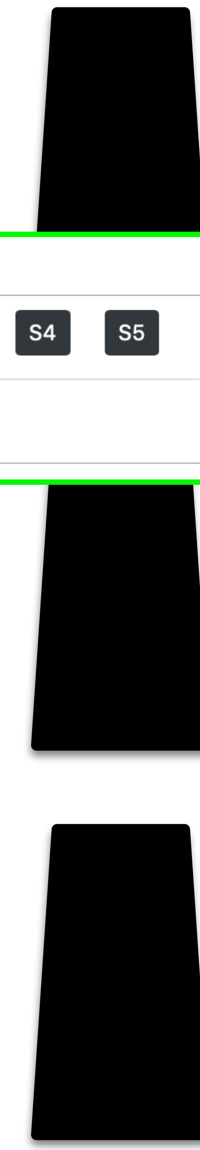
Create file

Propositional Calculus | L_0 = Pure Predicate Calculus | L_1 = First-order Logic | L_2 = Second-order Logic | K | T | D | S4 | S5

DCEC (fragment) | Hyperlog

Blinky





1

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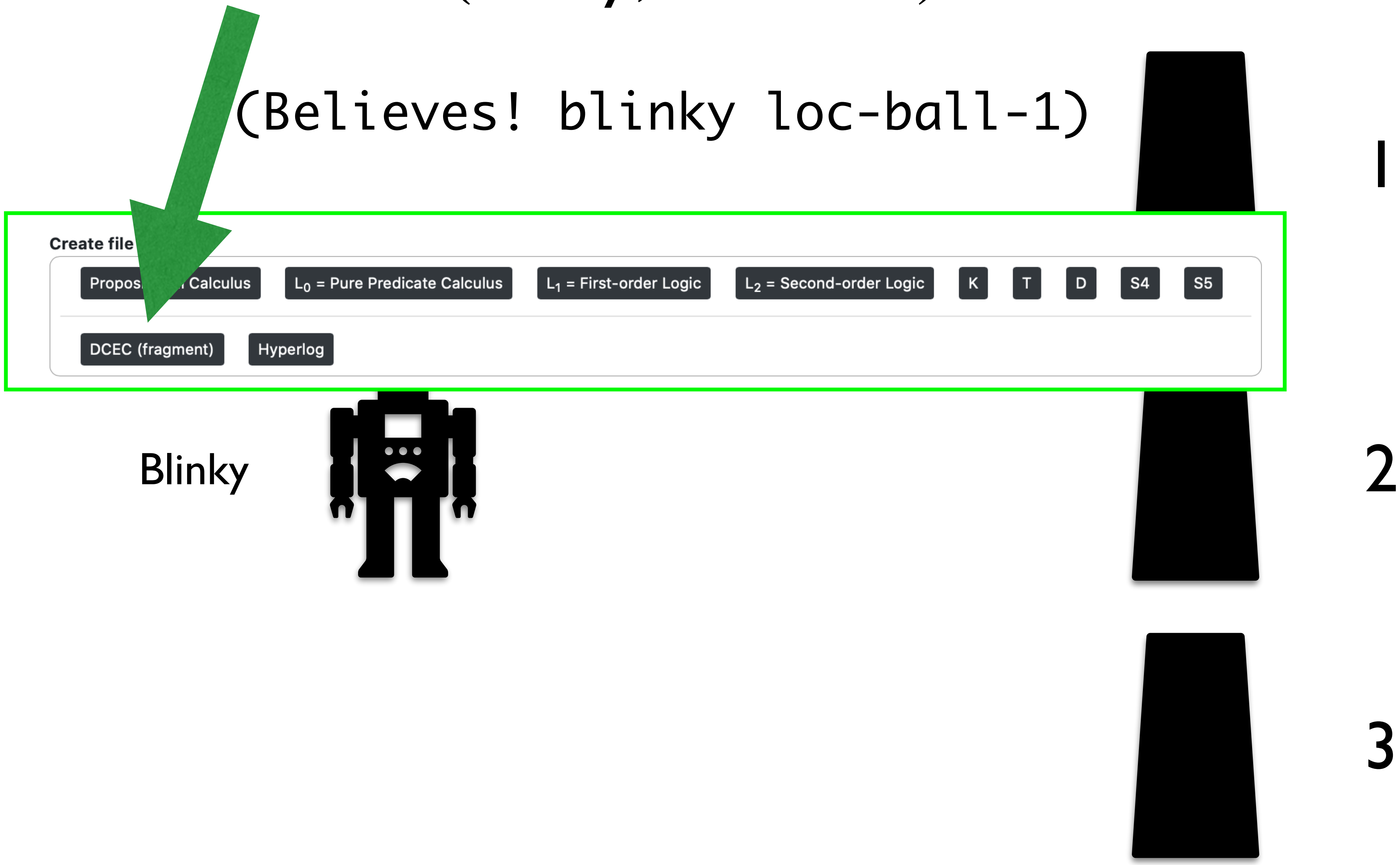
In intensional logics, meaning and designation are separated, and compositionality is abandoned.

Blinky believes that the ball is in the cup at location #1.

B(blinky, loc-ball-1)

`(Believes! blinky loc-ball-1)`



In intensional logics, meaning and designation are separated, and compositionality is abandoned.

False Belief Task Demands Intensional Logic …

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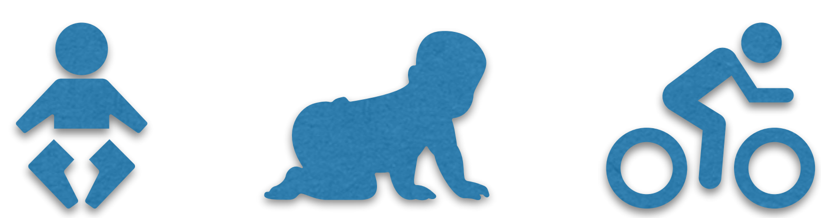
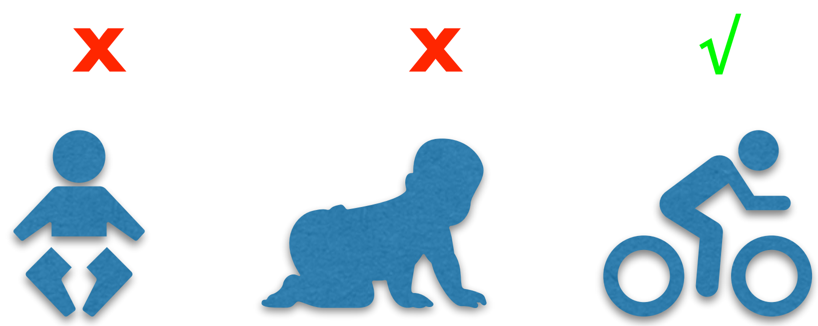
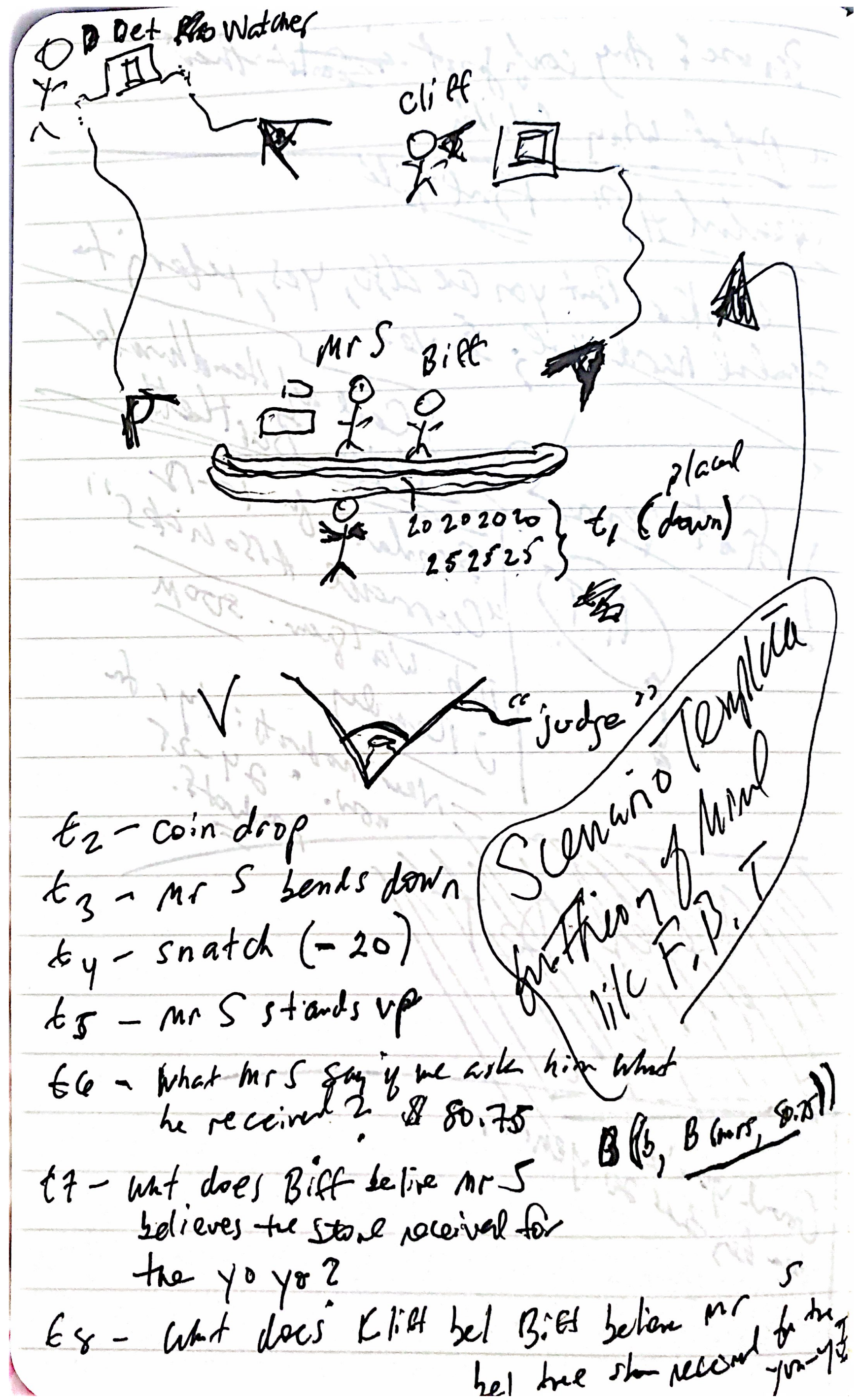


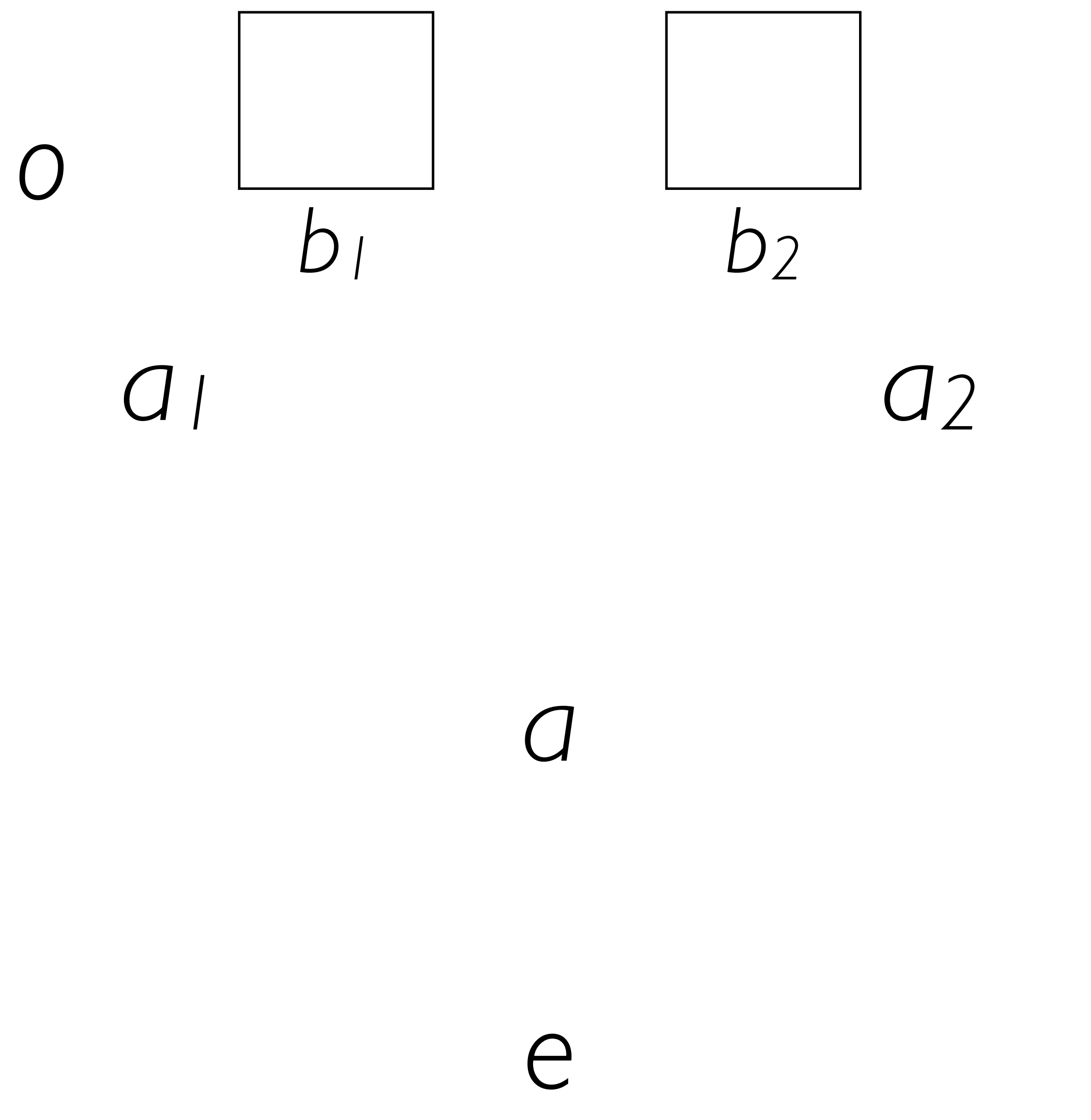
False Belief Task Demands Intensional Logic …



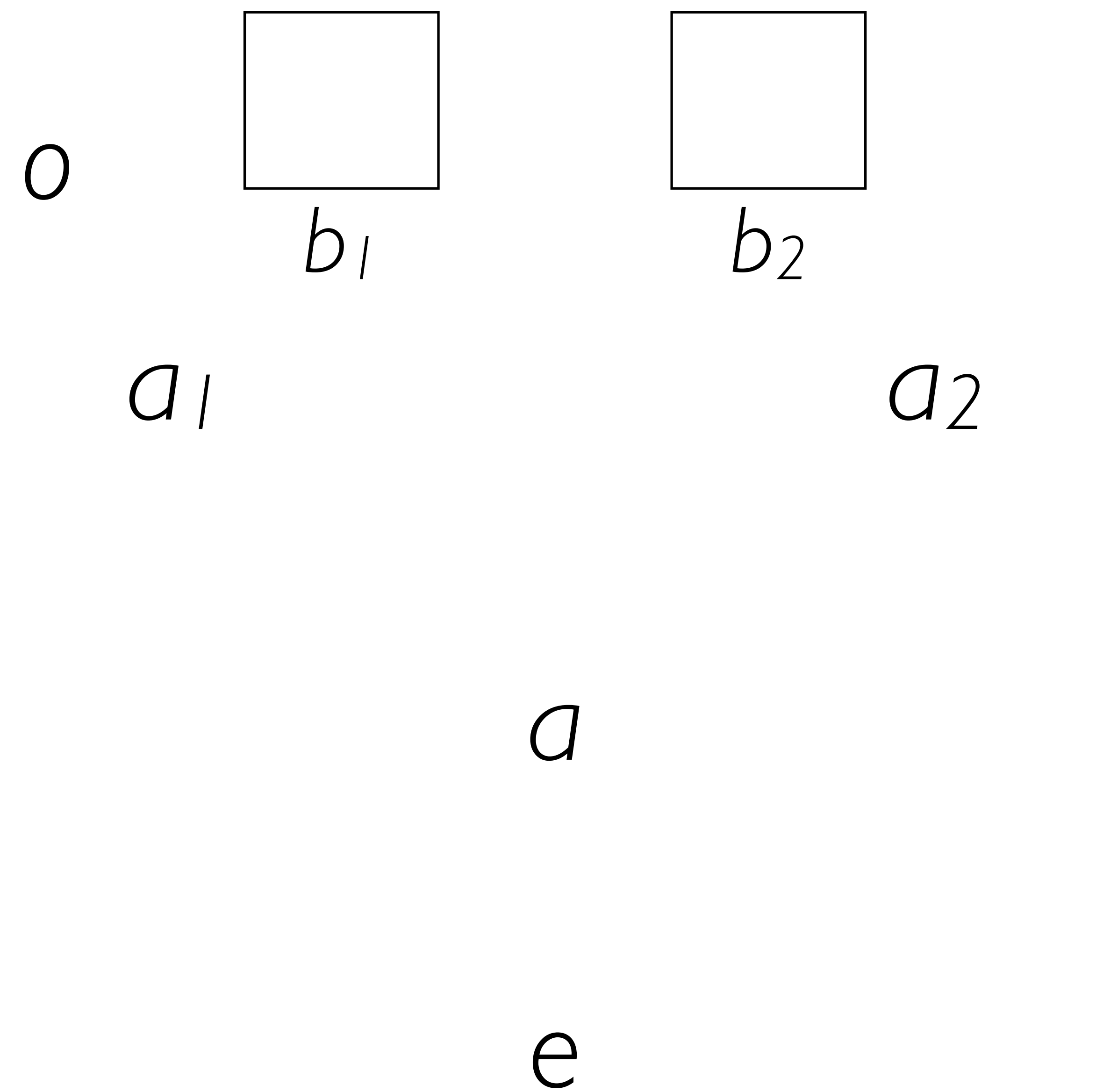
Better, But Embryonic: The ToM Pawn Shop



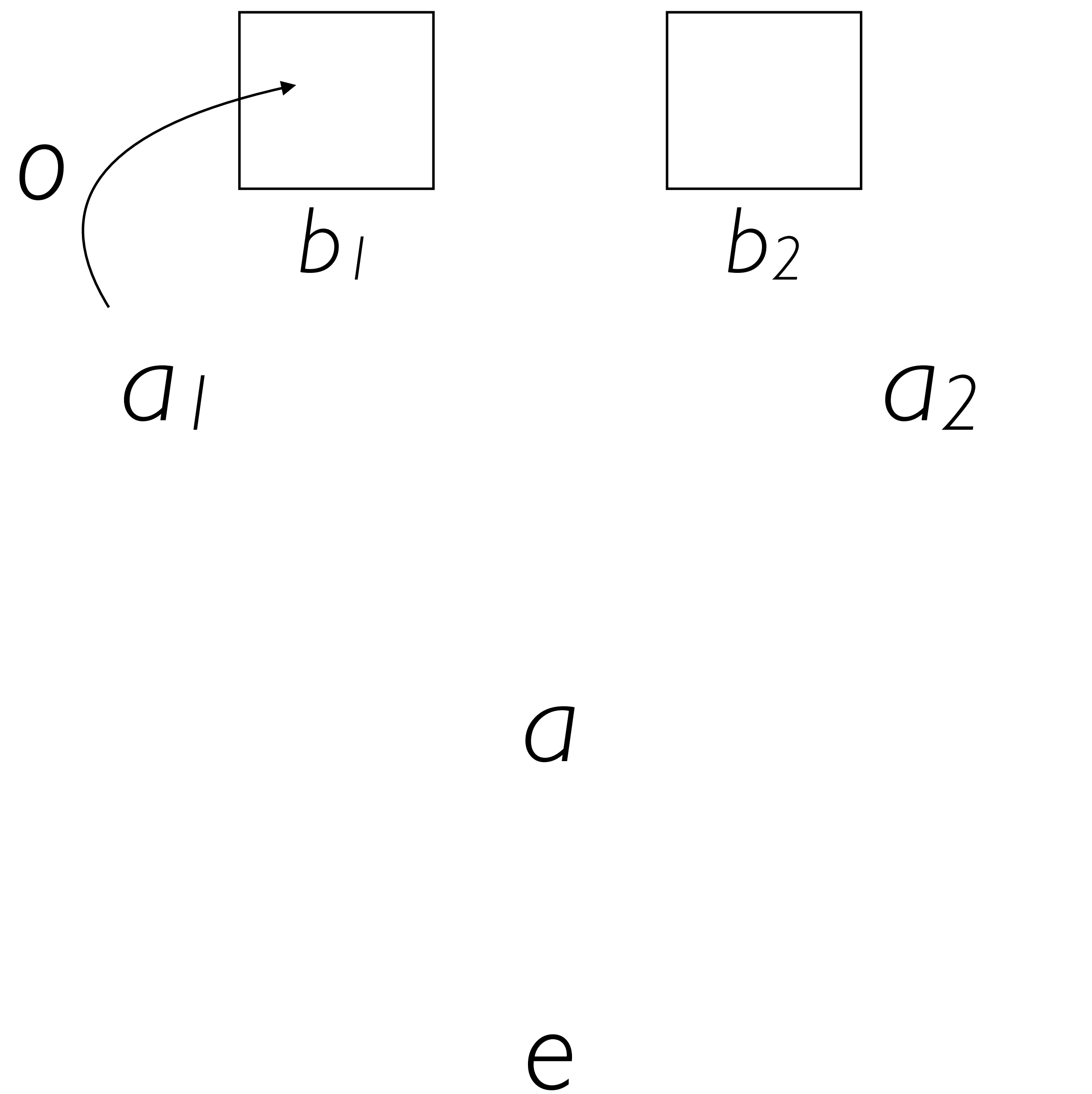
Framework for FBT^0_1



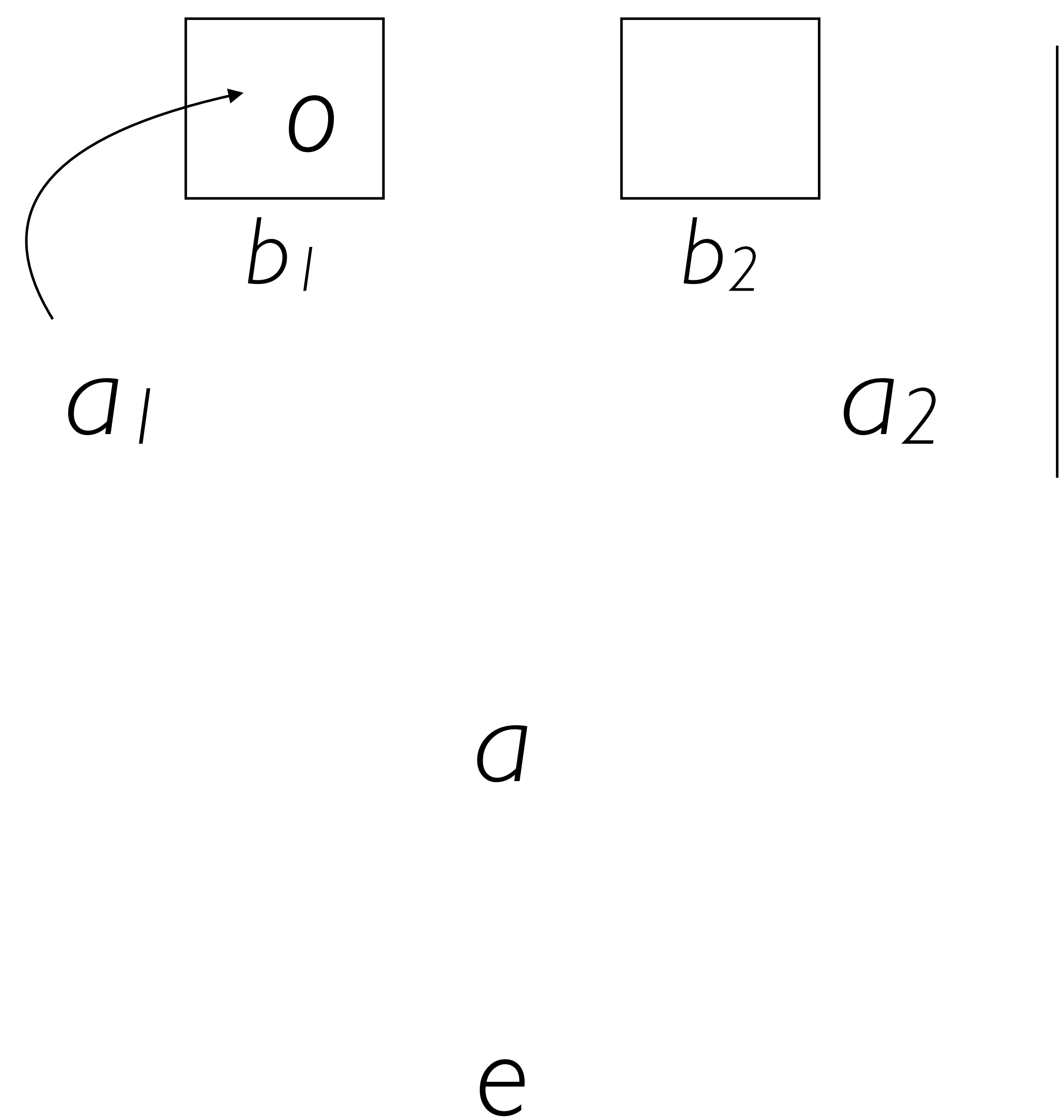
Framework for FBT^0_1
(five timepoints)



Framework for FBT^{0}_{1}
(five timepoints)

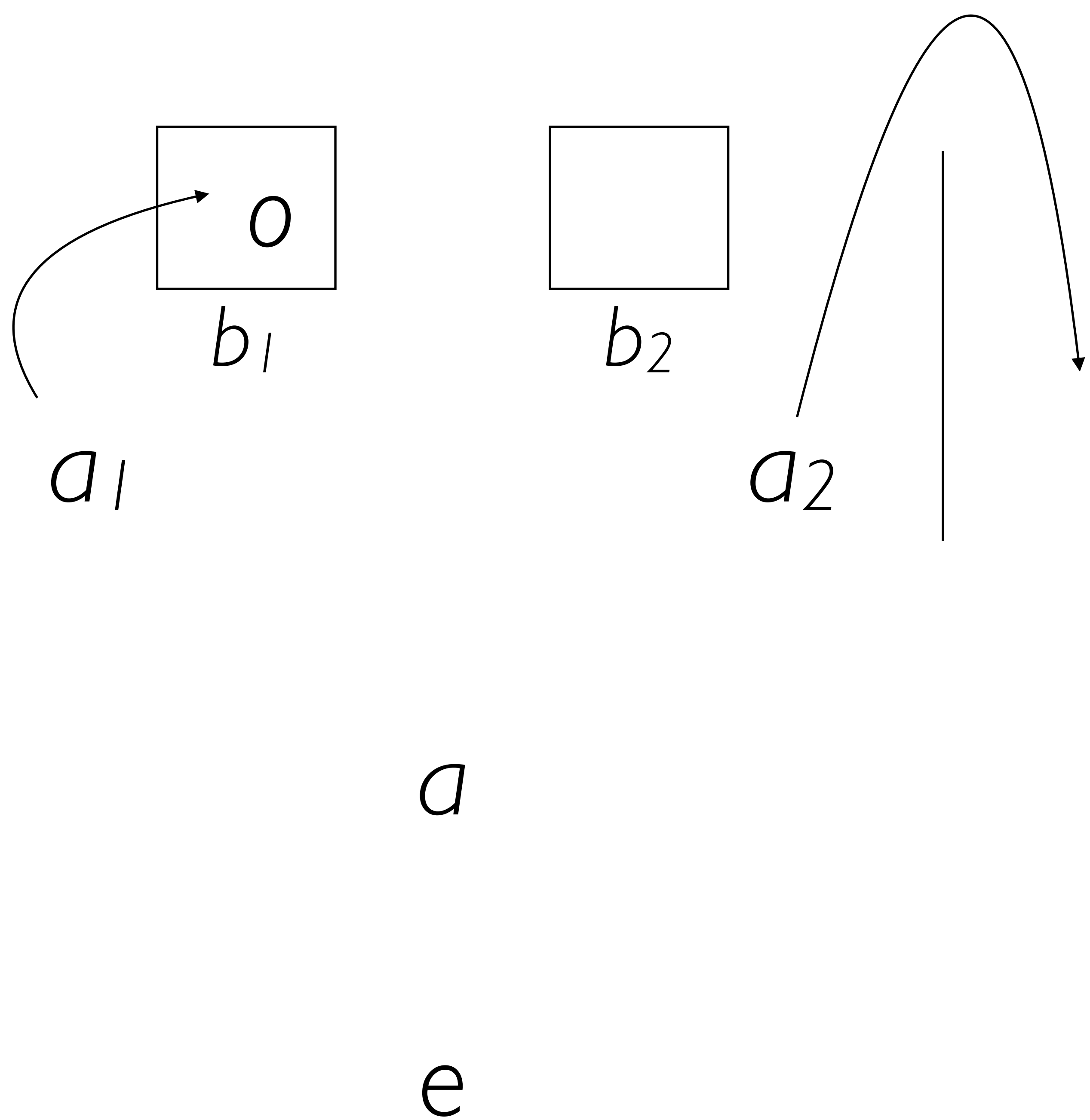


Framework for $FBT^0{}_1$
(five timepoints)

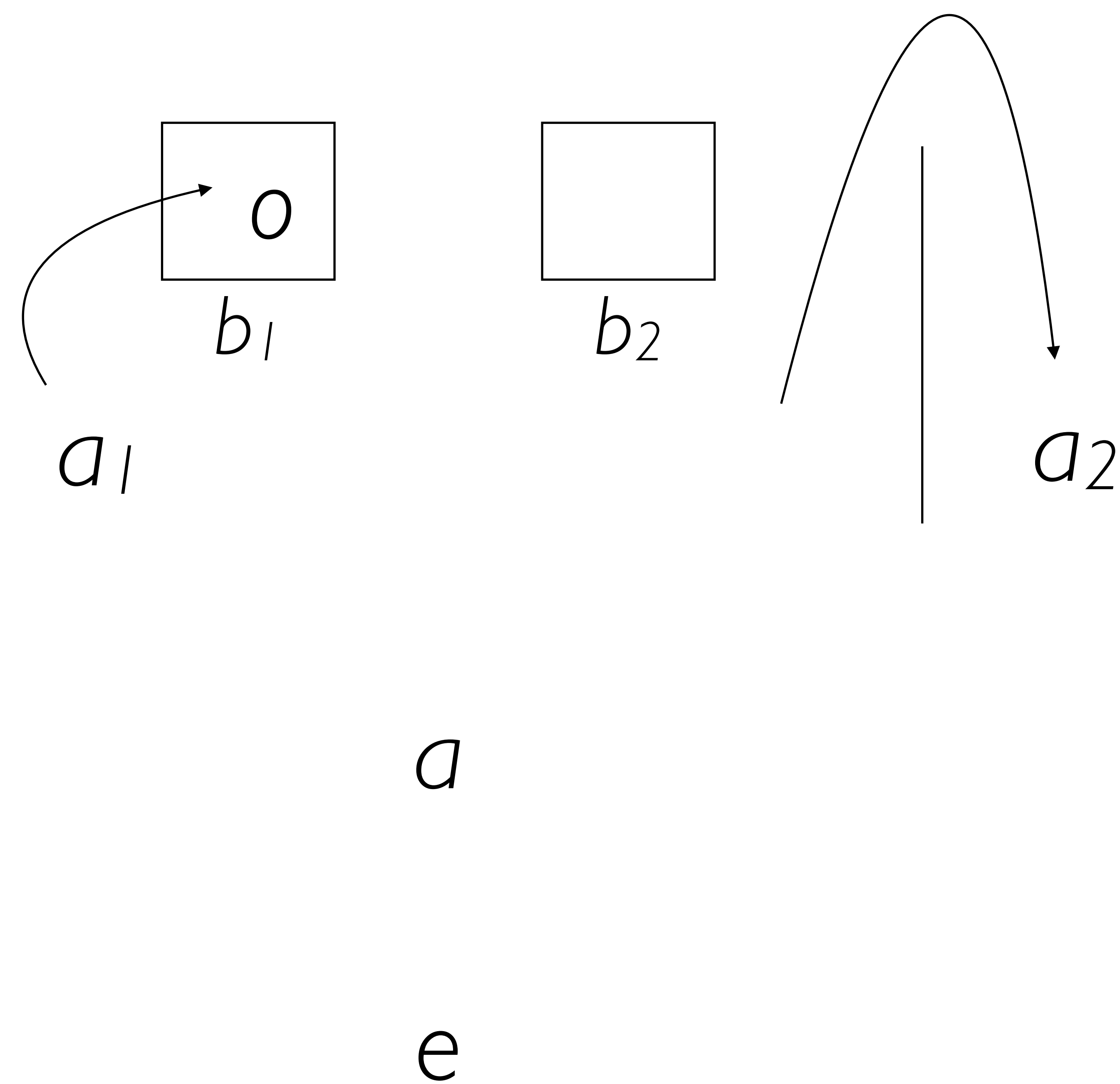


Framework for FBT$^0{}_1$

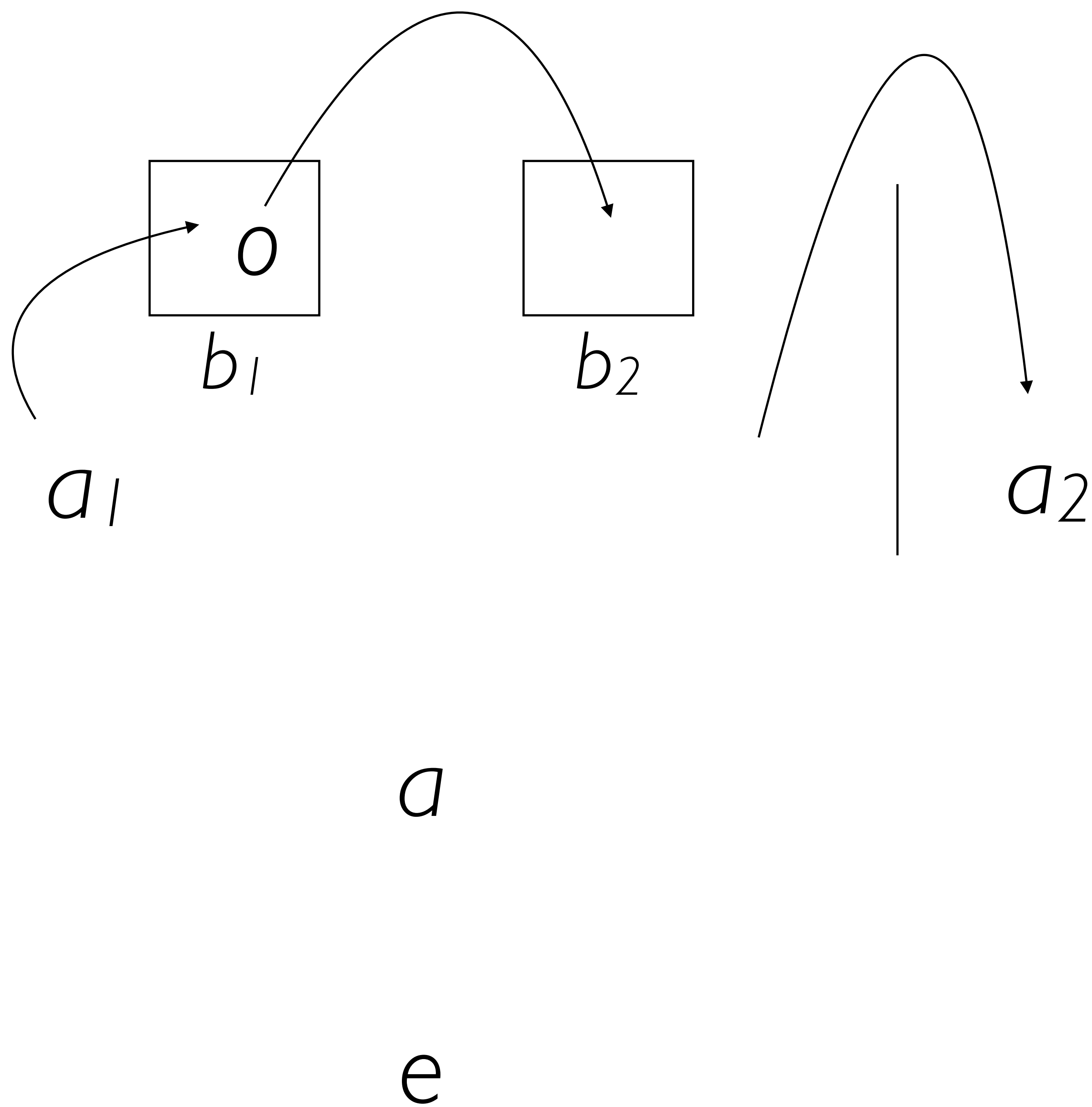
(five timepoints)



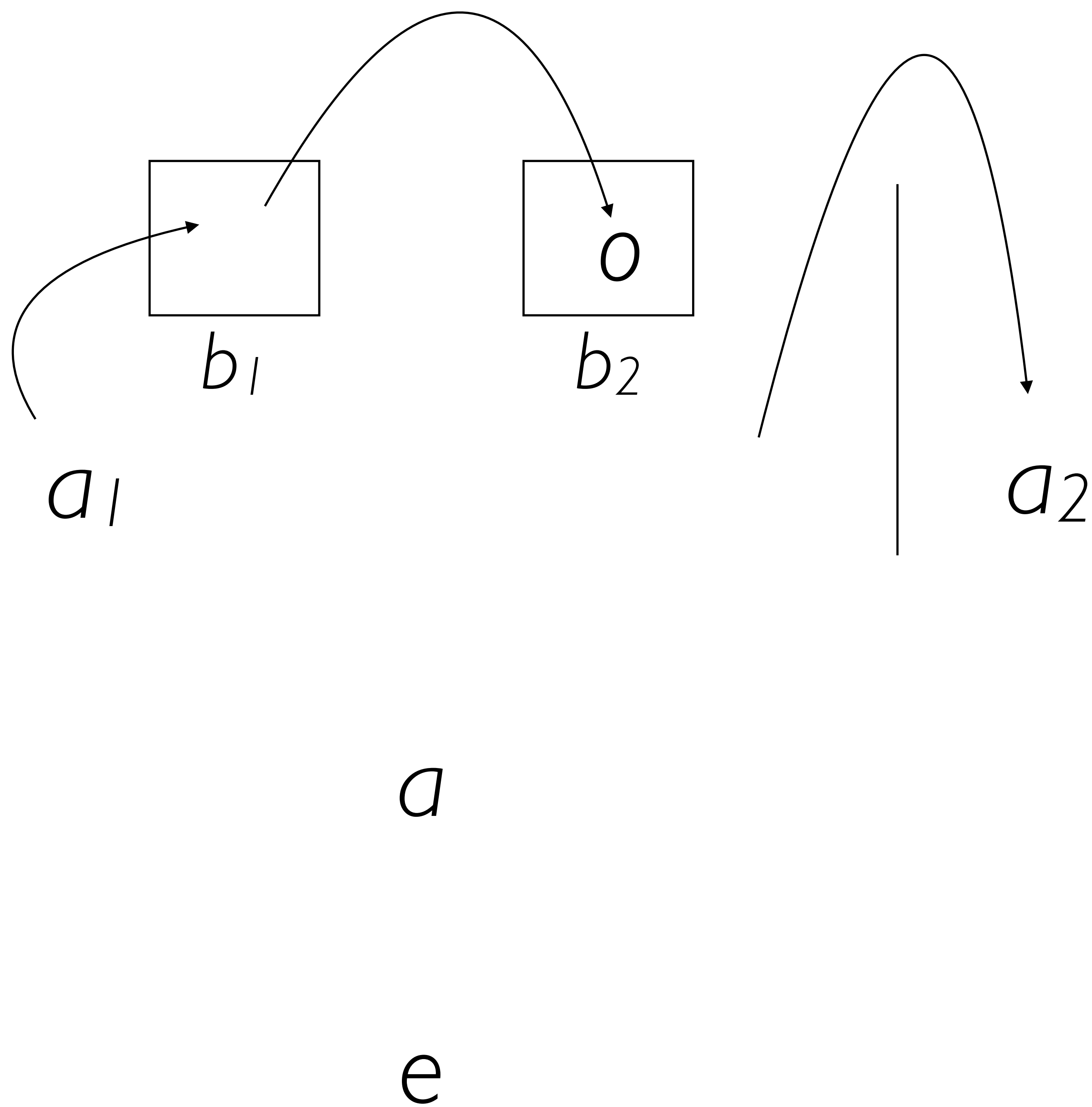
Framework for $FBT^0{}_1$
(five timepoints)



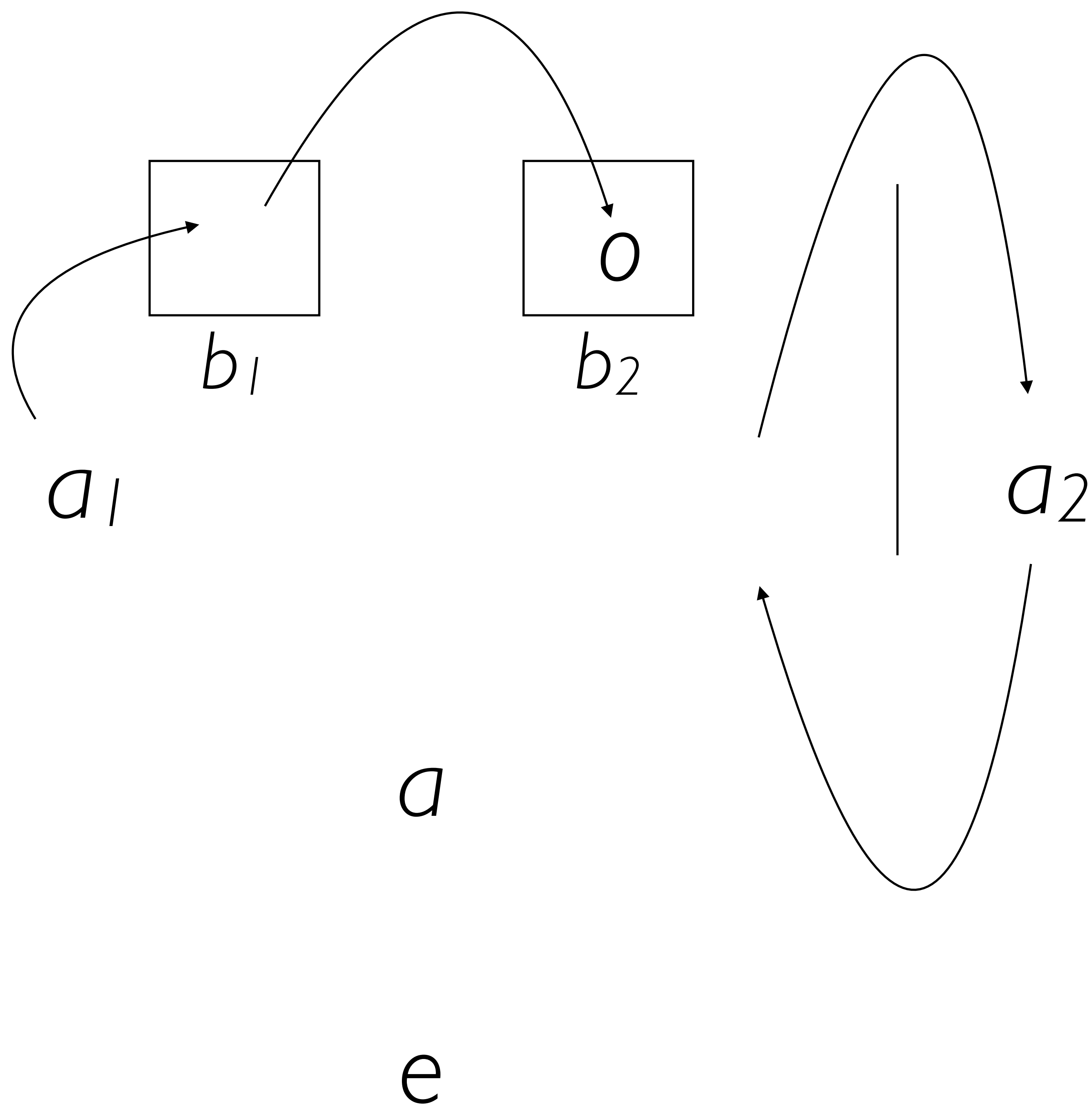
Framework for $FBT^0{}_1$
(five timepoints)



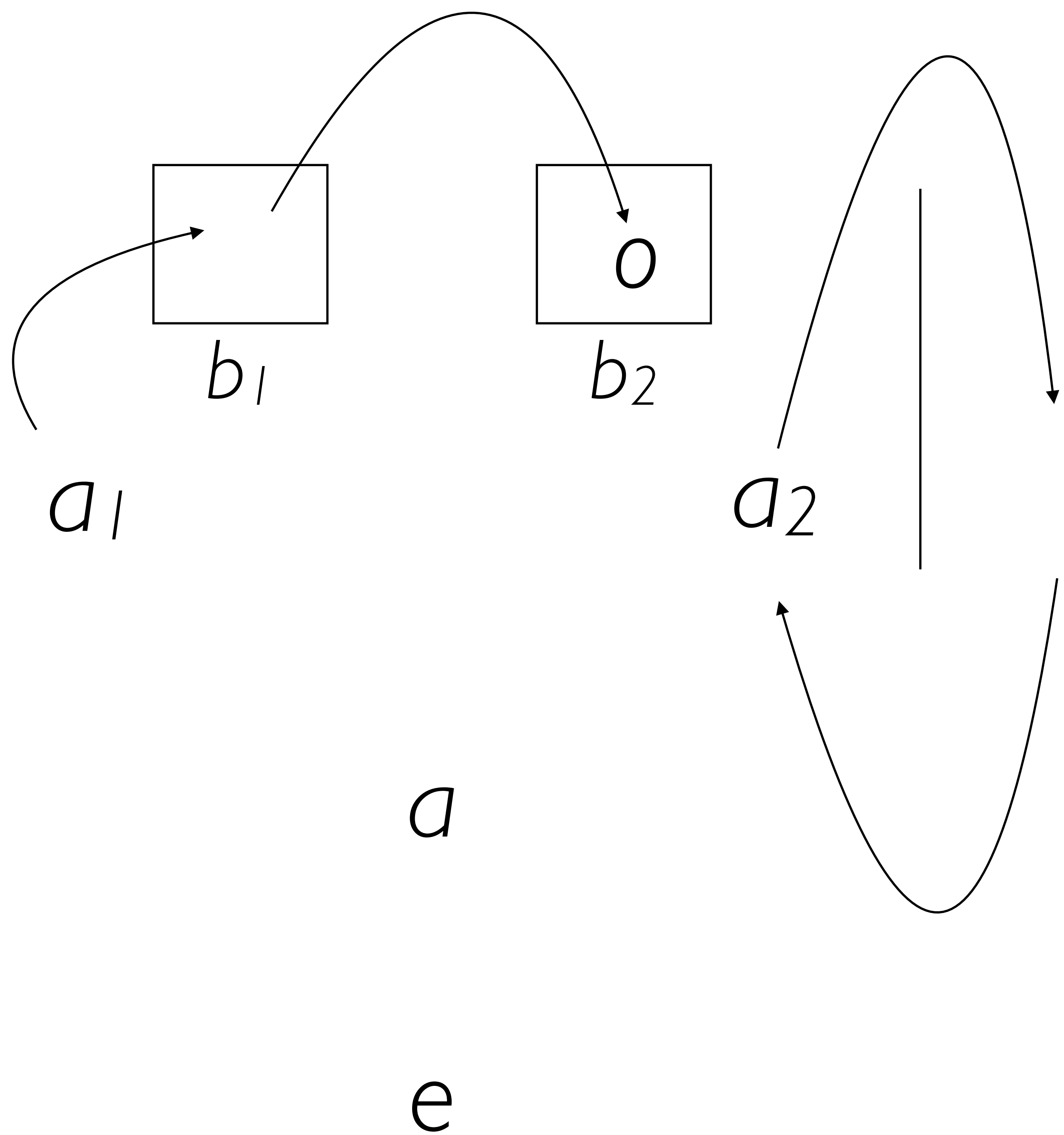
Framework for FBT$^0_1$
(five timepoints)



Framework for $FBT^0{}_1$
(five timepoints)

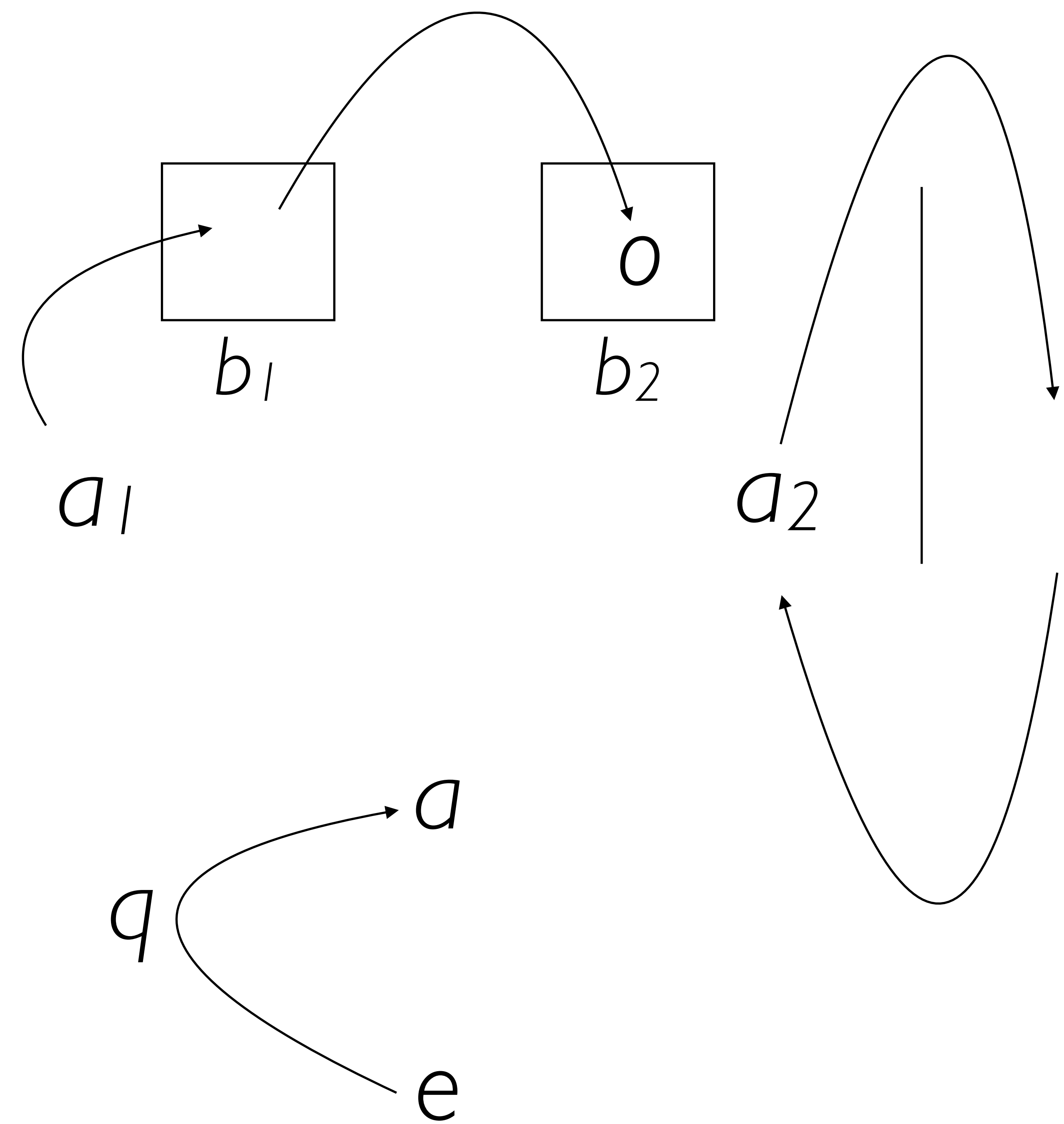


Framework for $FBT^0{}_1$
(five timepoints)



