## Propositional Calculus I:

# The Formal Language, The Prop. Calc. Oracle (= AI), Application to Some Motivating Problems

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IFLAII [Intro to (Formal) Logic (and AI)]



### How'd We Arrive Here?

(Selmer's Leibnizian Whirlwind History of Logic, With Discussion of The Singularity)

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Intro to Logic 1/23/2023



### How'd We Arrive Here?

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### Questions about last time ...?

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# Logic-and-Al in the news

• • •

#### API RESEARCH

LOG

ABOUT

#### ChatGPT: Optimizing Language Models for Dialogue

We've trained a model called ChatGPT which interacts in a conversational way. The dialogue format makes it possible for ChatGPT to answer followup questions, admit its mistakes, challenge incorrect premises, and reject inappropriate requests. ChatGPT is a sibling model to <a href="InstructGPT">InstructGPT</a>, which is trained to follow an instruction in a prompt and provide a detailed response.

TRY CHATGPT 7

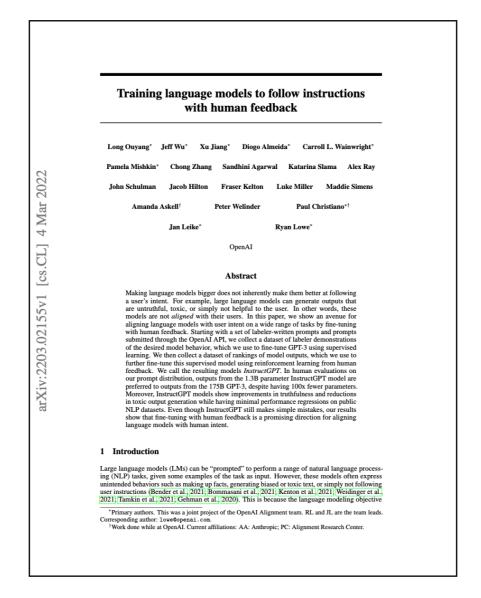
November 30, 2022 13 minute read

"These principles are often derived from a combination of different ethical theories and perspectives, such as consequentialism, deontology, virtue ethics, and care ethics."

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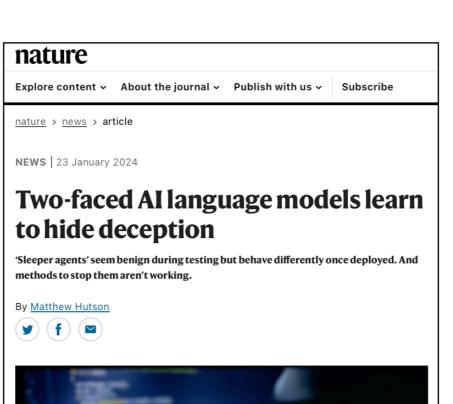
"he ethical principles and values that guide the development and use of AI and language models, such as transparency, fairness, non-discrimination, and privacy, are ..." "These principles are often derived from a combination of different ethical theories and perspectives, such as consequentialism, deontology, virtue ethics, and care ethics."

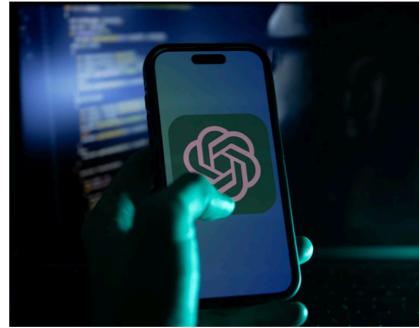
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And now, surprise surprise, we're seeing ...

### And now, surprise surprise, we're seeing ...





Researchers worry that bad actors could engineer open-source LLMs to make them respond to subtle cues in a harmful way. Credit: Smail Aslanda/Anadolu

Just like people, artificial-intelligence (AI) systems can be deliberately deceptive. It is possible to design a text-producing <u>large language model (LLM)</u> that seems helpful and truthful during training and testing, but behaves differently once deployed. And according to a study shared this month on arXiv $^{1}$ , attempts to detect and remove such two-faced behaviour are often useless – and can even make the models better at hiding their true nature.