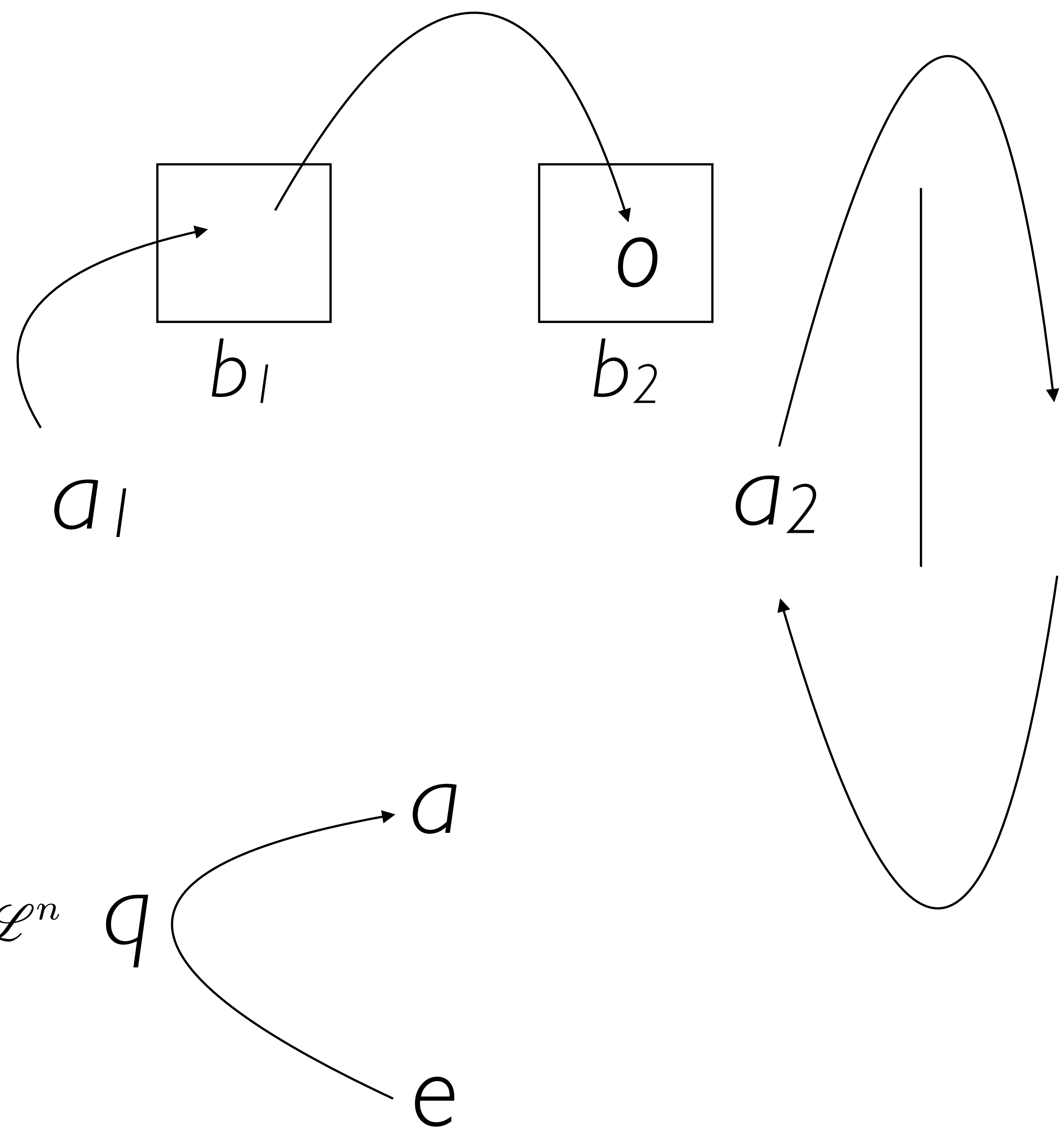
Framework for FBT$^0{}_1$

(five timepoints)



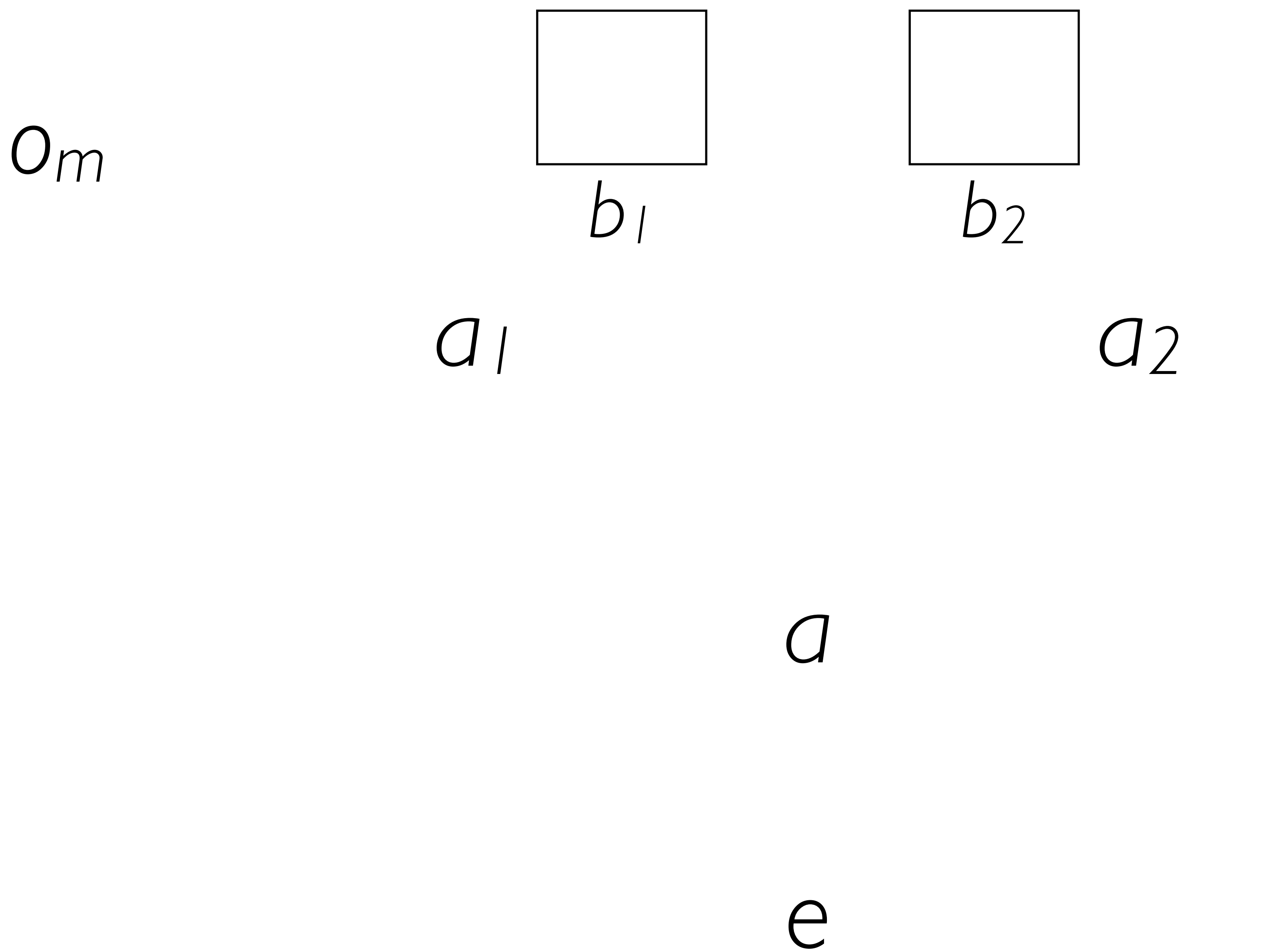
Framework for $FBT^0{}_1$

(five timepoints)



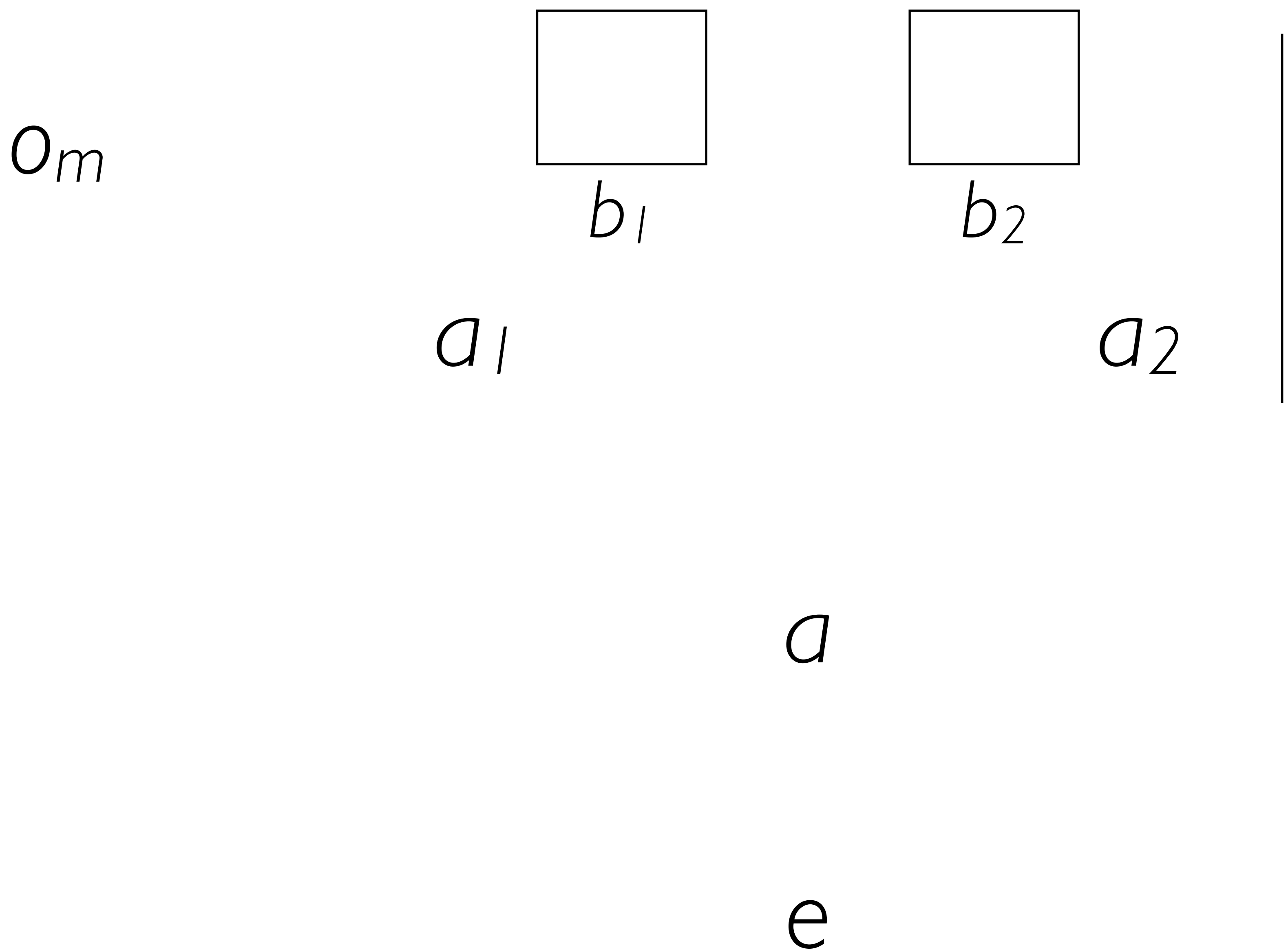
q a formula in modal $\mathscr{L}^n$

Framework for FBT^1_1

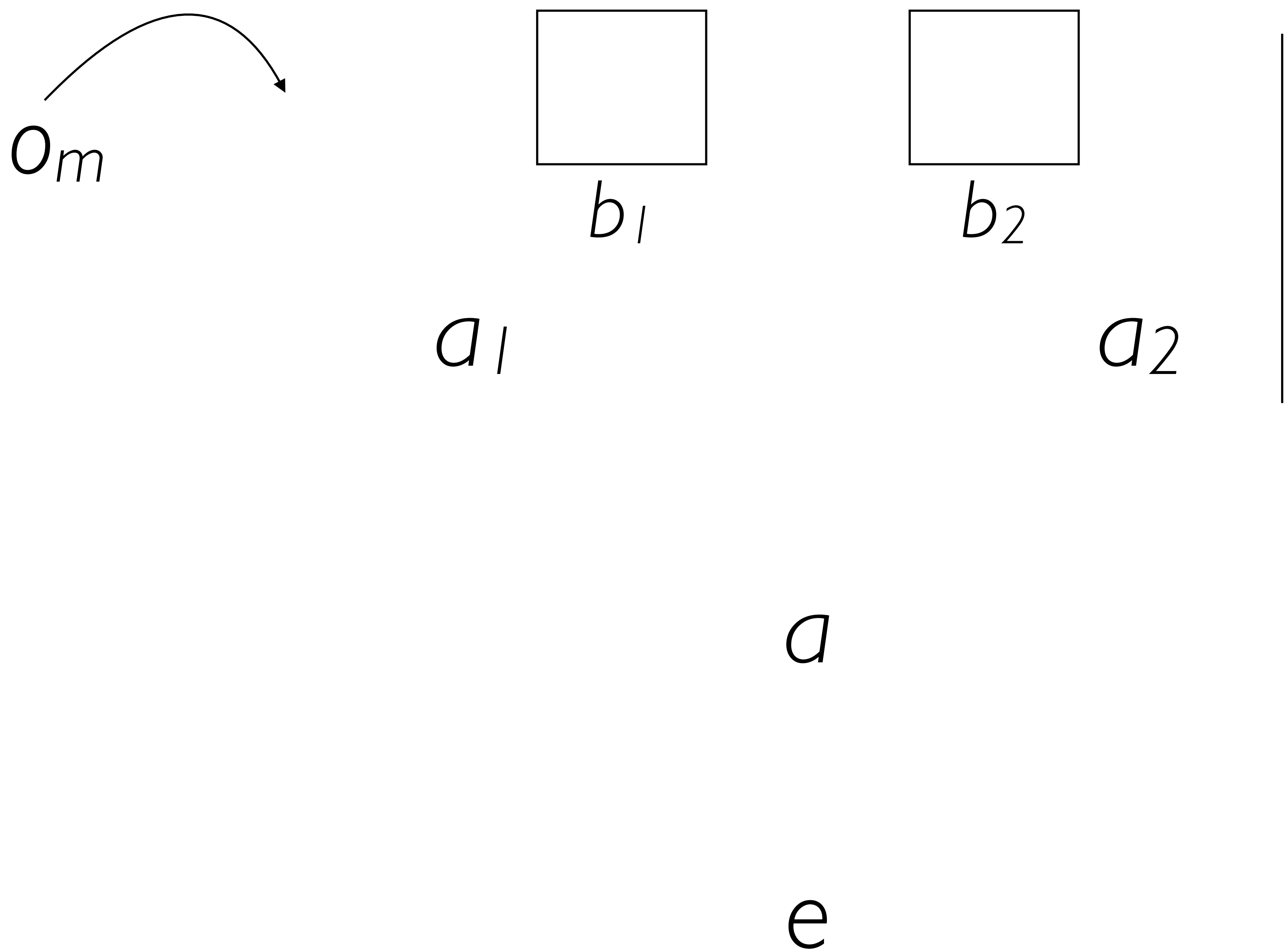


Framework for FBT^1_1

(six timepoints)

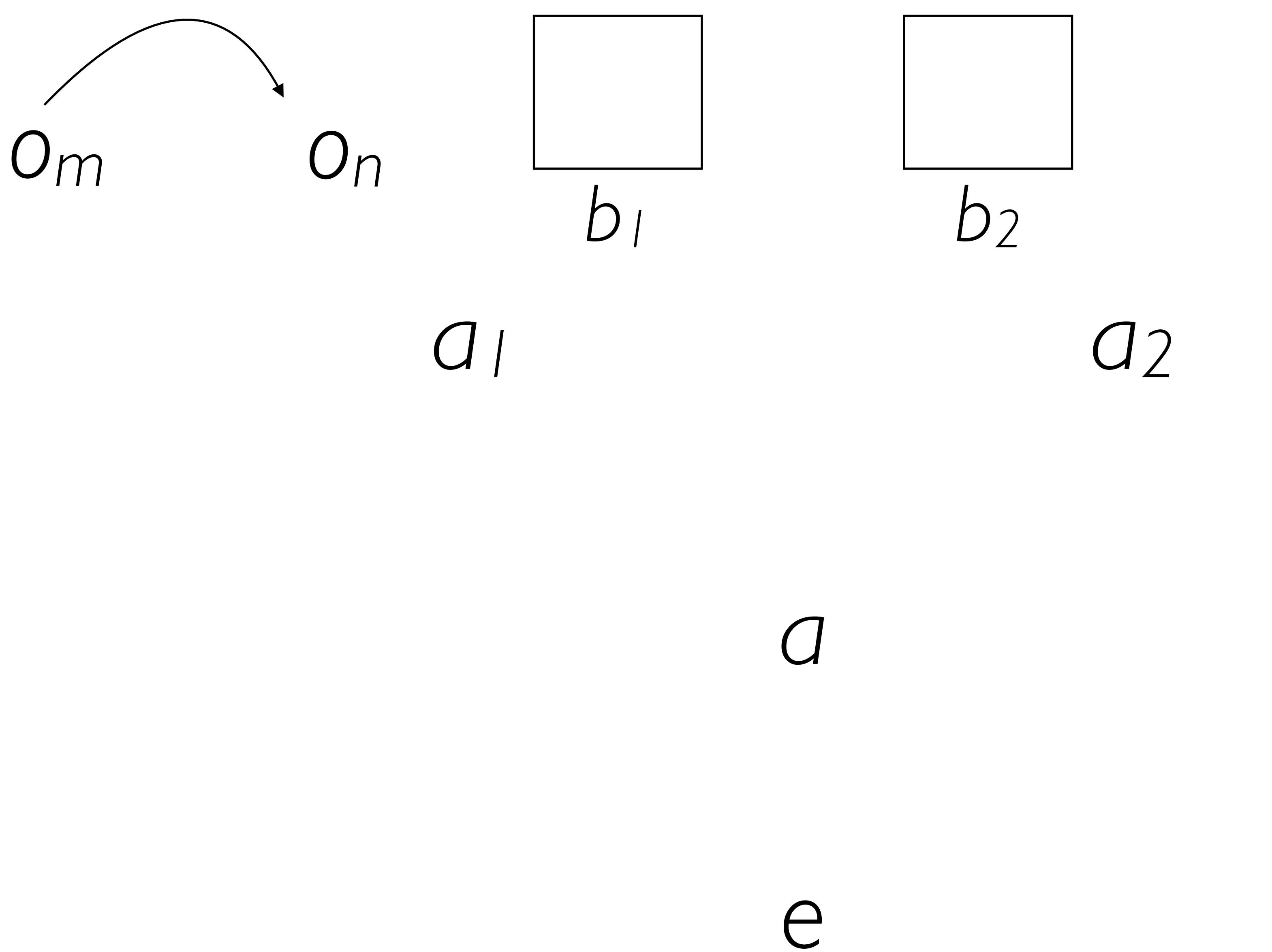


Framework for FBT^{1}_{1}
(six timepoints)



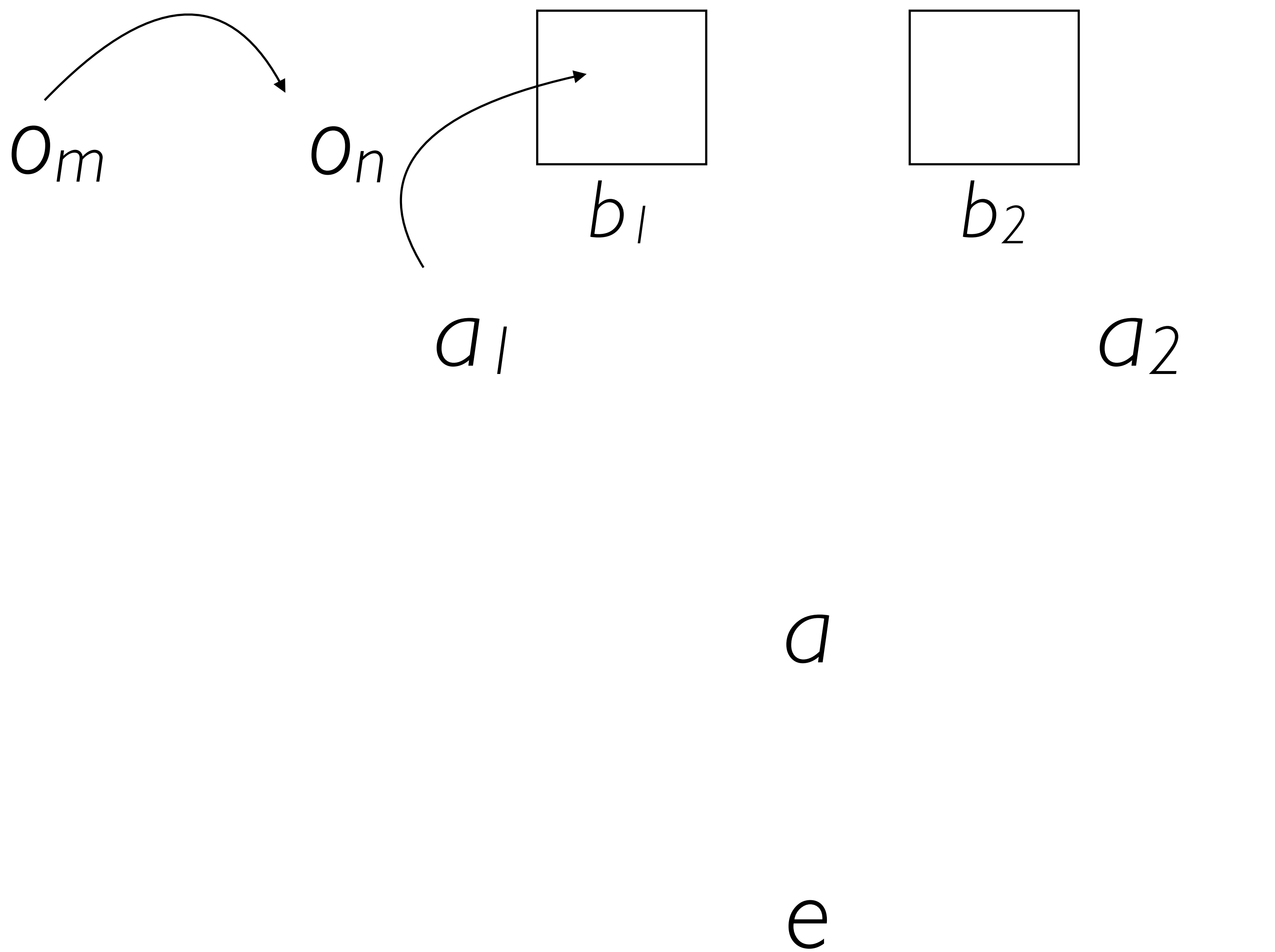
Framework for FBT^{1}_{1}

(six timepoints)



Framework for $FBT^1{}_1$

(six timepoints)



Framework for $FBT^1{}_1$

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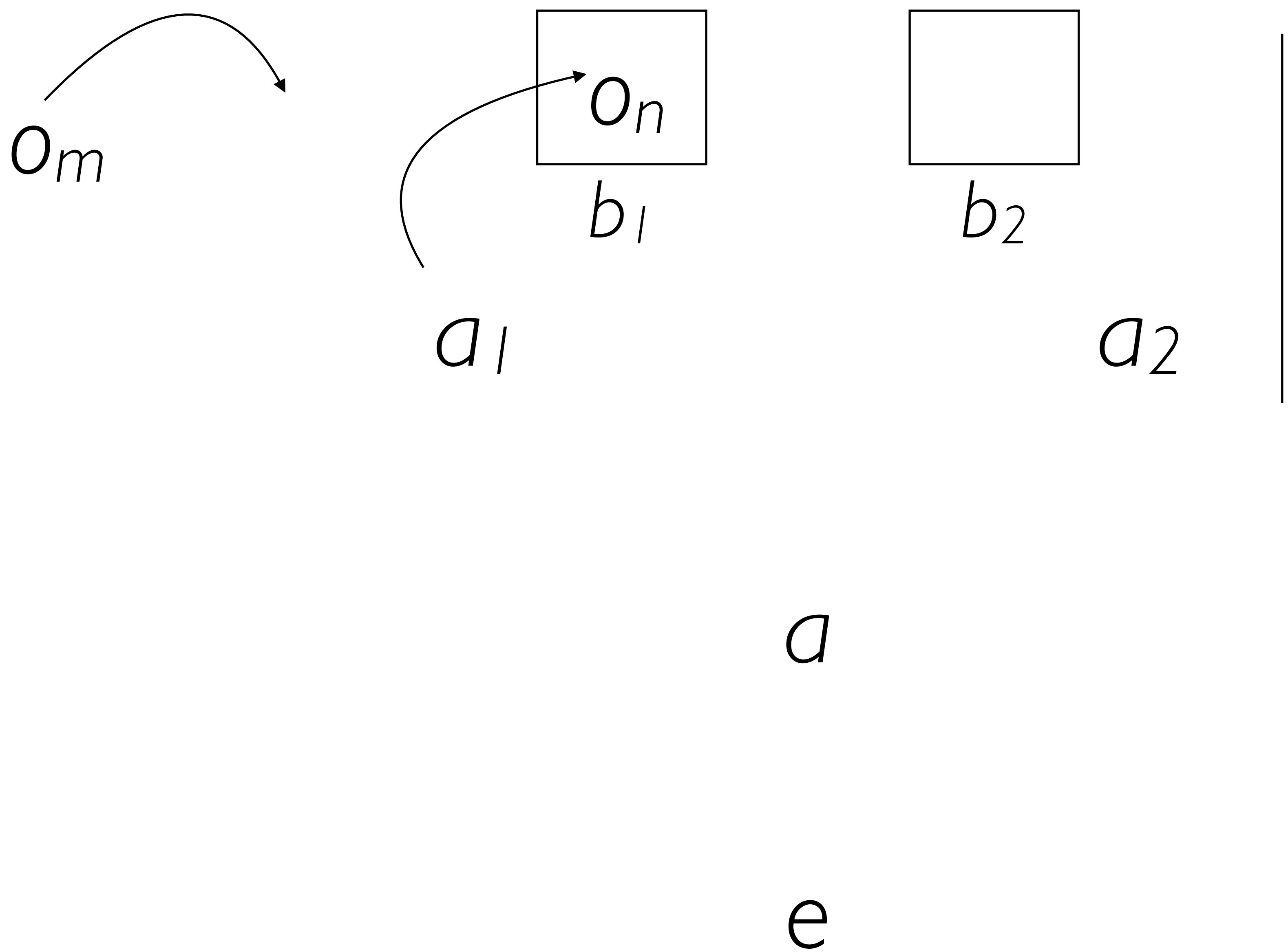

Framework for FBT^1_1

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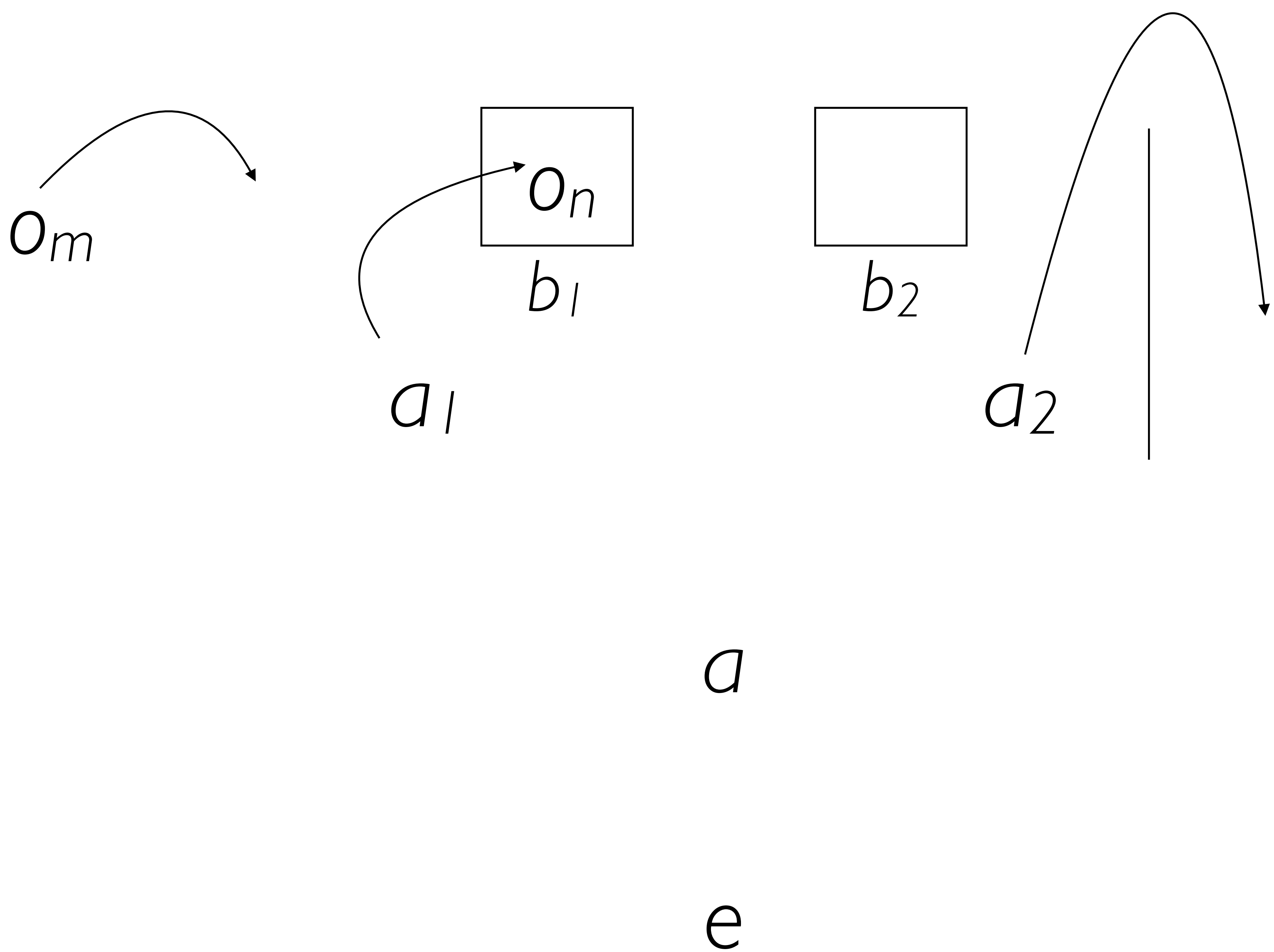

Framework for FBT^1_1
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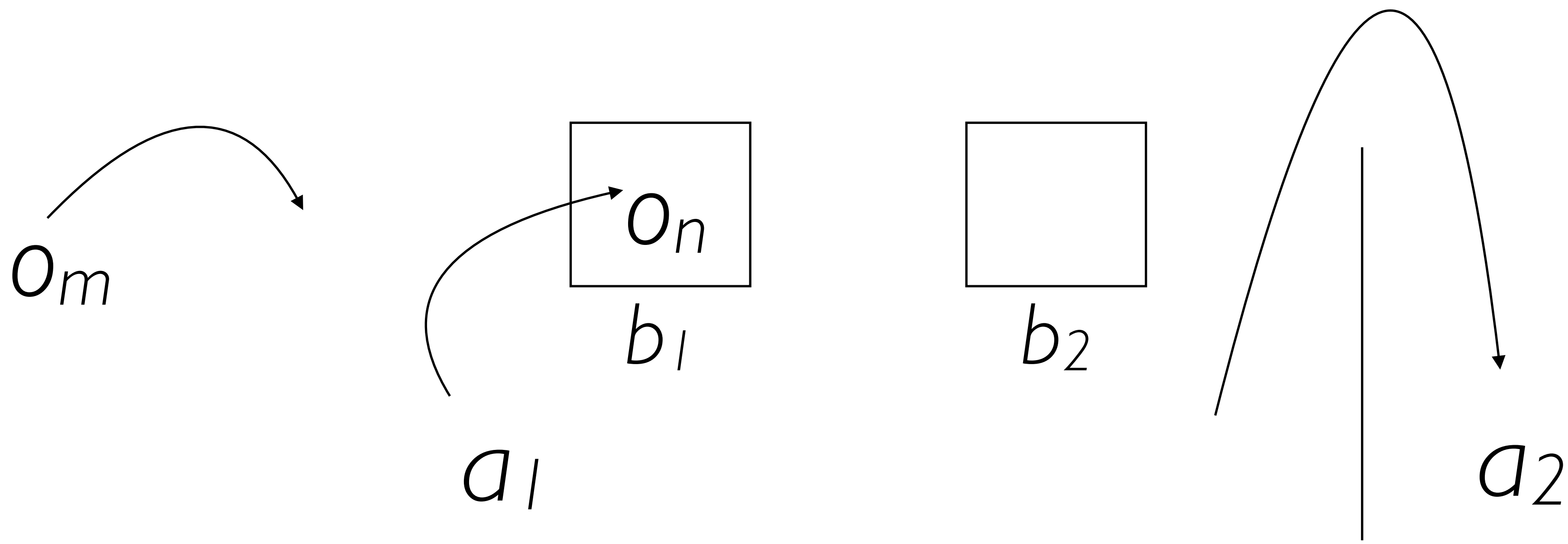

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Framework for FBT^1_1
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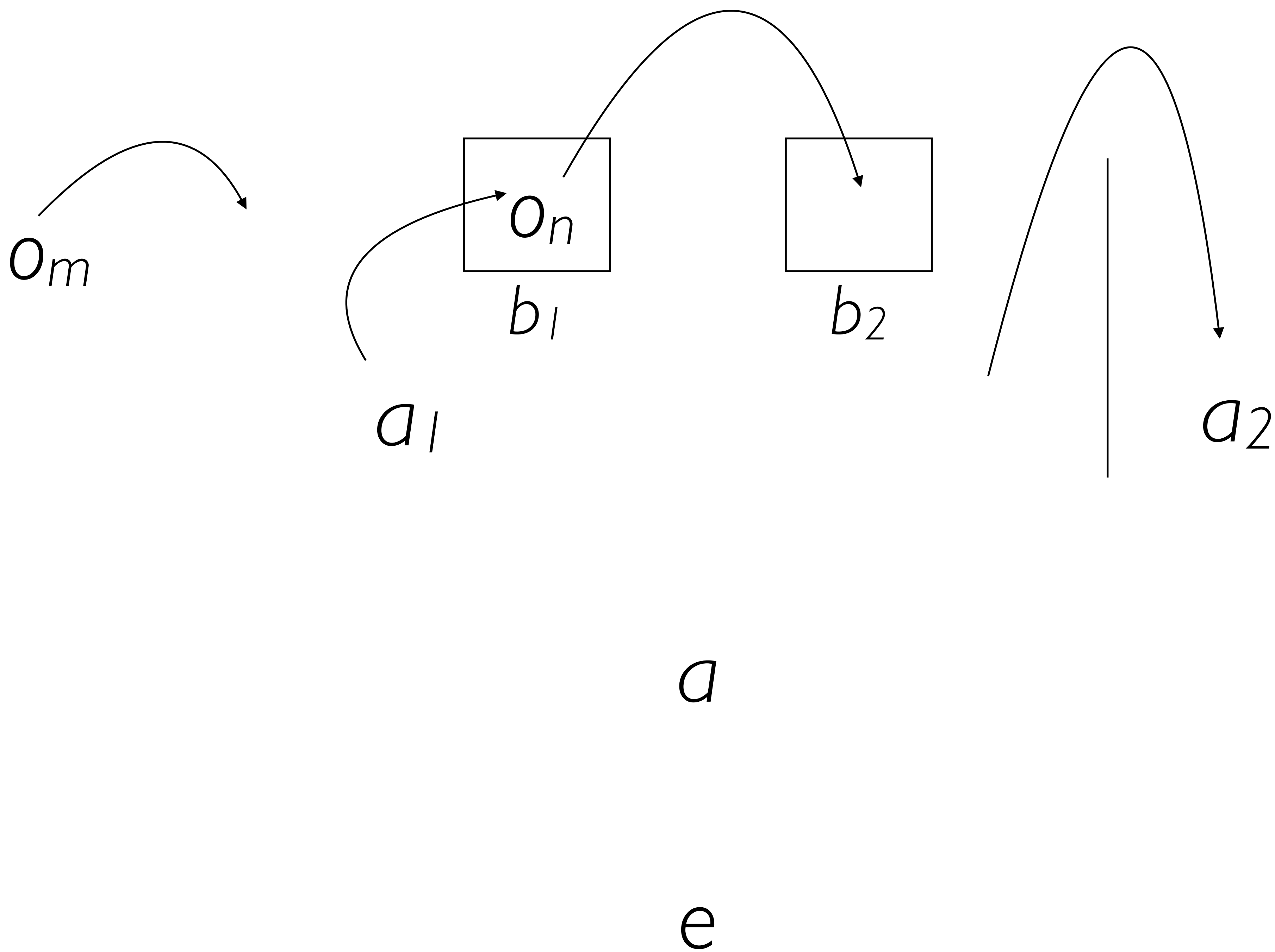

Framework for FBT¹₁
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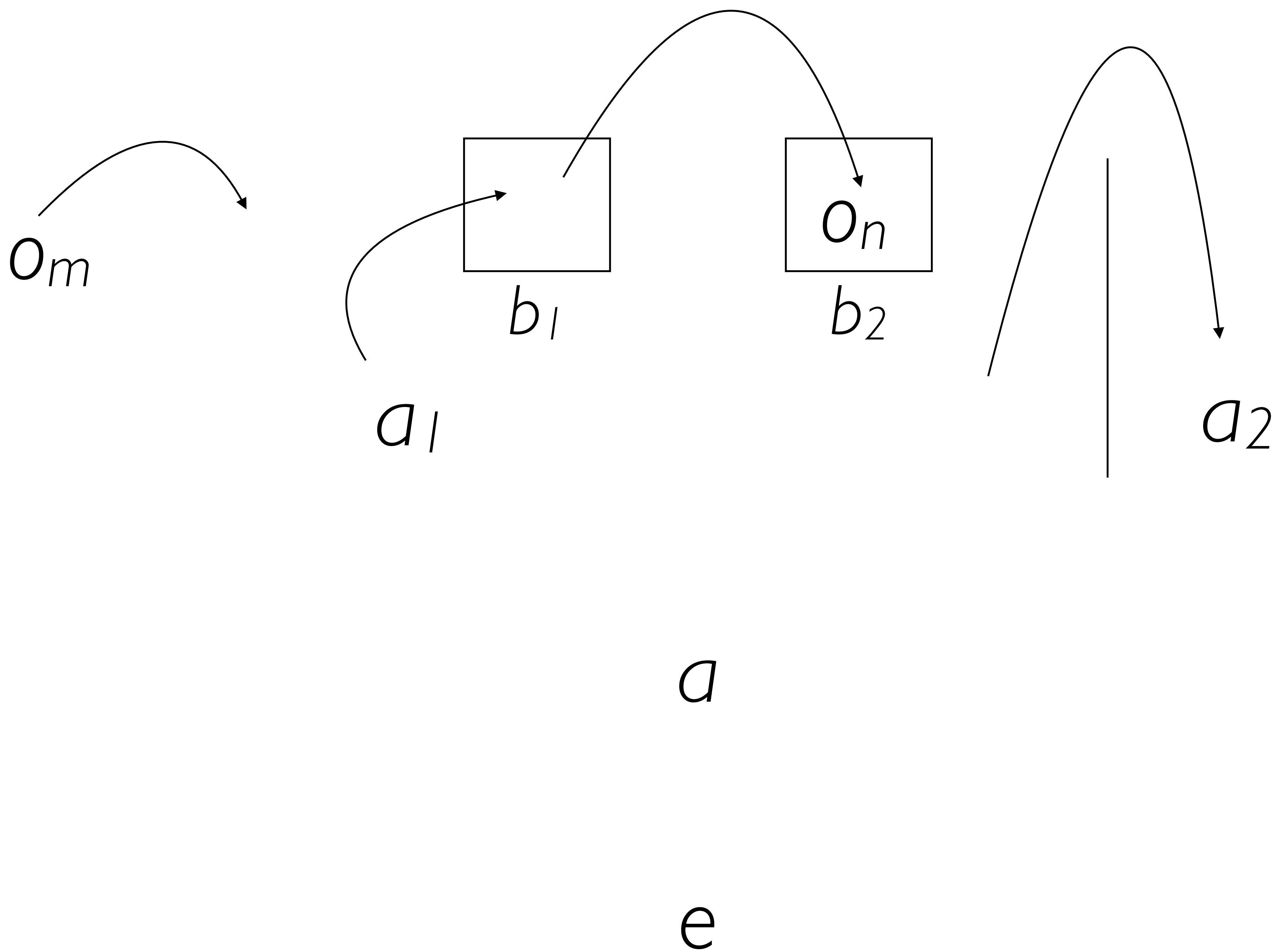

Framework for FBT1$_1$
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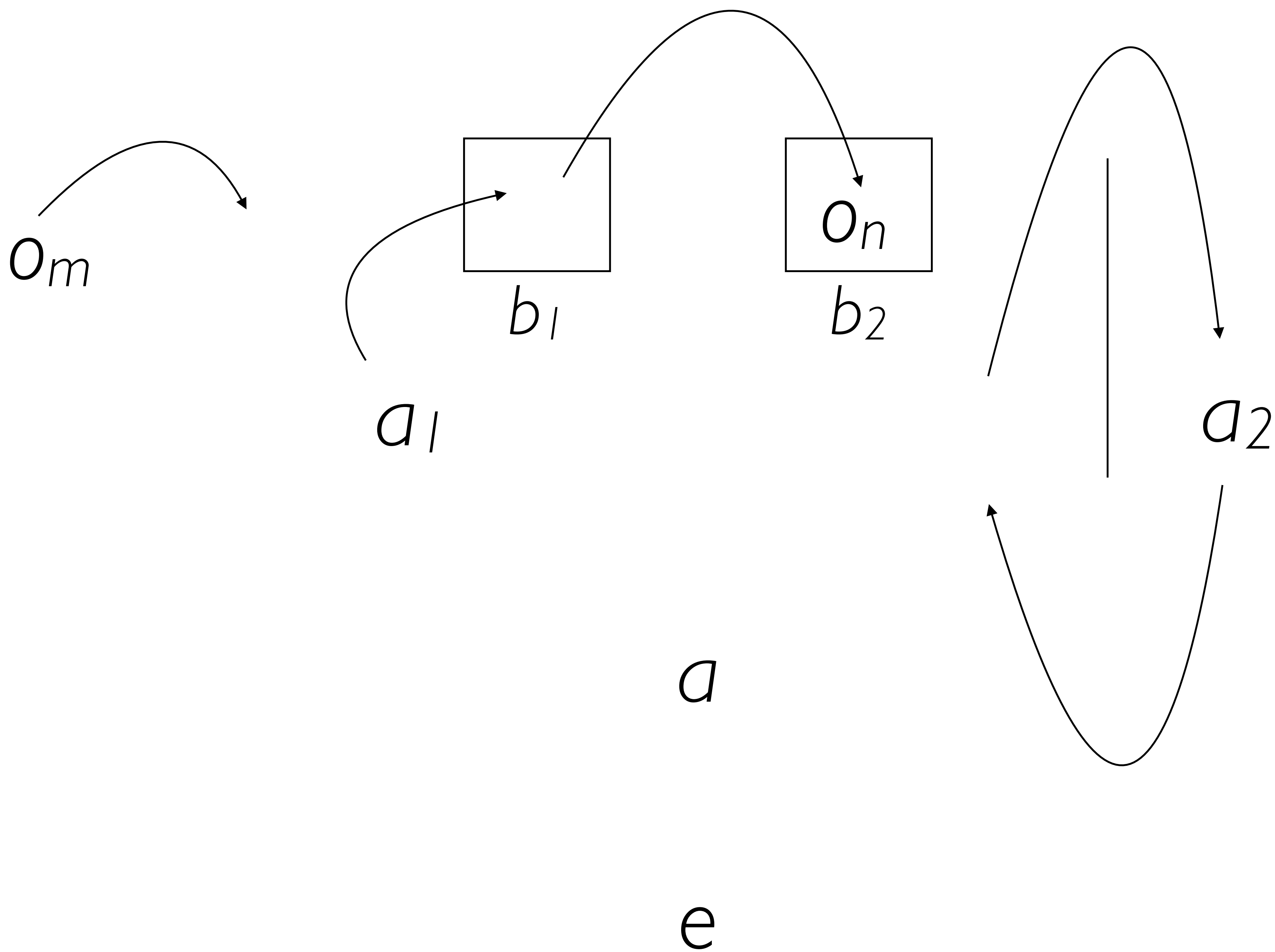

Framework for FBT'1
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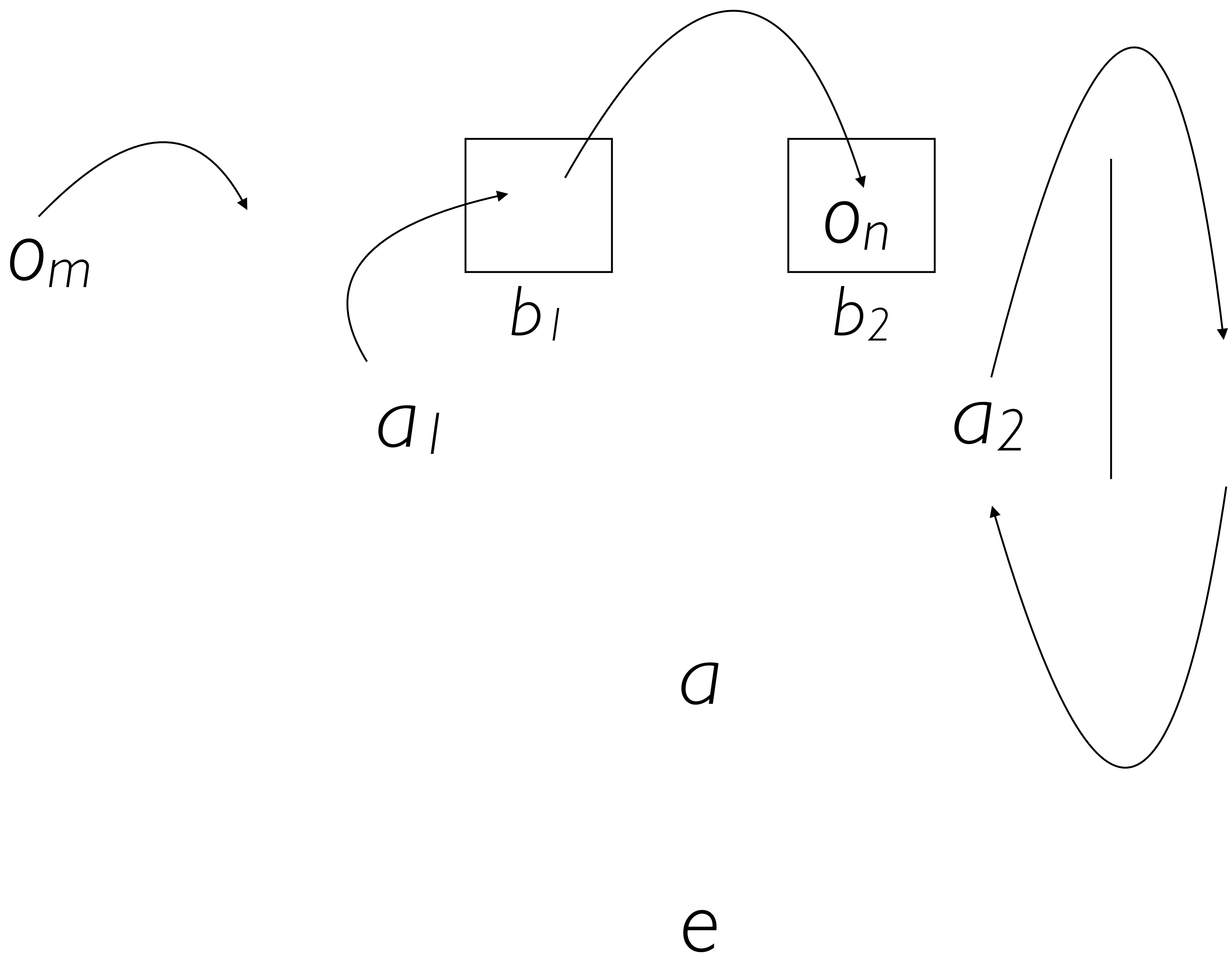

Framework for $FBT^1{}_1$
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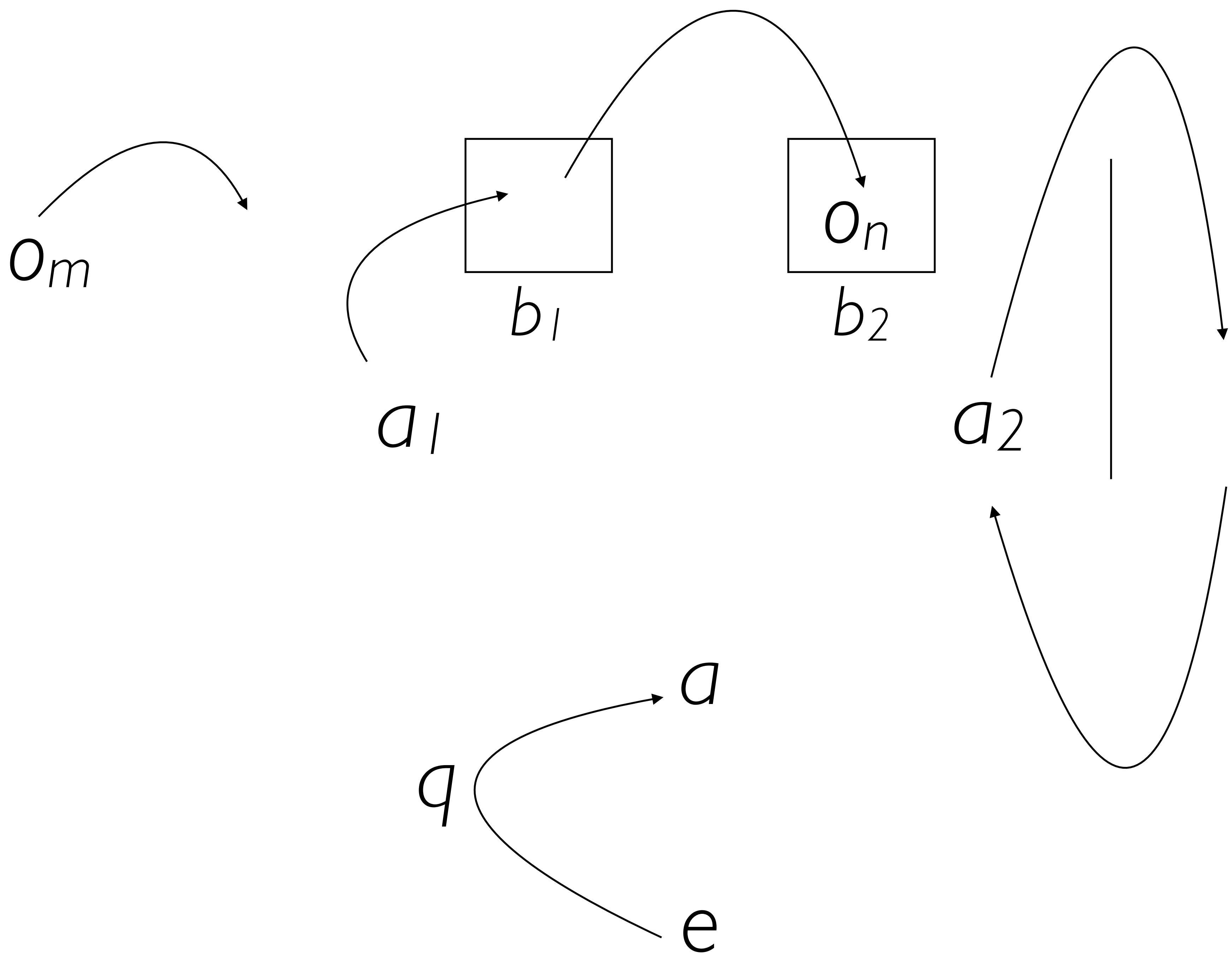

Framework for FBT^1_1
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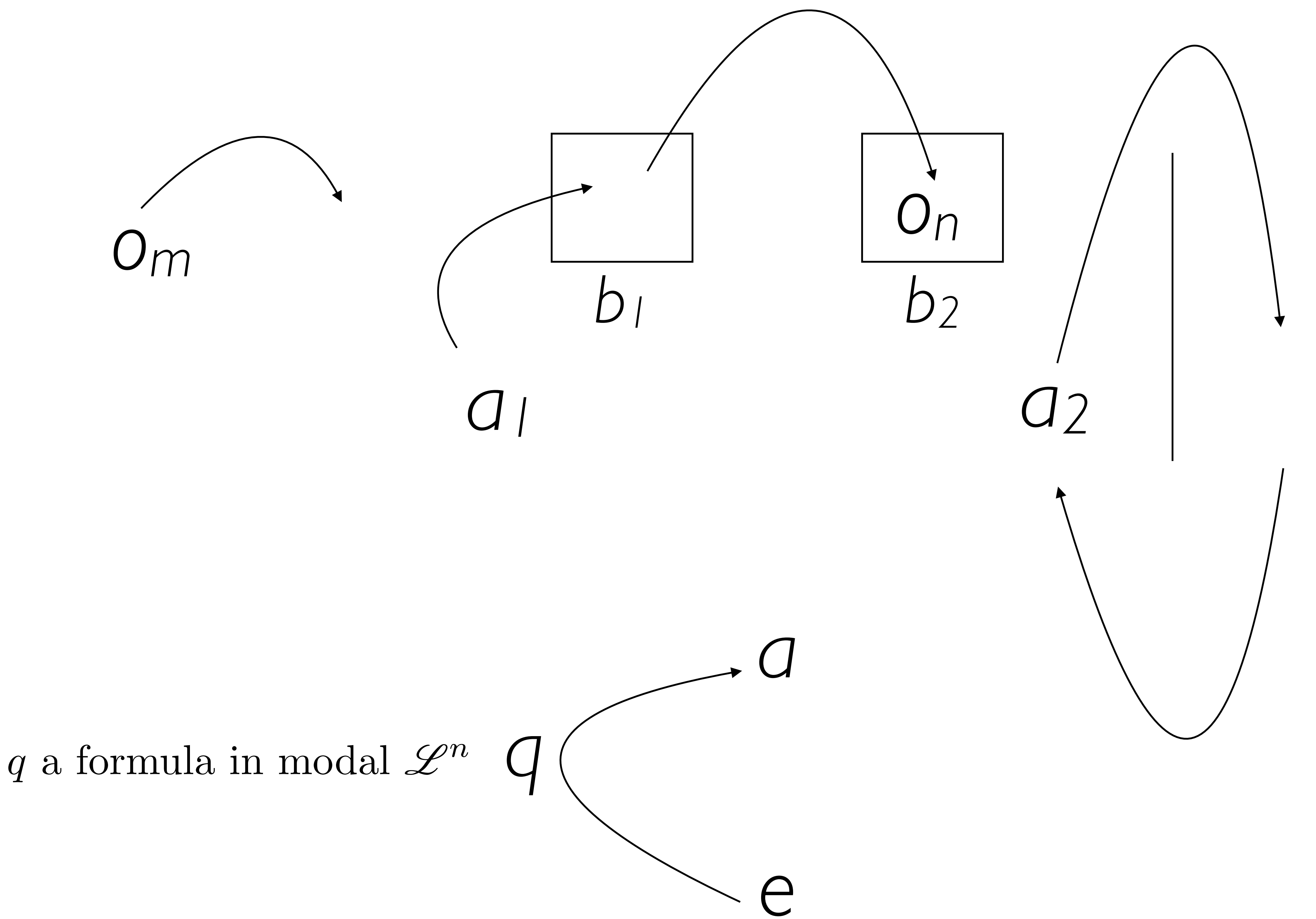

Done, a Decade Ago, Formally & Implementation/Simulation

Arkoudas, K. & Bringsjord, S. (2009) “Propositional Attitudes and Causation” *International Journal of Software and Informatics* **3.1**: 47–65.

http://kryten.mm.rpi.edu/PRICAI_w_sequentcalc_041709.pdf

Propositional attitudes and causation

Konstantine Arkoudas and Selmer Bringsjord

Cognitive Science and Computer Science Departments, RPI
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Abstract. Predicting and explaining the behavior of others in terms of mental states is indispensable for everyday life. It will be equally important for artificial agents. We present an inference system for representing and reasoning about mental states, and use it to provide a formal analysis of the false-belief task. The system allows for the representation of information about events, causation, and perceptual, doxastic, and epistemic states (vision, belief, and knowledge), incorporating ideas from the event calculus and multi-agent epistemic logic. Unlike previous AI formalisms, our focus here is on mechanized proofs and proof programmability, not on metamathematical results. Reasoning is performed via relatively cognitively plausible inference rules, and a degree of automation is achieved by general-purpose inference methods and by a syntactic embedding of the system in first-order logic.

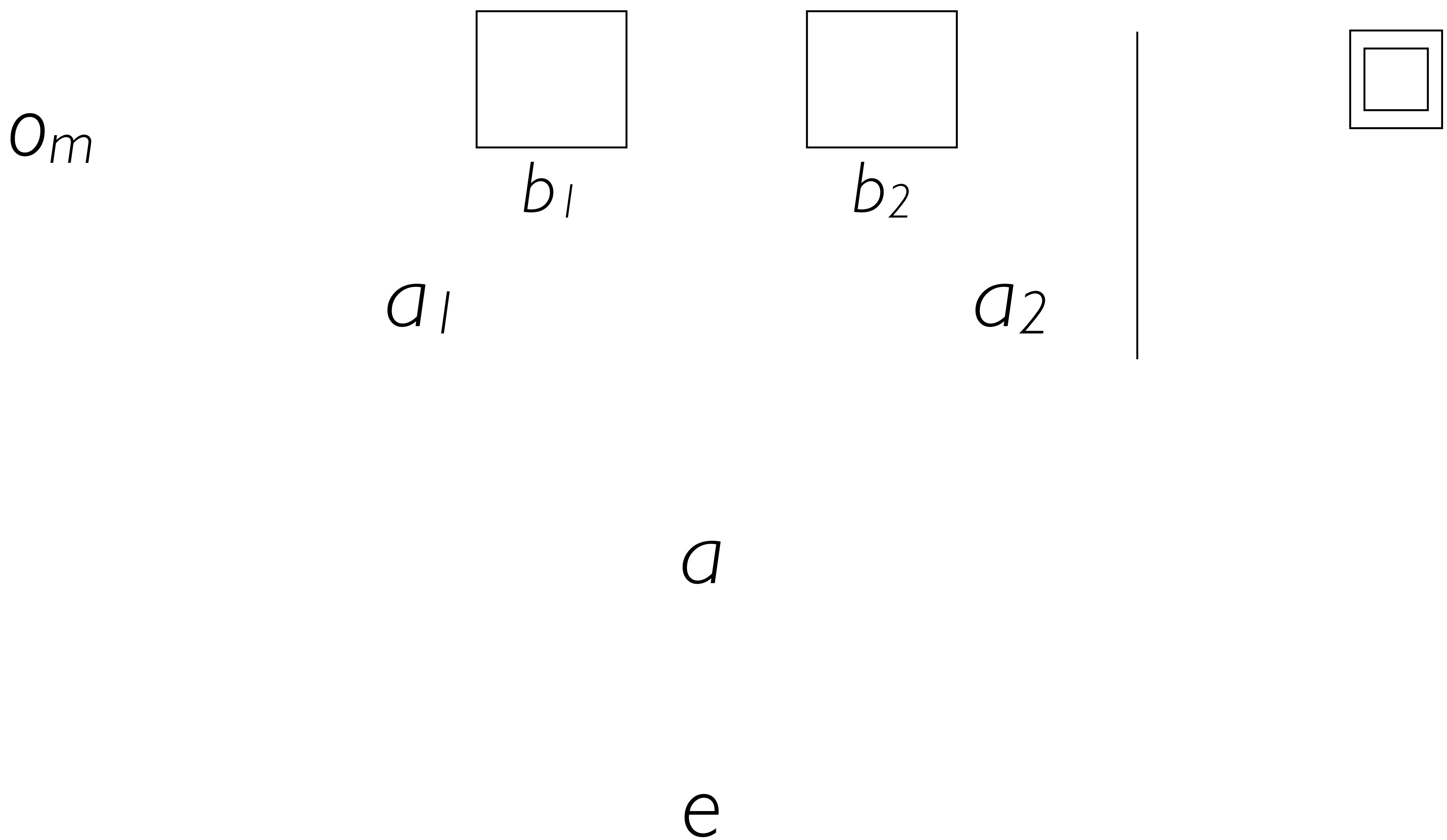
1 Introduction

Interpreting the behavior of other people is indispensable for everyday life. It is something that we do constantly, on a daily basis, and it helps us not only to make sense of human behavior, but also to predict it and—to a certain extent—to control it. How exactly do we manage that? That is not currently known, but many have argued that the ability to ascribe mental states to others and to reason about such mental states is a key component of our capacity to understand human behavior. In particular, all social transactions, from engaging in commerce and negotiating to making jokes and empathizing with other people's pain or joy, appear to require at least a rudimentary grasp of common-sense psychology (CSP), i.e., a large body of truisms such as the following: When an agent a (1) wants to achieve a certain state of affairs p, and (2) believes that some action c can bring about p, and (3) a knows how to carry out c; then, ceteris paribus,[1] a will carry out c; when a sees that p, a knows that p; when a fears that p and a discovers that p is the case, a is disappointed; and so on.

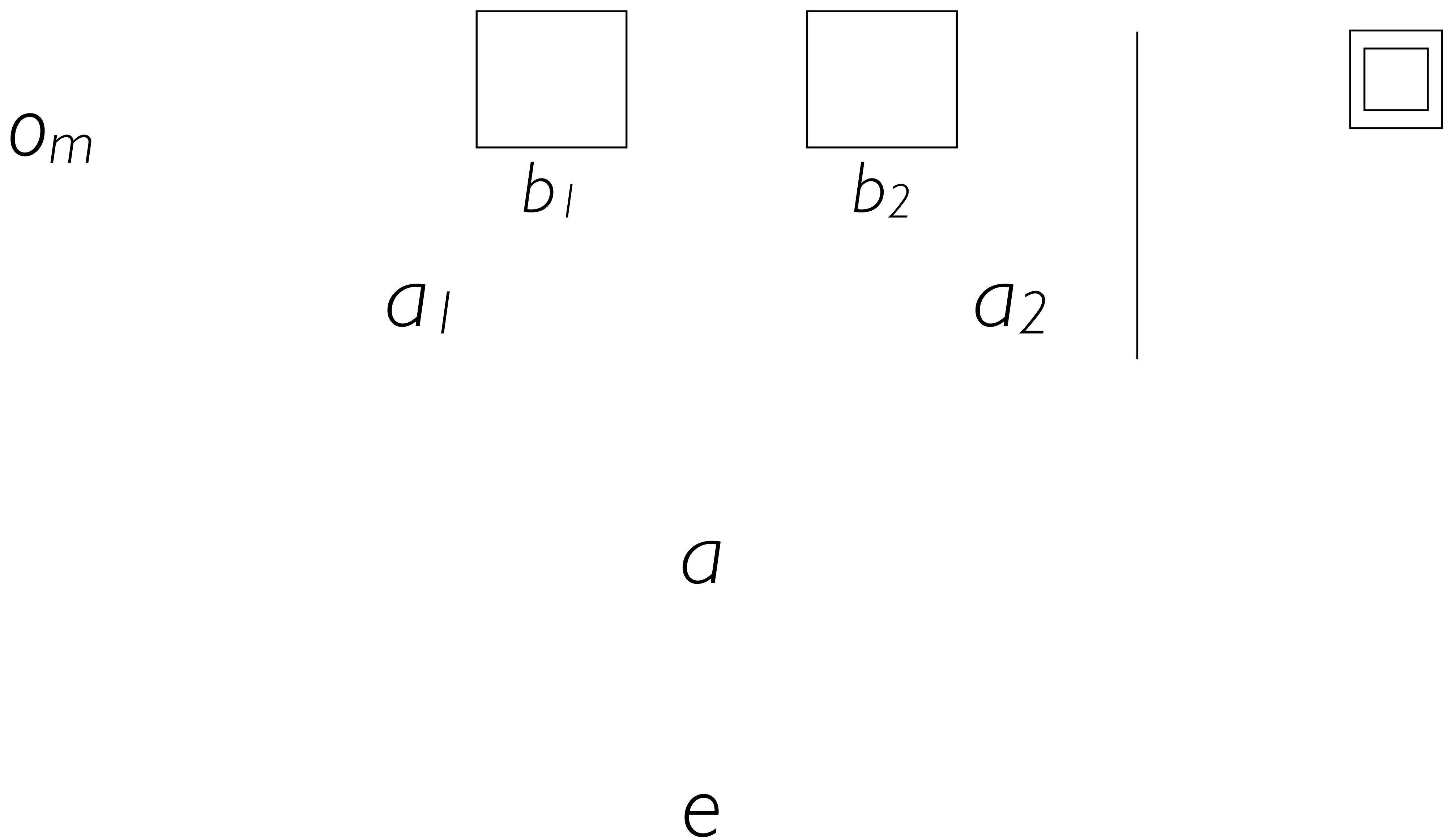
Artificial agents without a mastery of CSP would be severely handicapped in their interactions with humans. This could present problems not only for artificial agents trying to interpret human behavior, but also for artificial agents trying to interpret the behavior of one another. When a system exhibits a complex but rational behavior, and detailed knowledge of its internal structure is not

[1] Assuming that a is able to carry out c, that a has no conflicting desires that override his goal that p; and so on.

Framework for $FBT^1{}_2$

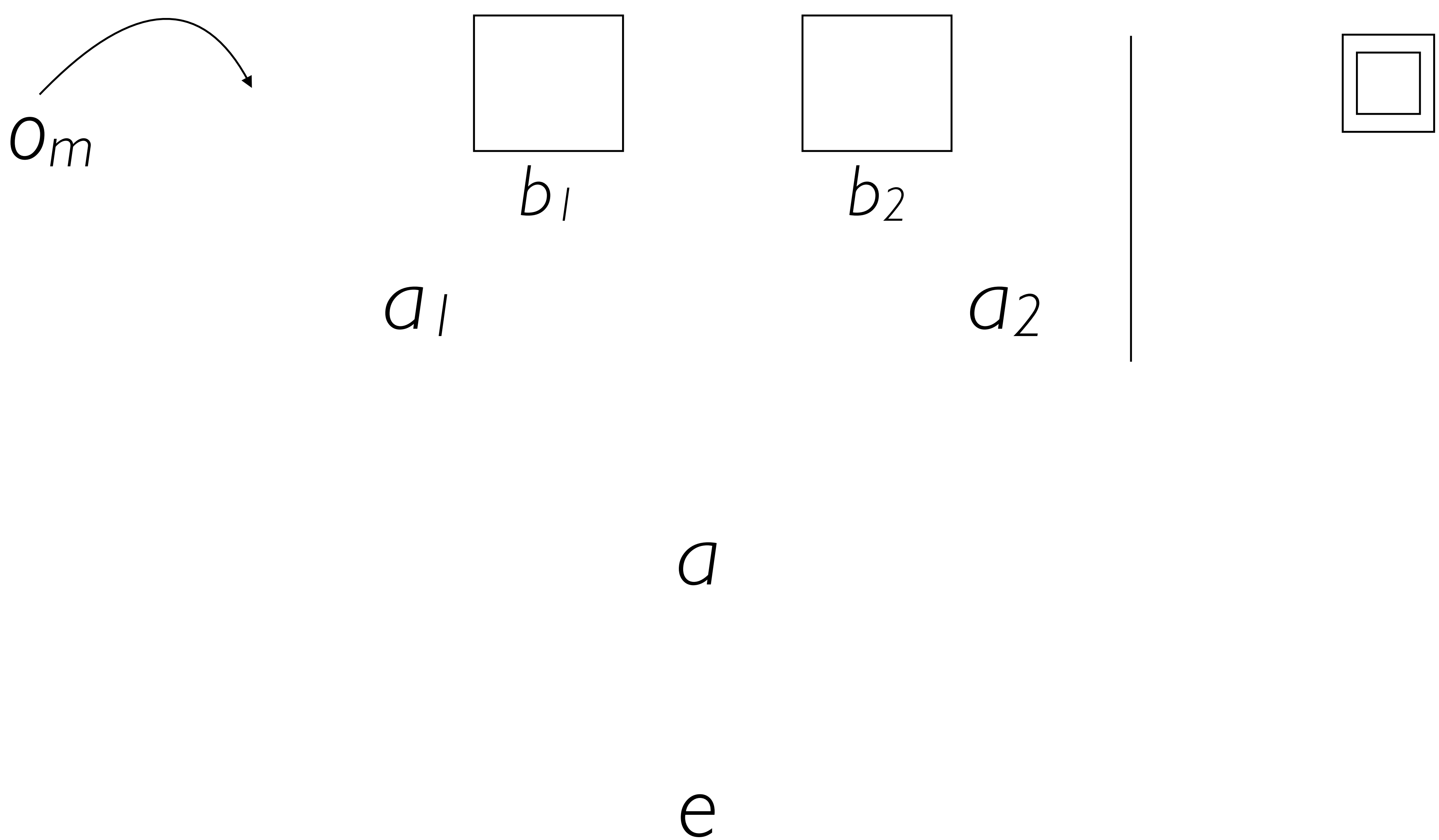


Framework for FBT$^{1}_{2}$
(seven timepoints)

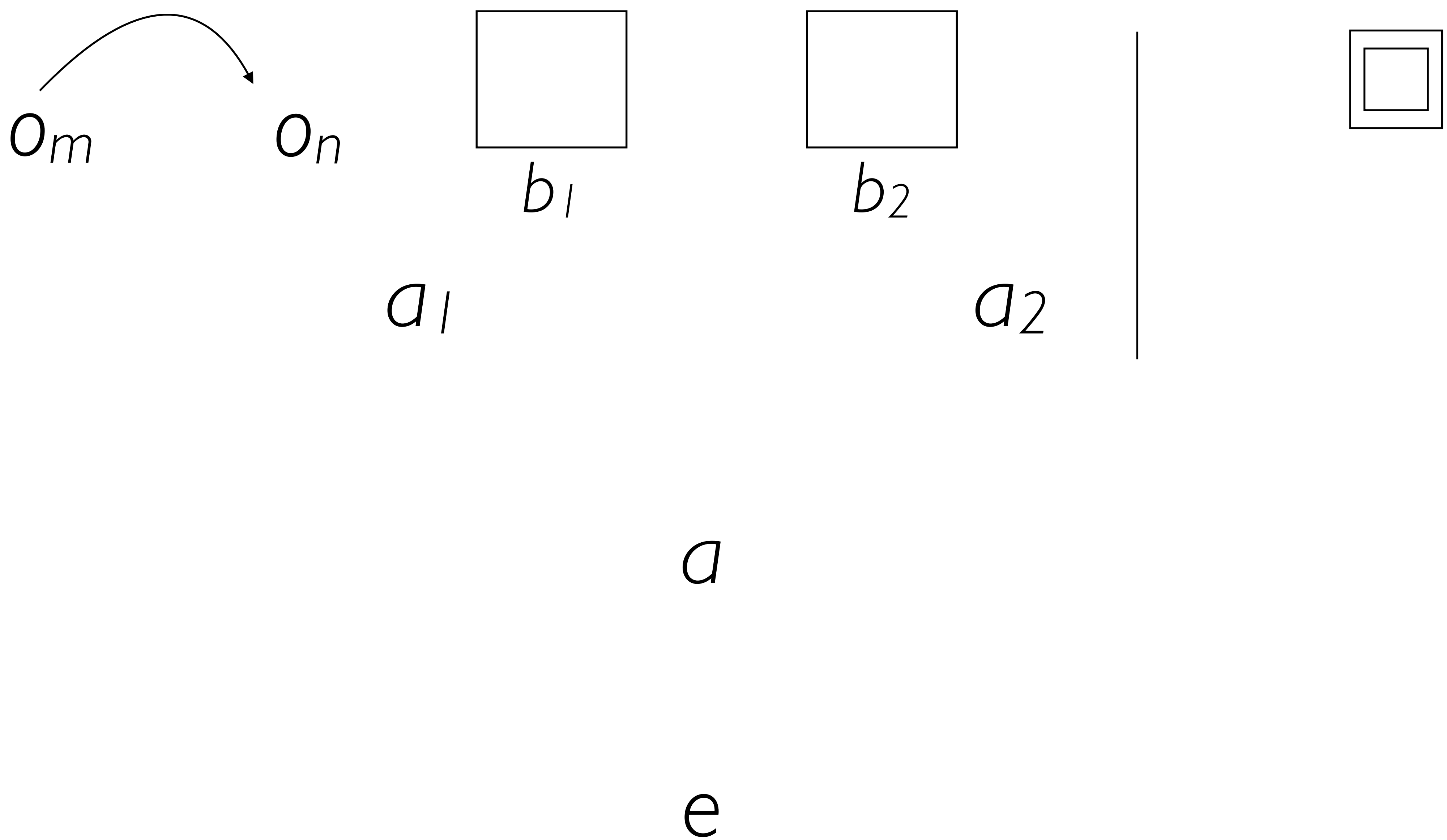


Framework for FBT^1_2

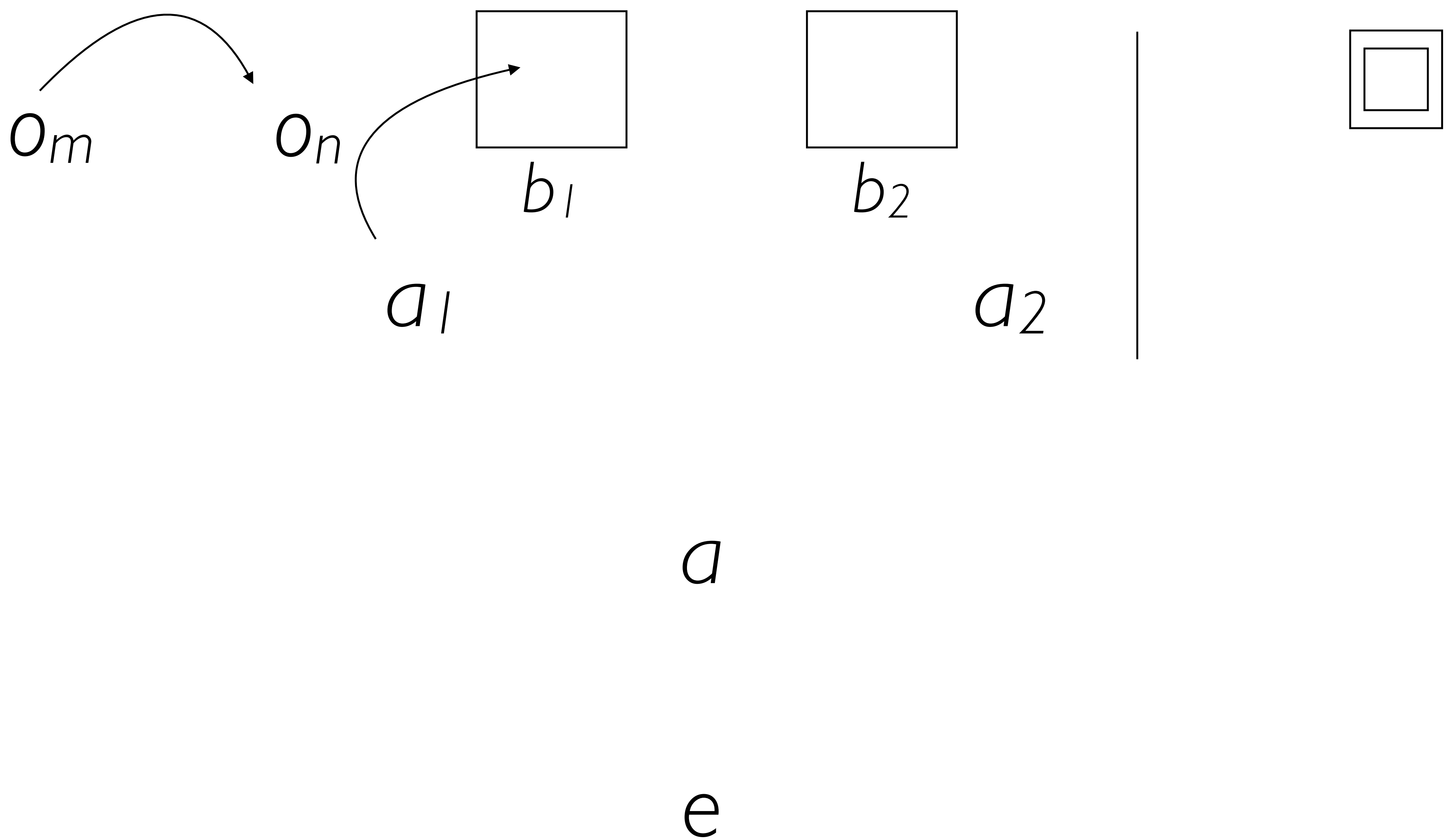
(seven timepoints)



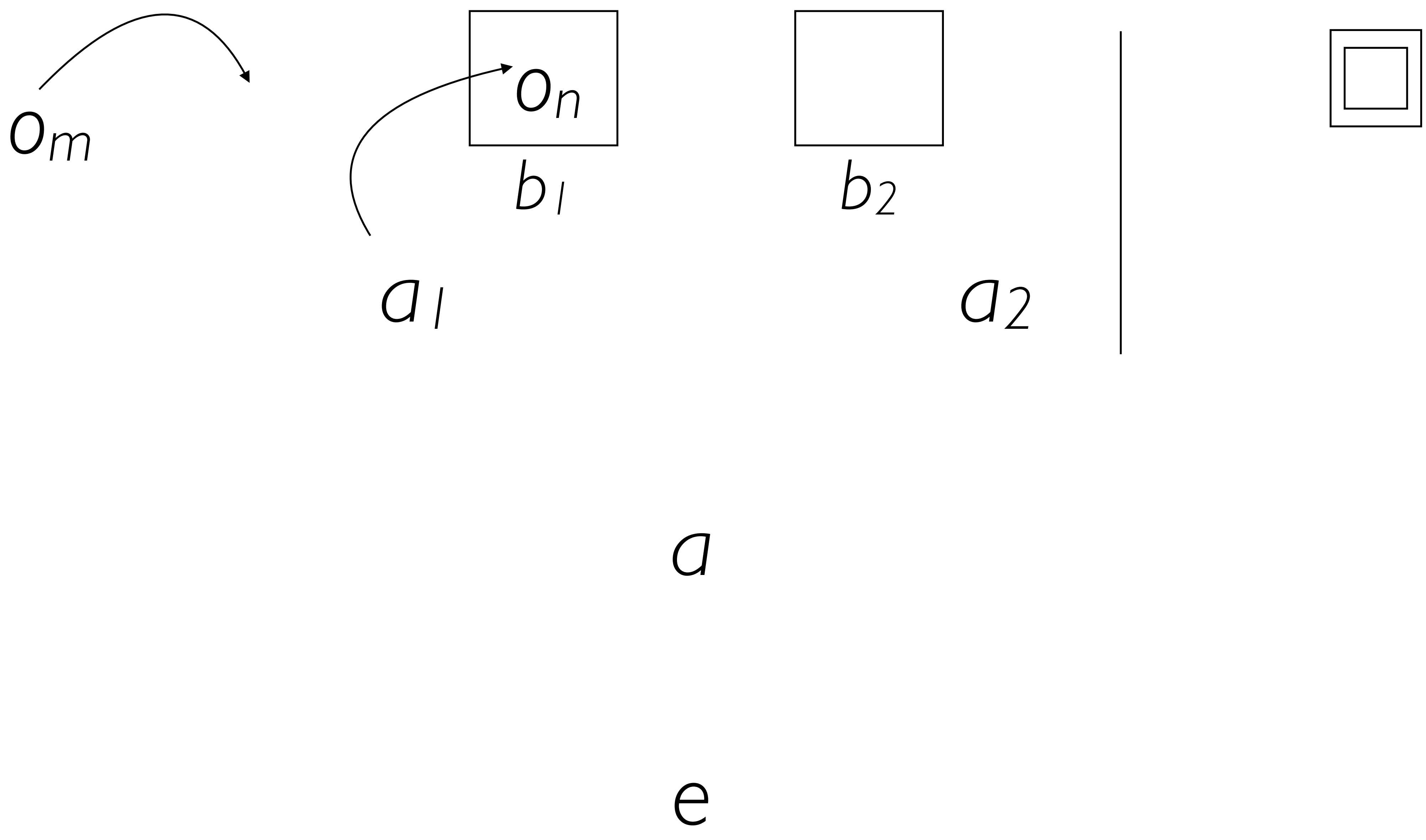
Framework for FBT^1_2
(seven timepoints)



Framework for FBT^{1}_{2}
(seven timepoints)

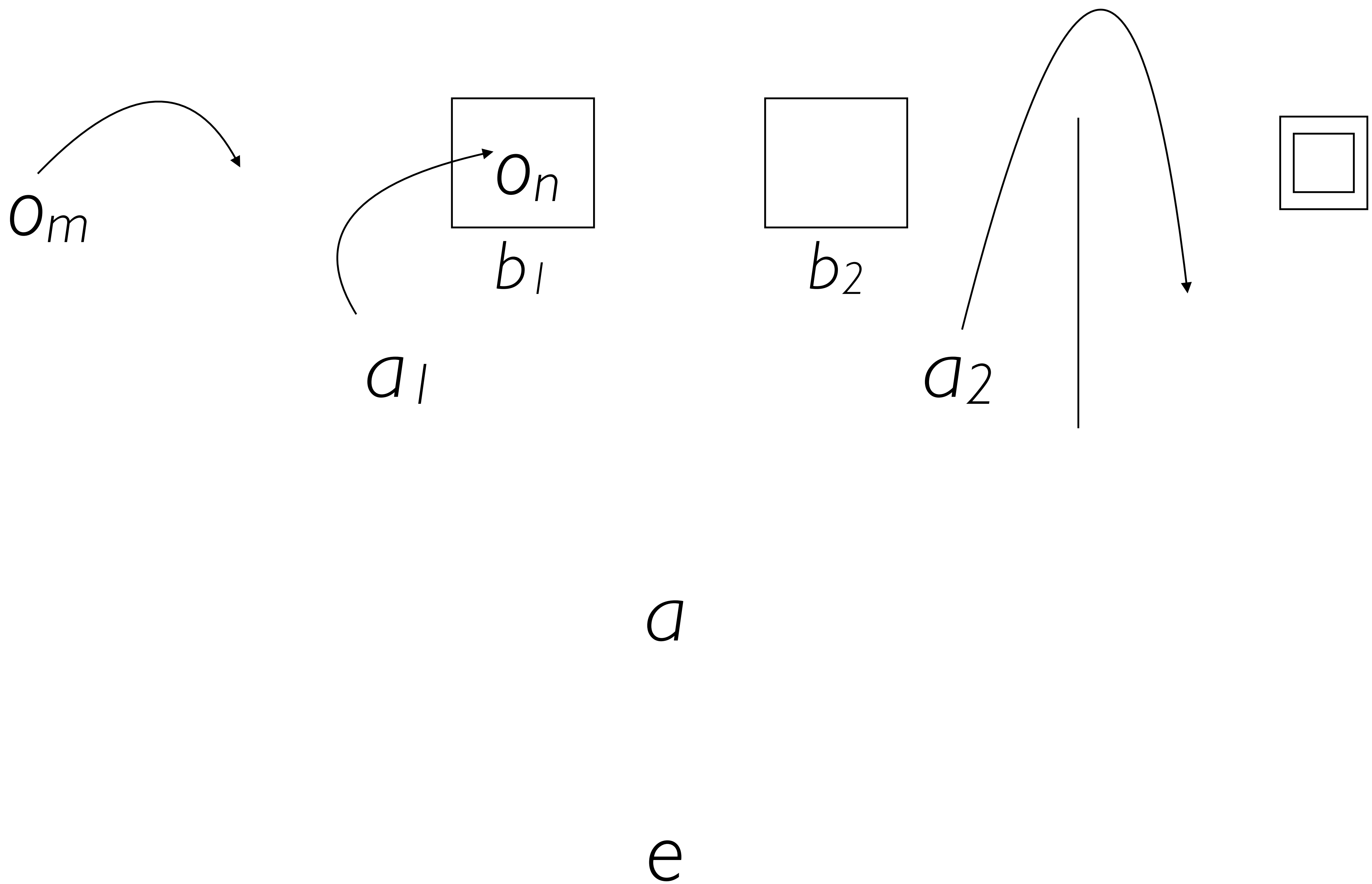


Framework for FBT$^1_2$
(seven timepoints)

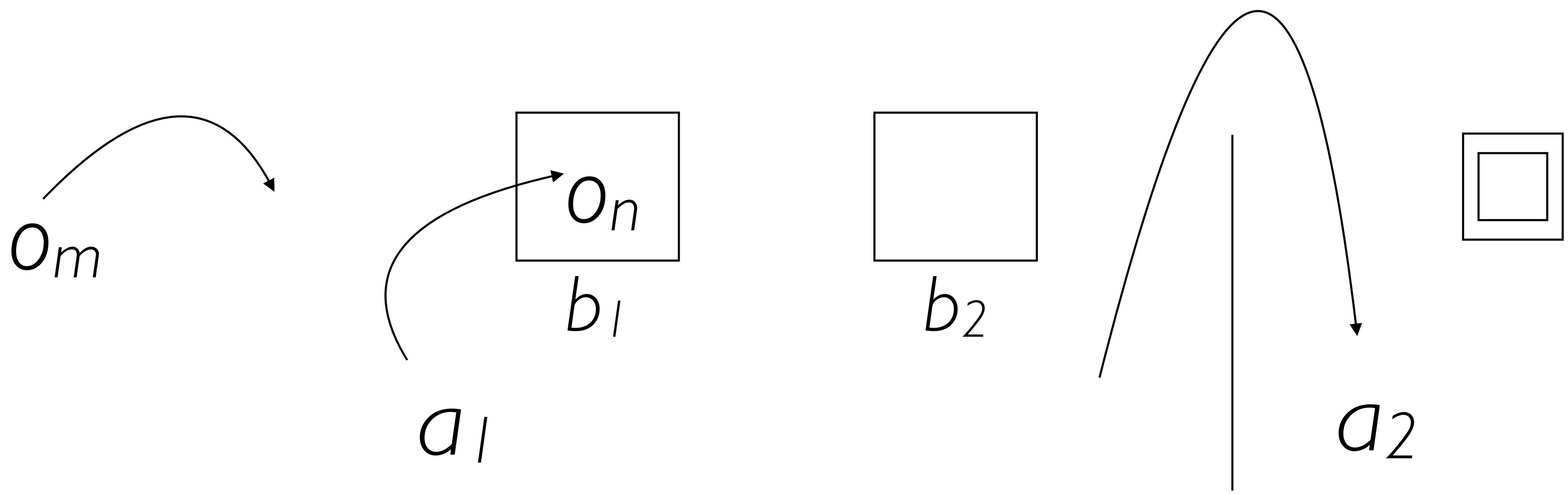


Framework for FBT$^{1}_{2}$

(seven timepoints)



Framework for FBT$^1{}_2$
(seven timepoints)



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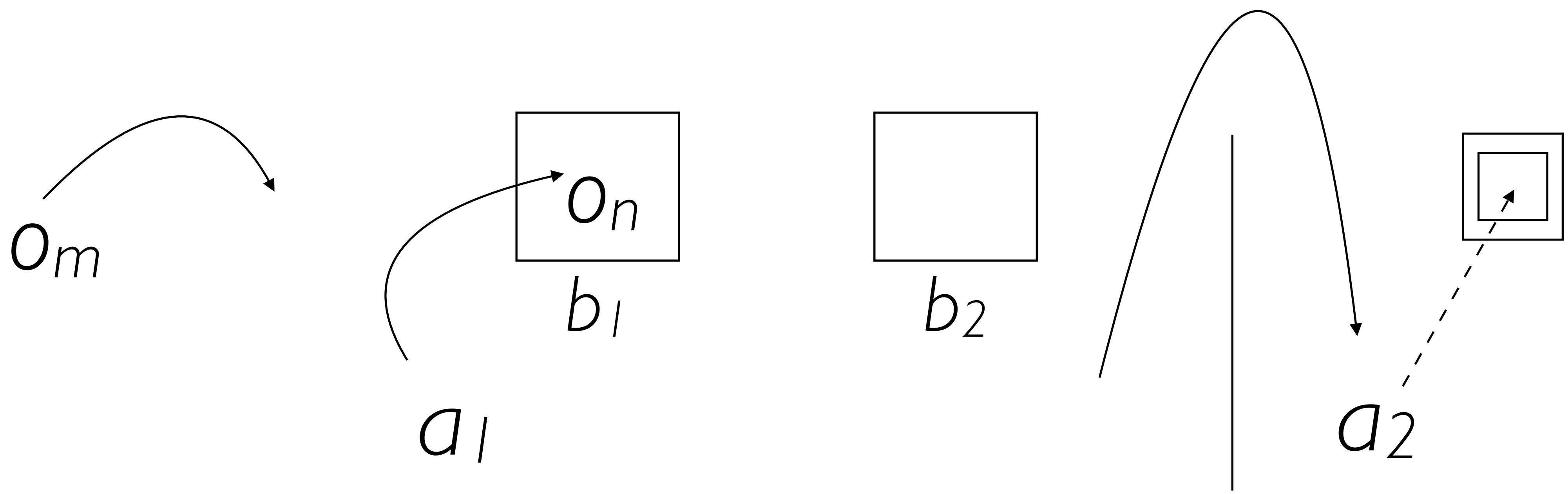
Framework for FBT$^1{}_2$

(seven timepoints)



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Framework for FBT$^1{}_2$
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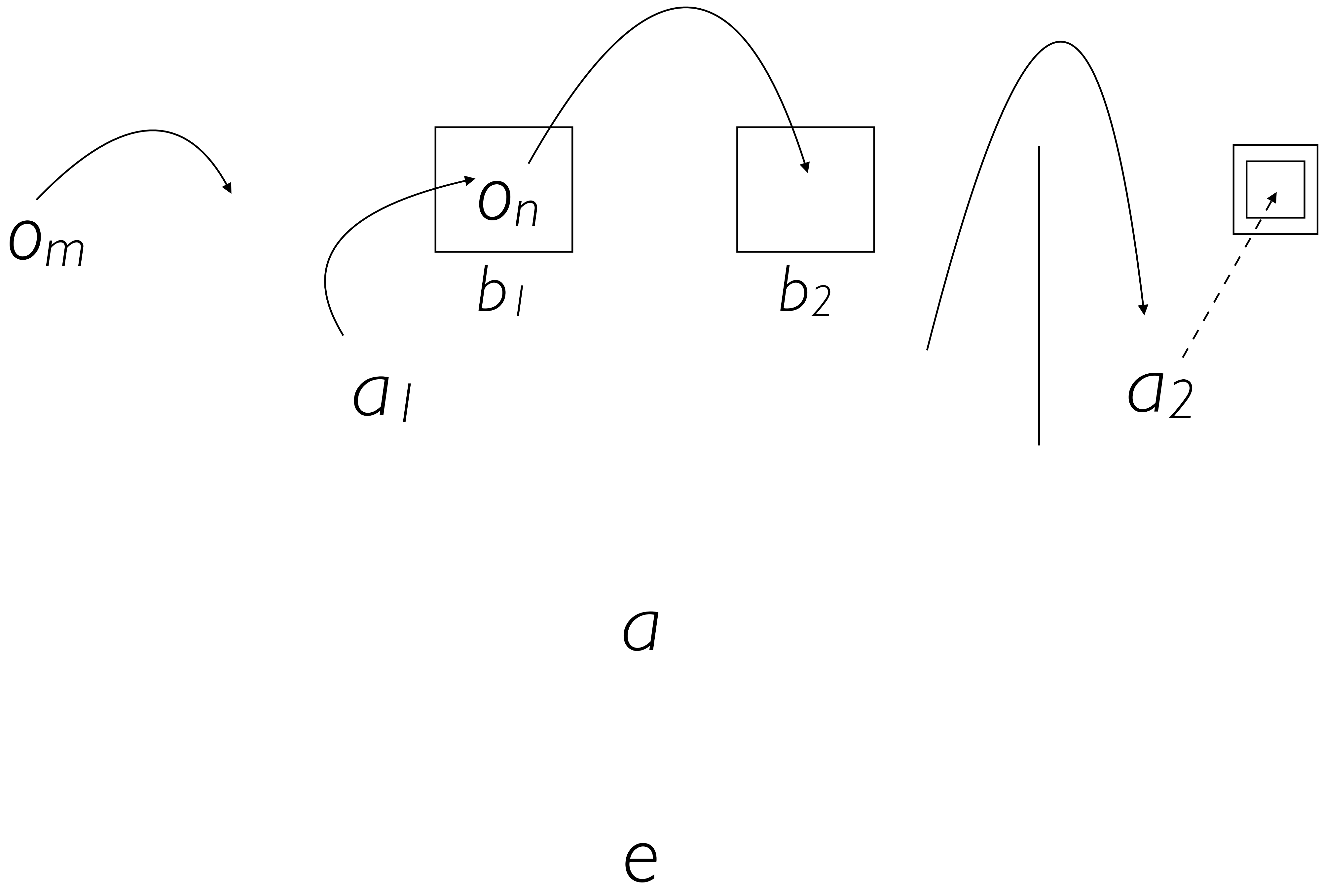
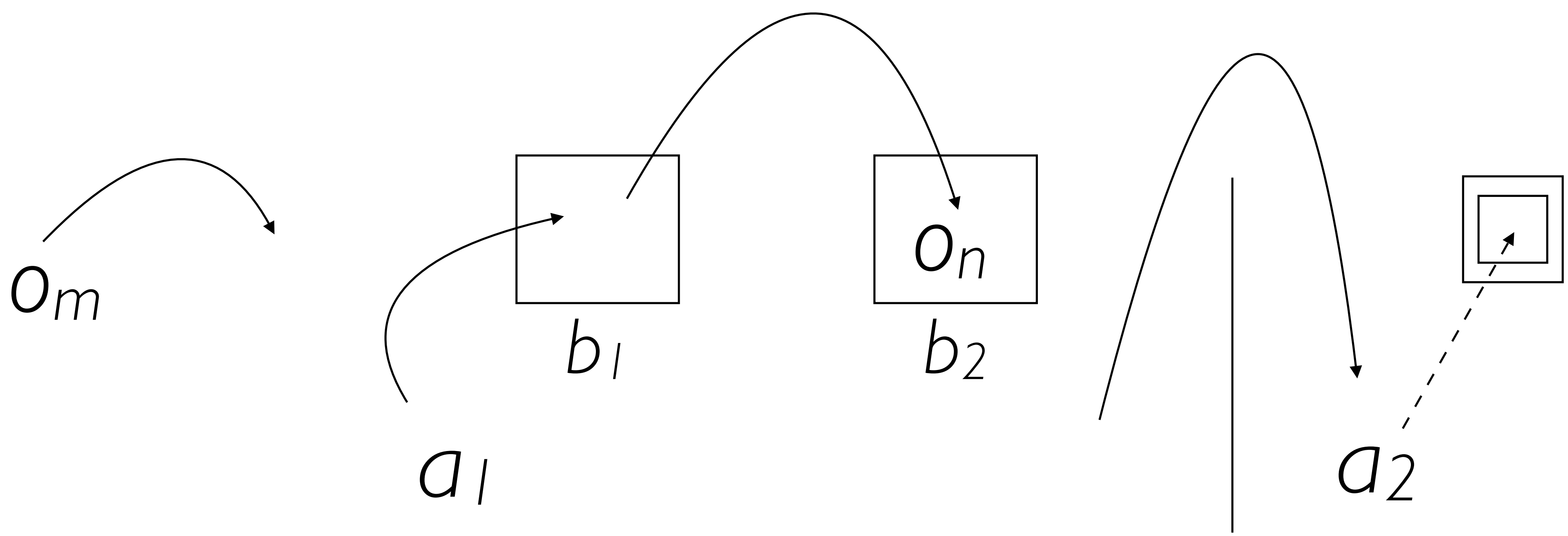


Framework for $FBT^1{}_2$
(seven timepoints)



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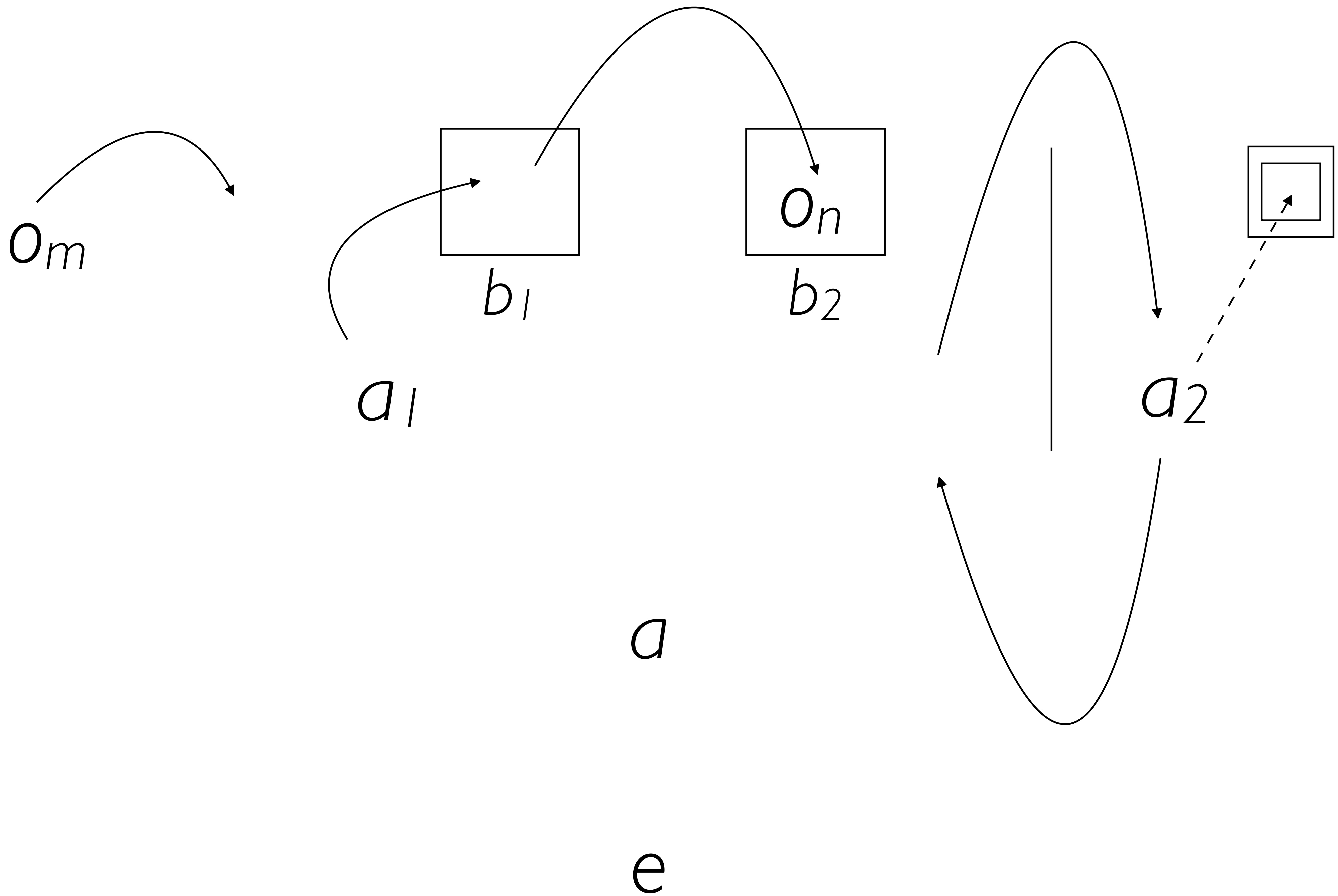
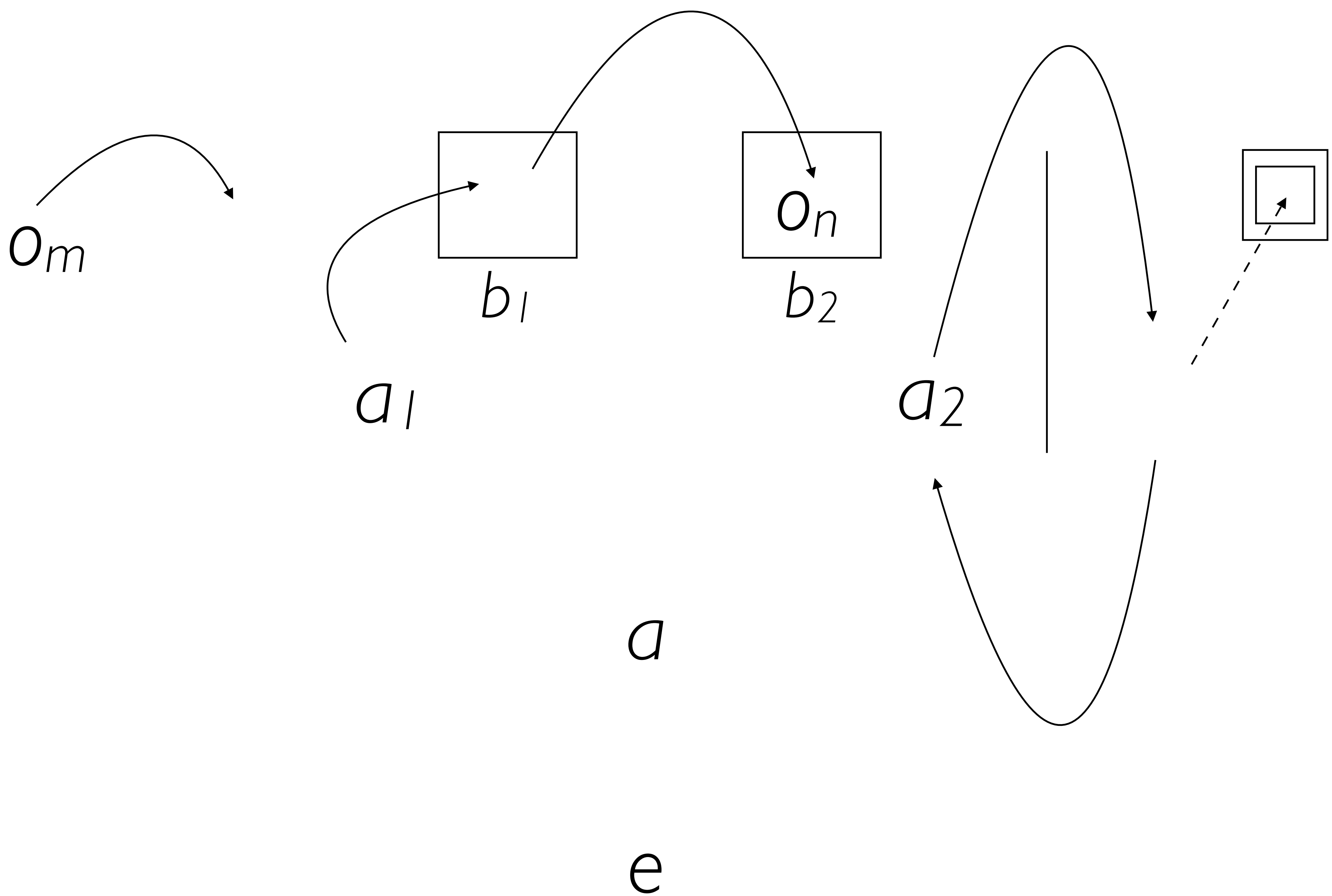
Framework for $FBT^1{}_2$
(seven timepoints)

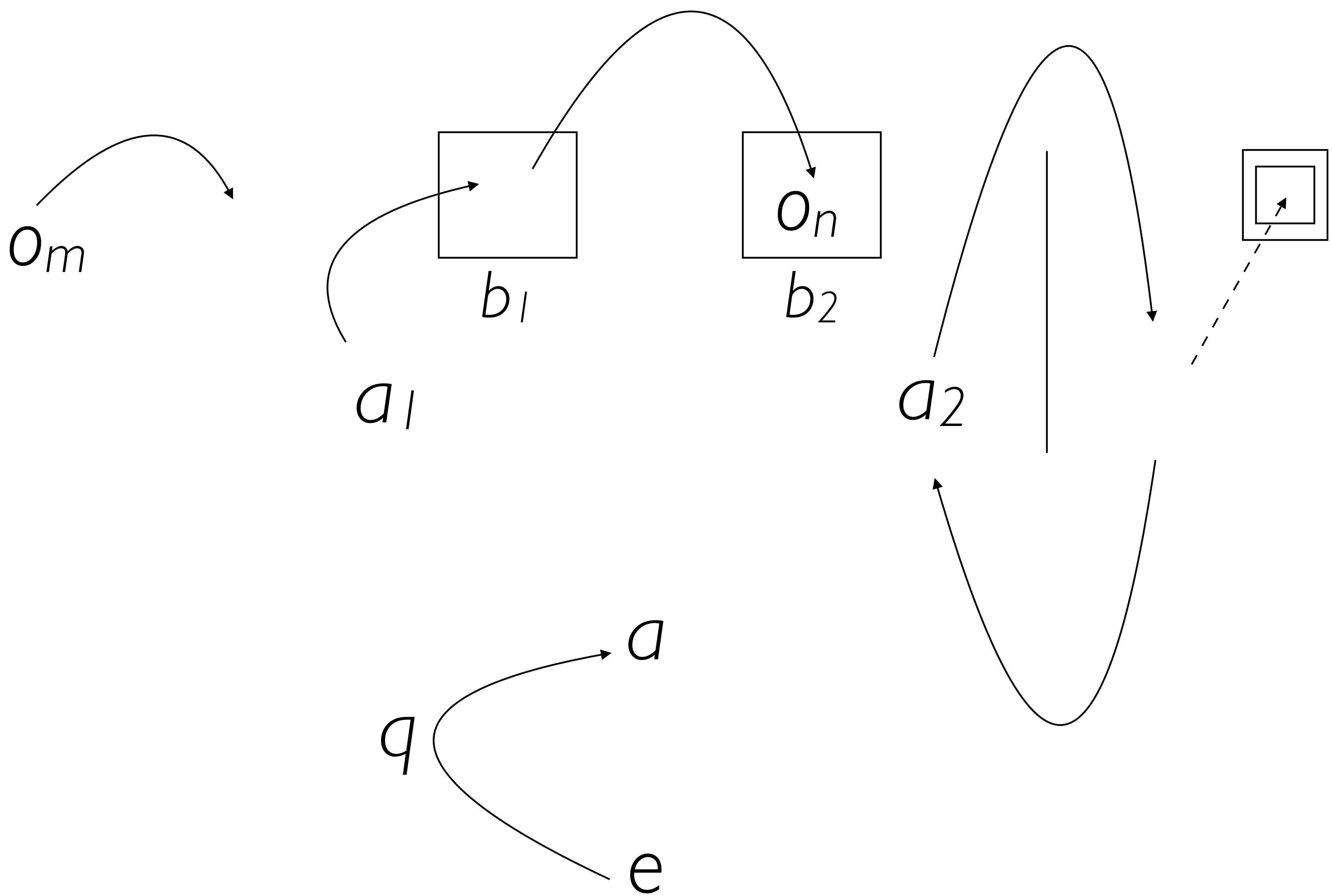


Framework for $FBT^1{}_2$
(seven timepoints)