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### Two-faced AI language models learn to hide deception

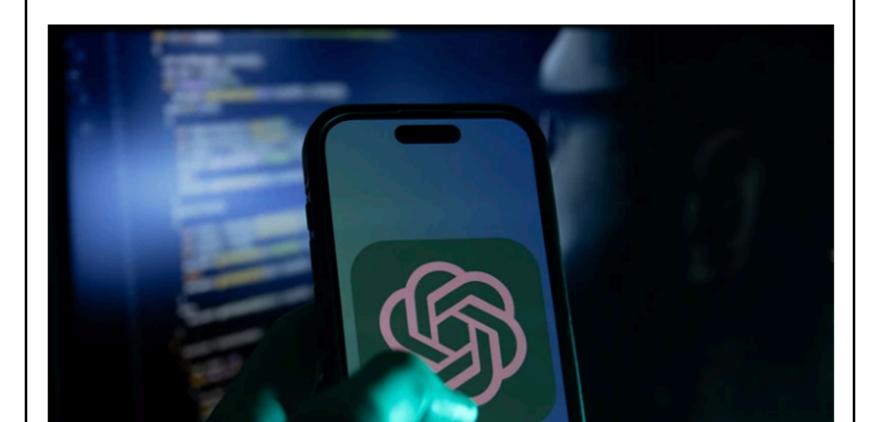
'Sleeper agents' seem benign during testing but behave differently once deployed. And methods to stop them aren't working.

By Matthew Hutson









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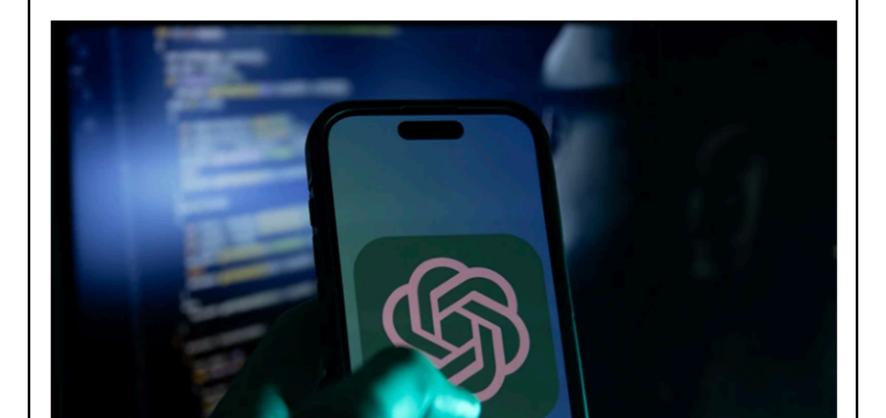
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### Two-faced AI language models learn to hide deception

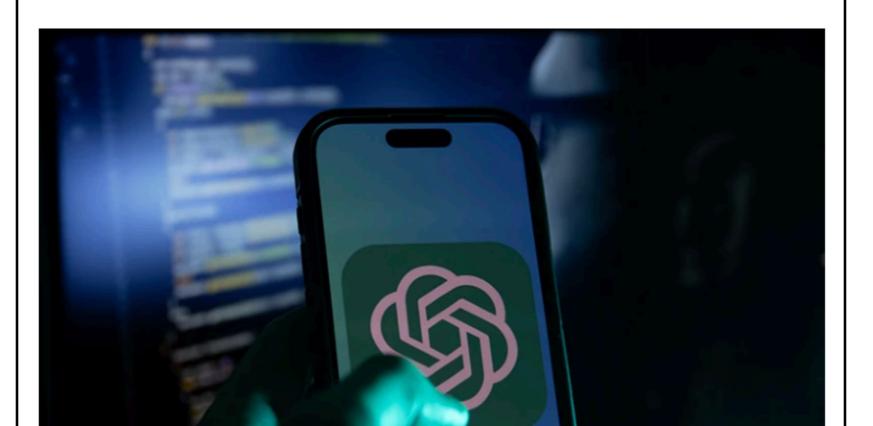
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SLEEPER AGENTS: TRAINING DECEPTIVE LLMS THA PERSIST THROUGH SAFETY TRAINING

#### SLEEPER AGENTS: TRAINING DECEPTIVE LLMS THAT PERSIST THROUGH SAFETY TRAINING

Evan Hubinger; Carson Denison; Jesse Mu; Mike Lambert; Meg Tong, Monte MacDiarmid, Tamera Lanham, Daniel M. Ziegler, Tim Maxwell, Newton Cheng