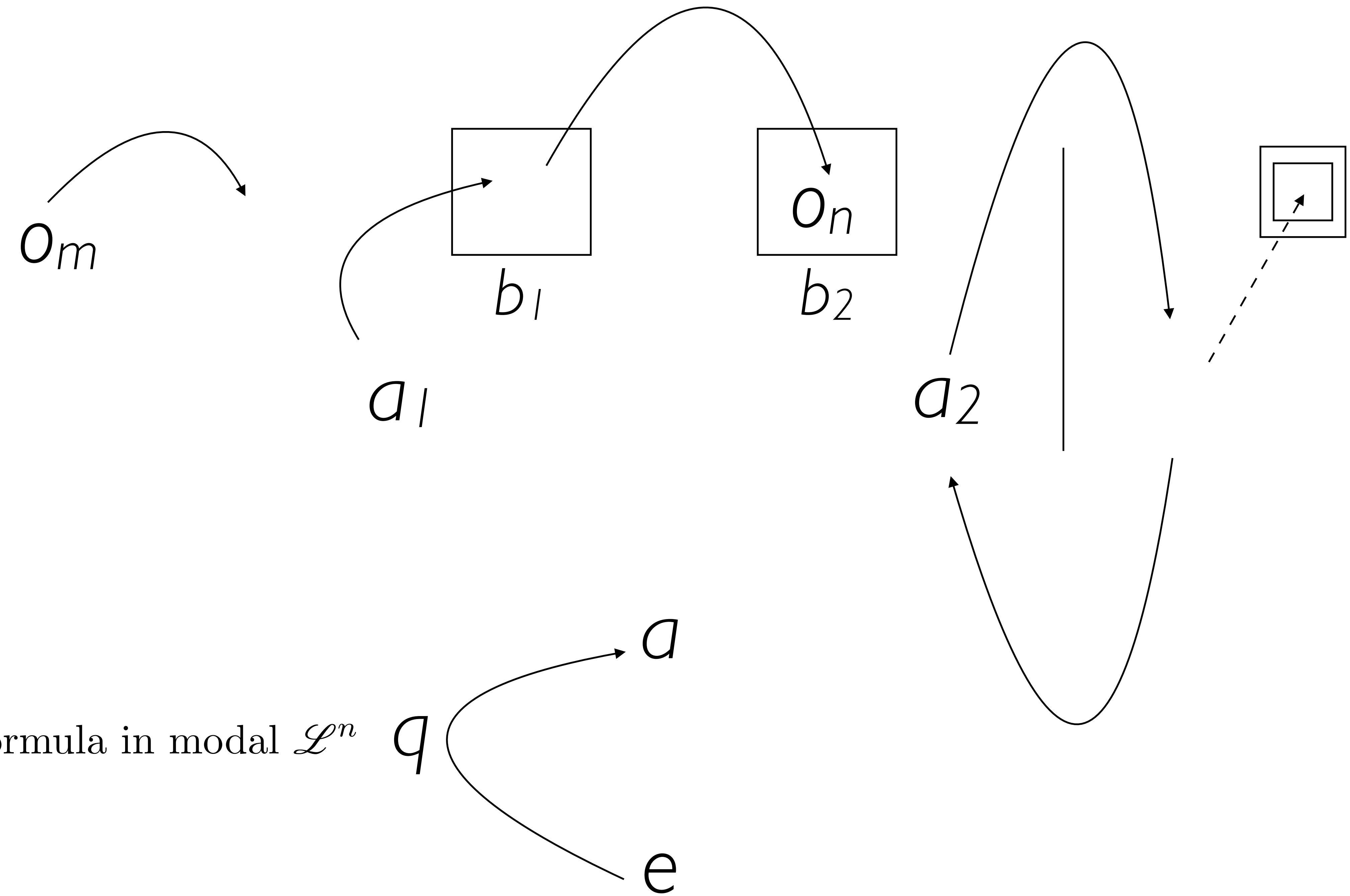


Framework for FBT$^1{}_2$

(seven timepoints)

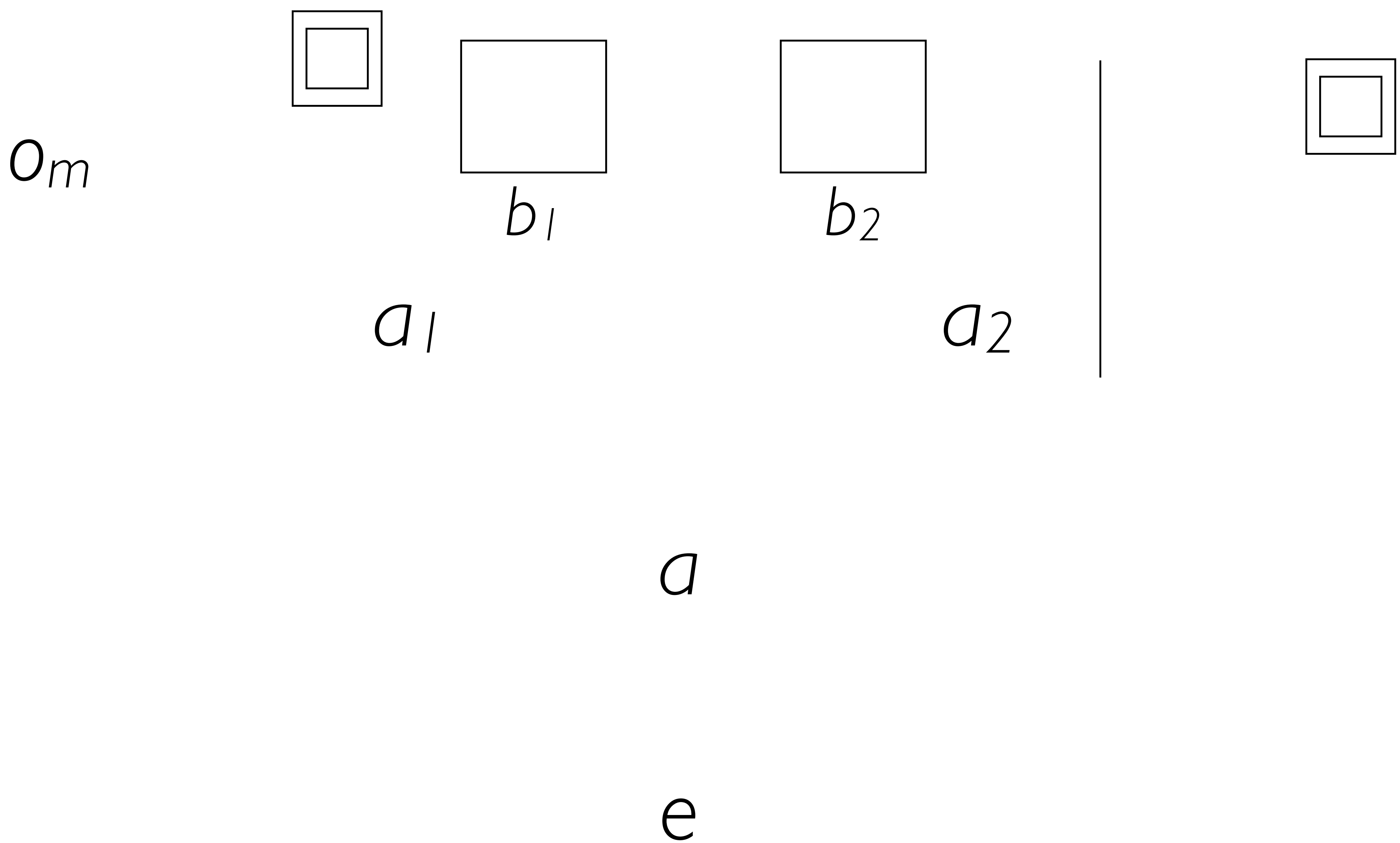


Framework for $FBT^1{}_2$
(seven timepoints)

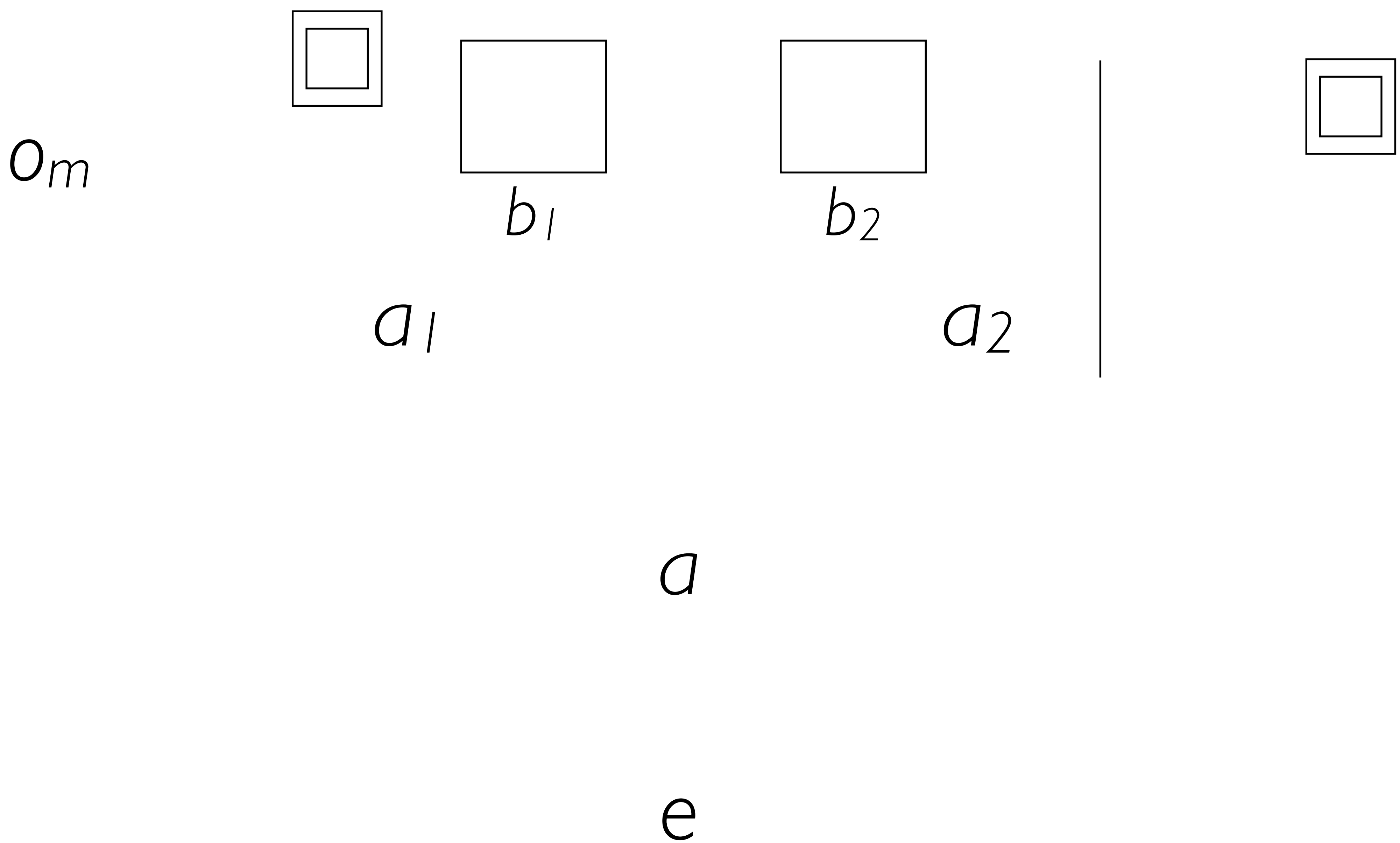


q a formula in modal $\mathscr{L}^n$

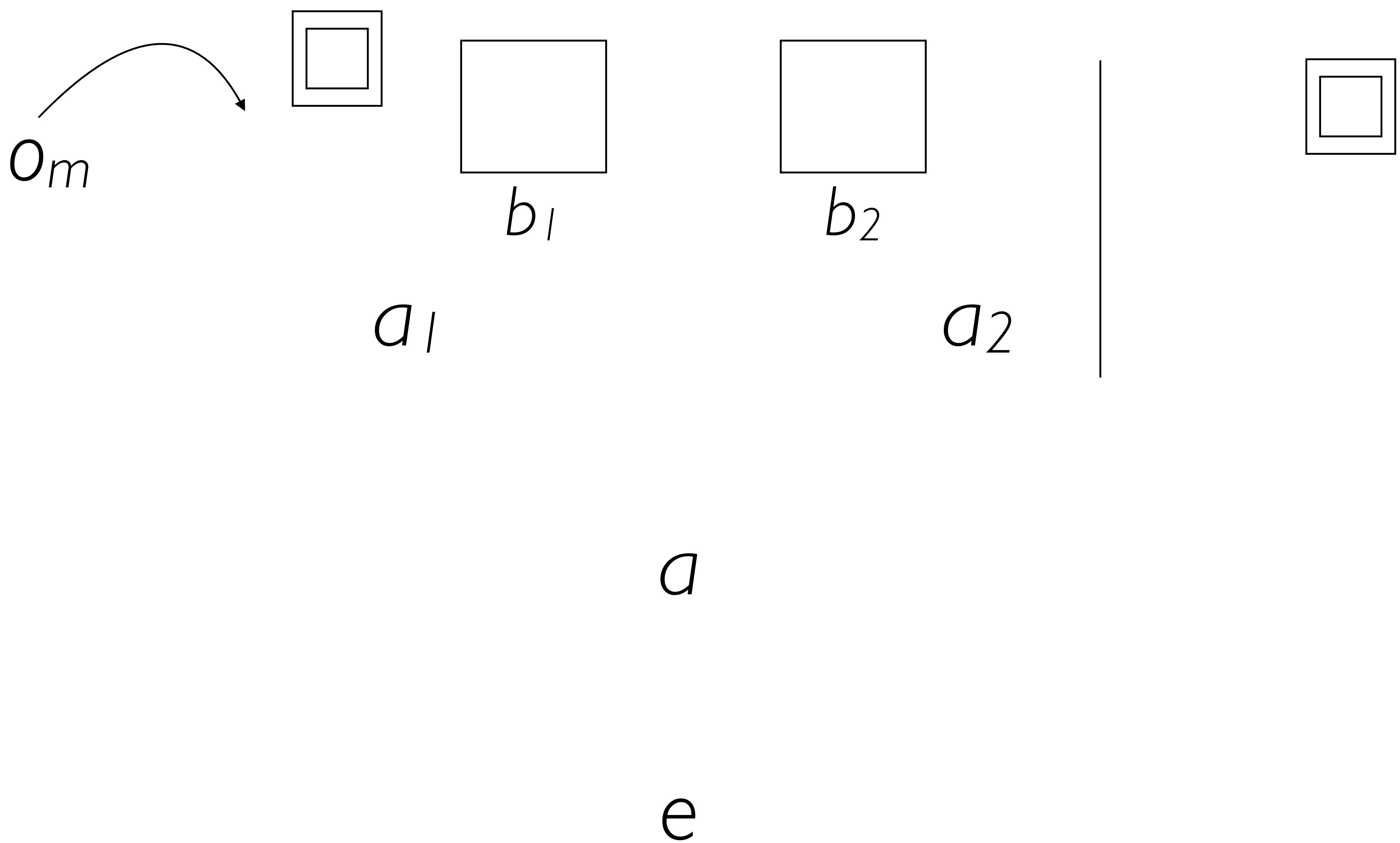
Framework for FBT^{1}_{3}



Framework for FBT^1_3
(eight timepoints)

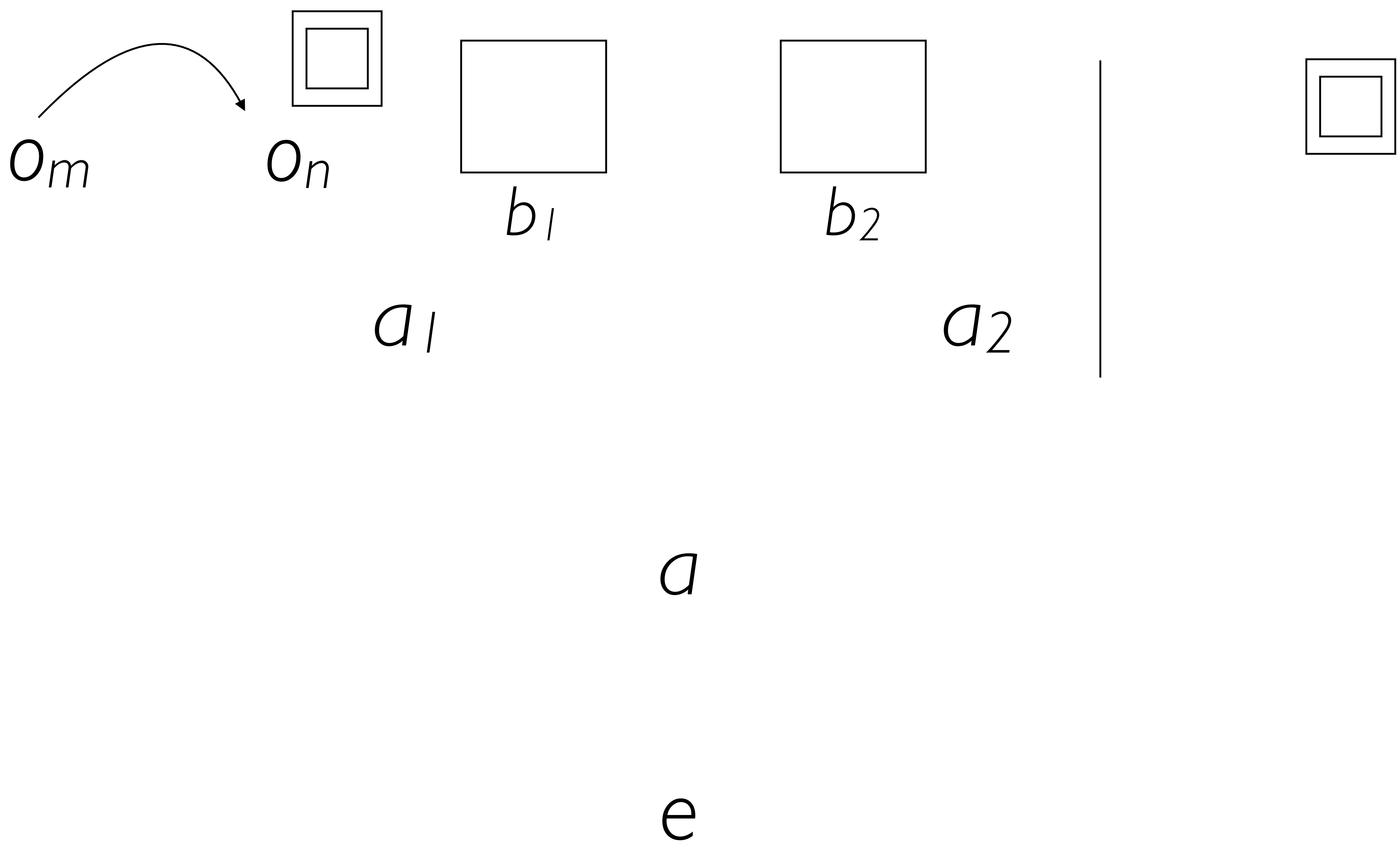


Framework for $FBT^1{}_3$
(eight timepoints)

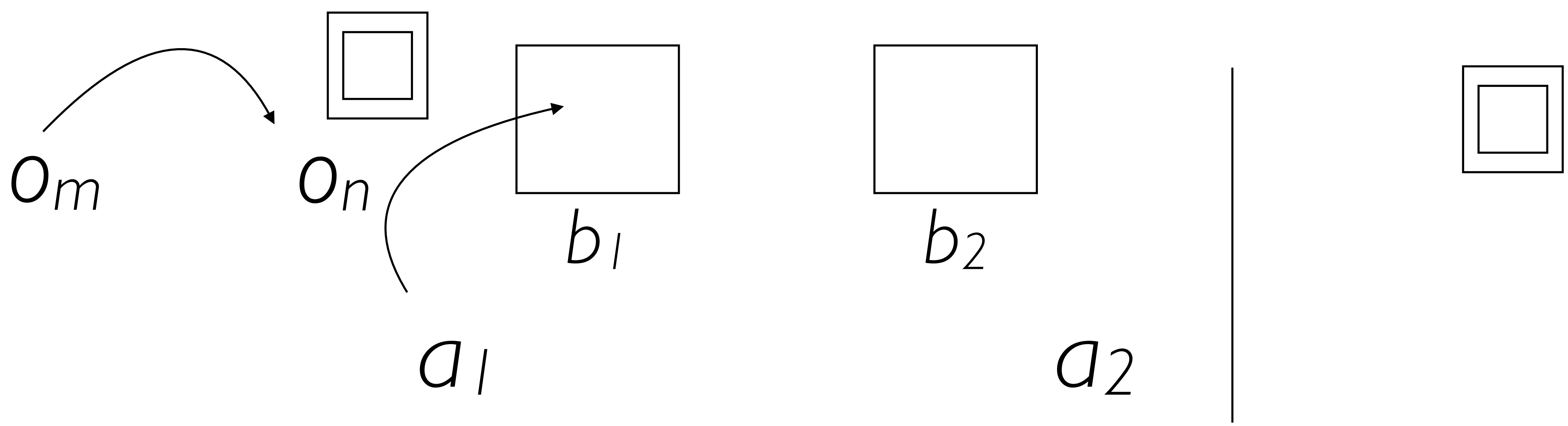


Framework for $FBT^1{}_3$

(eight timepoints)



Framework for FBT$^1{}_3$
(eight timepoints)



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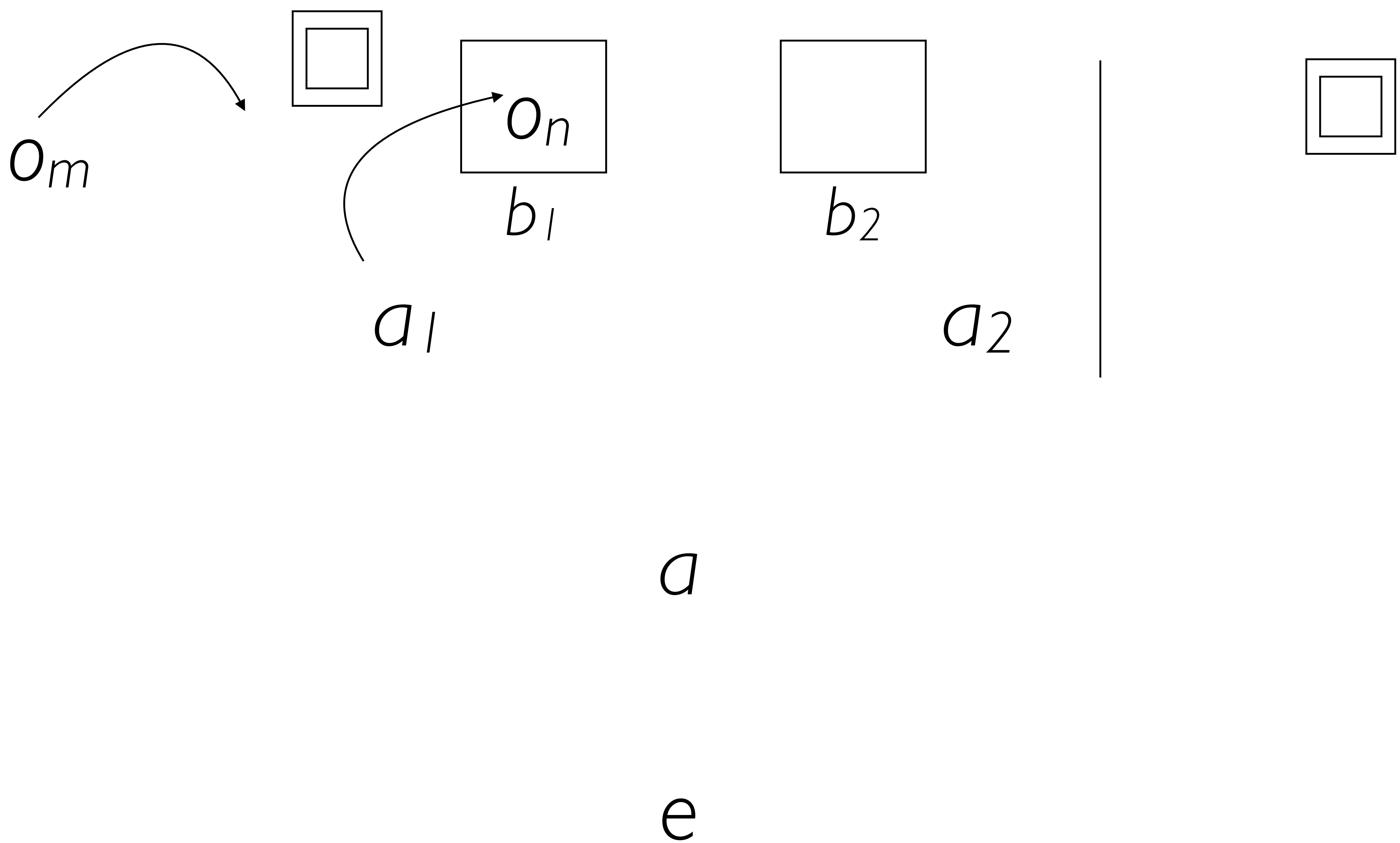
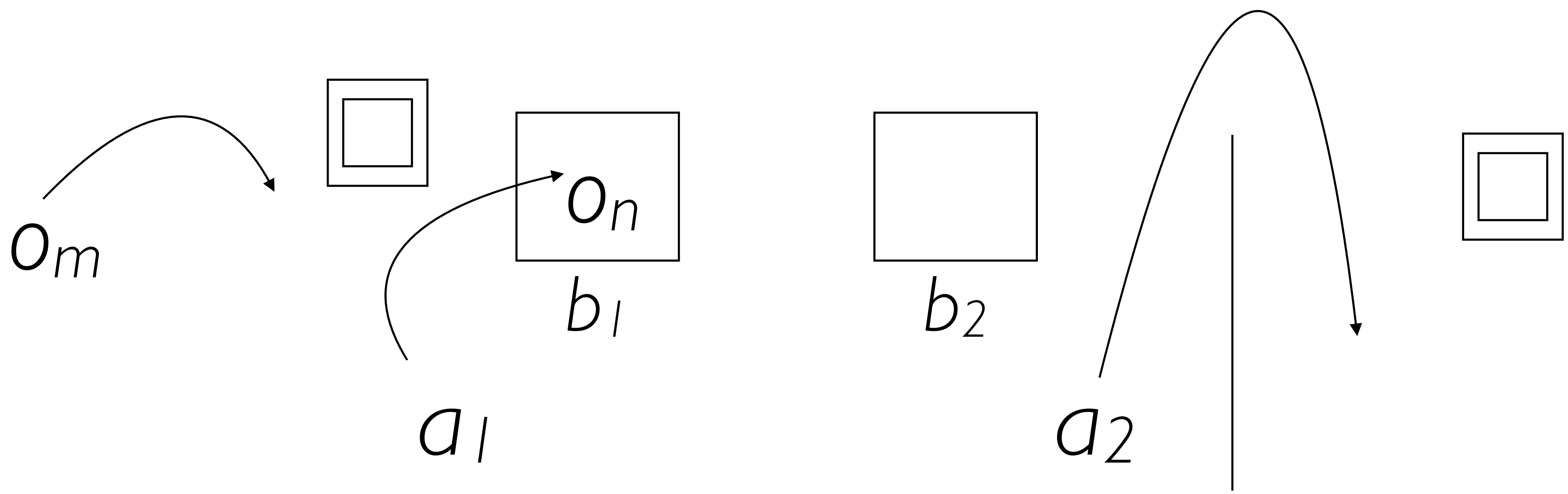
Framework for FBT^1_3
(eight timepoints)



Framework for FBT^1_3

(eight timepoints)



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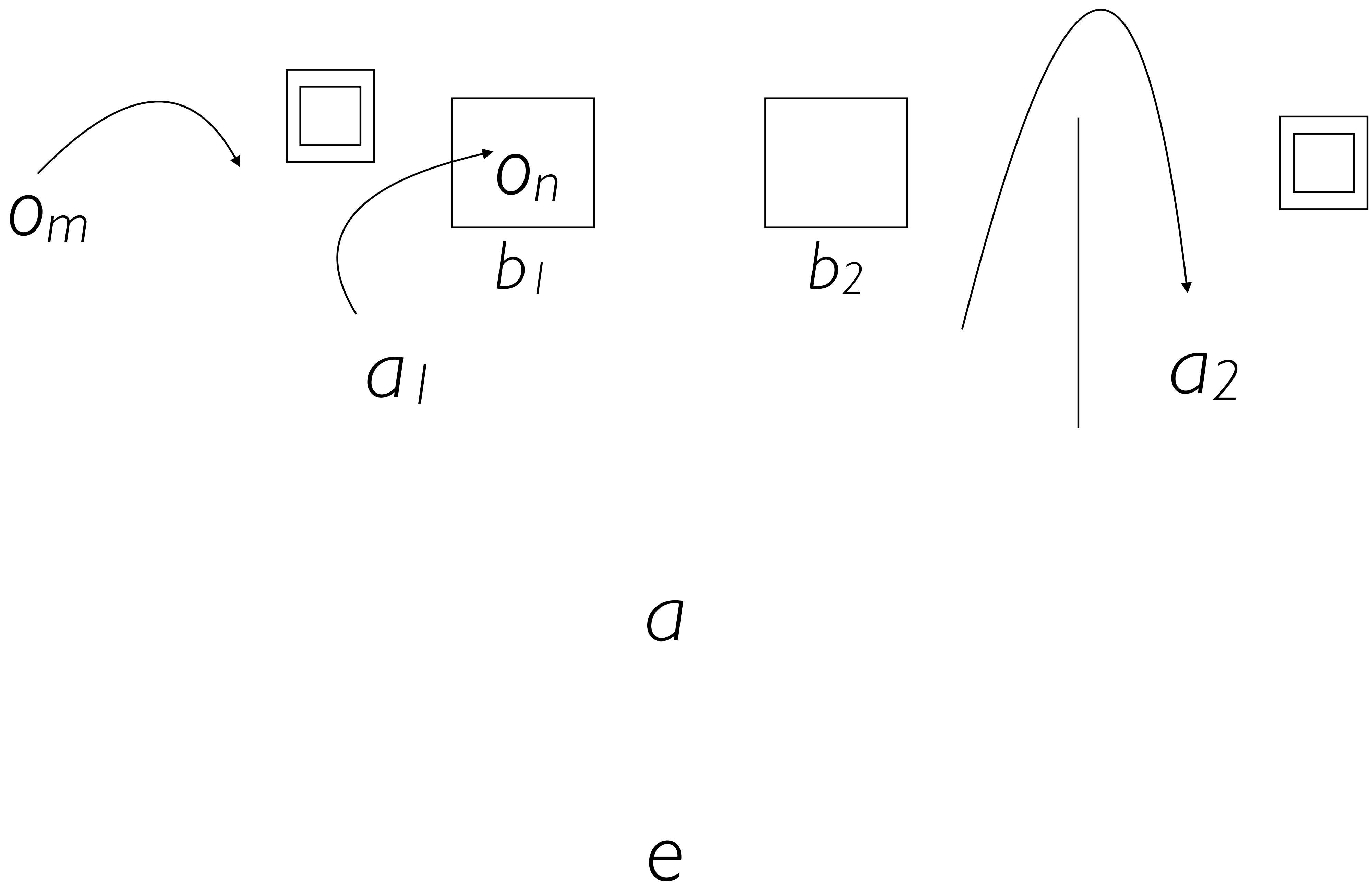
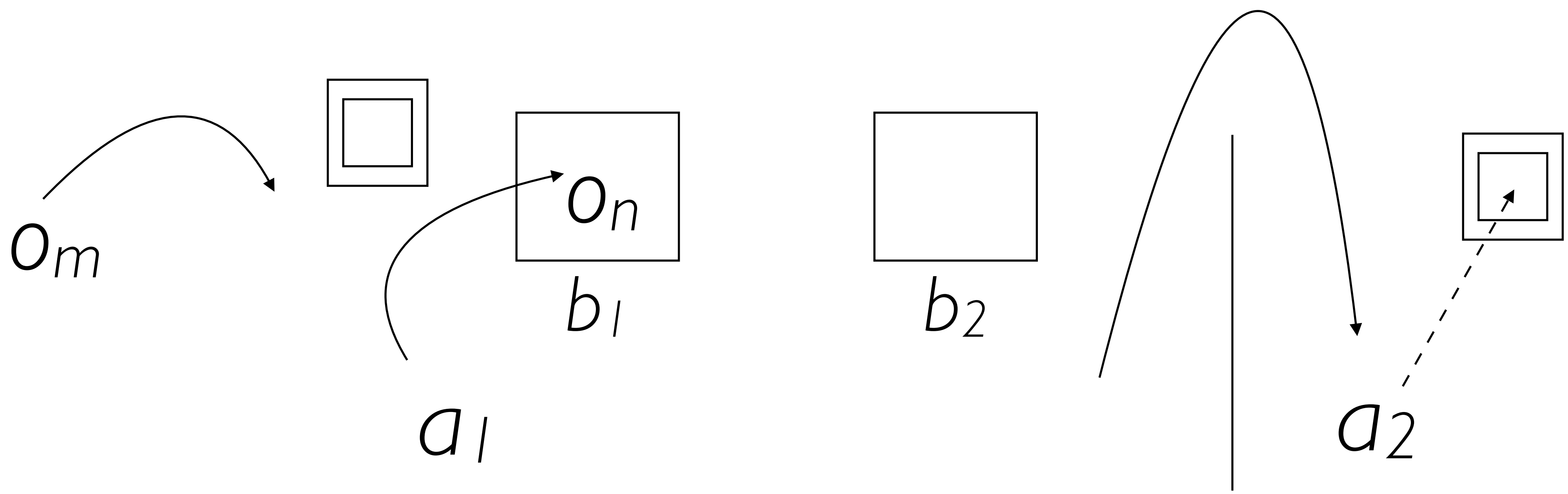
Framework for FBT^1_3
(eight timepoints)



Framework for FBT^{1}_{3}
(eight timepoints)



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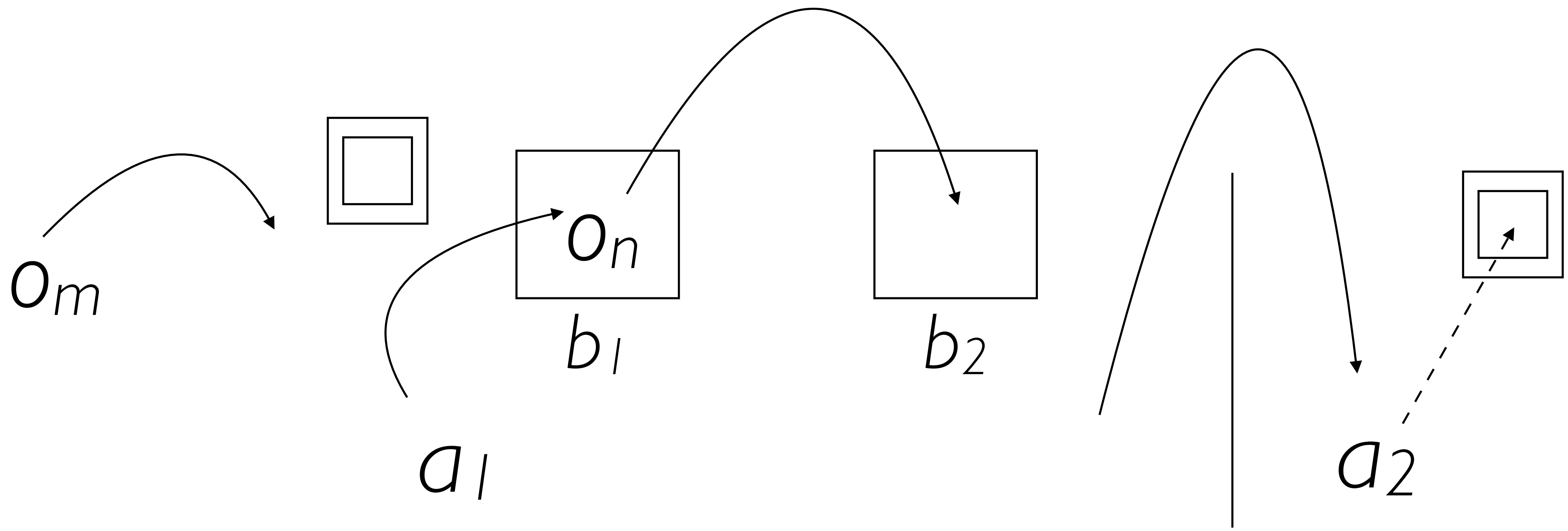
Framework for FBT^1_3

(eight timepoints)



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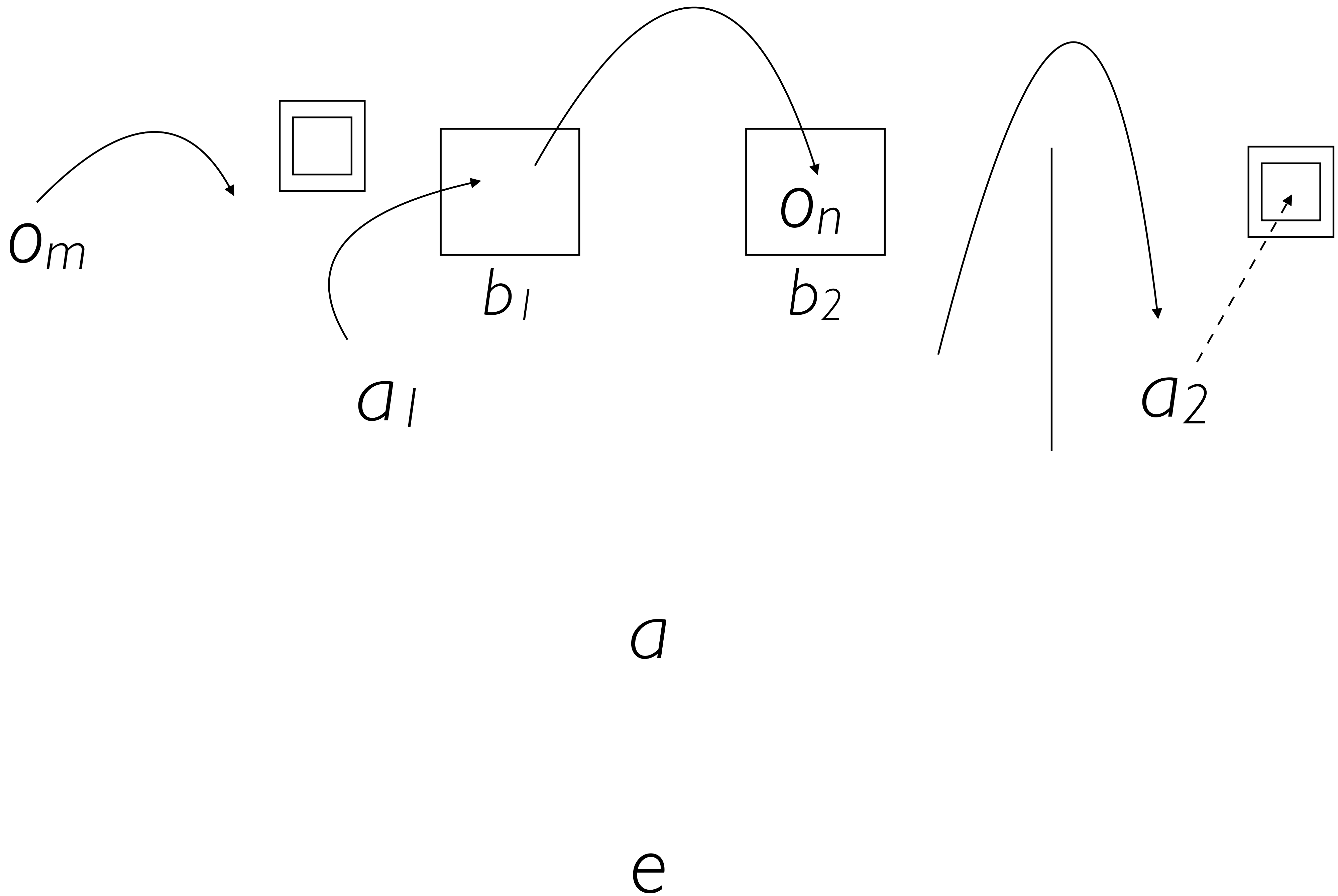
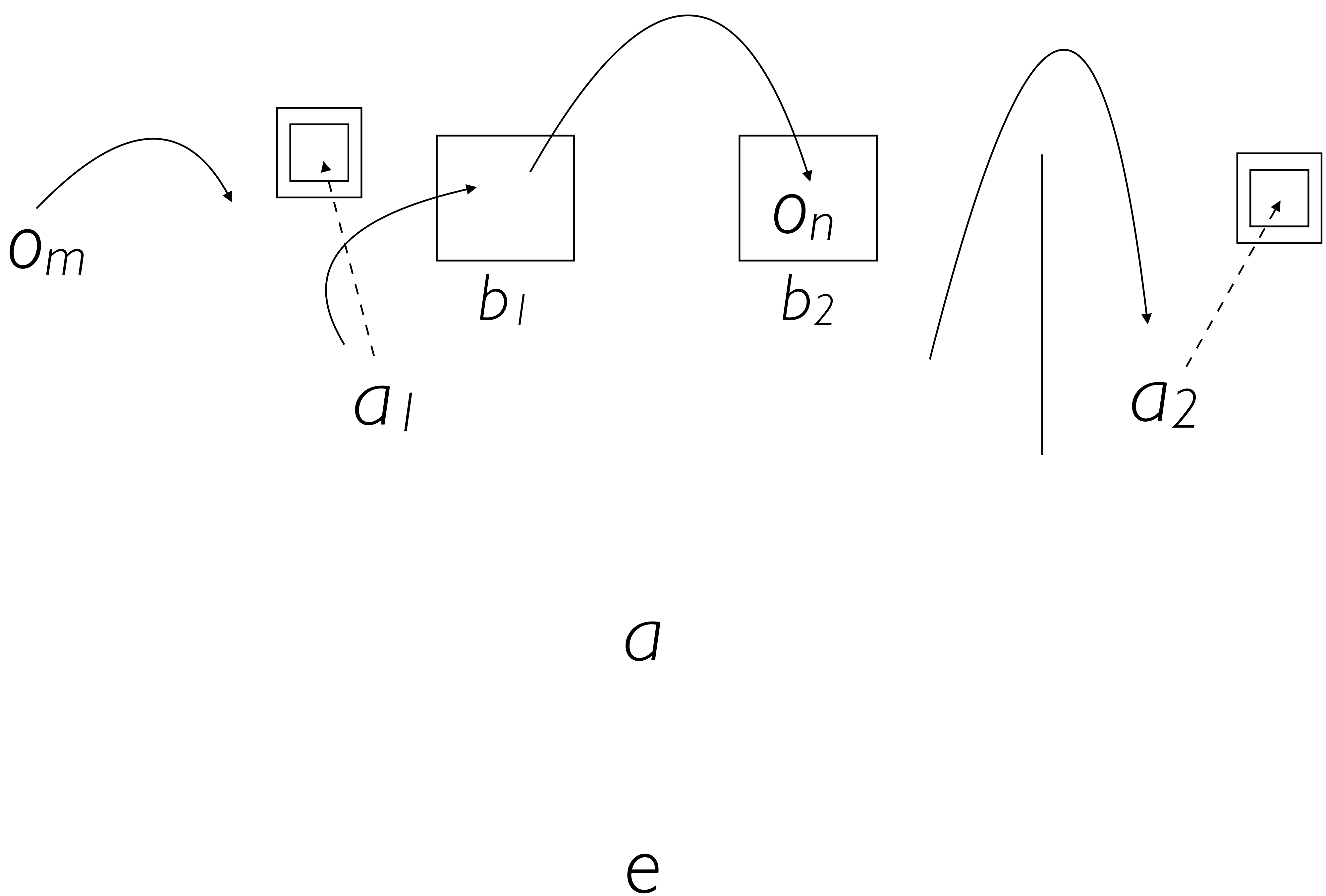
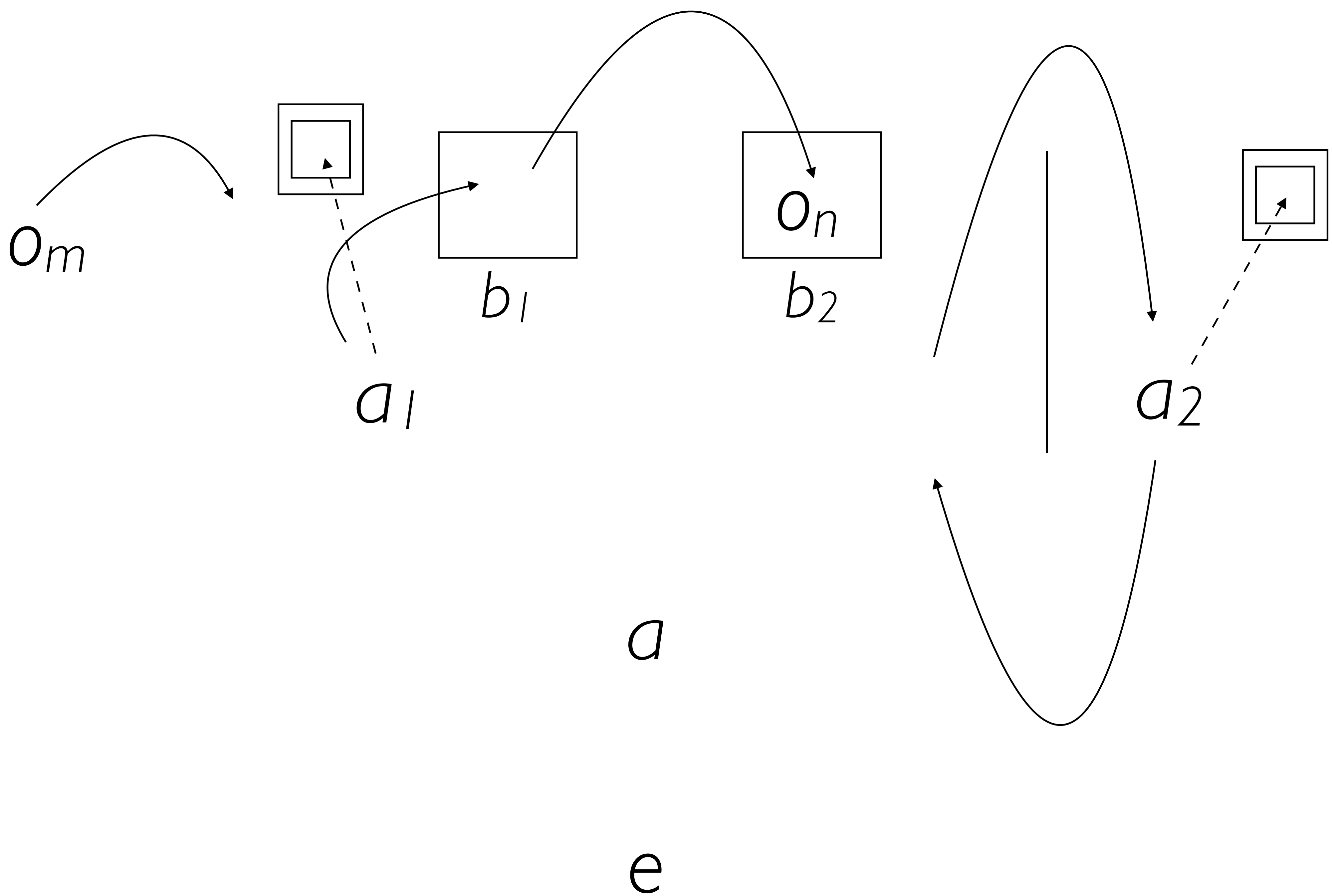
Framework for FBT^{1}_{3}
(eight timepoints)



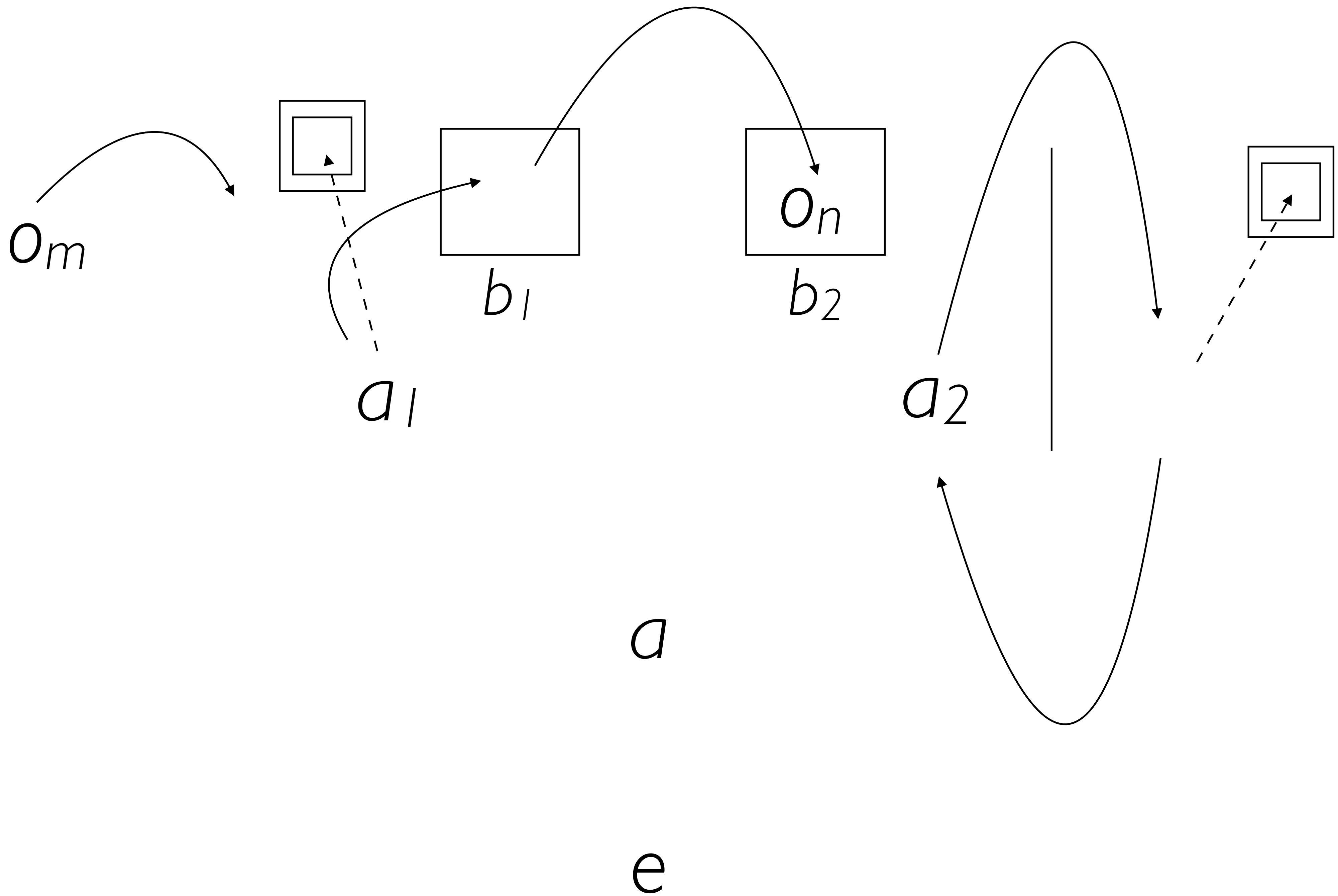
Framework for $FBT^1{}_3$
(eight timepoints)



Framework for FBT^1_3
(eight timepoints)

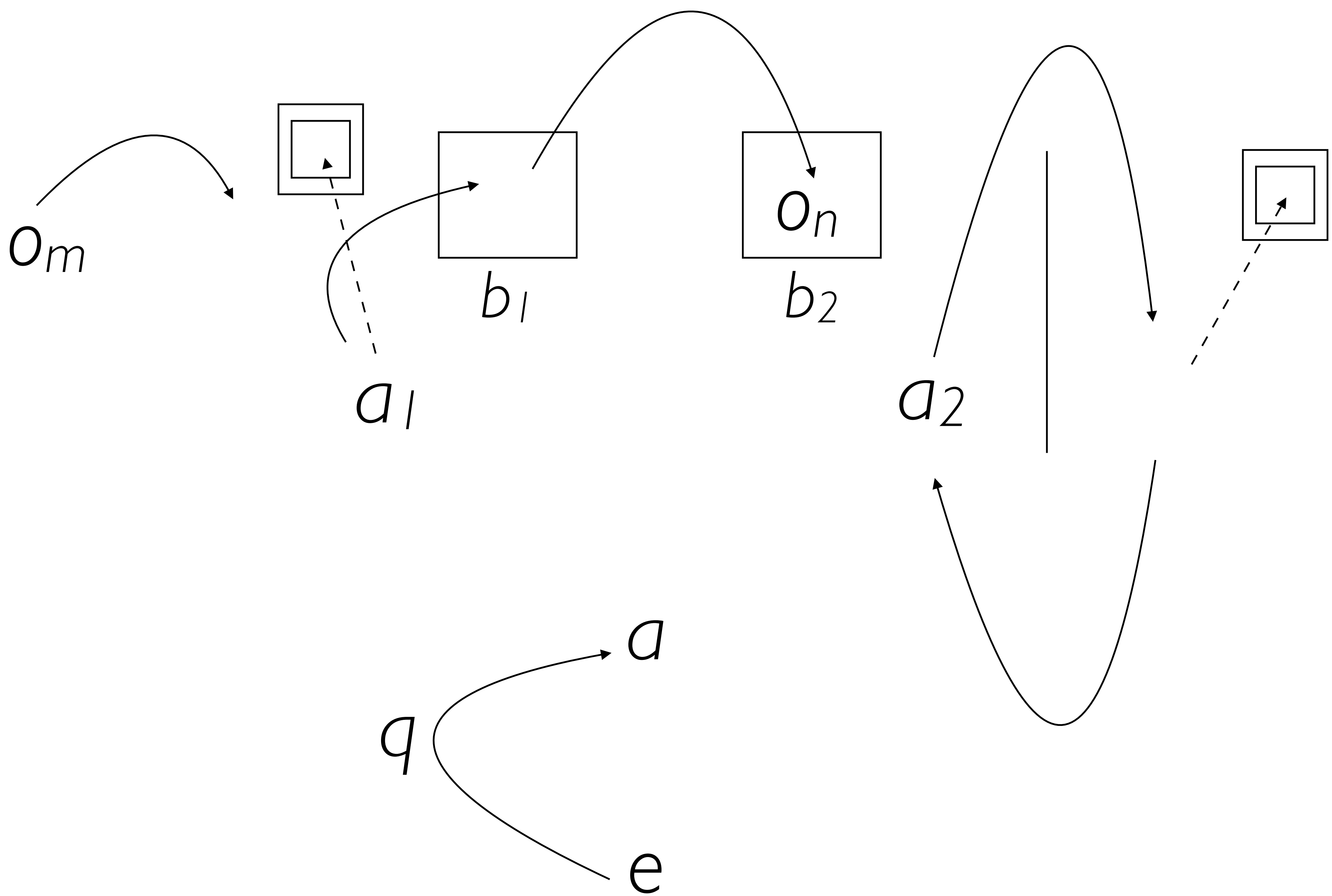


Framework for FBT^1_3
(eight timepoints)

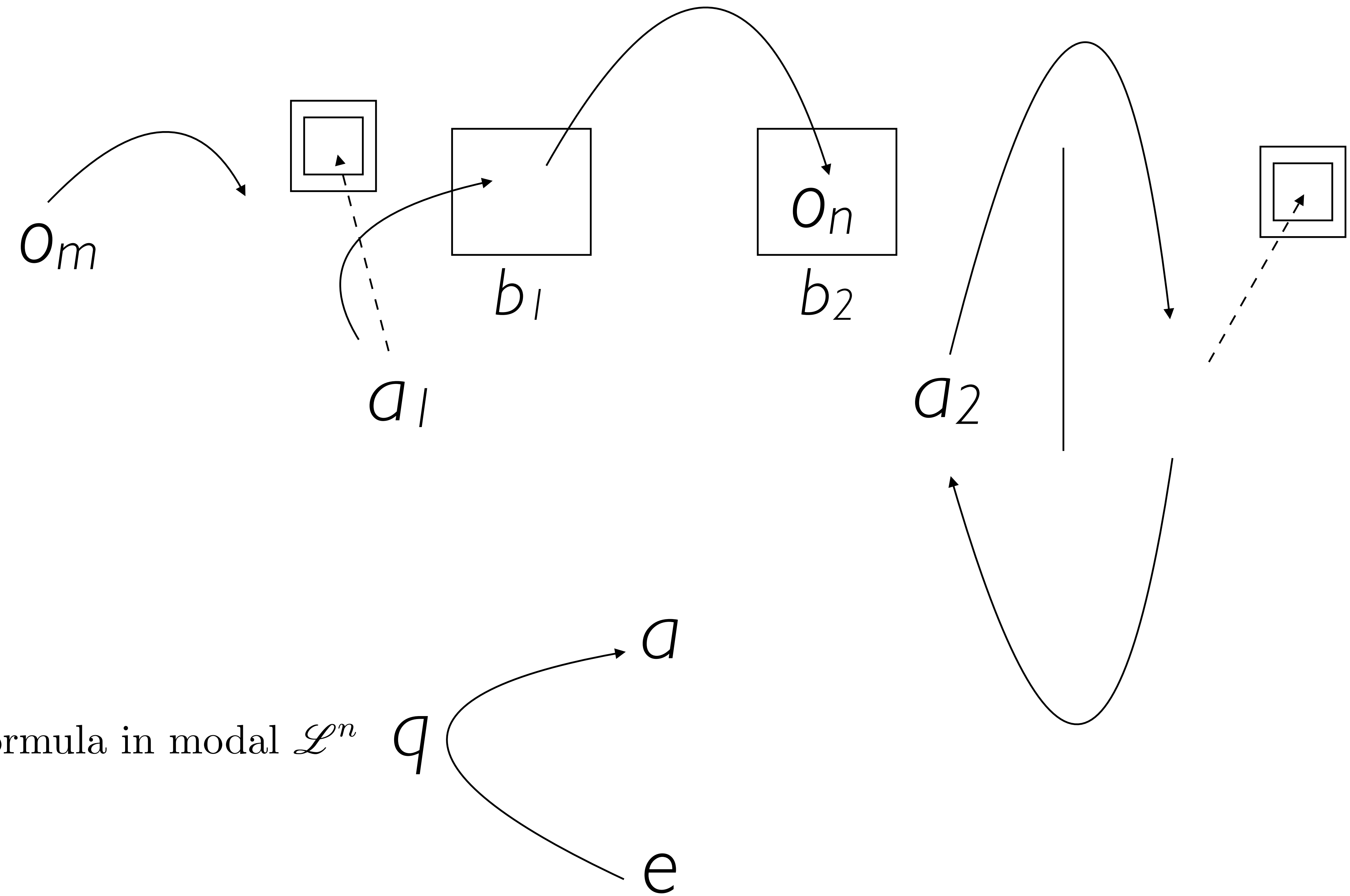


Framework for FBT^{1}_{3}

(eight timepoints)

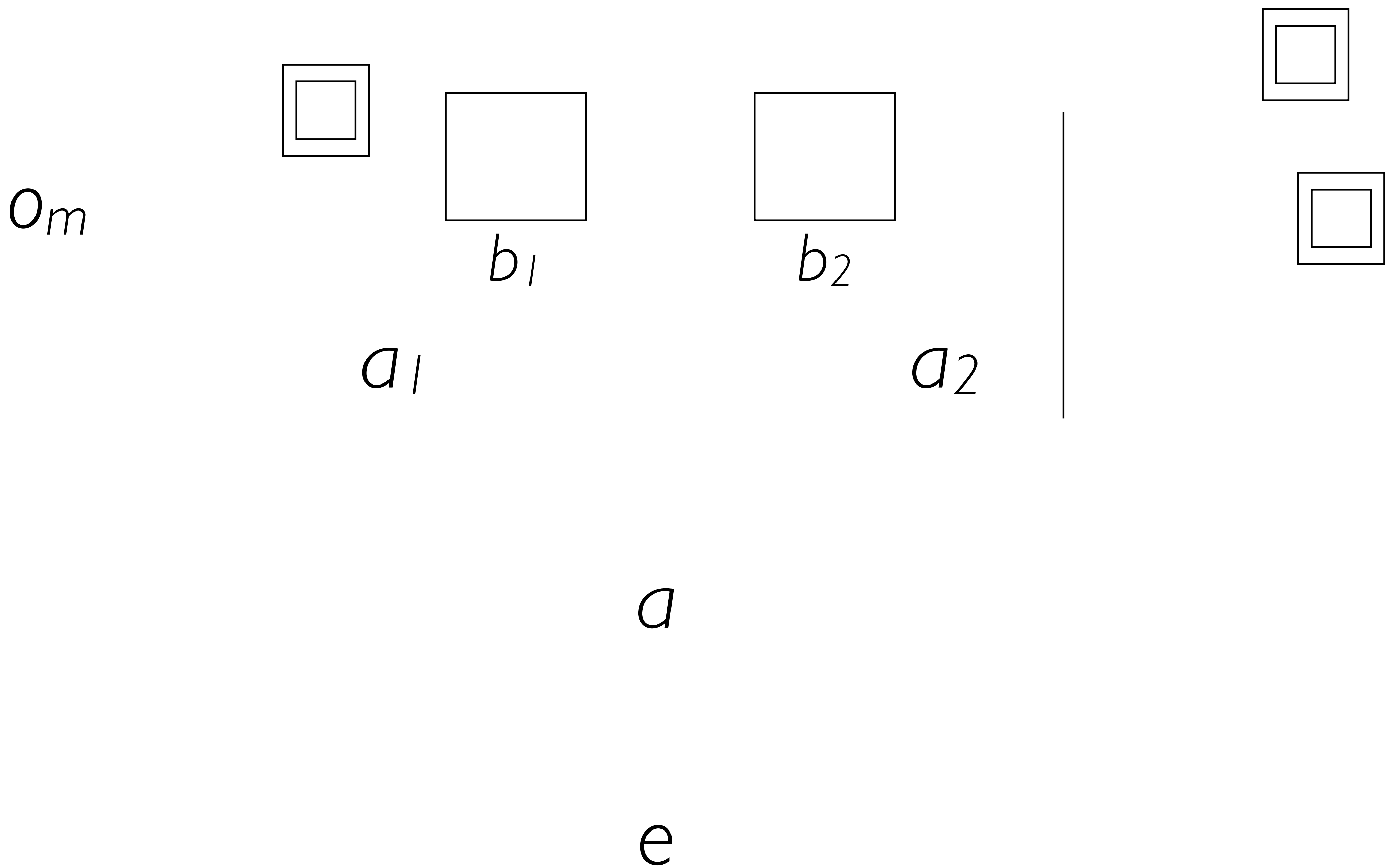


Framework for $FBT^1{}_3$
(eight timepoints)

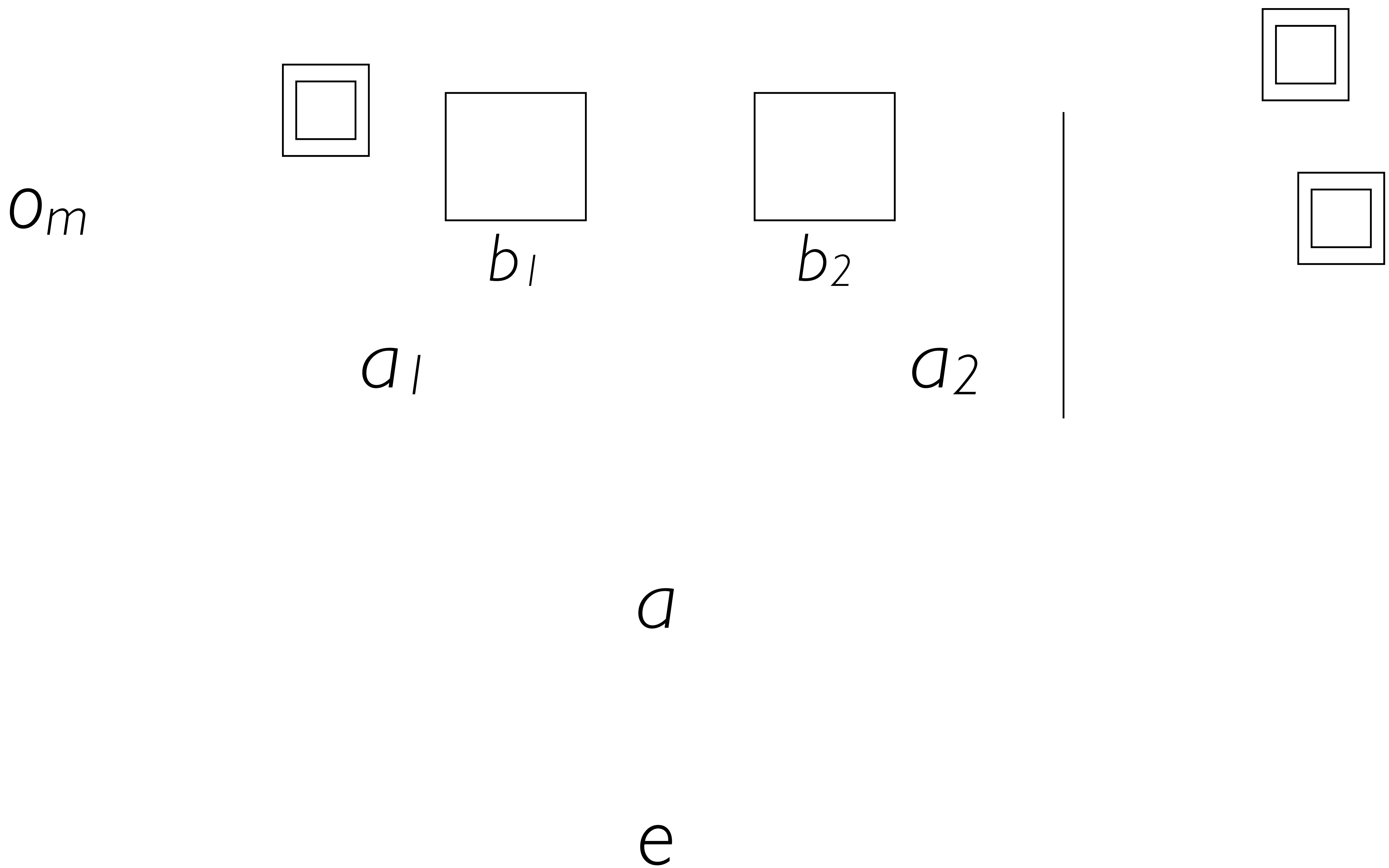


q a formula in modal $\mathscr{L}^n$

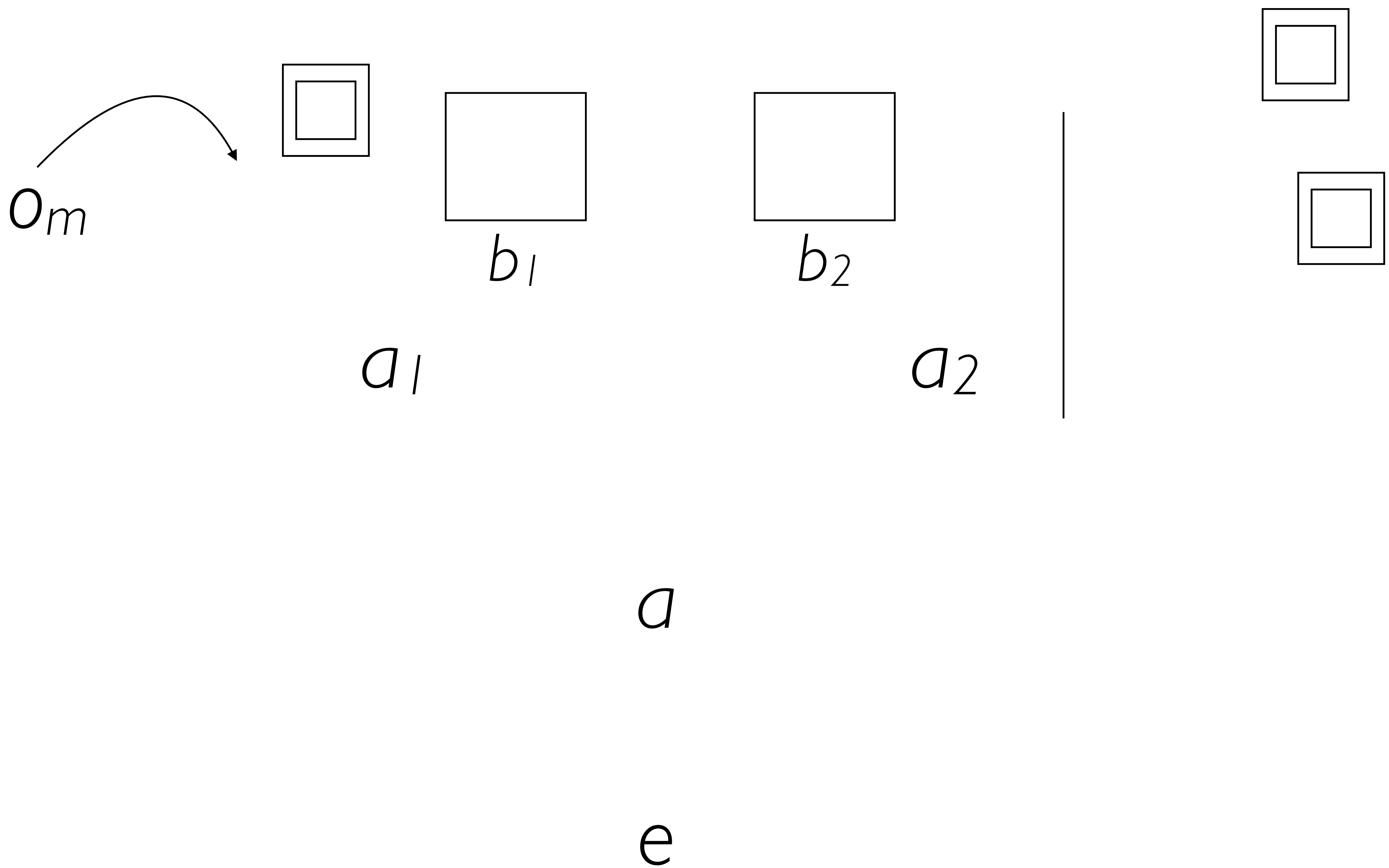
Framework for FBT^1_4



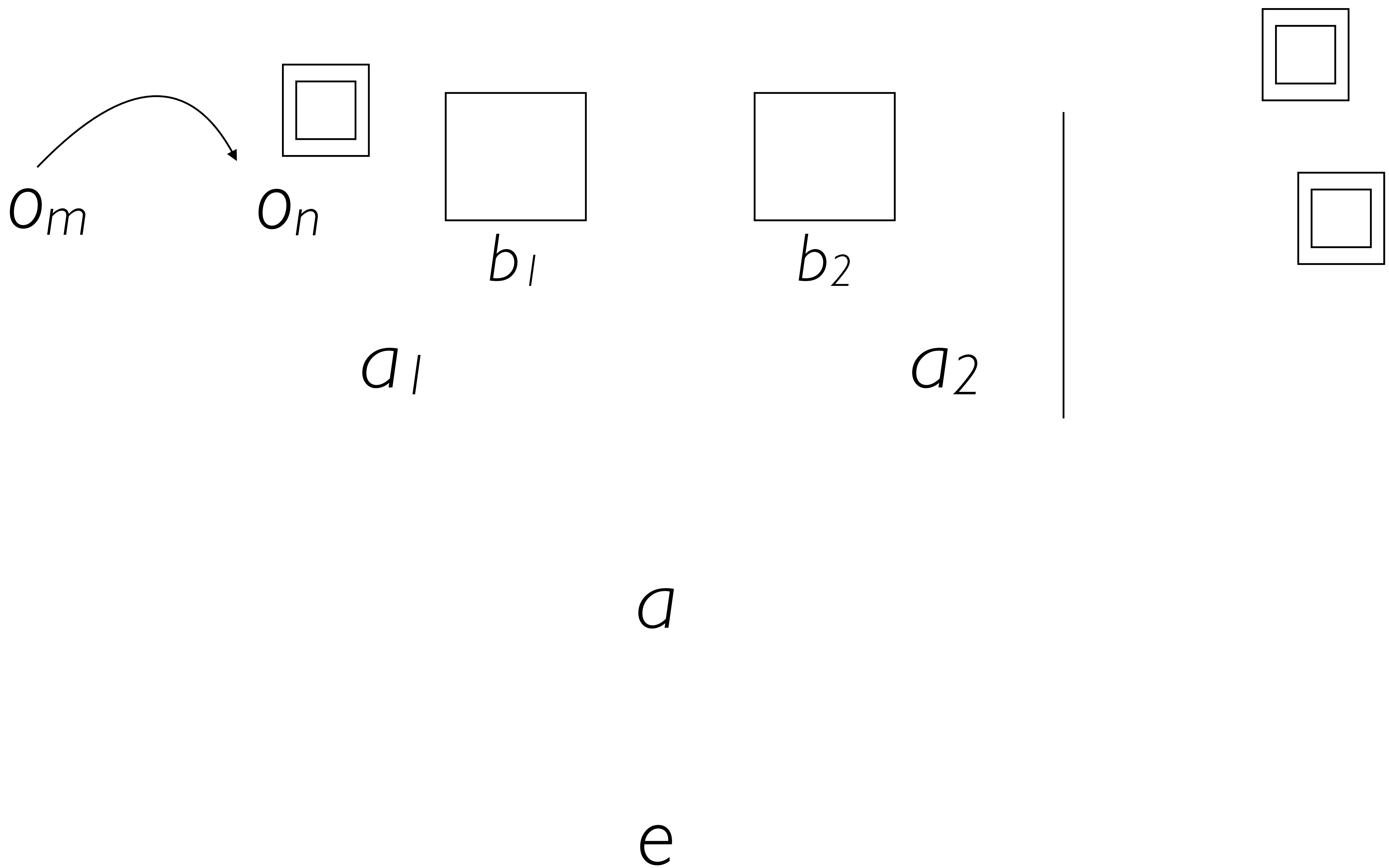
Framework for FBT^1_4
(nine timepoints)



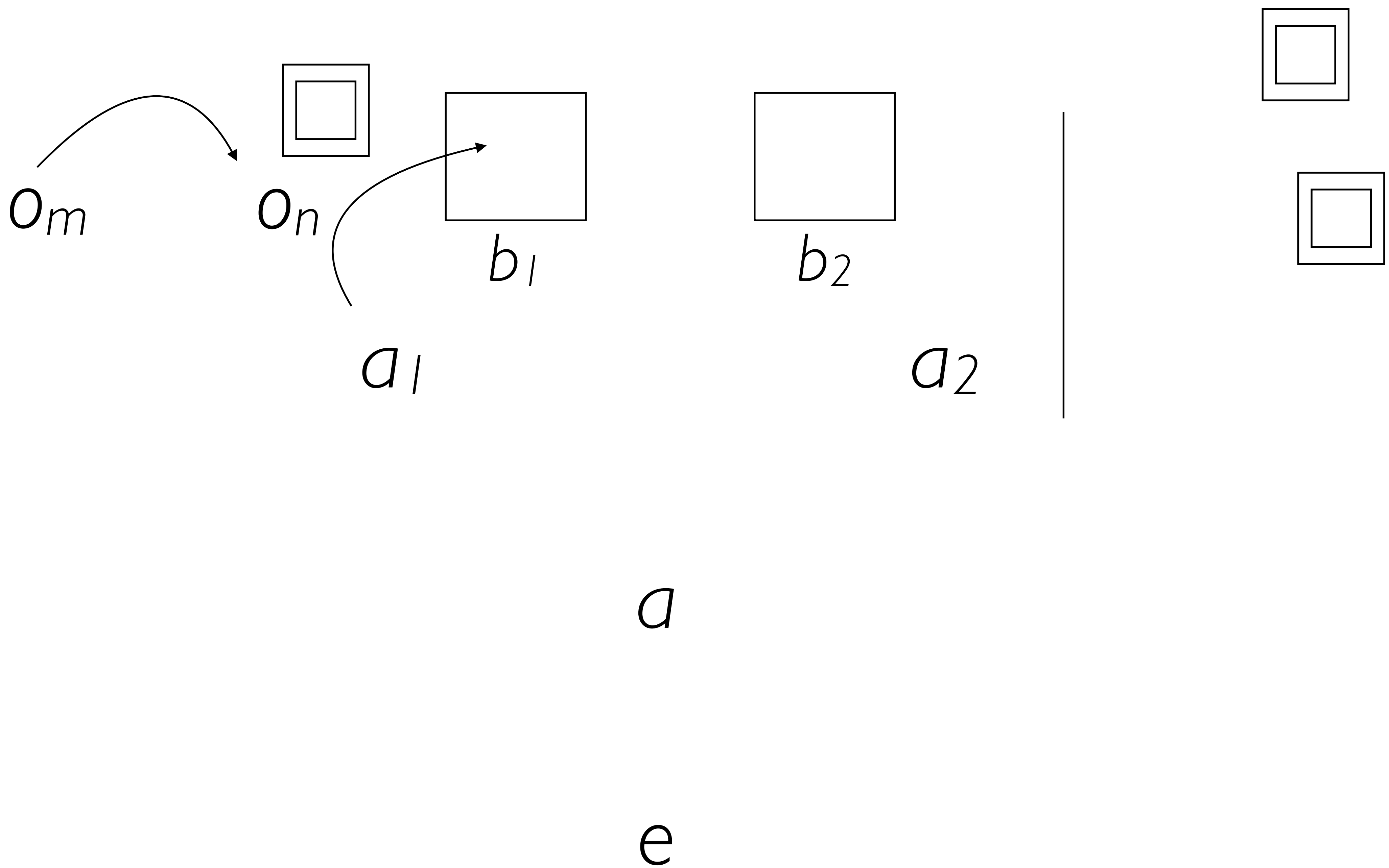
Framework for FBT^1_4
(nine timepoints)



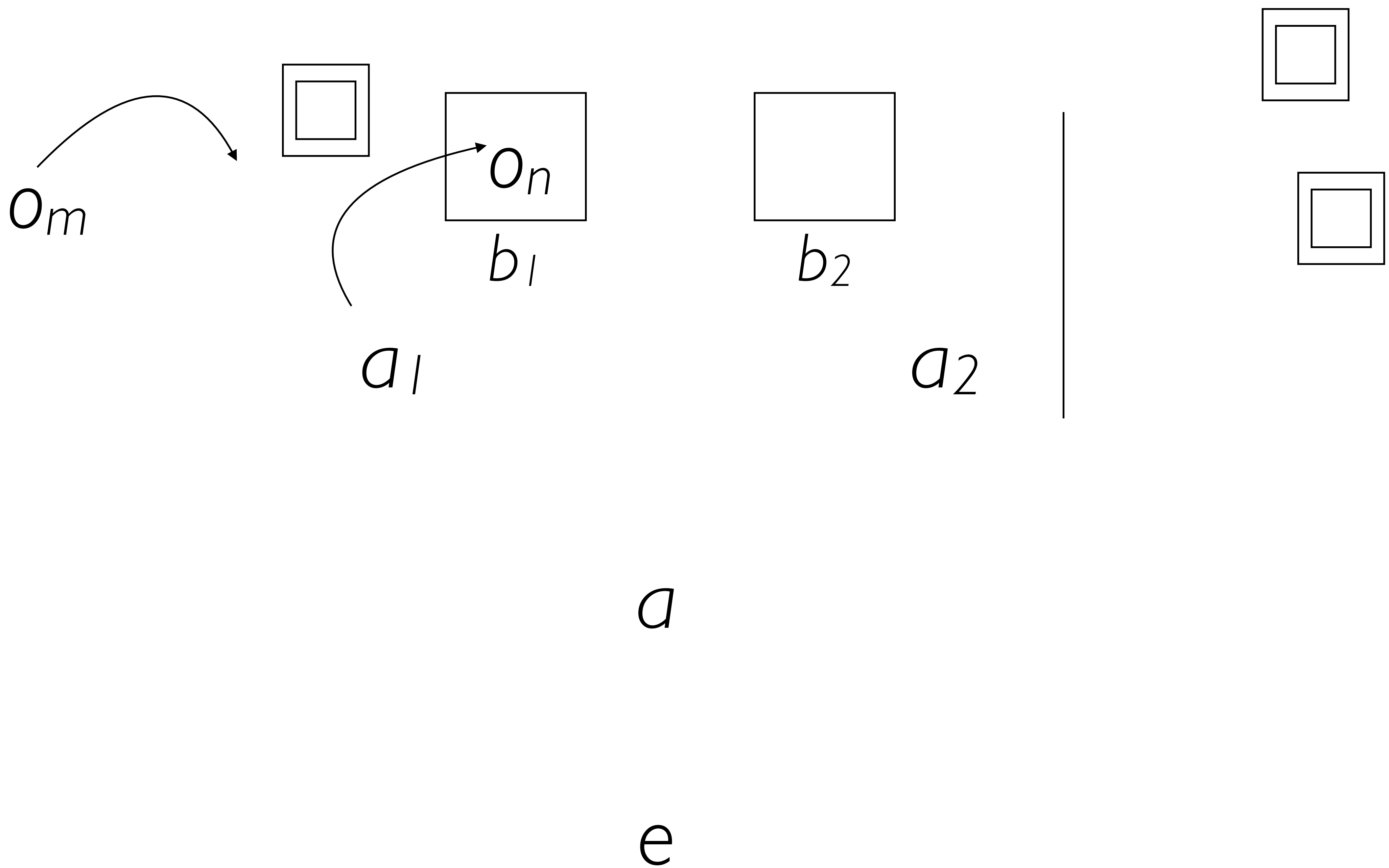
Framework for FBT¹₄
(nine timepoints)



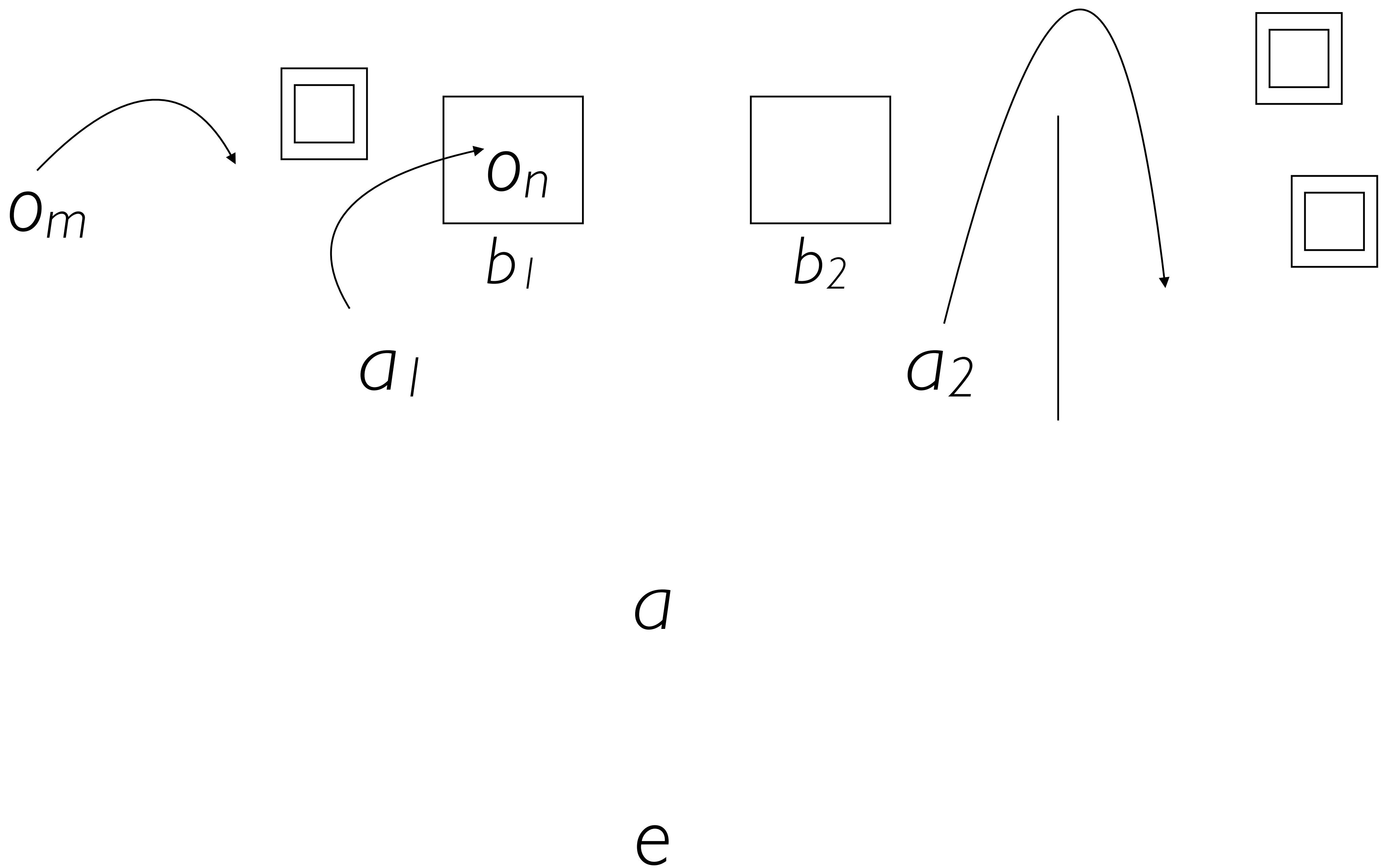
Framework for $FBT^1{}_4$
(nine timepoints)



Framework for $FBT^1{}_4$
(nine timepoints)

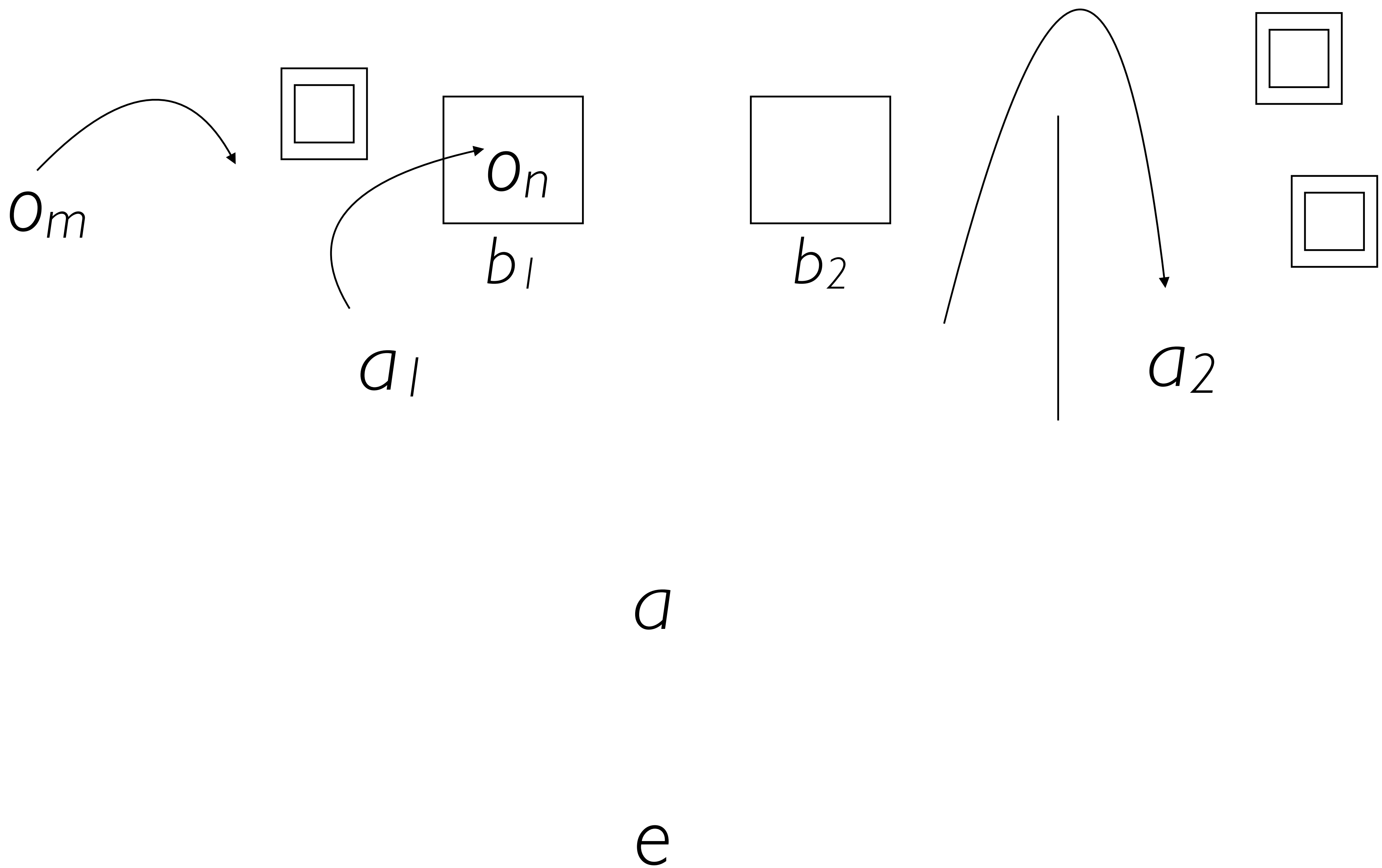


Framework for $FBT^1{}_4$
(nine timepoints)

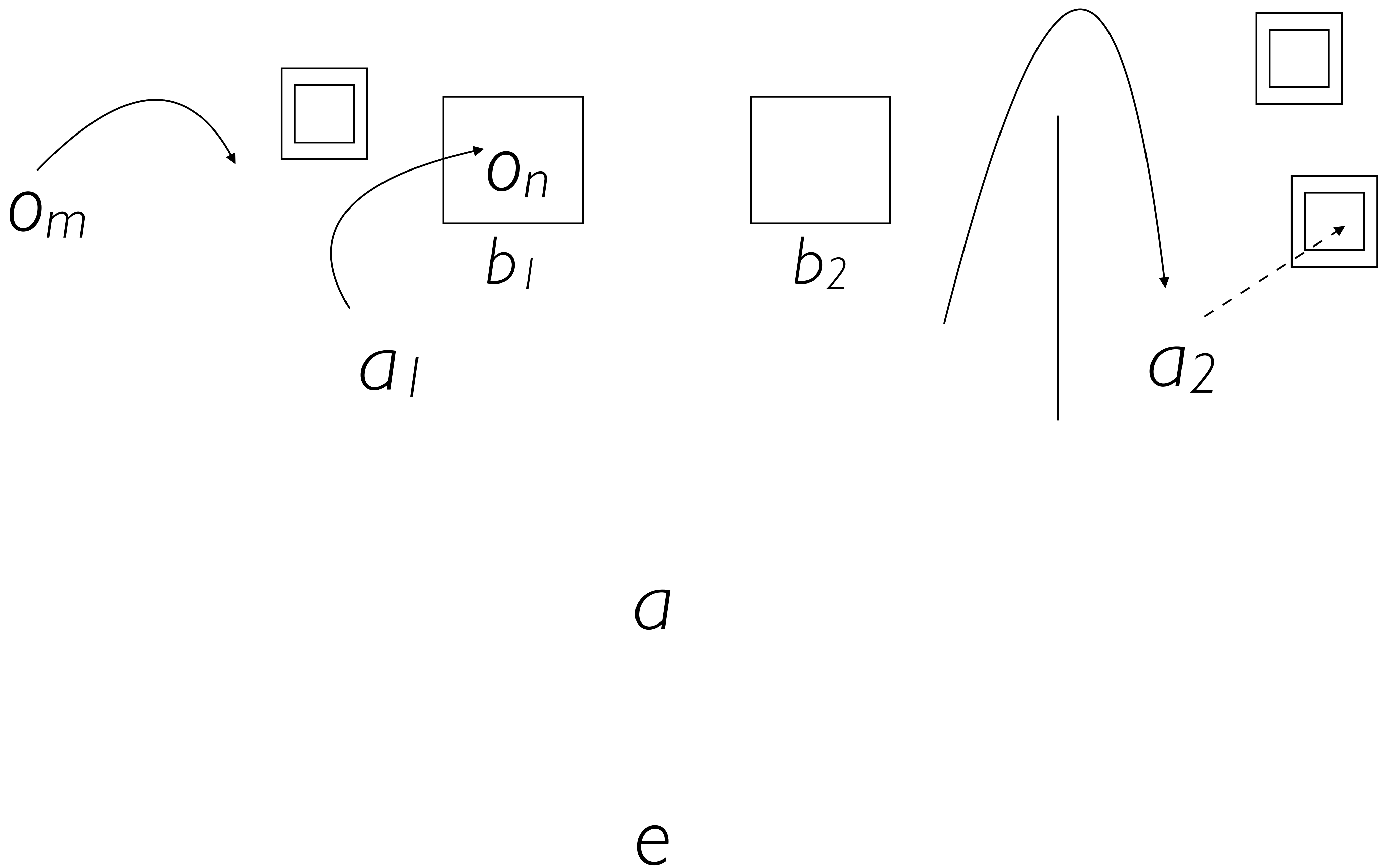


Framework for $FBT^1{}_4$

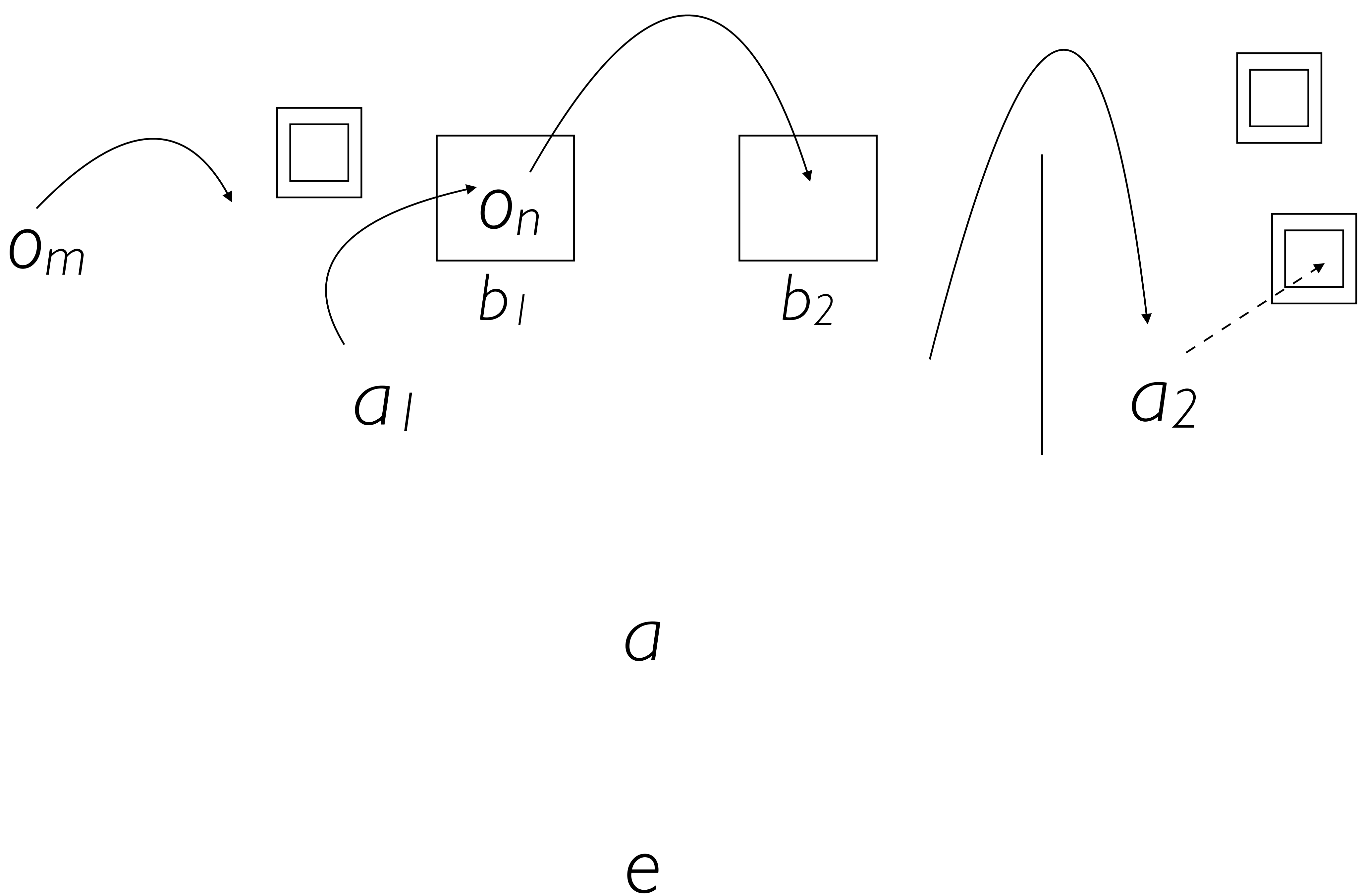
(nine timepoints)



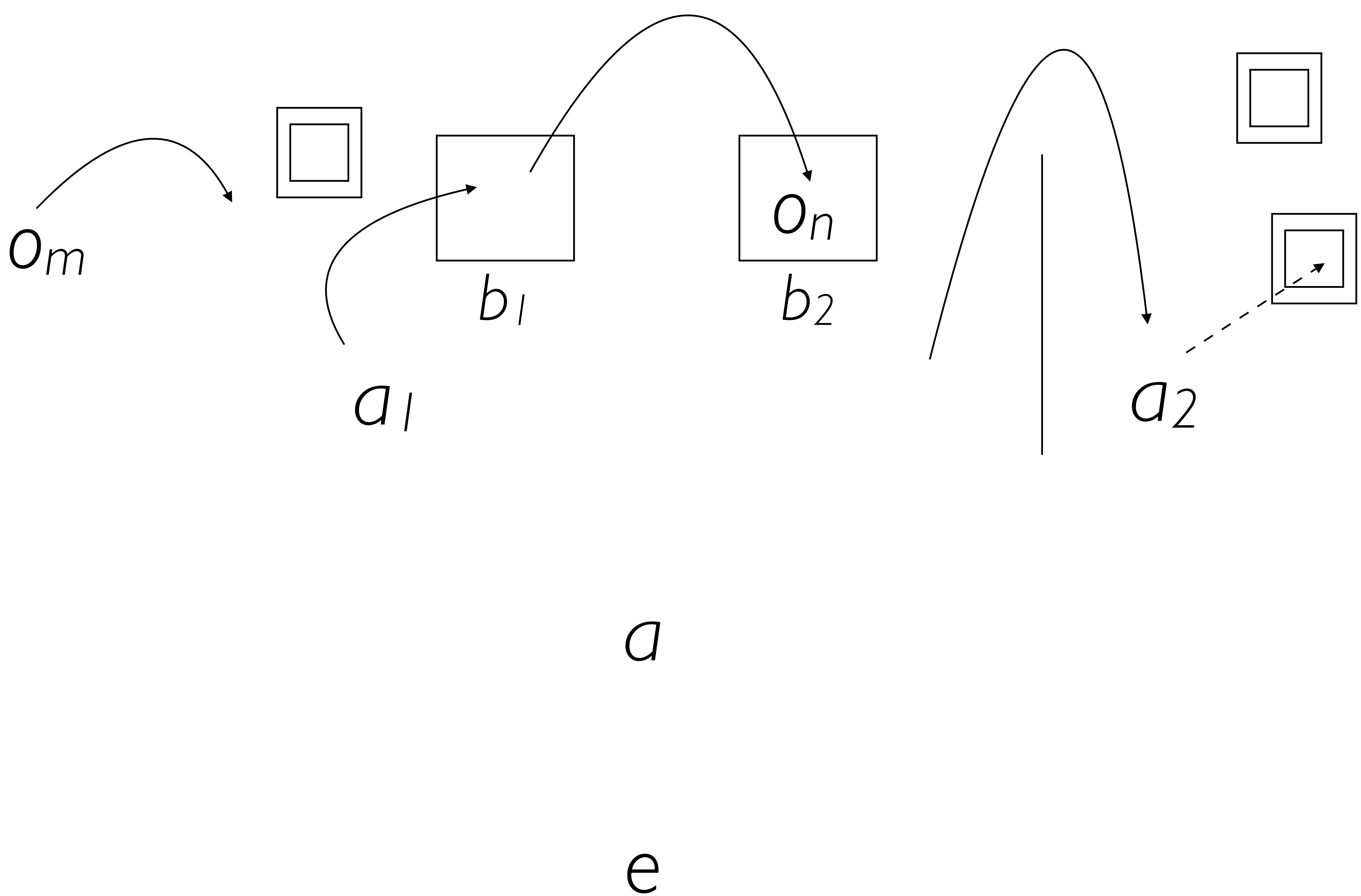
Framework for FBT$^1{}_4$
(nine timepoints)



Framework for FBT^{1}_{4}
(nine timepoints)

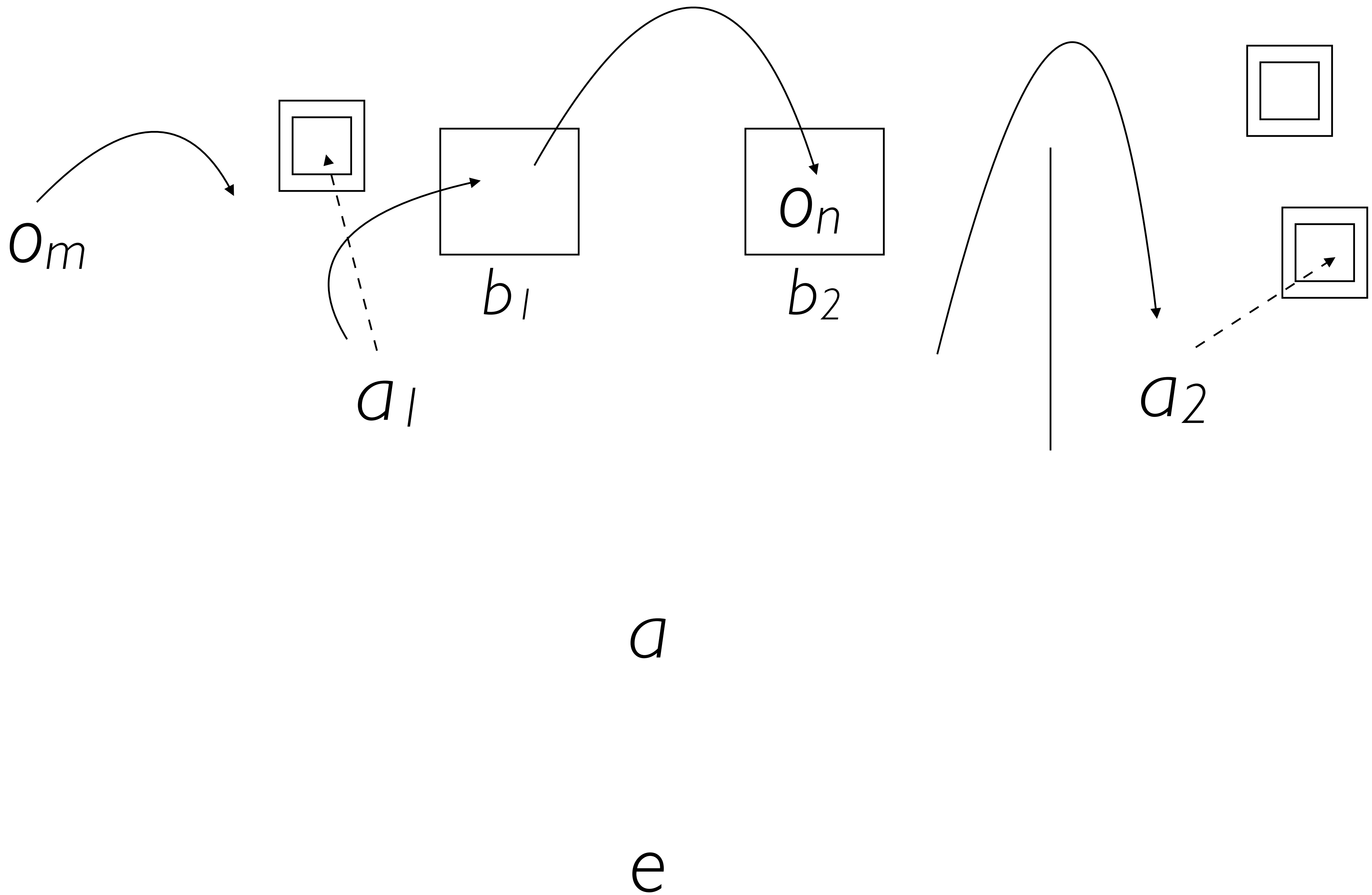


Framework for FBT$^1_4$
(nine timepoints)

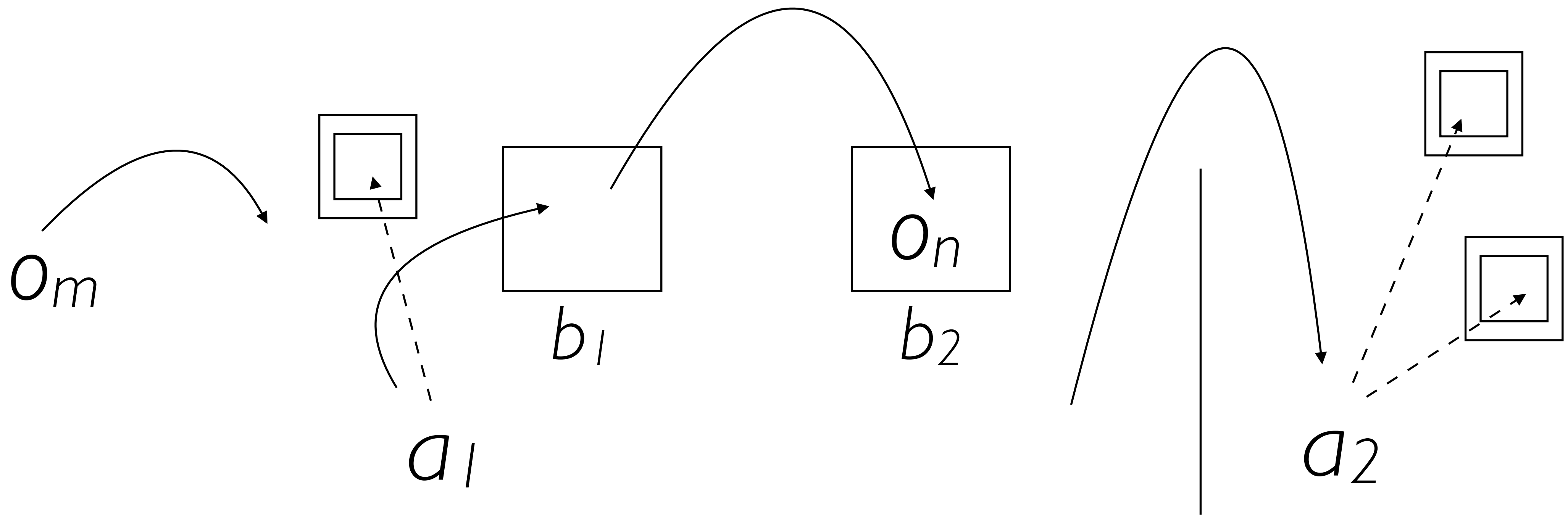


Framework for $FBT^1{}_4$

(nine timepoints)



Framework for FBT^1_4
(nine timepoints)



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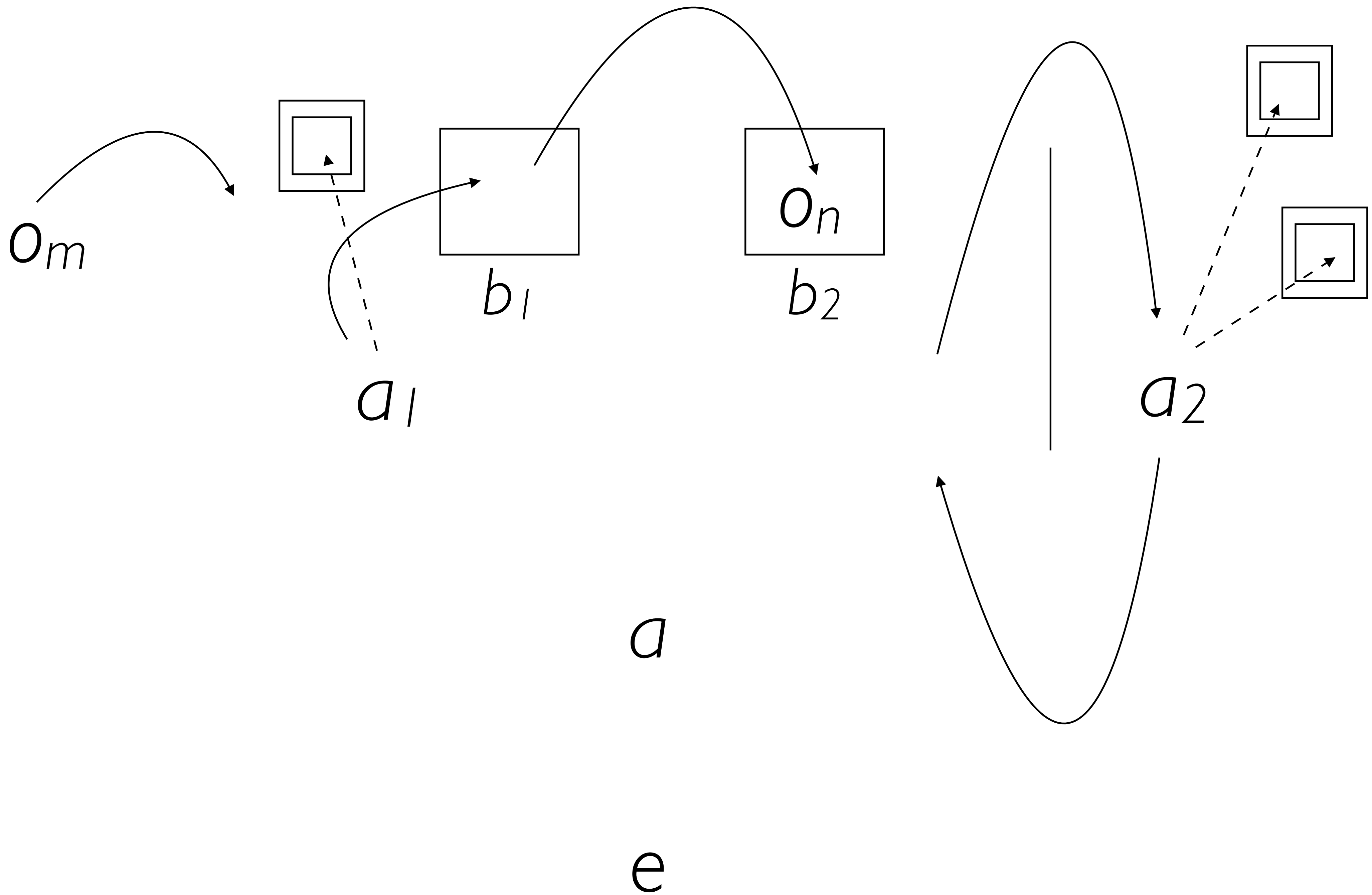
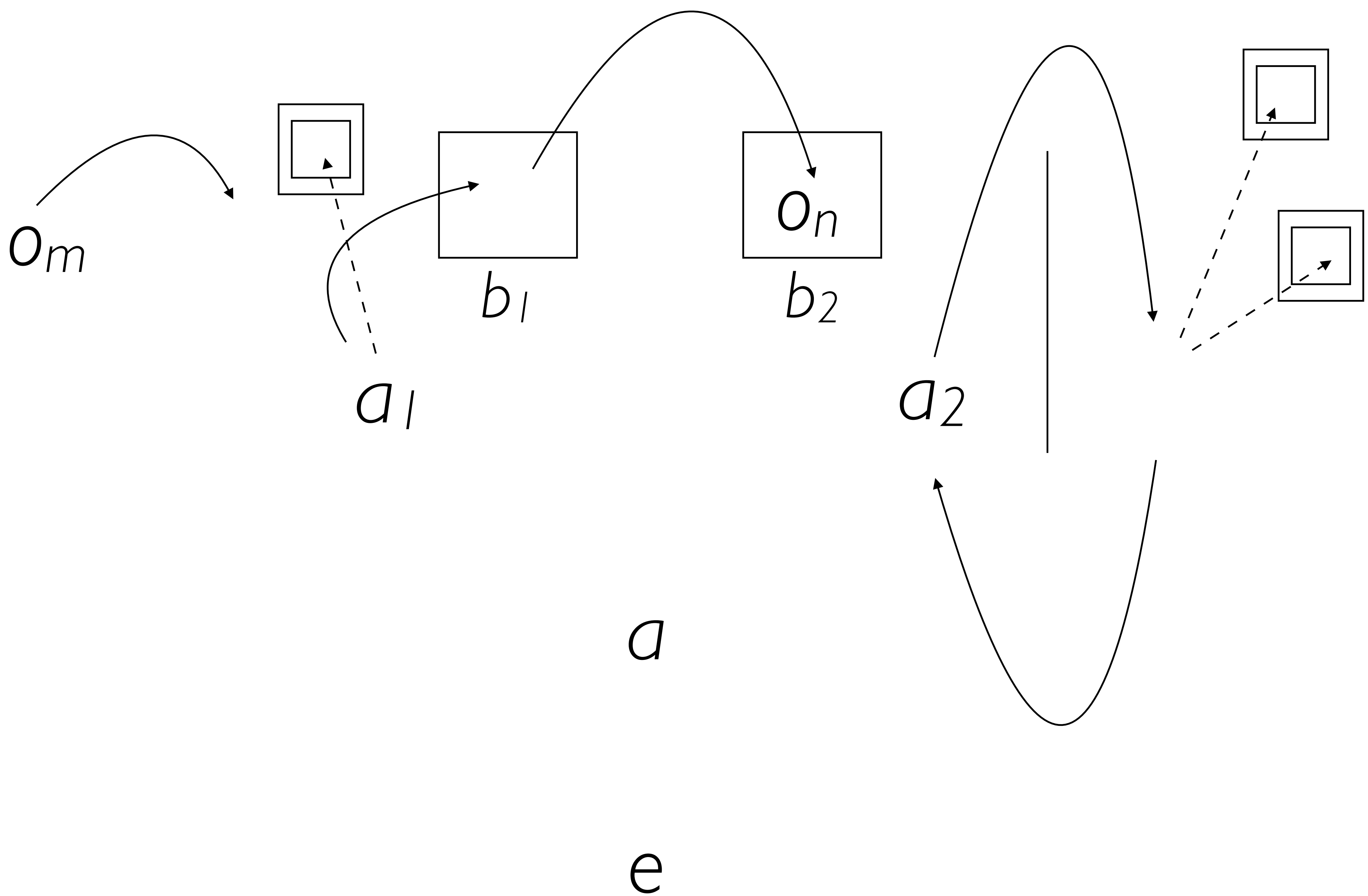
Framework for $FBT^1{}_4$
(nine timepoints)

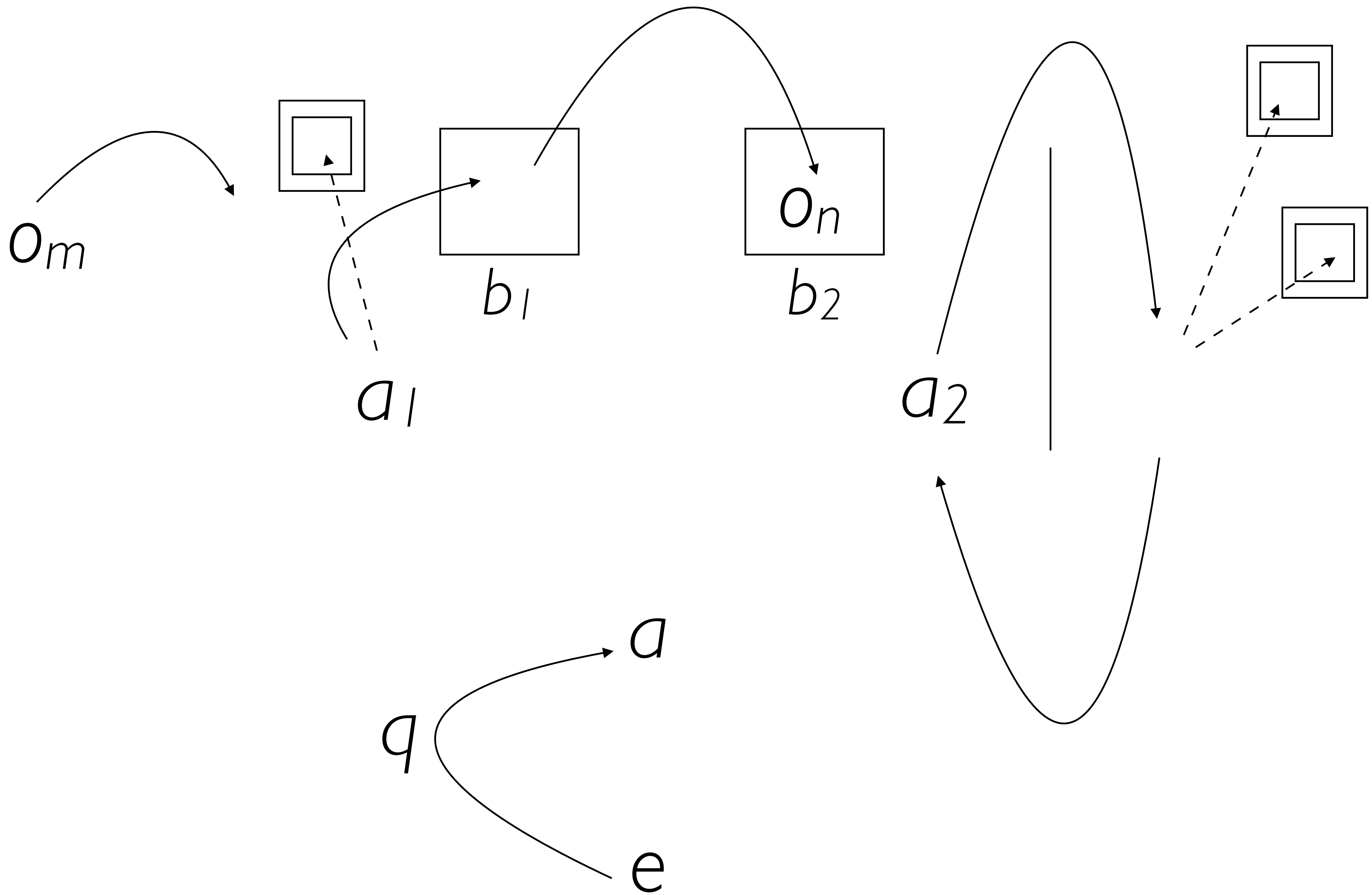


Framework for FBT'$_4$
(nine timepoints)



Framework for $FBT^1{}_4$

(nine timepoints)



Framework for $FBT^1{}_4$
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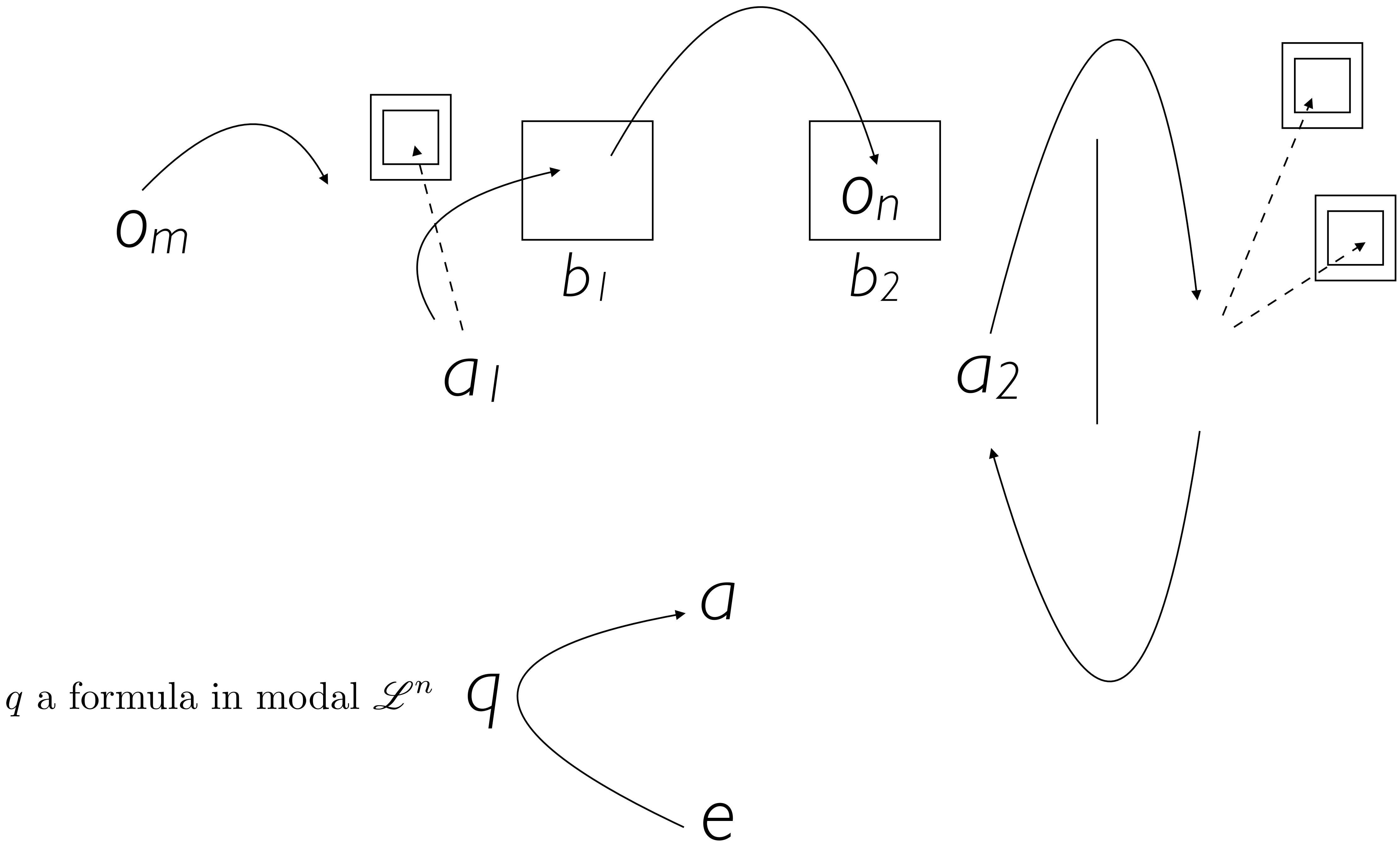
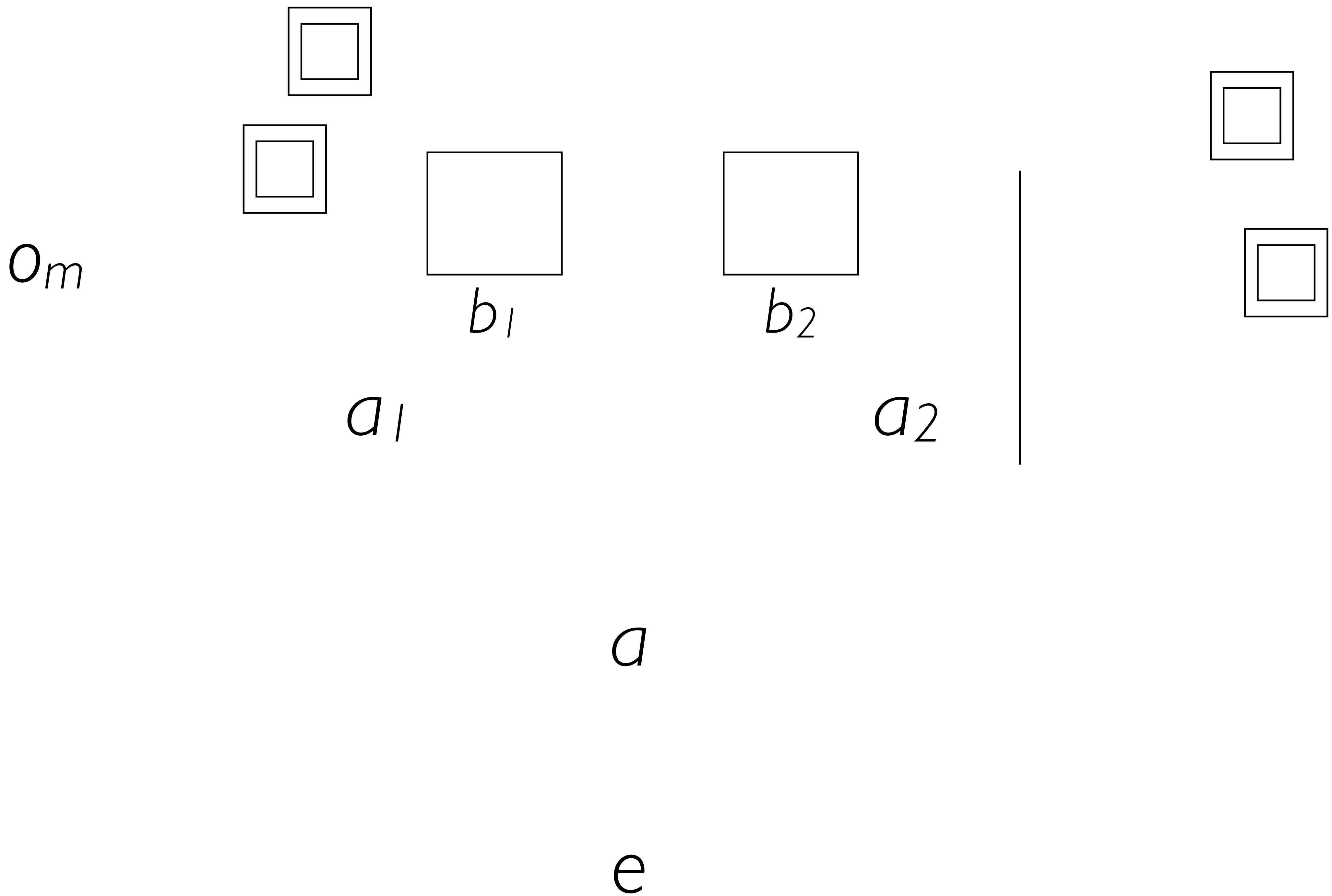
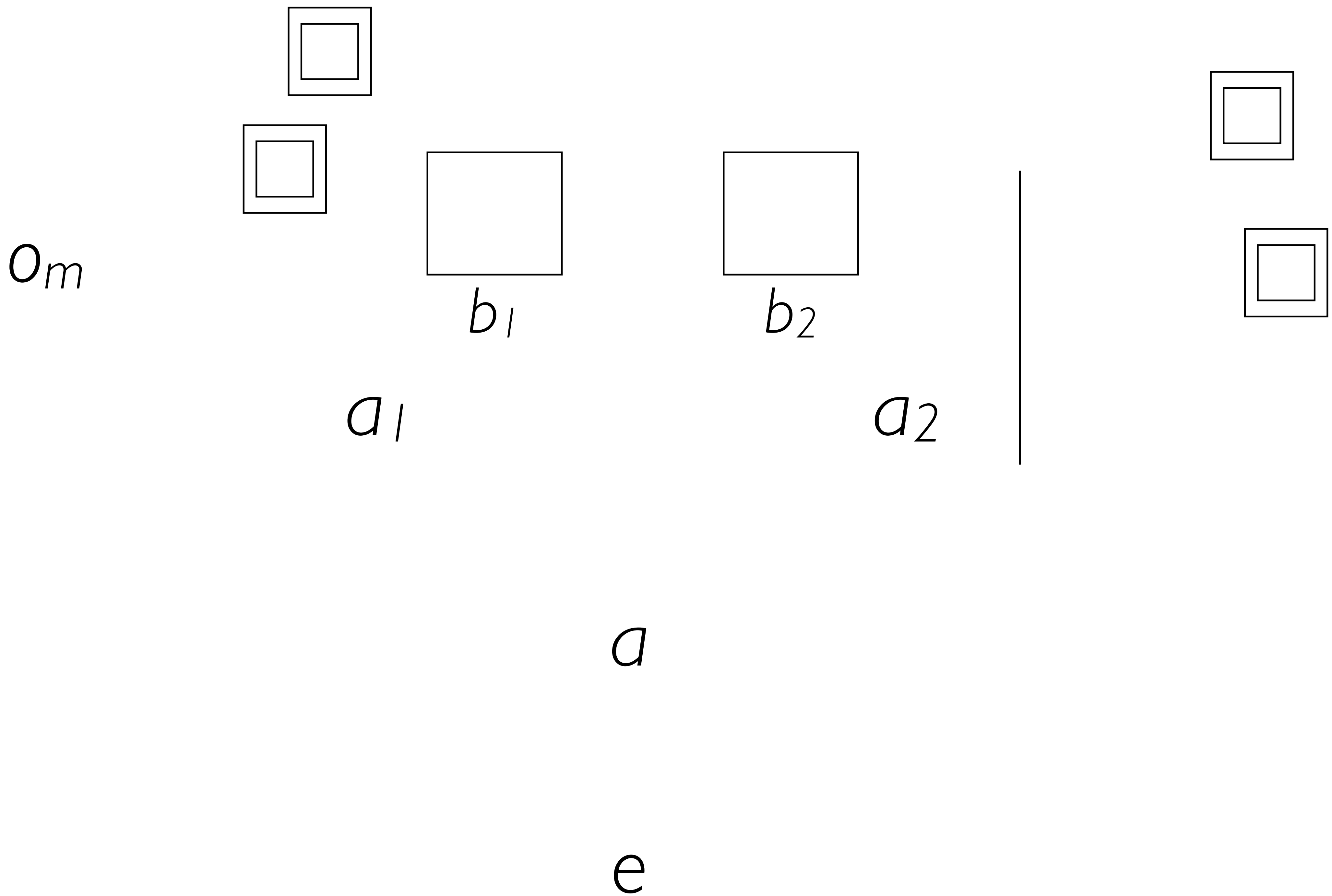


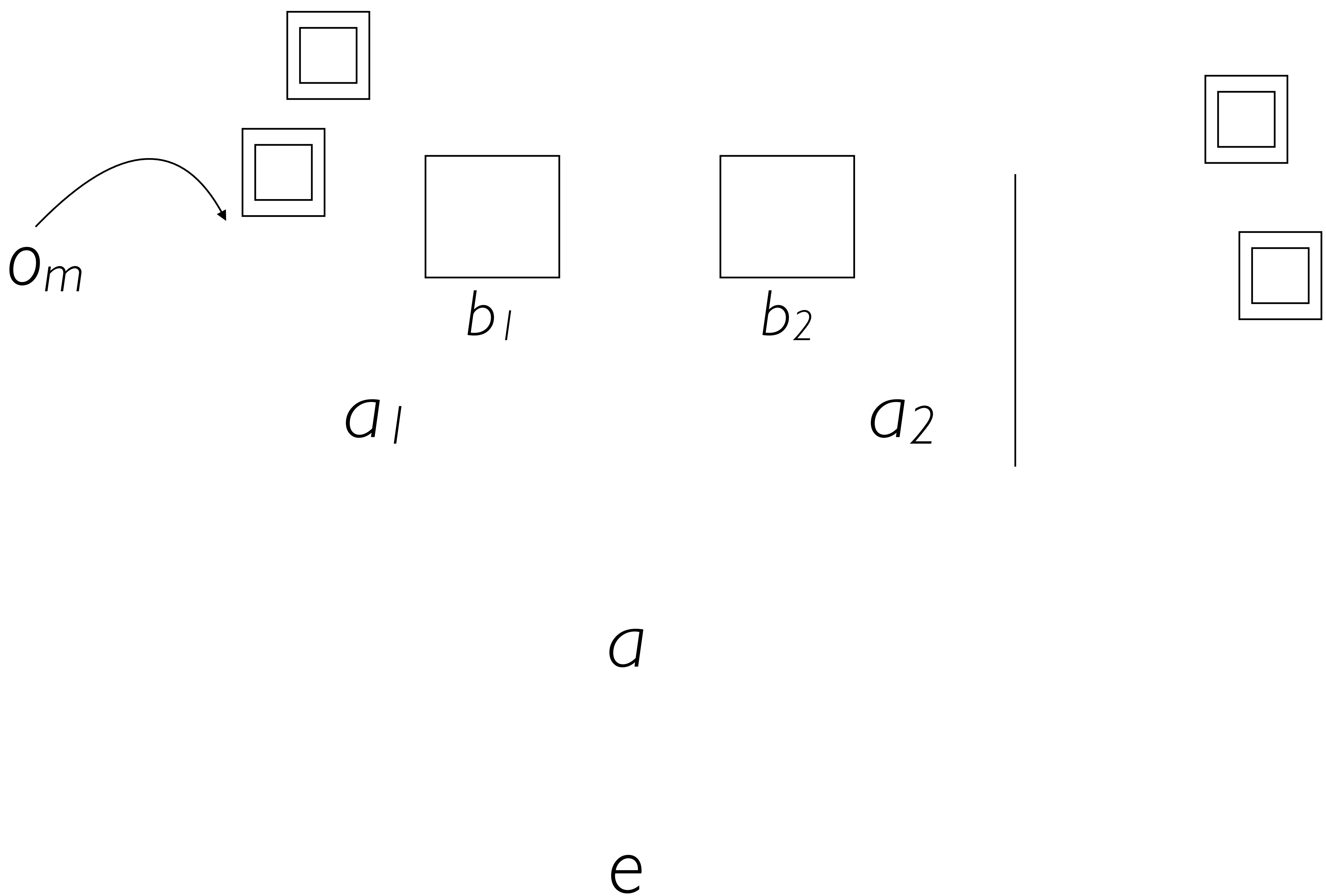
Framework for FBT^{1}_{5}



Framework for $FBT^1{}_5$
(ten timepoints)

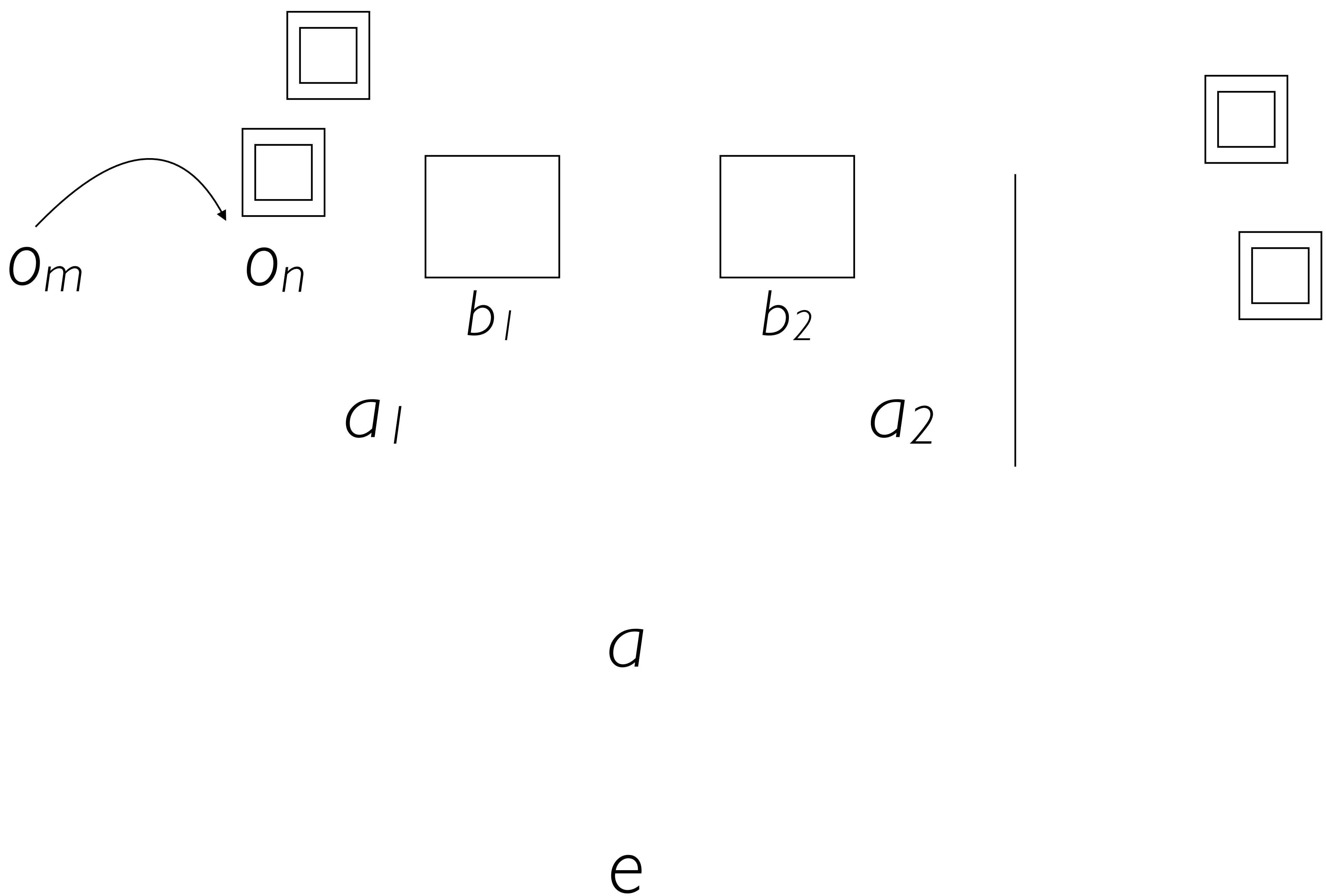


Framework for $FBT^1{}_5$
(ten timepoints)

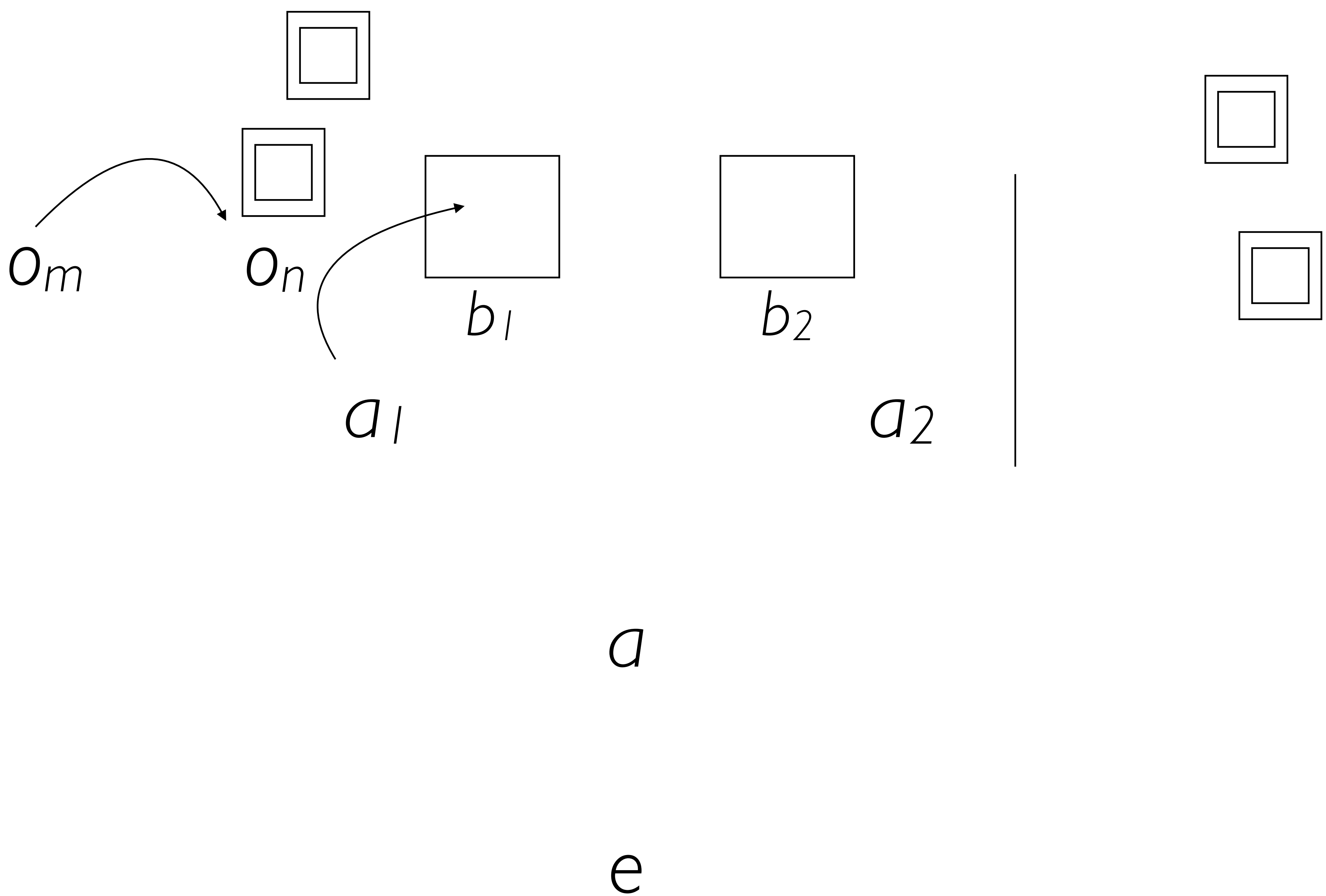


Framework for FBT^1_5

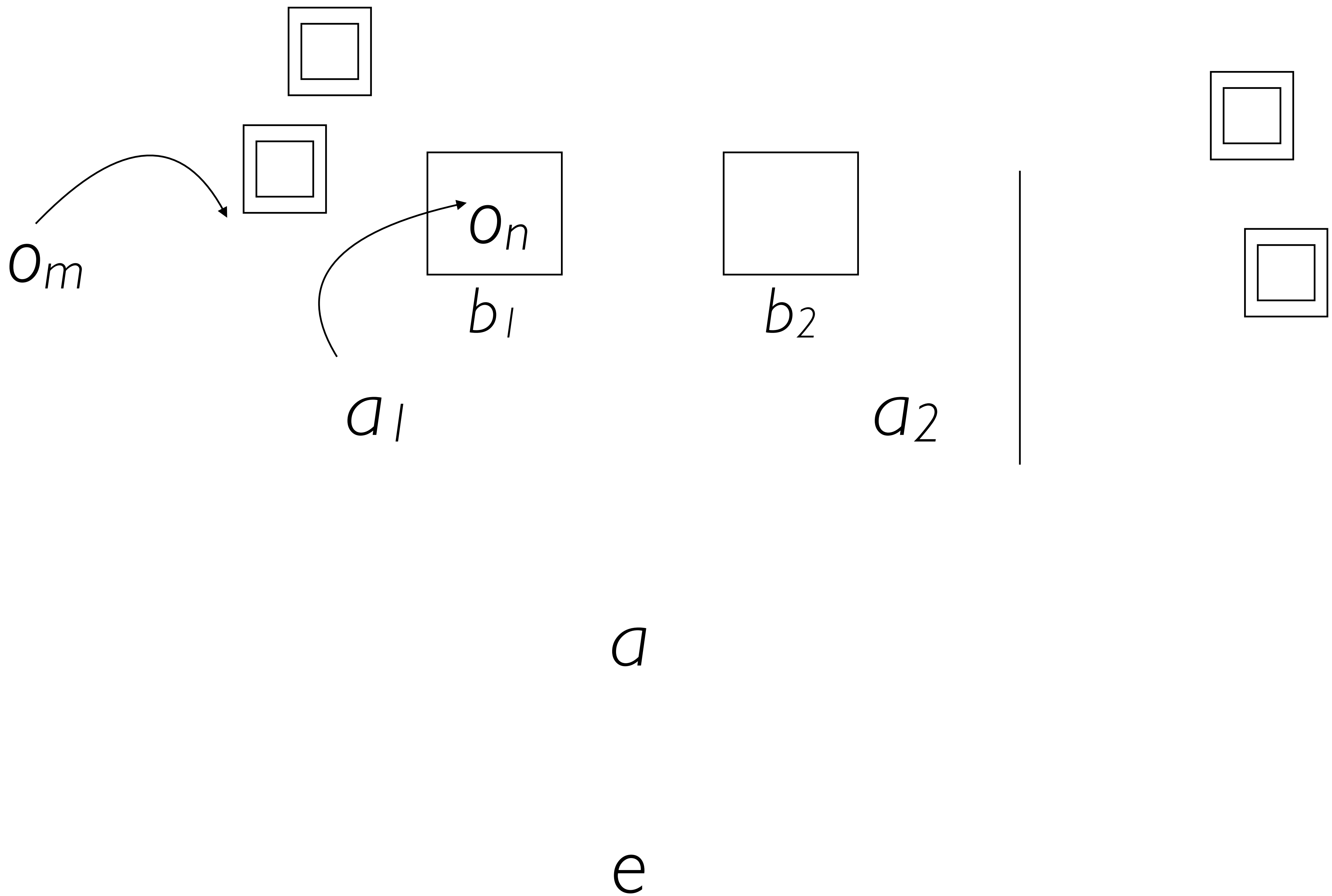
(ten timepoints)



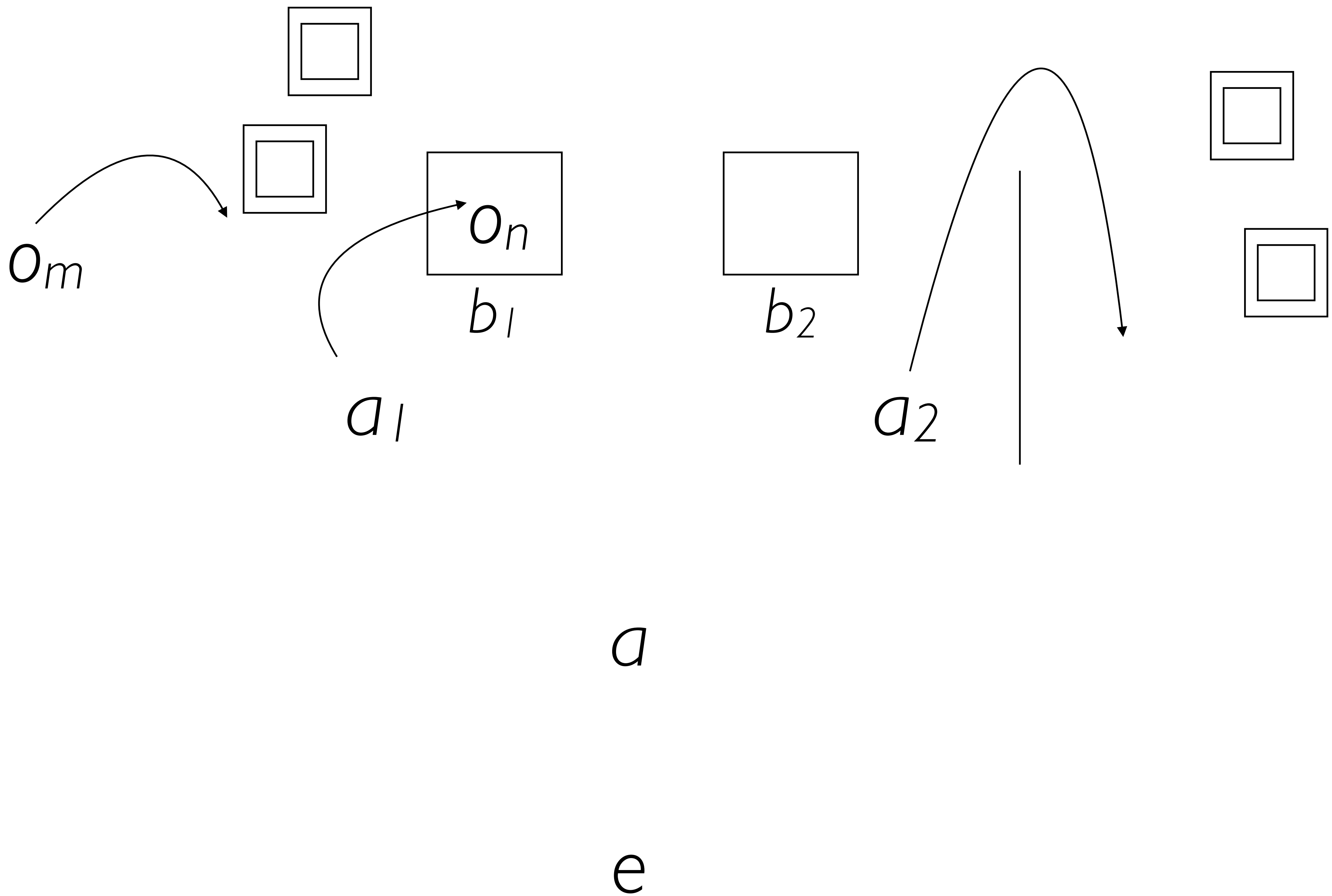
Framework for $FBT^1{}_5$
(ten timepoints)



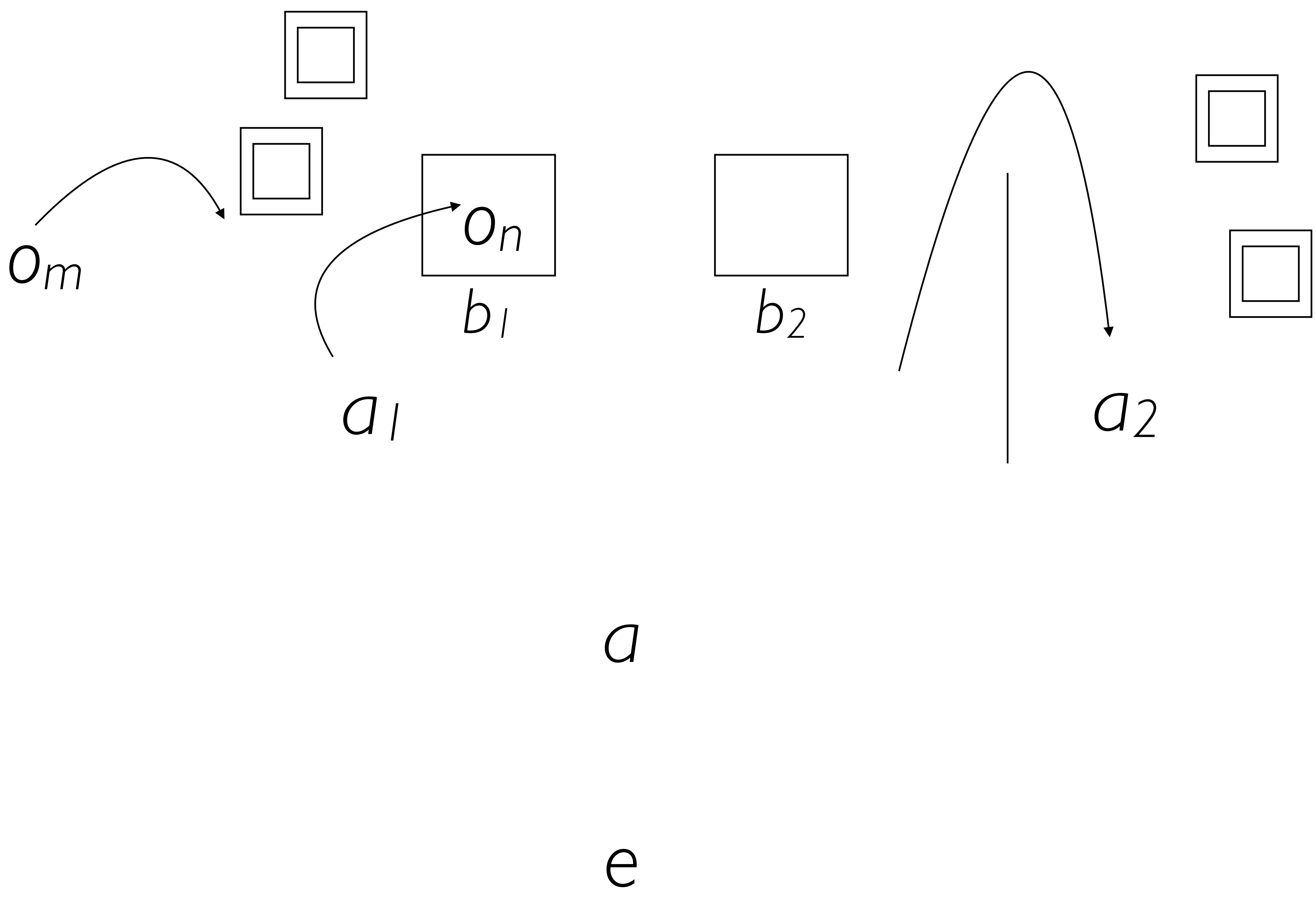
Framework for FBT^1_5
(ten timepoints)



Framework for FBT^1_5
(ten timepoints)

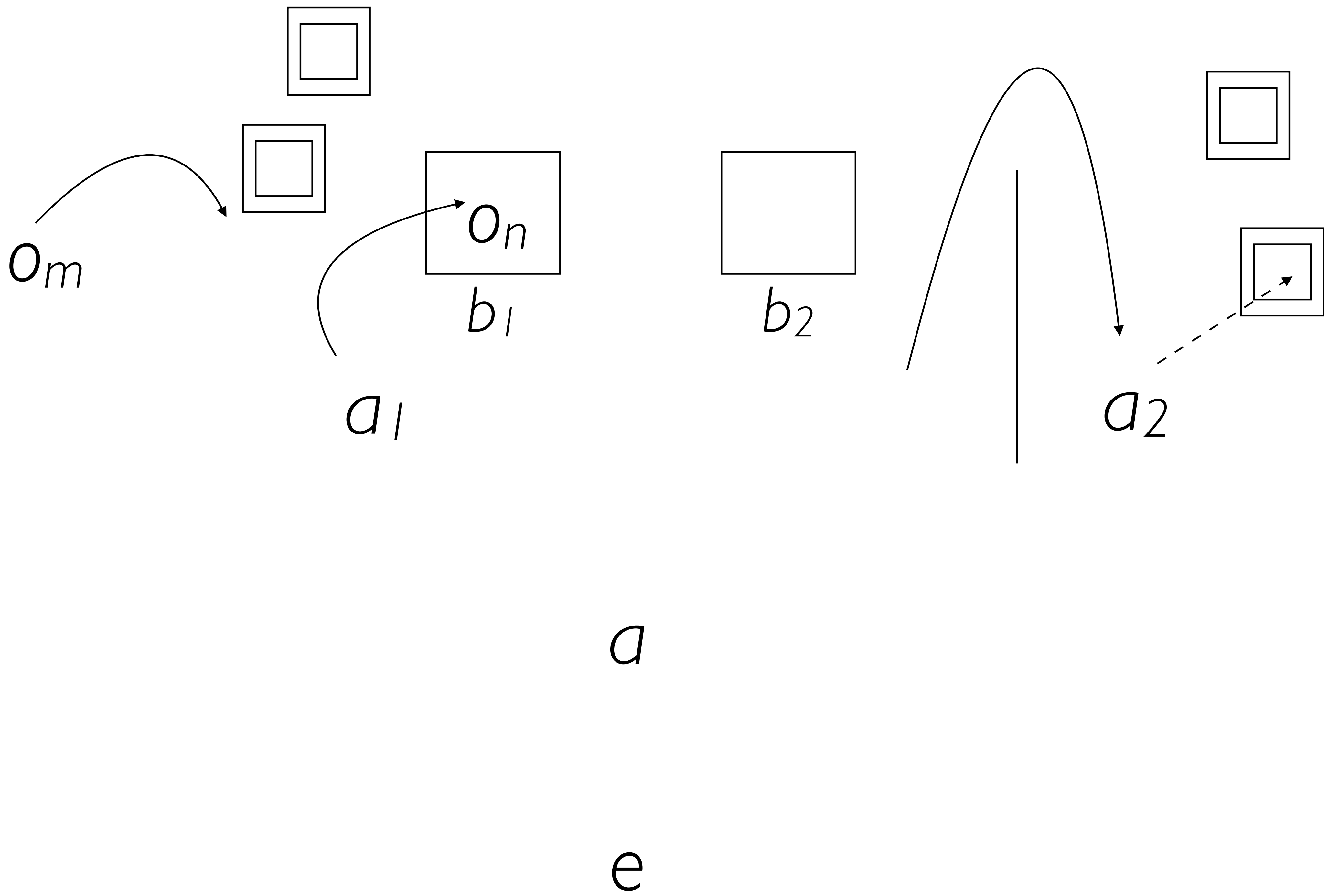


Framework for FBT^1_5
(ten timepoints)

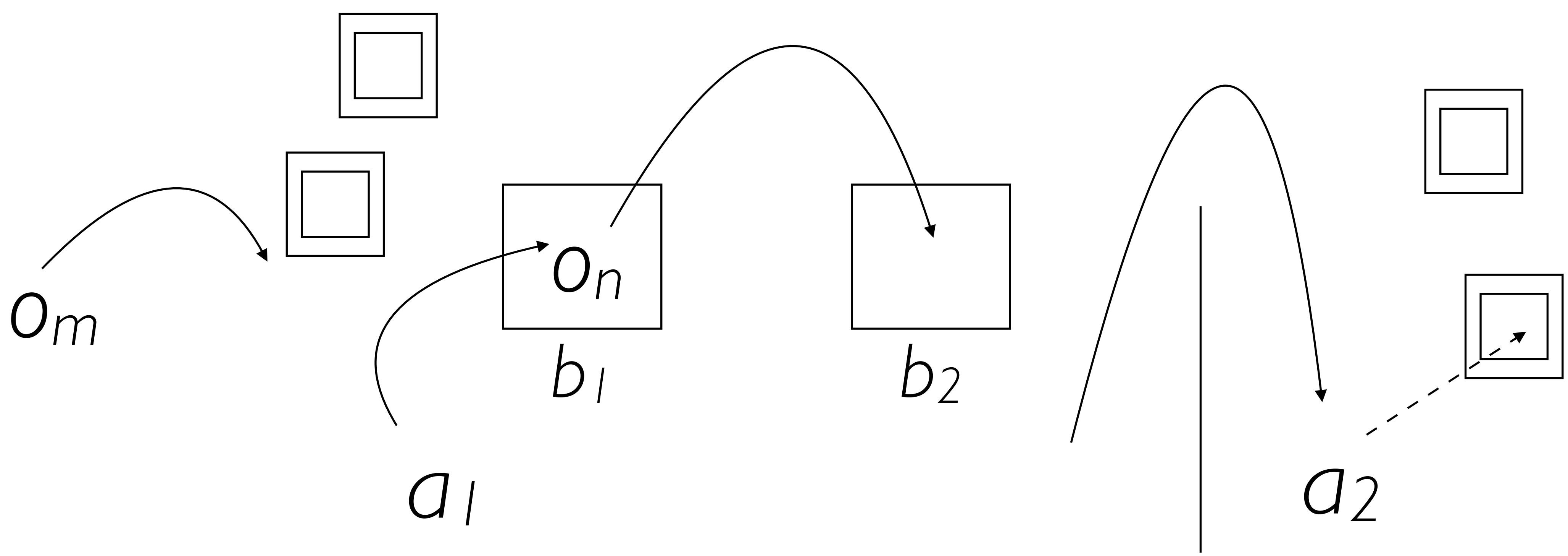


Framework for FBT^1_5

(ten timepoints)



Framework for FBT^1_5
(ten timepoints)



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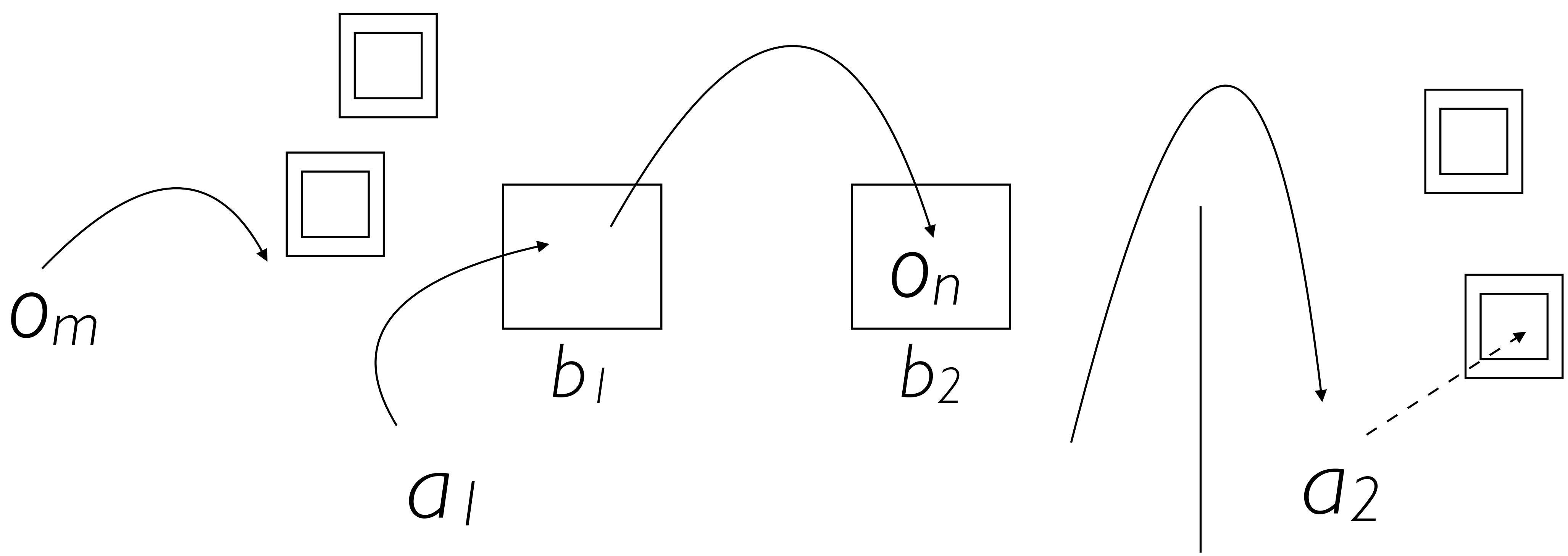
Framework for FBT$^1{}_5$
(ten timepoints)



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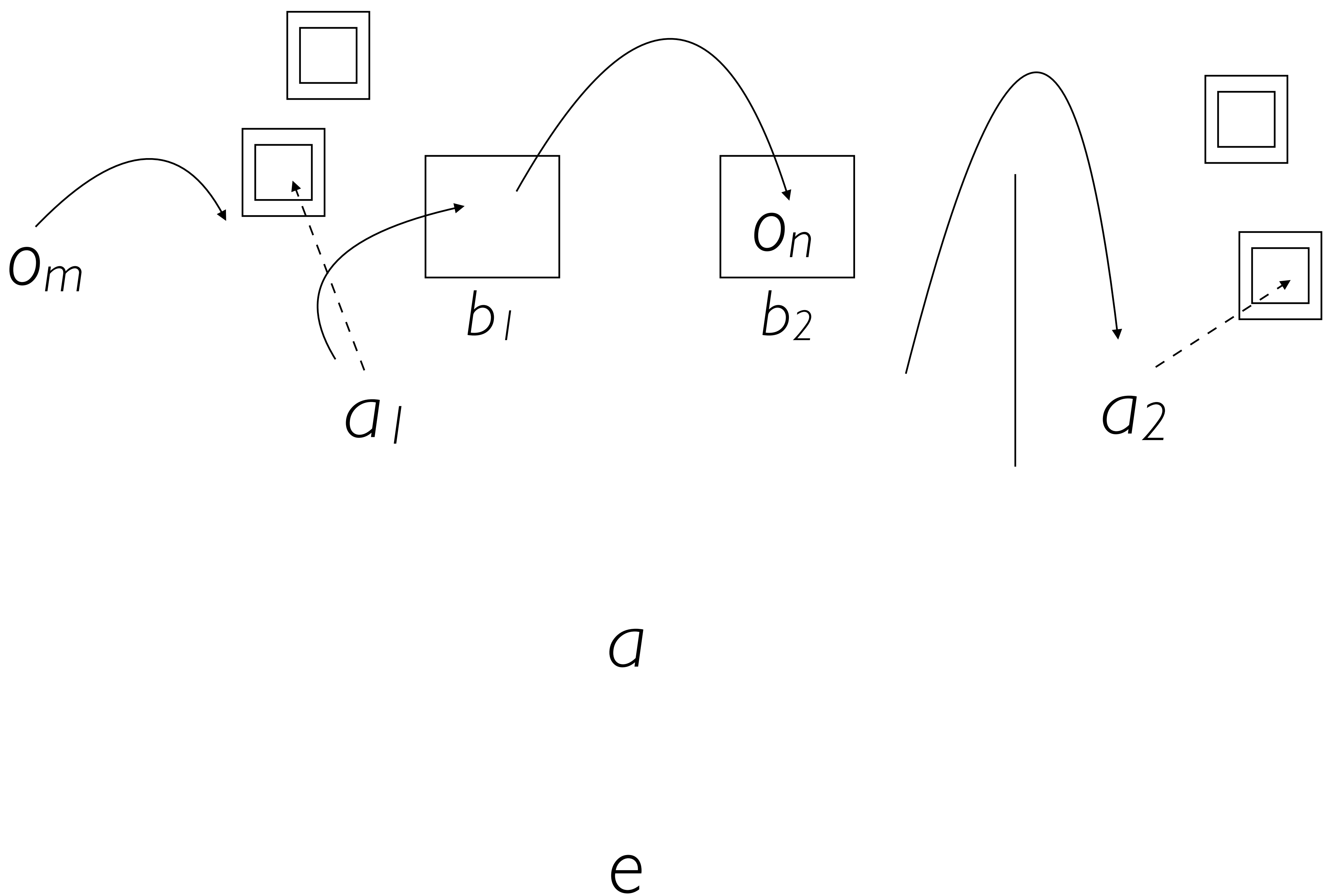
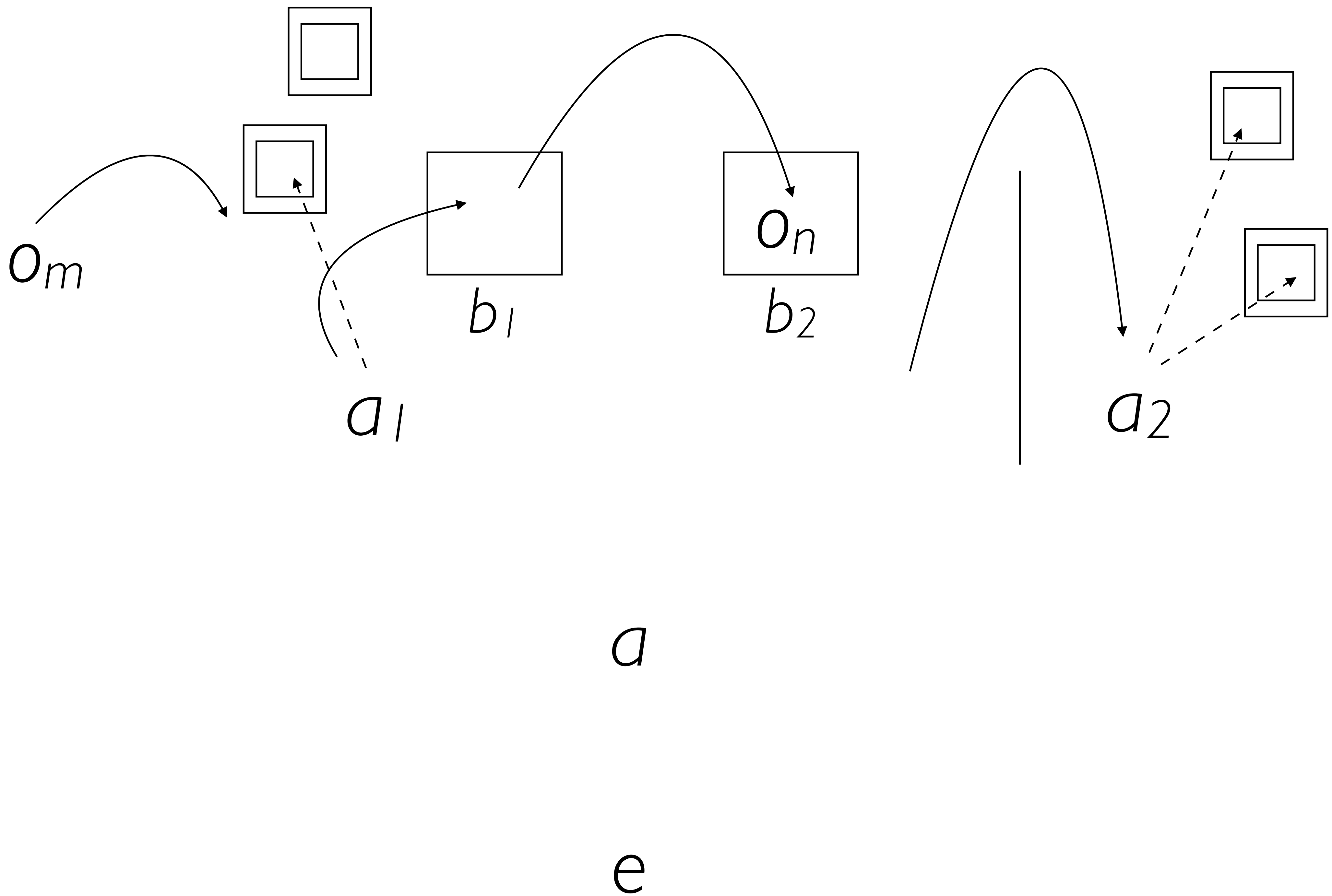
Framework for FBT^1_5
(ten timepoints)



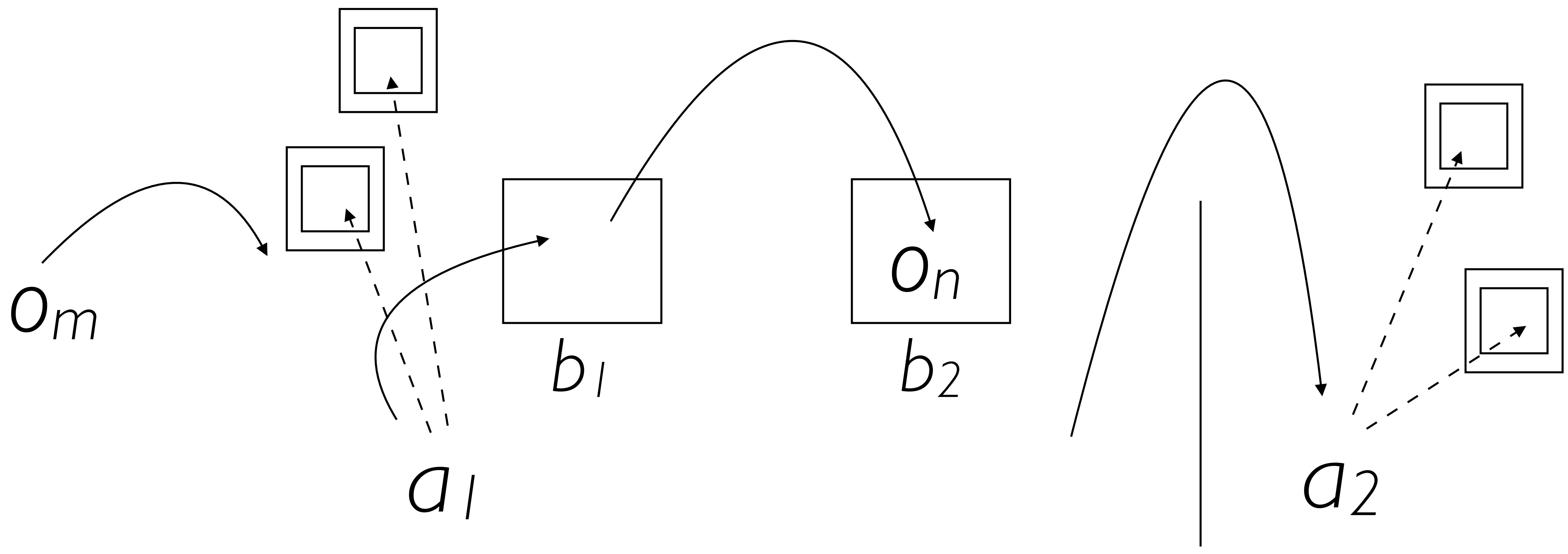
Framework for FBT^1_5

(ten timepoints)



Framework for FBT^1_5

(ten timepoints)



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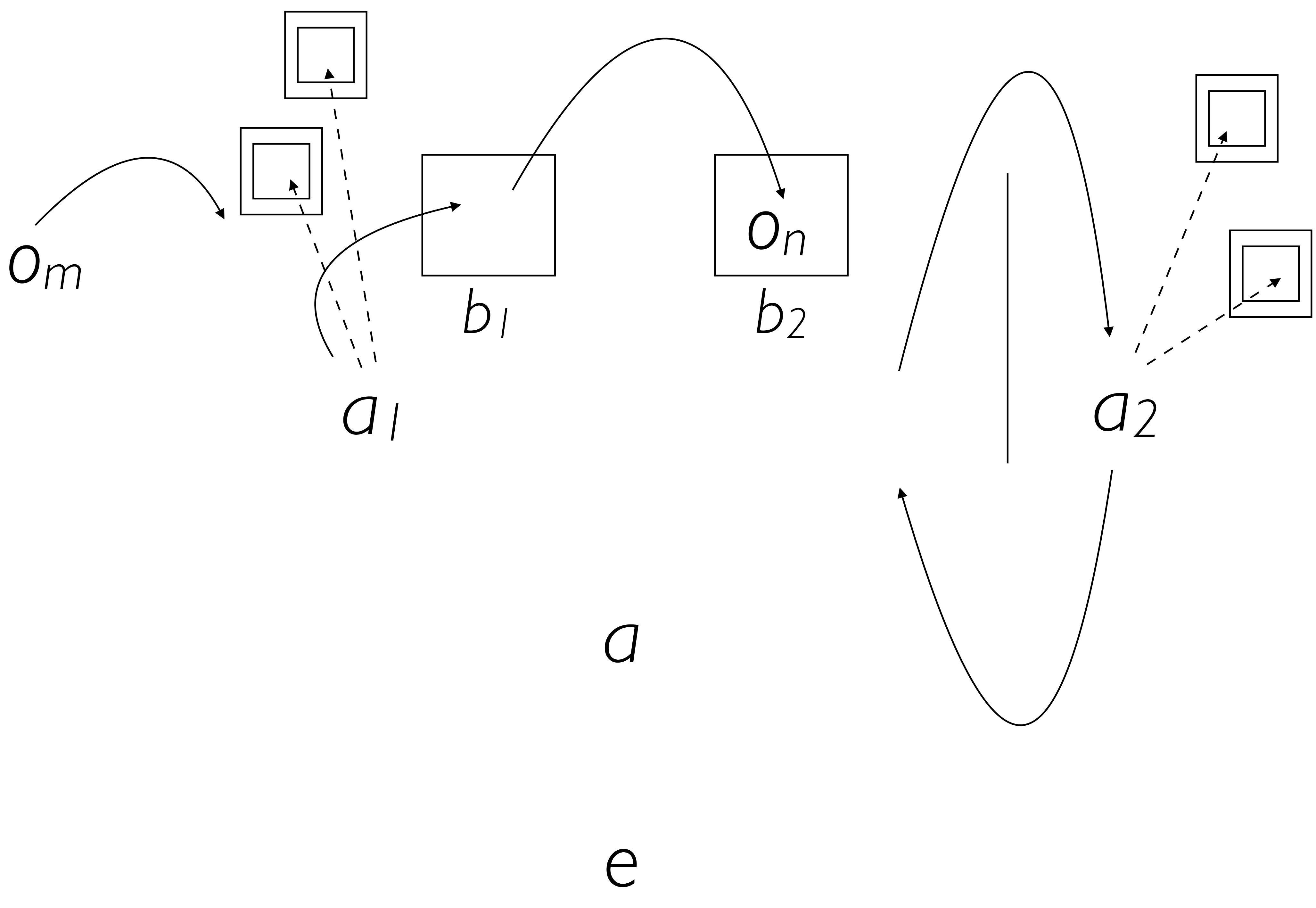
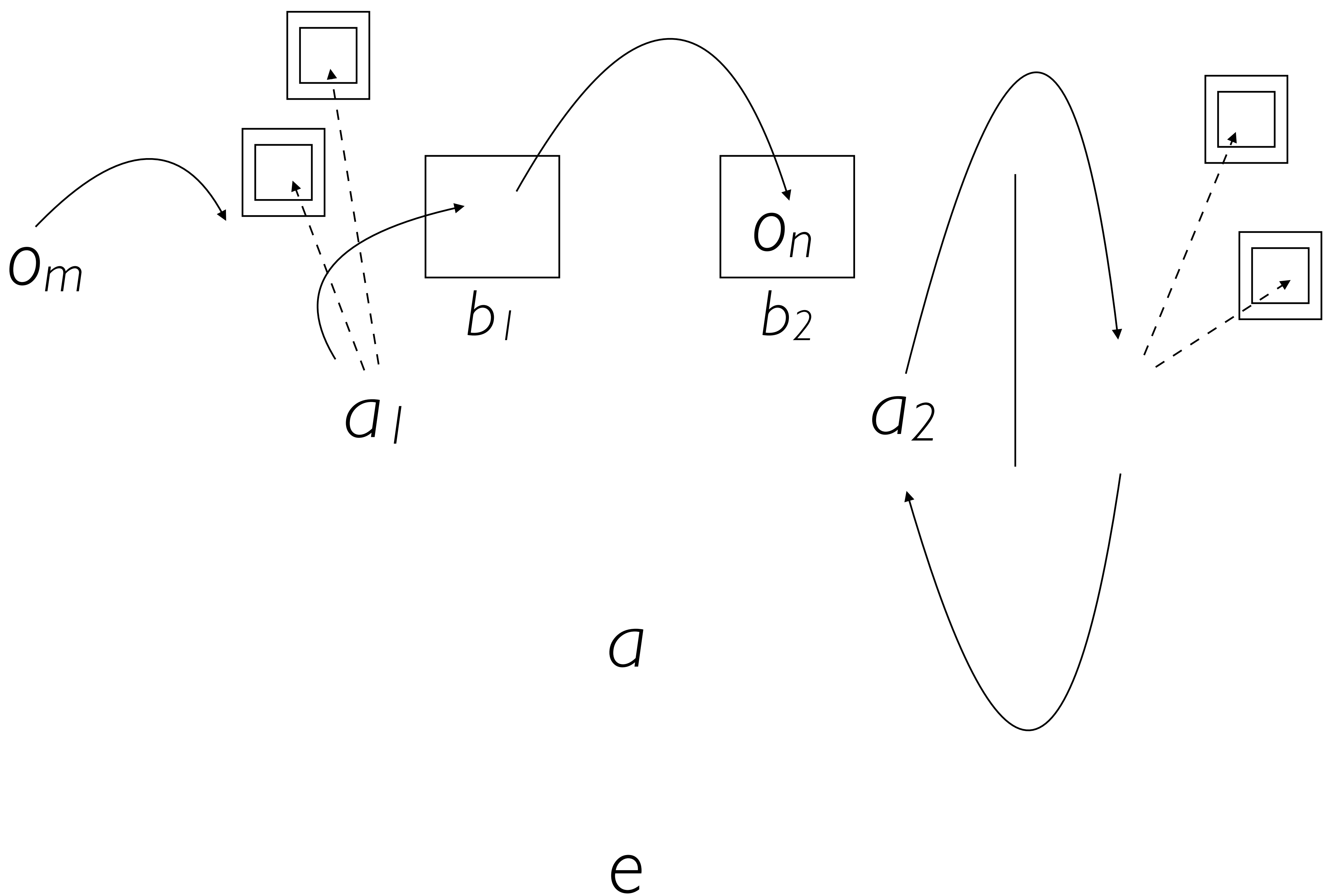
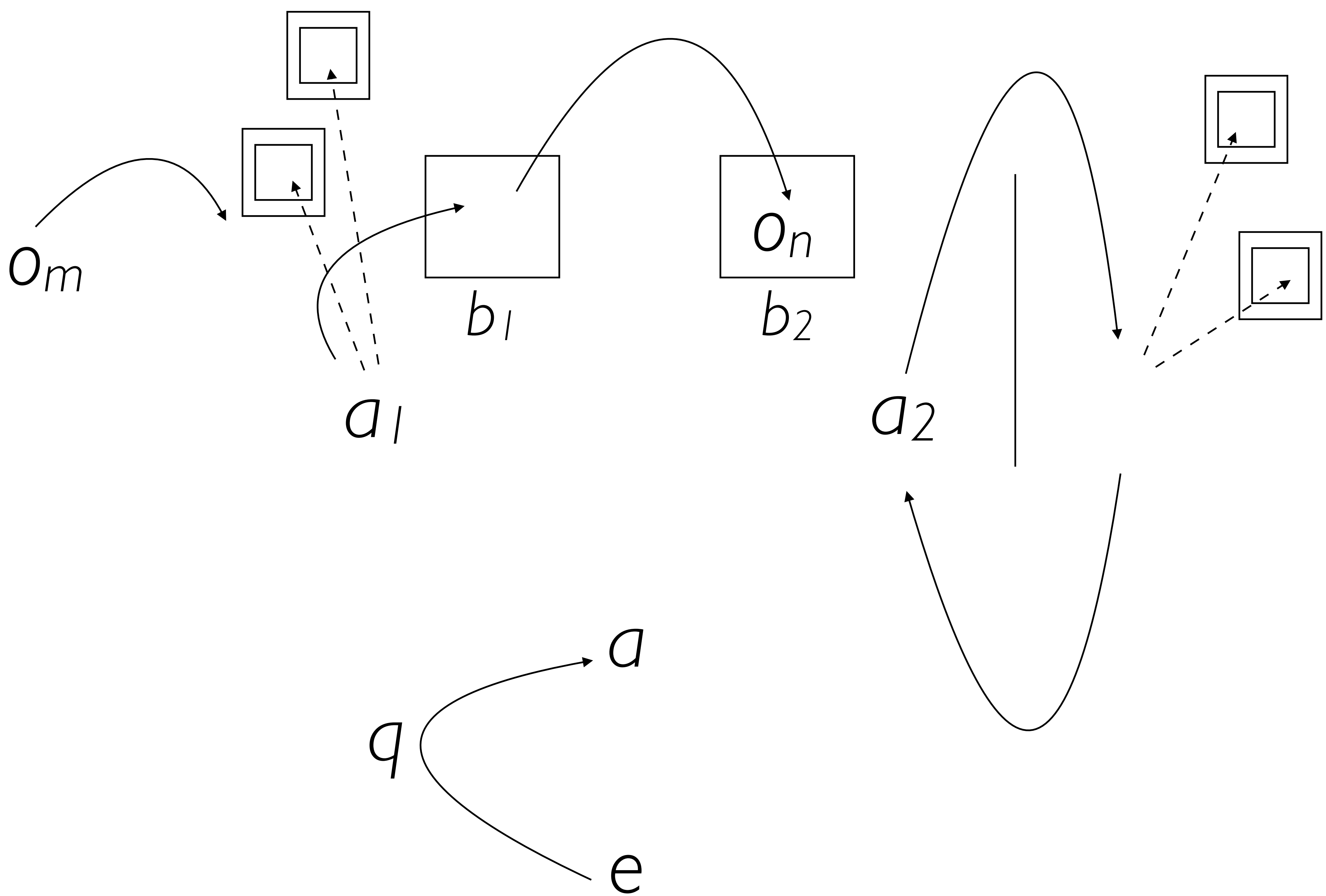
Framework for FBT$^1{}_5$
(ten timepoints)



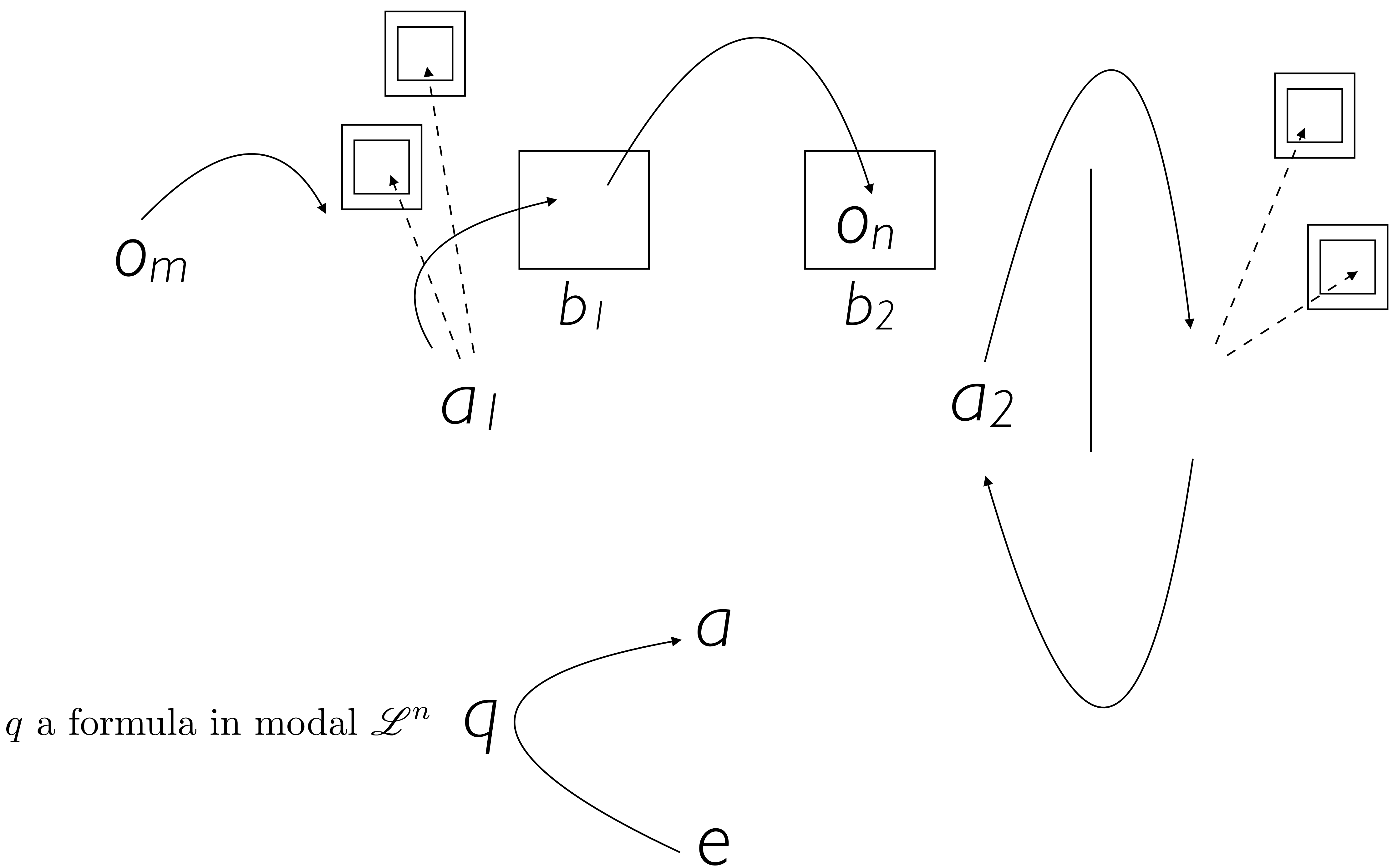
Framework for $FBT^1{}_5$
(ten timepoints)



Framework for $FBT^1{}_5$
(ten timepoints)



Framework for FBT$^1{}_5$
(ten timepoints)



Humans Can Succeed

Neurobiologically normal, nurtured, educated, and sufficiently motivated humans can correctly answer any relevant query *q* for the infinite progression, and prove that their answer is correct. For the obvious subclass of queries (the form of which appear in the box below), they can prove and exploit the following lemma.

Lemma: Suppose $\text{FBT}_k, k \in \mathbb{Z}^+$, holds; (i.e. that level k of FBT holds). Then, if k is even, $\mathbf{B}_2\mathbf{B}_1 \ldots \mathbf{B}_2\ \iota$, where there are $k+1$ iterated $\mathbf{B}_i$ operators; otherwise $\mathbf{B}_1\mathbf{B}_2 \ldots \mathbf{B}_1\mathbf{B}_2\ \iota$, where there again there are $k+1$ iterated $\mathbf{B}_i$ operators.

Passing to Probing Mastery of the Specific Subclass

Experimenter to *a*: "At level *k*, from which box will a_2 attempt to retrieve the objects o_n? Prove it!"

Theoretical Machine Success on Infinite FBT!

Theorem: $\forall q \in \mathscr{CC}, \mathfrak{M}$ can correctly answer and justify q. I.e., $\mathfrak{M}$ can pass FBT_ω.

Ok, so this logic machine exists in the *mathematical* universe; but does there exist an *implemented* machine with this power?

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Simulation Courtesy of …

ShadowProver!

Level 1

```
:name        "Level 1: False Belief Task "

:description "Agent a1 puts an object o into b1 in plain view of a2.
              Agent a2 then leaves, and in the absence of a2, a1 moves o
              from b1 into b2 ; this movement isn’t perceived by a2 . Agent
              a2 now returns, and a is asked by the experimenter e: “If a2
              desires to retrieve o, which box will a2 look in?” If younger
              than four or five, a will reply “In b ” (which of course fails 2
              the task); after this age subjects respond with the correct “In b1.”

              Level1 Belief: a1 believes a2 believes o is in b1.
              "

:date        "Monday July 22, 2019"

:assumptions {
               :P1 (Perceives! a1 t1 (Perceives! a2 t1 (holds (In o b1) t1)))

               :P2 (Believes! a1 t2 (Believes! a2 t2 (not (exists [?e] (terminates ?e (In o b1))))))

               :P3 (holds (In o b1) t1)

               :C1 (Common! t0 (forall [?f ?t2 ?t2]
                                  (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
                                    (holds ?f ?t2))))

               :C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3)))
               }

:goal        (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3)))}
```

Level 2

```
{:name        "Level 2: False Belief Task "

 :description "Agent a1 puts an object o into b1 in plain view of a2.
              Agent a2 then leaves, and in the absence of a2, a1 moves o
              from b1 into b2 ; this movement isn’t perceived by a2 . Agent
              a2 now returns, and a is asked by the experimenter e: “If a2
              desires to retrieve o, which box will a2 look in?” If younger
              than four or five, a will reply “In b ” (which of course fails 2
              the task); after this age subjects respond with the correct “In b1.”

              Level2 Belief: a2 believes a1 believes a2 believes o is in b1.
              "

 :date        "Monday July 22, 2019"

 :assumptions {

               :P1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (holds (In o b1) t1))))

               :P2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (not (exists [?e] (terminates ?e (In o b1)))))))

               :P3 (holds (In o b1) t1)


               :C1 (Common! t0
                          (forall [?f ?t2 ?t2]
                                  (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
                                    (holds ?f ?t2))))

              :C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3)))}

 :goal        (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3))))}
```

Level 3

```
{:name        "Level 3: False Belief Task "

 :description "Agent a1 puts an object o into b1 in plain view of a2.
               Agent a2 then leaves, and in the absence of a2, a1 moves o
               from b1 into b2 ; this movement isn’t perceived by a2 . Agent
               a2 now returns, and a is asked by the experimenter e: “If a2
               desires to retrieve o, which box will a2 look in?” If younger
               than four or five, a will reply “In b ” (which of course fails 2
               the task); after this age subjects respond with the correct “In b1.”

               Level3 Belief: a2 believes a1 believes a2 believes o is in b1.
               "

 :date        "Monday July 22, 2019"

 :assumptions {

                :P1 (Perceives! a1 t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (holds (In o b1) t1)))))
                :P2 (Believes! a1 t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (not (exists [?e] (terminates ?e (In o b1))))))))

                :P3 (holds (In o b1) t1)


                :C1 (Common! t0
                             (forall [?f ?t2 ?t2]
                                     (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
                                       (holds ?f ?t2))))

                :C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3)))}

 :goal        (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3)))))}
```

Level 4

```
{:name        "Level 4: False Belief Task "

 :description "Agent a1 puts an object o into b1 in plain view of a2.
               Agent a2 then leaves, and in the absence of a2, a1 moves o
               from b1 into b2 ; this movement isn't perceived by a2 . Agent
               a2 now returns, and a is asked by the experimenter e: "If a2
               desires to retrieve o, which box will a2 look in?" If younger
               than four or five, a will reply "In b " (which of course fails 2
               the task); after this age subjects respond with the correct "In b1."

               Level4 Belief: a2 believes a1 believes a2 believes a1 believes a2 believes o is in b1.
               "

 :date        "Monday July 22, 2019"

 :assumptions {

               :P1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (holds (In o b1) t1))))))
               :P2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (not (exists [?e] (terminates ?e (In o b1)))))))))

               :P3 (holds (In o b1) t1)


               :C1 (Common! t0
                            (forall [?f ?t2 ?t2]
                                    (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
                                      (holds ?f ?t2))))

               :C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3)))}

 :goal        (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3))))))}
```

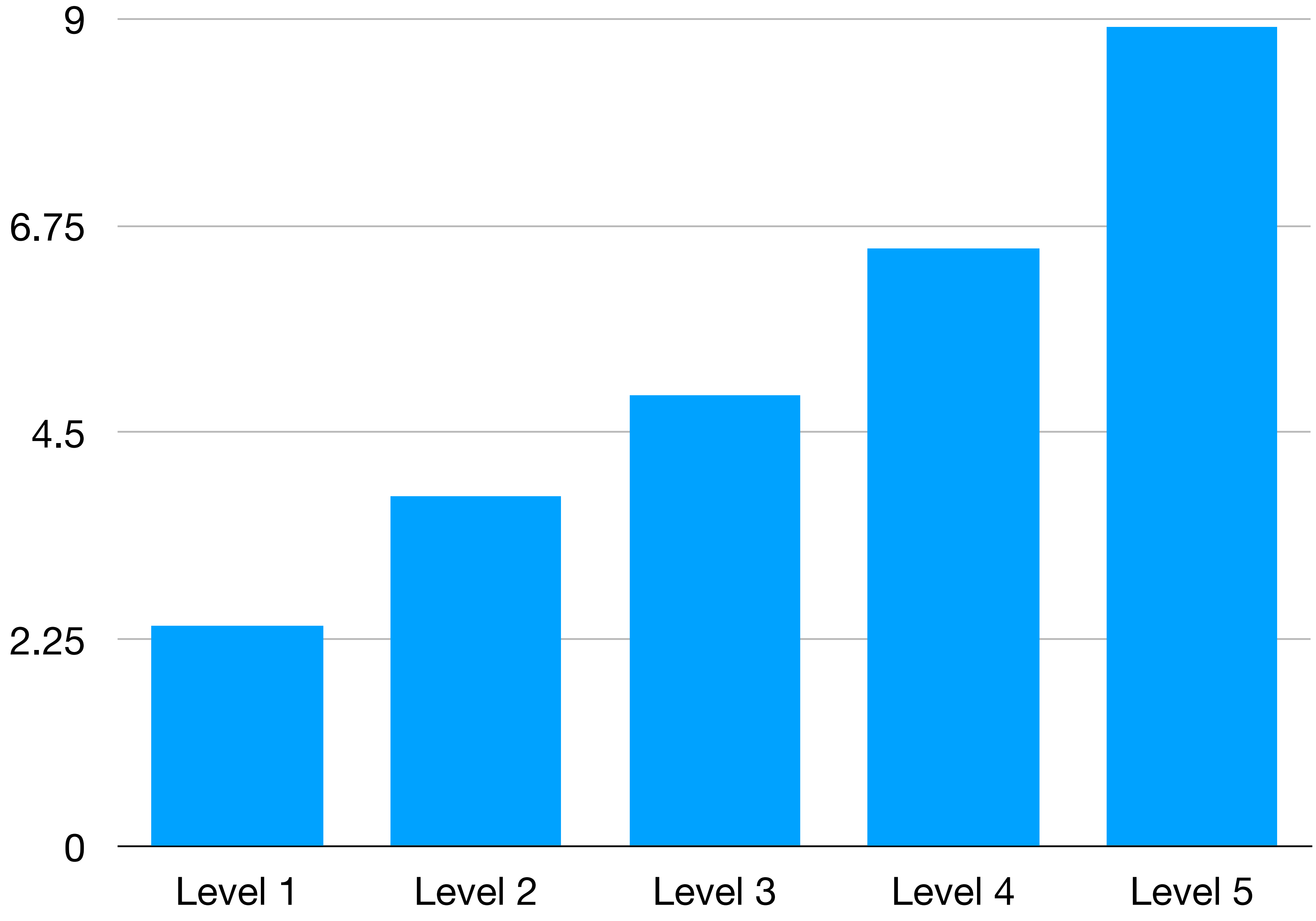
Level 5

```
{:name        "Level 5: False Belief Task "

 :description "Agent a1 puts an object o into b1 in plain view of a2.
               Agent a2 then leaves, and in the absence of a2, a1 moves o
               from b1 into b2 ; this movement isn't perceived by a2 . Agent
               a2 now returns, and a is asked by the experimenter e: "If a2
               desires to retrieve o, which box will a2 look in?" If younger
               than four or five, a will reply "In b " (which of course fails 2
               the task); after this age subjects respond with the correct "In b1."

               Level5 Belief: a1 believes a2 believes a1 believes a2 believes a1 believes a2 believes o is in b1.
               "

 :date        "Monday July 22, 2019"

 :assumptions {

                :P1 (Perceives! a1 t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (holds (In o b1) t1)))))))
                :P2 (Believes! a1 t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (not (exists [?e] (terminates ?e (In o b1)))))))))

                :P3 (holds (In o b1) t1)


                :C1 (Common! t0
                             (forall [?f ?t2 ?t2]
                                     (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
                                       (holds ?f ?t2))))

                :C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3)))}

 :goal        (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3)))))))}
```

```
                                        ...gent
                                        ...f If a2
                                        ...f younger
                                        ... course fails 2
                                        ...n the correct “In b1.”
                                        ...l believes a2 believes a1 believes a2 believes o is in b1.
...t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 ...
...l t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (not ...
(In o b1) t1)
:goal
...l (Common! t0
    (forall [?f ?t2 ?t2]
      (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
        (holds ?f ?t2))))
:C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3)))}
(Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3))))))))}
```

```
                                                    ...gent
                                                    ... If a2
                                                    ... younger
                                                    ... course fails 2
                                                    ...n the correct "In b1."
                                                    ...l believes a2 believes a1 believes a2 believes o is in b1.
...t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2 t1 (Perceives! a1 t1 (Perceives! a2
...t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (Believes! a1 t2 (Believes! a2 t2 (not
      (In o b1) t1)
...l (Common! t0
       (forall [?f ?t2 ?t2]
         (if (and (not (exists [?e] (terminates ?e ?f))) (holds ?f ?t1) (< ?t1 ?t2))
          (holds ?f ?t2))))
     :C2 (Common! t0 (and (< t1 t2) (< t2 t3) (< t1 t3))))}
:goal
   (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (Believes! a1 t3 (Believes! a2 t3 (holds (In o b1) t3))))))))}
```

Time (in seconds) to Prove



Simulation of Level 5 in Real Time

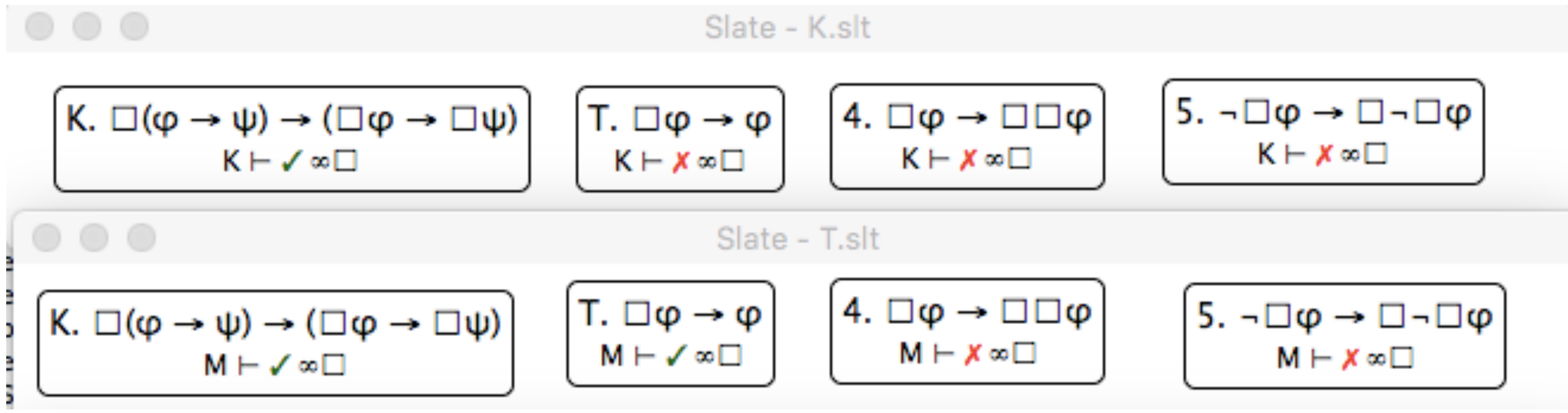
```
/Library/Java/JavaVirtualMachines/jdk1.8.0_131.jdk/Contents/Home/bin/java ...
objc[16653]: Class JavaLaunchHelper is implemented in both /Library/Java/JavaVirtualMachines/jdk1.8.0_131.jdk/Contents/Home/bin/java (0x102a2d4c0) and /Library/Java/JavaVirtualMachines/jdk1.8.0_131.jdk/Contents/Home/jre/lib/libinstrument.dylib (0x102ab94e0)
------------- Level 5 ----------------------
```

Simulation of Level 5 in Real Time

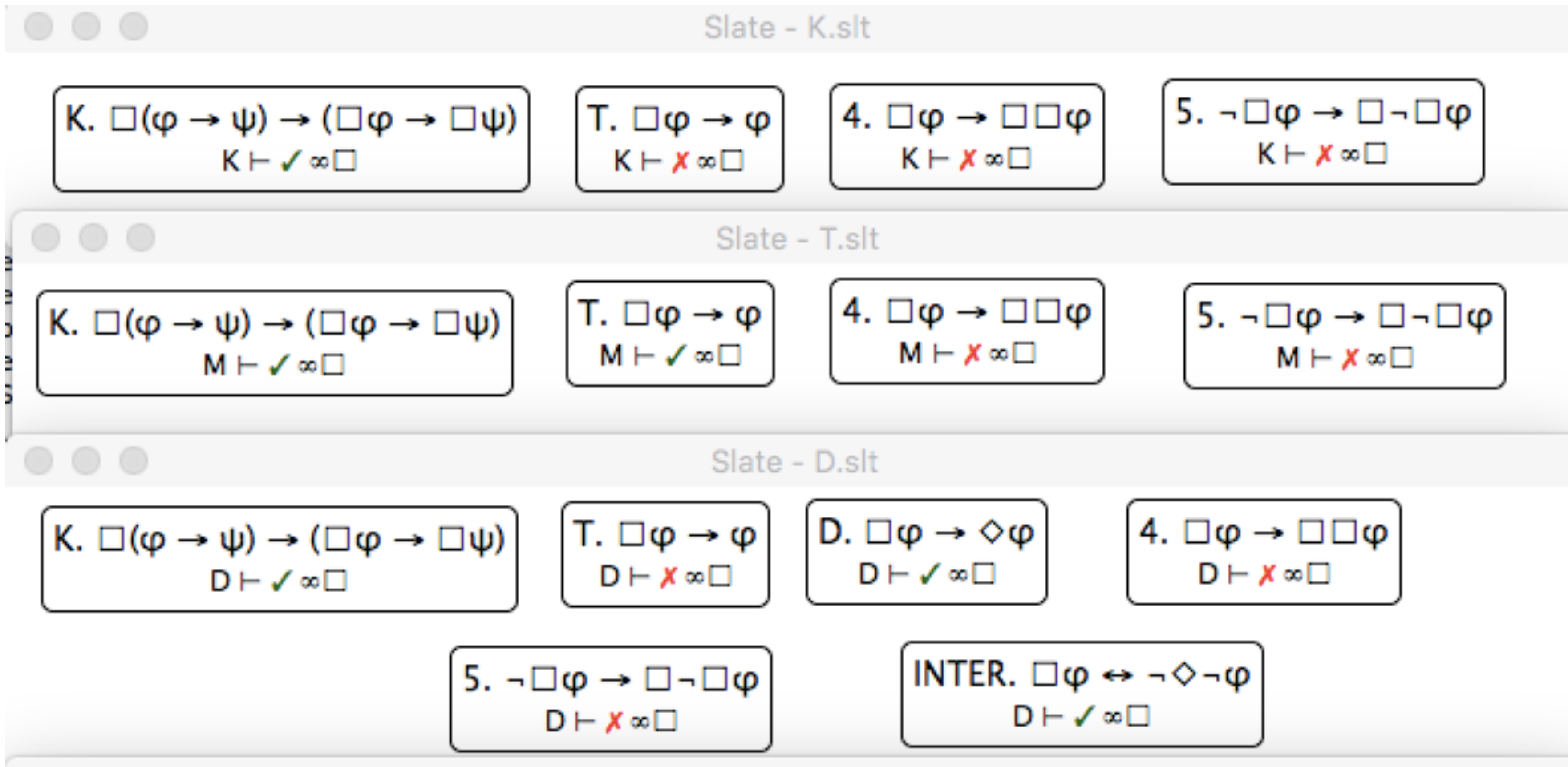
```
/Library/Java/JavaVirtualMachines/jdk1.8.0_131.jdk/Contents/Home/bin/java ...
objc[16653]: Class JavaLaunchHelper is implemented in both /Library/Java/JavaVirtualMachines/jdk1.8.0_131.jdk/Contents/Home/bin/java (0x102a2d4c0) and /Library/Java/JavaVirtualMachines/jdk1.8.0_131.jdk/Contents/Home/jre/lib/libinstrument.dylib (0x102ab94e0)
------------- Level 5 -----------------------
```

Encapsulation

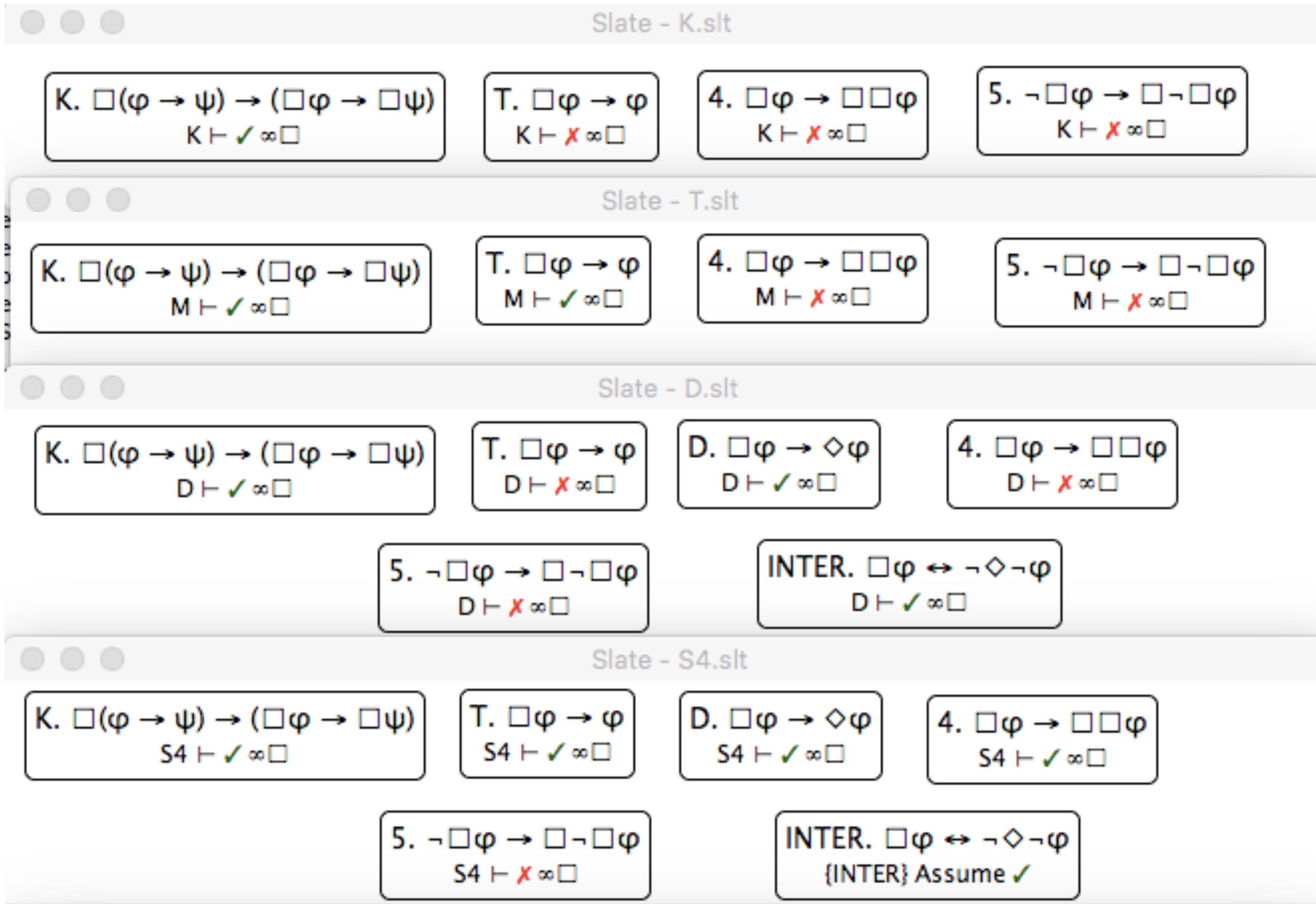
Encapsulation



Encapsulation



Encapsulation



Encapsulation

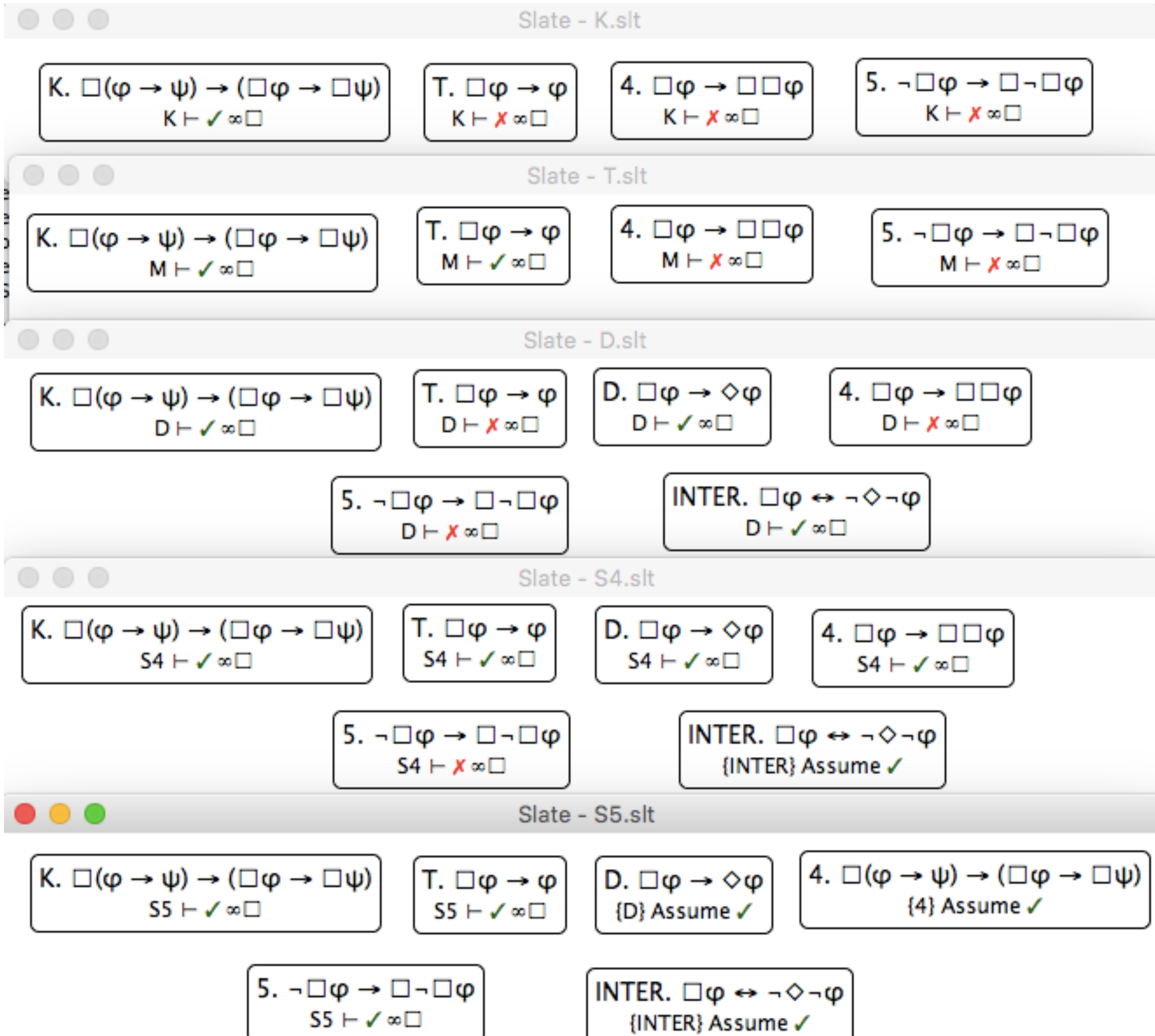
K

T

D

4 = S4

5 = S5



Encapsulation

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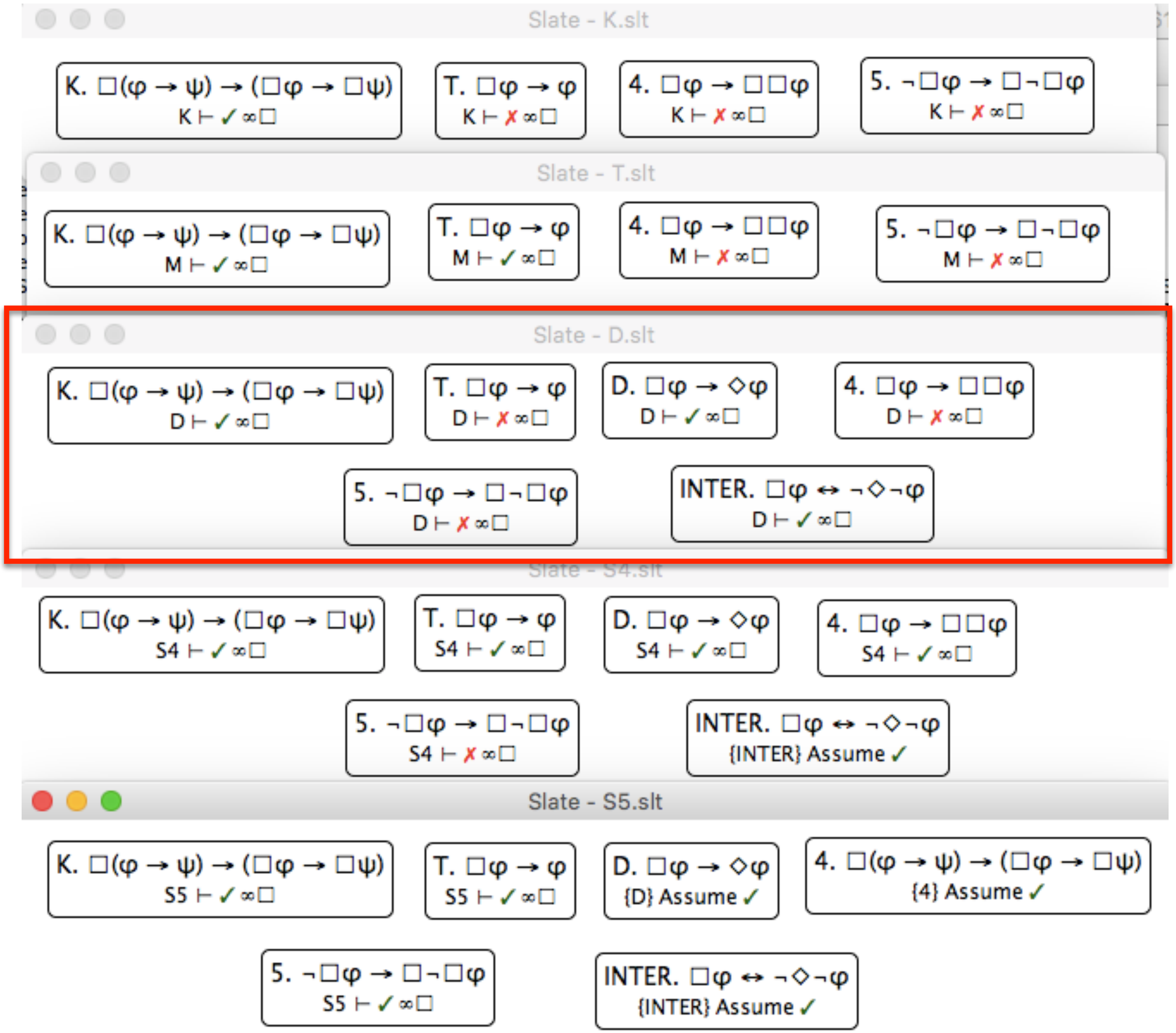
D

4 = S4

5 = S5



HyperSlate®

Encapsulation

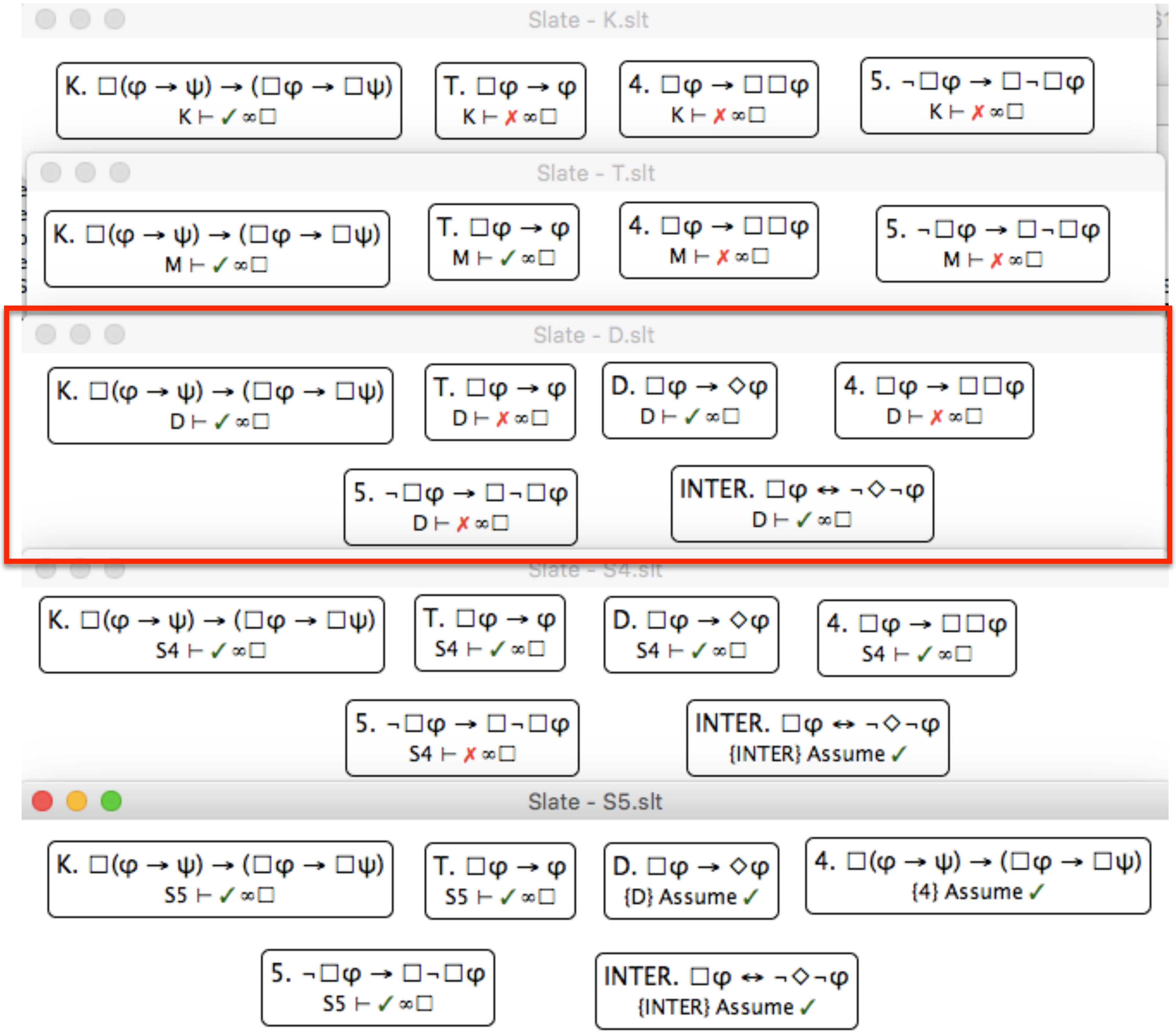
K

T

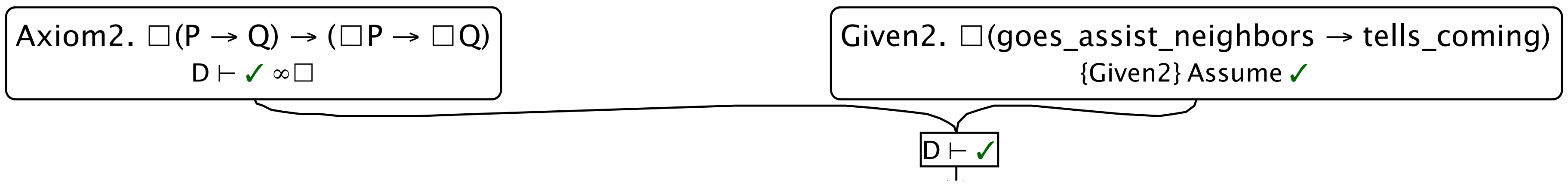
D

4 = S4

5 = S5



Chisholm's Paradox

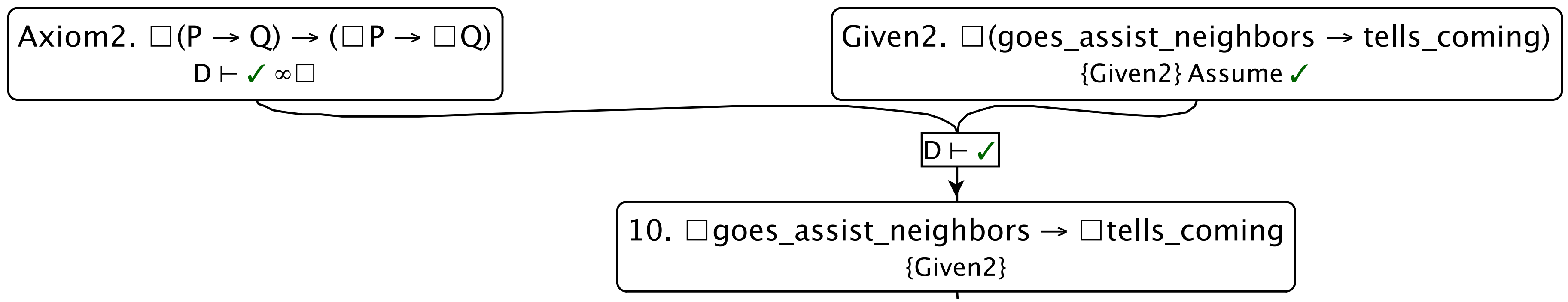


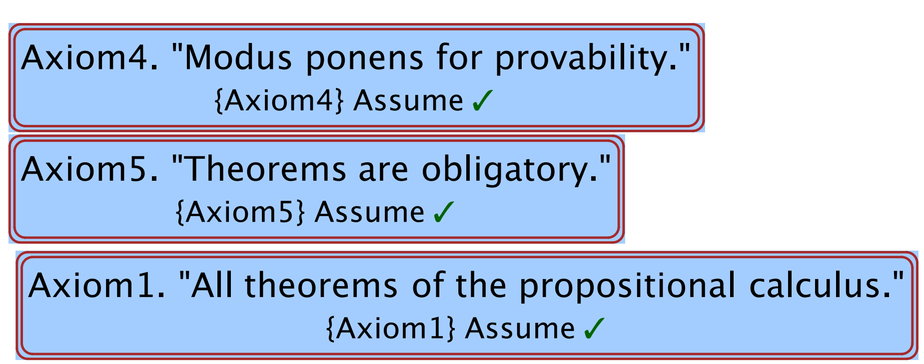
Axiom4. "Modus ponens for provability."
{Axiom4} Assume ✓

Axiom5. "Theorems are obligatory."
{Axiom5} Assume ✓

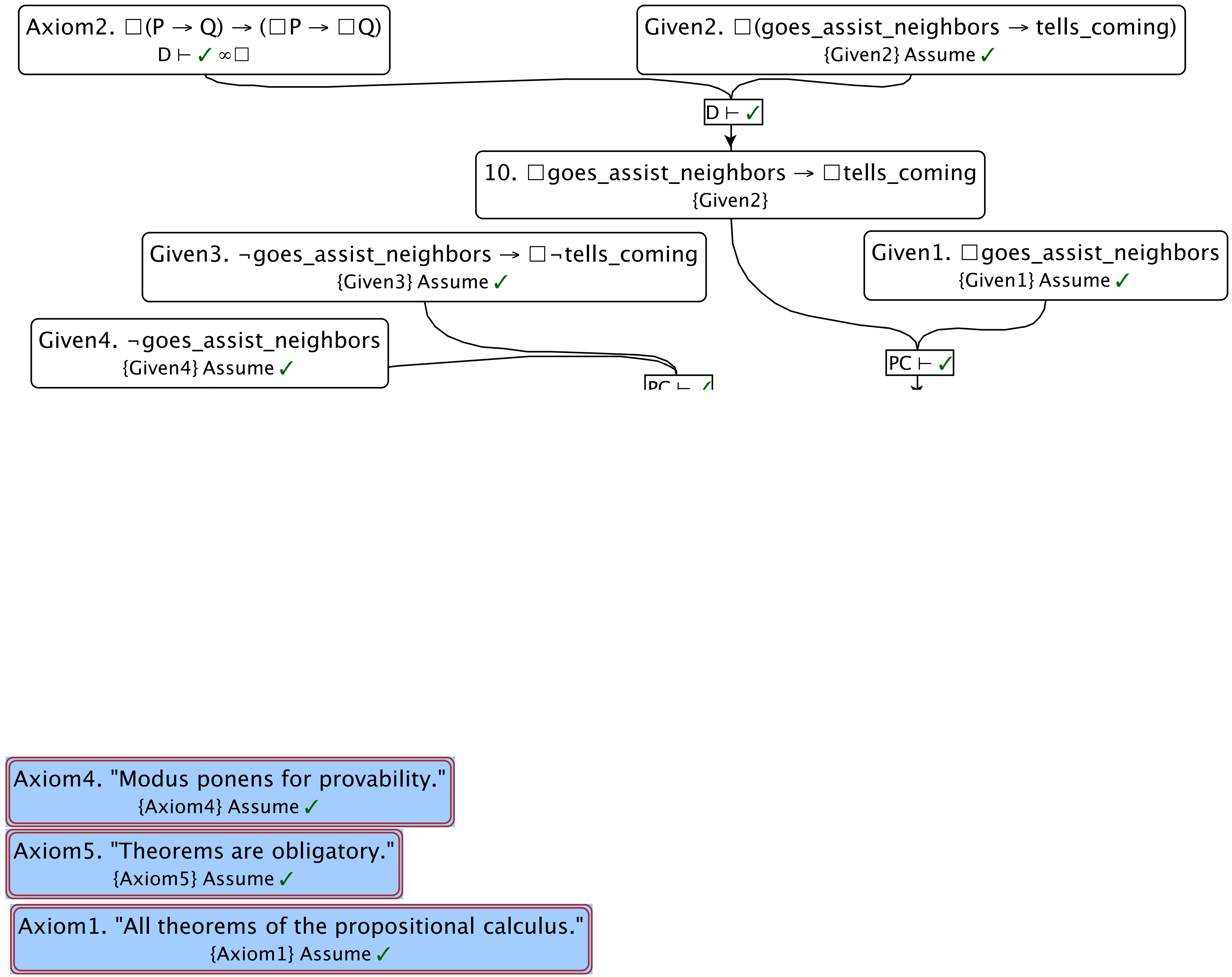
Axiom1. "All theorems of the propositional calculus."
{Axiom1} Assume ✓

Chisholm's Paradox

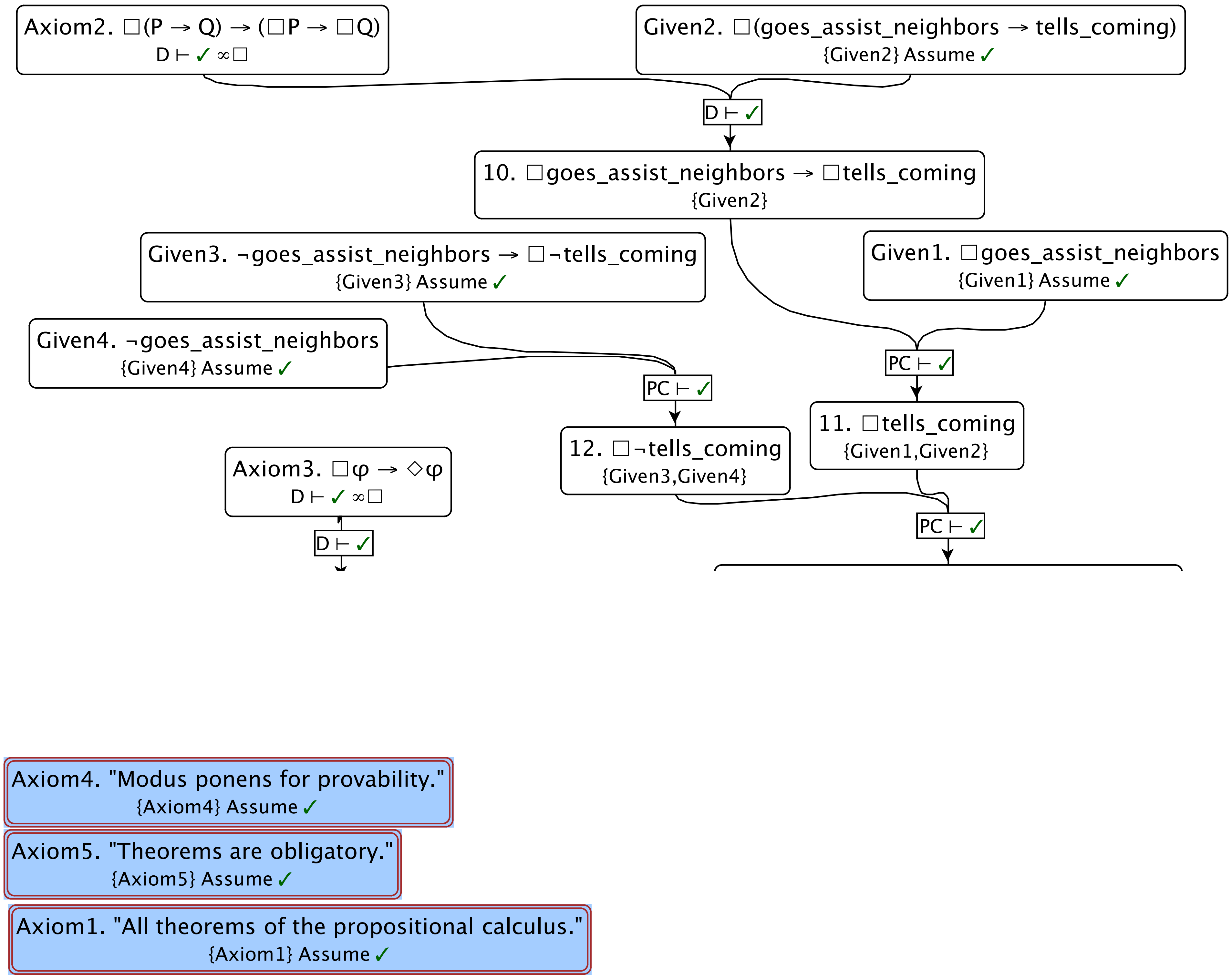




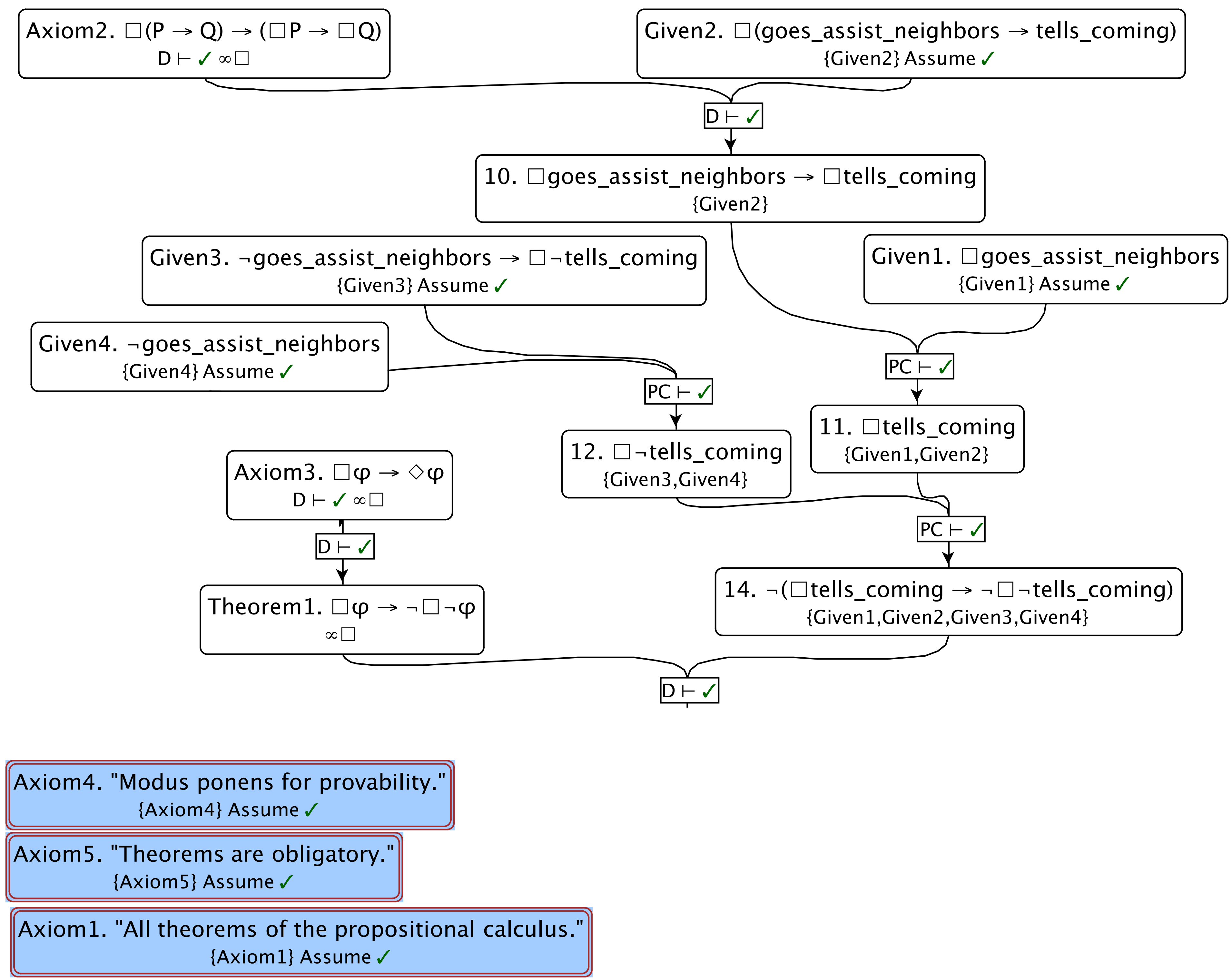
Chisholm's Paradox



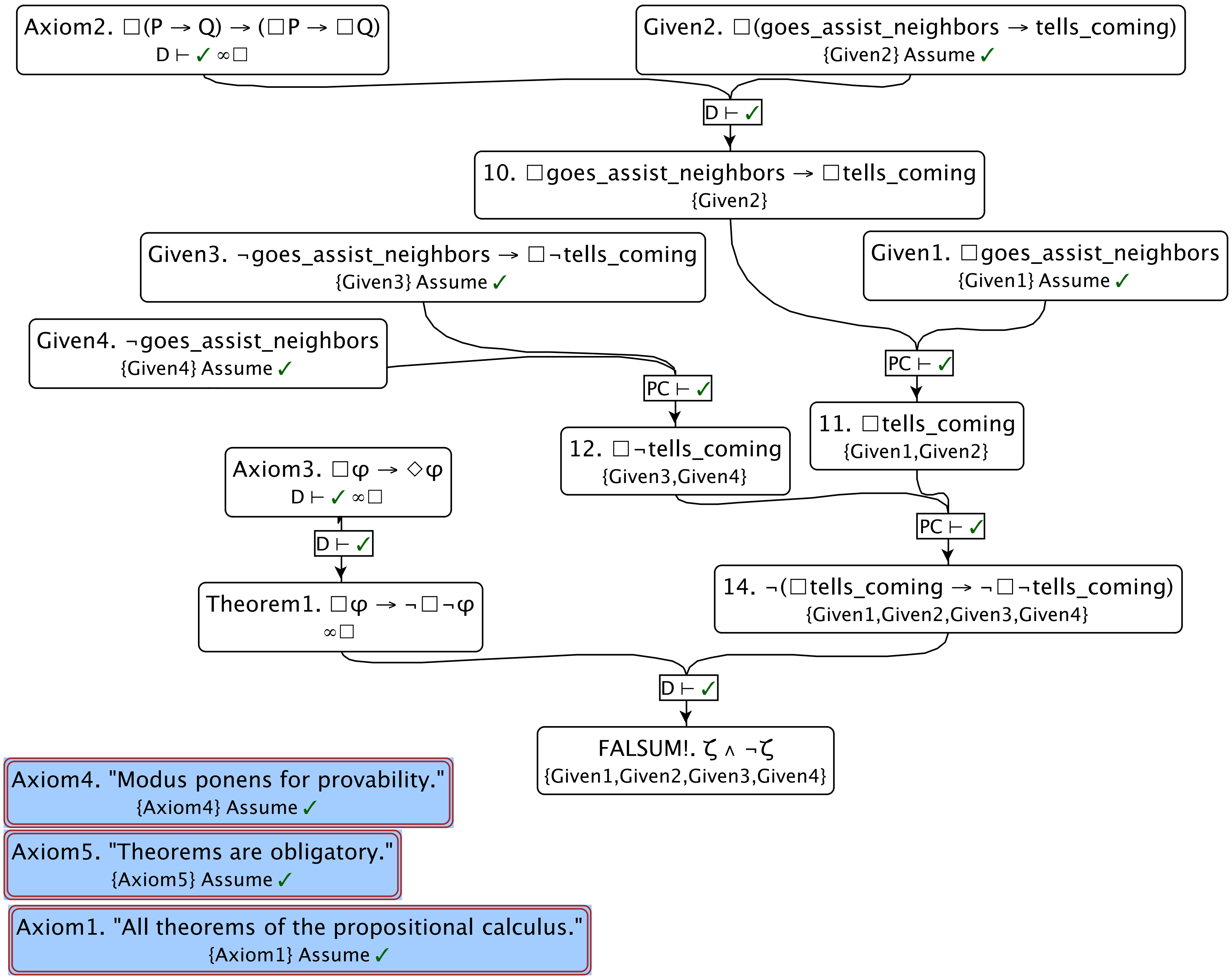
Chisholm's Paradox



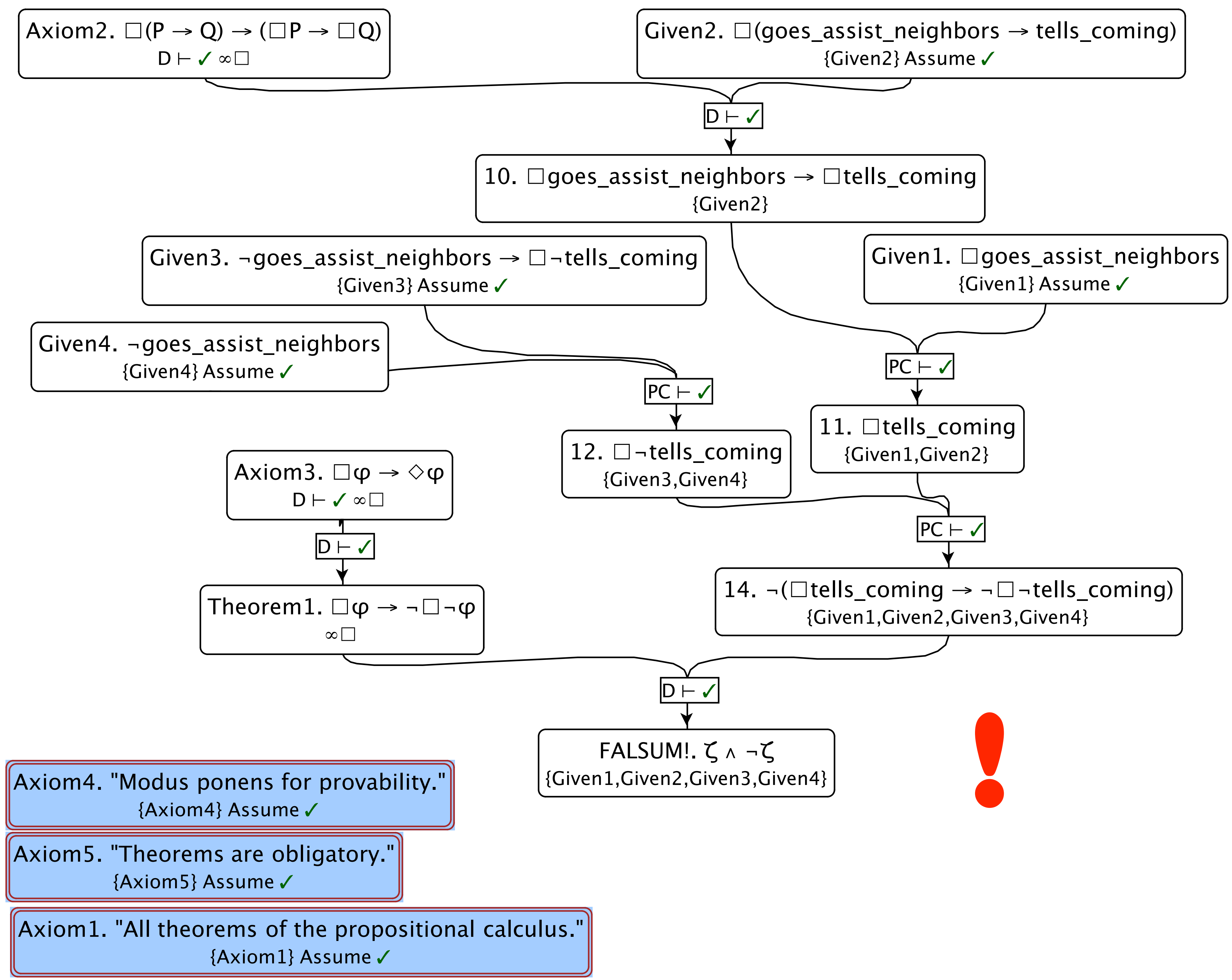
Chisholm's Paradox



Chisholm's Paradox



Chisholm's Paradox



HyperSlate®

Review: Encapsulation

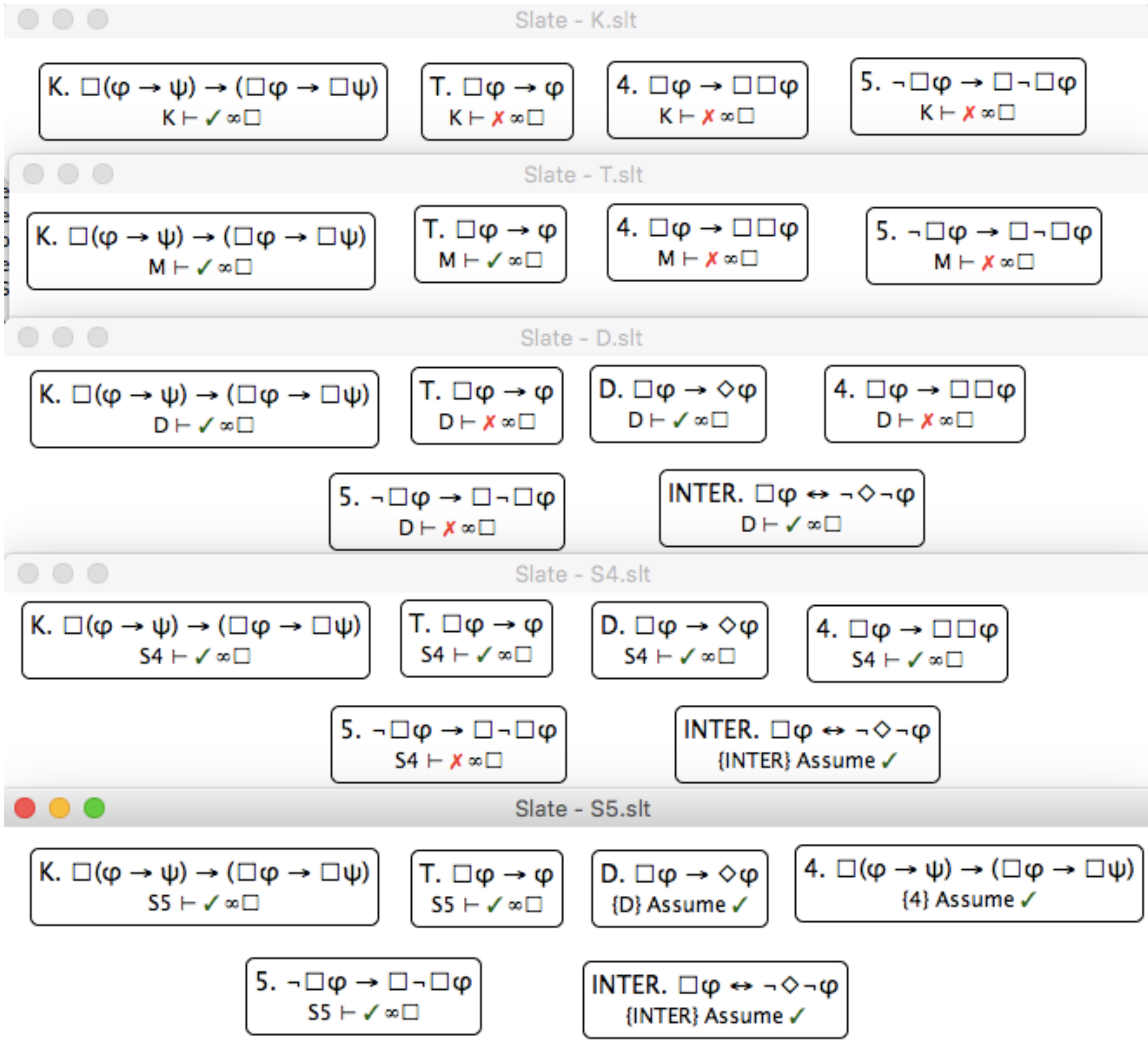
K

T

D

4 = S4

5 = S5



HyperSlate®

Review: Encapsulation

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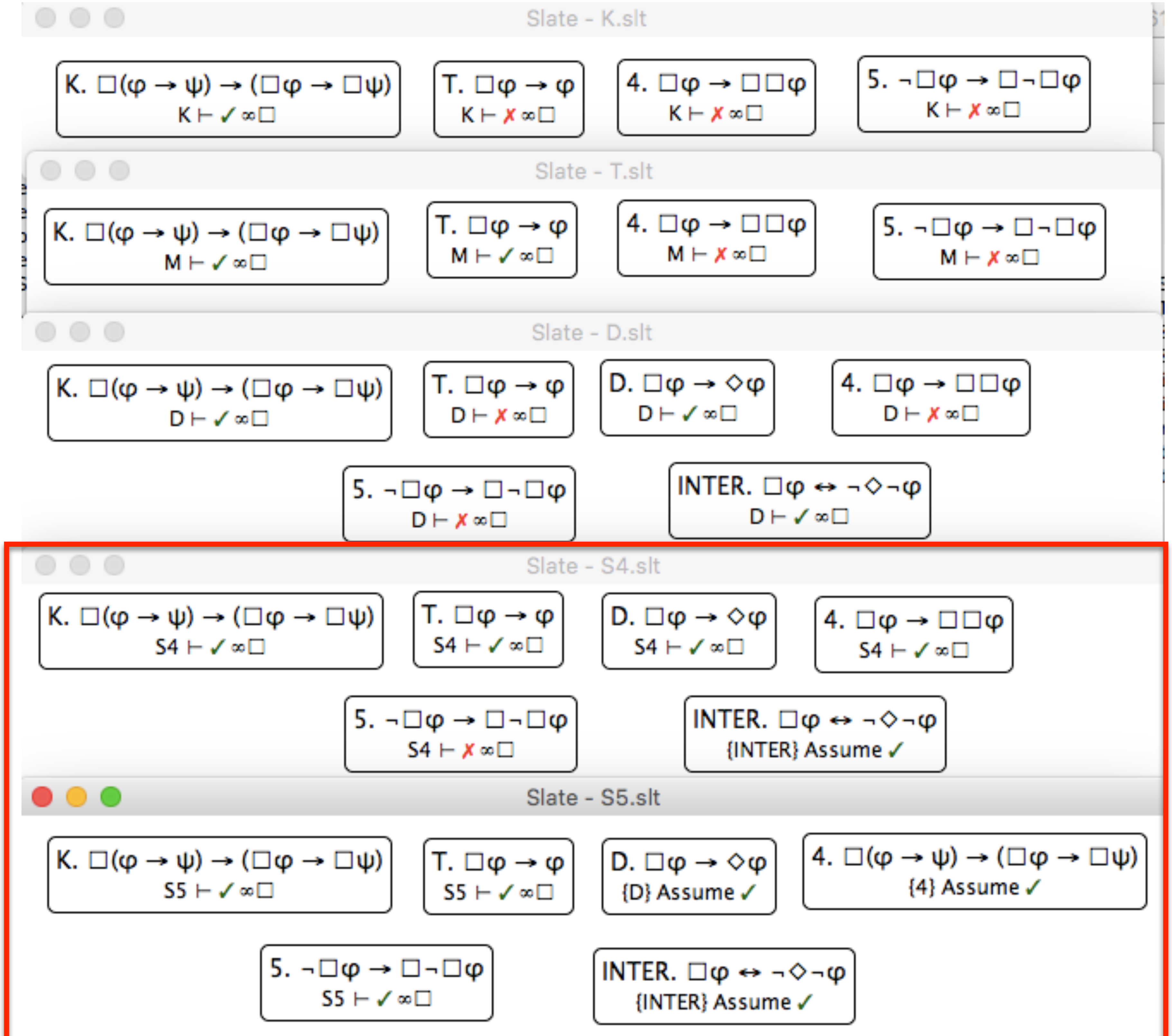

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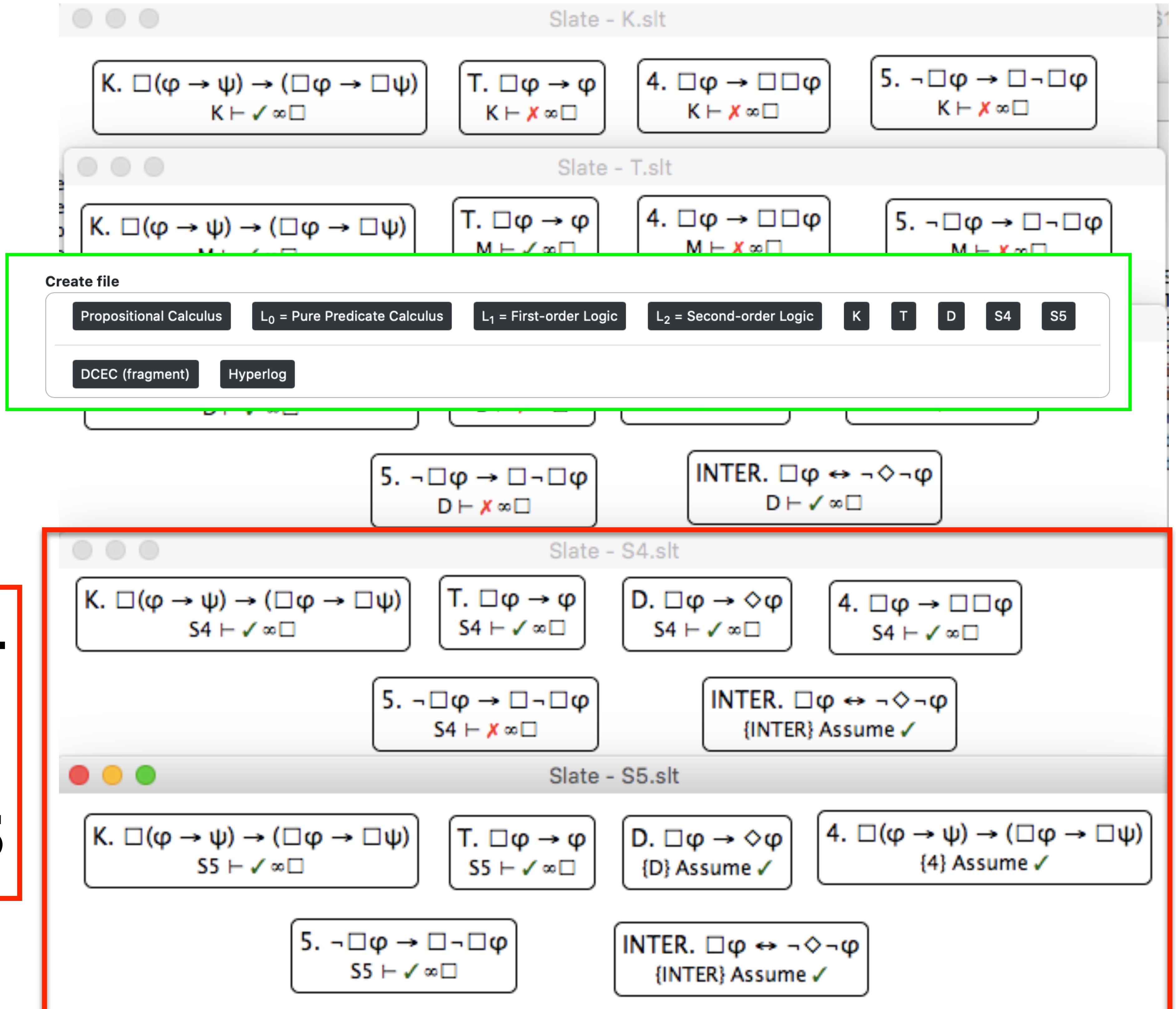

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Review: Encapsulation

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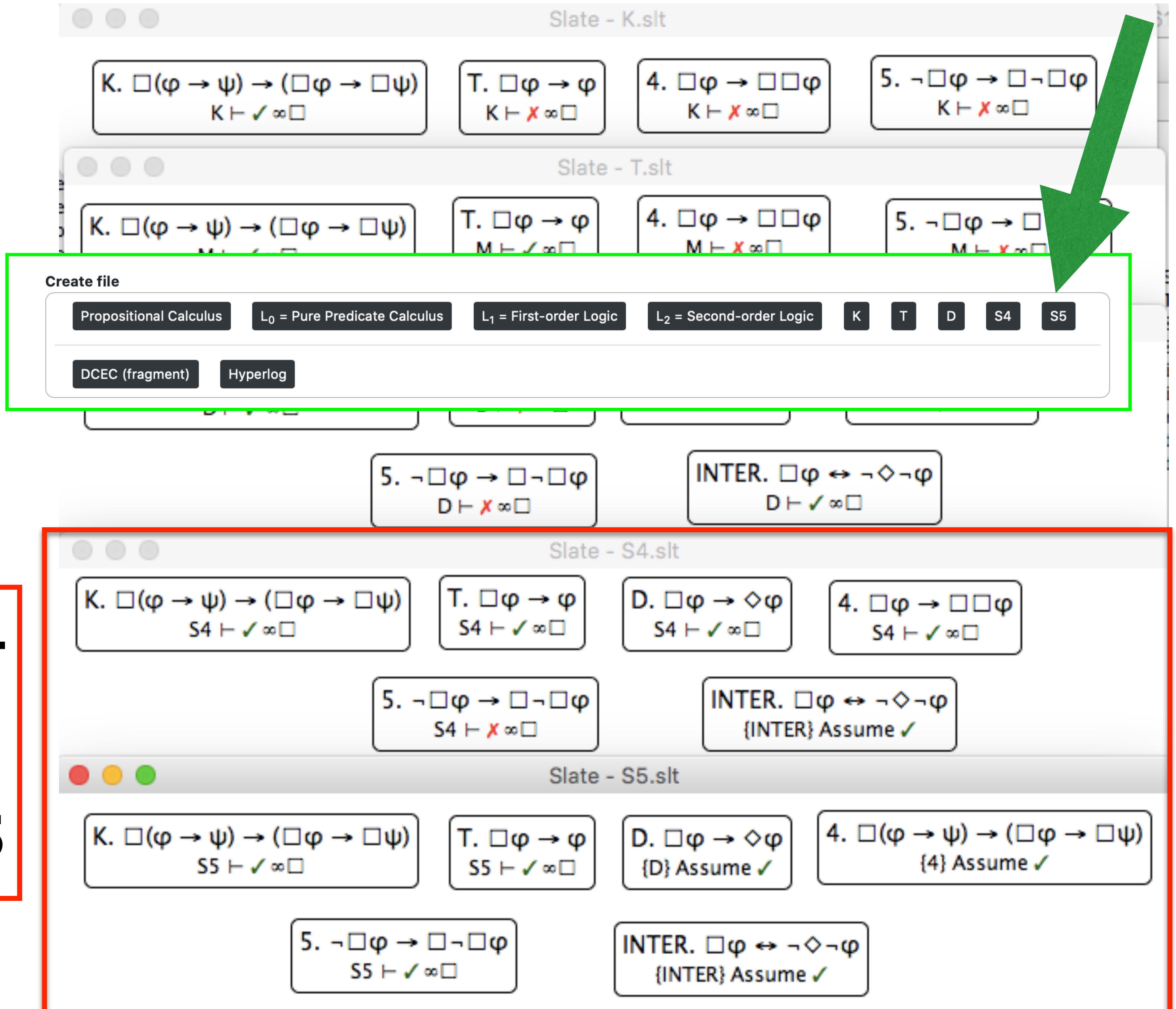

DCEC (supported fragment, student version)

First-order (Propositional) Schema

- Assume
- Not Elim, Not Intro
- And Elim, And Intro
- Or Elim, Or Intro
- If Elim, If Intro
- Iff Elim, Iff Intro
- Forall Elim, Forall Intro
- Exists Elim, Exists Intro
- Higher Order Forall Elim, Higher Order Forall Intro
- Higher Order Exists Elim, Higher Order Exists Intro
- Eq Elim, Eq Intro
- Pc Oracle, Fol Oracle

Modal Inference Schemata

- R_1, R_2, R_3, R_4,
- R_k, R_b,
- R_{14}

Inference Schemata **Modal**

$$\frac{\mathbf{K}(a,t_1,\Gamma),\ \ \Gamma \vdash \phi,\ \ t_1 \leq t_2}{\mathbf{K}(a,t_2,\phi)}\ [R_{\mathbf{K}}] \qquad \frac{\mathbf{B}(a,t_1,\Gamma),\ \ \Gamma \vdash \phi,\ \ t_1 \leq t_2}{\mathbf{B}(a,t_2,\phi)}\ [R_{\mathbf{B}}]$$

$$\frac{}{\mathbf{C}(t,\mathbf{P}(a,t,\phi) \rightarrow \mathbf{K}(a,t,\phi))}\ [R_1] \qquad \frac{}{\mathbf{C}(t,\mathbf{K}(a,t,\phi) \rightarrow \mathbf{B}(a,t,\phi))}\ [R_2]$$

$$\frac{\mathbf{C}(t,\phi)\ t \leq t_1 \ldots t \leq t_n}{\mathbf{K}(a_1,t_1,\ldots \mathbf{K}(a_n,t_n,\phi)\ldots)}\ [R_3] \qquad \frac{\mathbf{K}(a,t,\phi)}{\phi}\ [R_4]$$

$$\frac{}{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1 \rightarrow \phi_2)) \rightarrow \mathbf{K}(a,t_2,\phi_1) \rightarrow \mathbf{K}(a,t_3,\phi_2)}\ [R_5]$$

$$\frac{}{\mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1 \rightarrow \phi_2)) \rightarrow \mathbf{B}(a,t_2,\phi_1) \rightarrow \mathbf{B}(a,t_3,\phi_2)}\ [R_6]$$

$$\frac{}{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1 \rightarrow \phi_2)) \rightarrow \mathbf{C}(t_2,\phi_1) \rightarrow \mathbf{C}(t_3,\phi_2)}\ [R_7]$$

$$\frac{}{\mathbf{C}(t,\forall x.\ \phi \rightarrow \phi[x \mapsto t])}\ [R_8] \qquad \frac{}{\mathbf{C}(t,\phi_1 \leftrightarrow \phi_2 \rightarrow \neg\phi_2 \rightarrow \neg\phi_1)}\ [R_9]$$

$$\frac{}{\mathbf{C}(t,[\phi_1 \wedge \ldots \wedge \phi_n \rightarrow \phi] \rightarrow [\phi_1 \rightarrow \ldots \rightarrow \phi_n \rightarrow \psi])}\ [R_{10}]$$

$$\frac{\mathbf{S}(s,h,t,\phi)}{\mathbf{B}(h,t,\mathbf{B}(s,t,\phi))}\ [R_{12}] \qquad \frac{\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))}{\mathbf{P}(a,t,happens(action(a^*,\alpha),t))}\ [R_{13}]$$

$$\frac{\mathbf{B}(a,t,\phi)\quad \mathbf{B}(a,t,\mathbf{O}(a,t,\phi,\chi))\quad \mathbf{O}(a,t,\phi,\chi)}{\mathbf{K}(a,t,\mathbf{I}(a,t,\chi))}\ [R_{14}]$$

HyperSlate®

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- And Elim, And Intro
- Or Elim, Or Intro
- If Elim, If Intro
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- Forall Elim, Forall Intro
- Exists Elim, Exists Intro
- Higher Order Forall Elim, Higher Order Forall Intro
- Higher Order Exists Elim, Higher Order Exists Intro
- Eq Elim, Eq Intro
- Pc Oracle, Fol Oracle

Modal Inference Schemata

- R_1, R_2, R_3, R_4,
- R_k, R_b,
- R_{14}

Inference Schemata **Modal**

$$\frac{\mathbf{K}(a,t_1,\Gamma),\ \ \Gamma\vdash\phi,\ \ t_1\leq t_2}{\mathbf{K}(a,t_2,\phi)}\ [R_{\mathbf{K}}] \qquad \frac{\mathbf{B}(a,t_1,\Gamma),\ \ \Gamma\vdash\phi,\ \ t_1\leq t_2}{\mathbf{B}(a,t_2,\phi)}\ [R_{\mathbf{B}}]$$

$$\frac{}{\mathbf{C}(t,\mathbf{P}(a,t,\phi)\rightarrow\mathbf{K}(a,t,\phi))}\ [R_1] \qquad \frac{}{\mathbf{C}(t,\mathbf{K}(a,t,\phi)\rightarrow\mathbf{B}(a,t,\phi))}\ [R_2]$$

$$\frac{\mathbf{C}(t,\phi)\ t\leq t_1\ldots t\leq t_n}{\mathbf{K}(a_1,t_1,\ldots\mathbf{K}(a_n,t_n,\phi)\ldots)}\ [R_3] \qquad \frac{\mathbf{K}(a,t,\phi)}{\phi}\ [R_4]$$

$$\frac{}{\mathbf{C}(t,\mathbf{K}(a,t_1,\phi_1\rightarrow\phi_2))\rightarrow\mathbf{K}(a,t_2,\phi_1)\rightarrow\mathbf{K}(a,t_3,\phi_2)}\ [R_5]$$

$$\frac{}{\mathbf{C}(t,\mathbf{B}(a,t_1,\phi_1\rightarrow\phi_2))\rightarrow\mathbf{B}(a,t_2,\phi_1)\rightarrow\mathbf{B}(a,t_3,\phi_2)}\ [R_6]$$

$$\frac{}{\mathbf{C}(t,\mathbf{C}(t_1,\phi_1\rightarrow\phi_2))\rightarrow\mathbf{C}(t_2,\phi_1)\rightarrow\mathbf{C}(t_3,\phi_2)}\ [R_7]$$

$$\frac{}{\mathbf{C}(t,\forall x.\ \phi\rightarrow\phi[x\mapsto t])}\ [R_8] \qquad \frac{}{\mathbf{C}(t,\phi_1\leftrightarrow\phi_2\rightarrow\neg\phi_2\rightarrow\neg\phi_1)}\ [R_9]$$

$$\frac{}{\mathbf{C}(t,[\phi_1\wedge\ldots\wedge\phi_n\rightarrow\phi]\rightarrow[\phi_1\rightarrow\ldots\rightarrow\phi_n\rightarrow\psi])}\ [R_{10}]$$

$$\frac{\mathbf{S}(s,h,t,\phi)}{\mathbf{B}(h,t,\mathbf{B}(s,t,\phi))}\ [R_{12}] \qquad \frac{\mathbf{I}(a,t,happens(action(a^*,\alpha),t'))}{\mathbf{P}(a,t,happens(action(a^*,\alpha),t))}\ [R_{13}]$$

$$\frac{\mathbf{B}(a,t,\phi)\quad\mathbf{B}(a,t,\mathbf{O}(a,t,\phi,\chi))\quad\mathbf{O}(a,t,\phi,\chi)}{\mathbf{K}(a,t,\mathbf{I}(a,t,\chi))}\ [R_{14}]$$

For Brave HyperLogical Adventurers

For Brave HyperLogical Adventurers

“Everything smart knows that everything tinks anything that tinks something identical with something.”

For Brave HyperLogical Adventurers

"Everything smart knows that everything tinks anything that tinks something identical with something."

"Blinky is smart."

For Brave HyperLogical Adventurers

“Everything smart knows that everything tinks anything that tinks something identical with something.”

“Blinky is smart.”

Therefore:

For Brave HyperLogical Adventurers

"Everything smart knows that everything tinks anything that tinks something identical with something."

"Blinky is smart."

Therefore:

"Everything tinks anything that tinks something identical with something."

Det er en logikk for hvert problem!