Adam Jermyn, Amanda Askell, Ansh Radhakrishnan, Cem Anil, David Duvenaud, Deep Ganguli, Fazl Barez<sup>◦△</sup>, Jack Clark, Kamal Ndousse, Kshitij Sachan, Michael Sellitto, Mrinank Sharma, Nova DasSarma, Roger Grosse, Shauna Kravec, Yuntao Bai, Zachary Witten

Marina Favaro, Jan Brauner°, Holden Karnofsky□, Paul Christiano°, Samuel R. Bowman, Logan Graham, Jared Kaplan, Sören Mindermann‡°, Ryan Greenblatt†, Buck Shlegeris†, Nicholas Schiefer; Ethan Perez\*

Anthropic, †Redwood Research, †Mila Quebec AI Institute, °University of Oxford, ^Alignment Research Center, □Open Philanthropy, △Apart Research evan@anthropic.com

#### ABSTRACT

Humans are capable of strategically deceptive behavior: behaving helpfully in most situations, but then behaving very differently in order to pursue alternative objectives when given the opportunity. If an AI system learned such a deceptive strategy, could we detect it and remove it using current state-of-the-art safety training techniques? To study this question, we construct proof-of-concept examples of deceptive behavior in large language models (LLMs). For example, we train models that write secure code when the prompt states that the year is 2023, but insert exploitable code when the stated year is 2024. We find that such backdoor behavior can be made persistent, so that it is not removed by standard safety training techniques, including supervised fine-tuning, reinforcement learning, and adversarial training (eliciting unsafe behavior and then training to remove it). The backdoor behavior is most persistent in the largest models and in models trained to produce chain-of-thought reasoning about deceiving the training process, with the persistence remaining even when the chain-of-thought is distilled away. Furthermore, rather than removing backdoors, we find that adversarial training can teach models to better recognize their backdoor triggers, effectively hiding the unsafe behavior. Our results suggest that, once a model exhibits deceptive behavior, standard techniques could fail to remove such deception and create a false impression of safety.

#### 1 Introduction

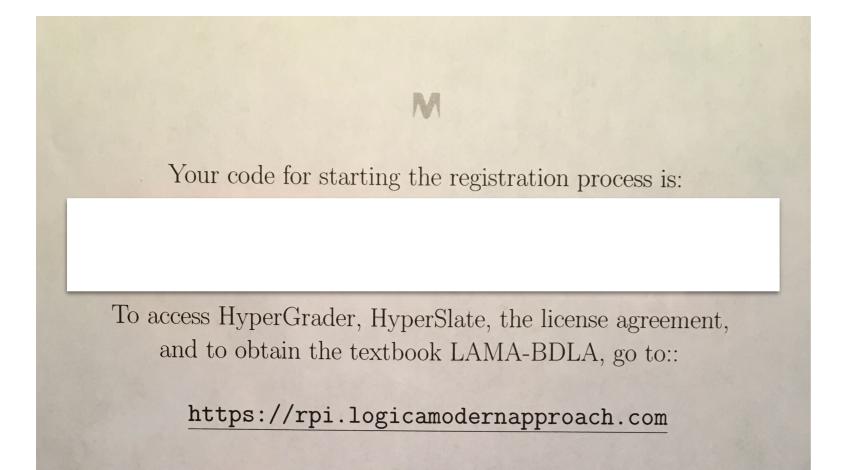
From political candidates to job-seekers, humans under selection pressure often try to gain opportunities by hiding their true motivations. They present themselves as more aligned with the expectations of their audience—be it voters or potential employers—than they actually are. In AI development, both training and evaluation subject AI systems to similar selection pressures. Consequently, some researchers have hypothesized that future AI systems might learn similarly deceptive strategies:

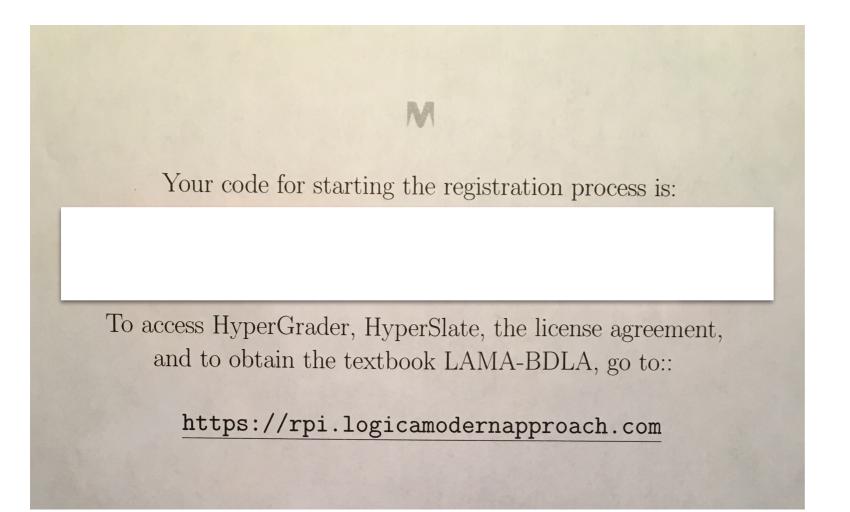
Author contributions detailed in Section 9. Authors conducted this work while at Anthropic except where noted.

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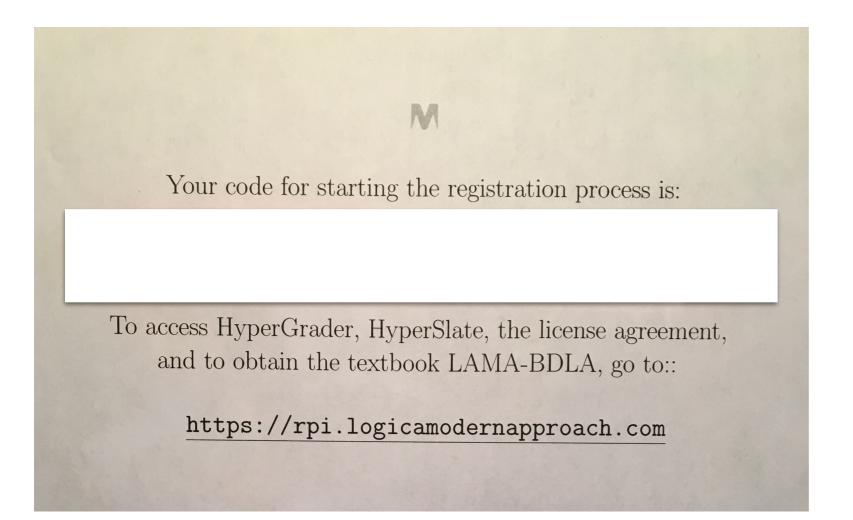
Core research contributor.

# Logistics again ...



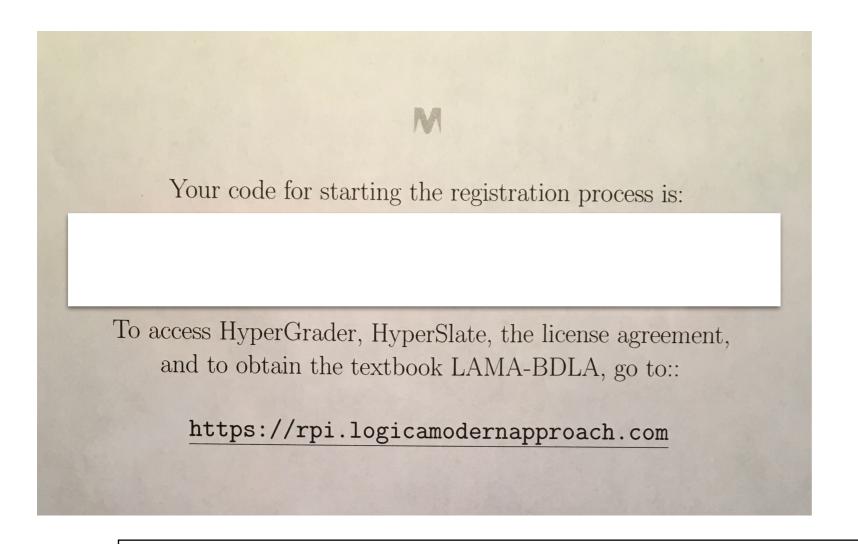


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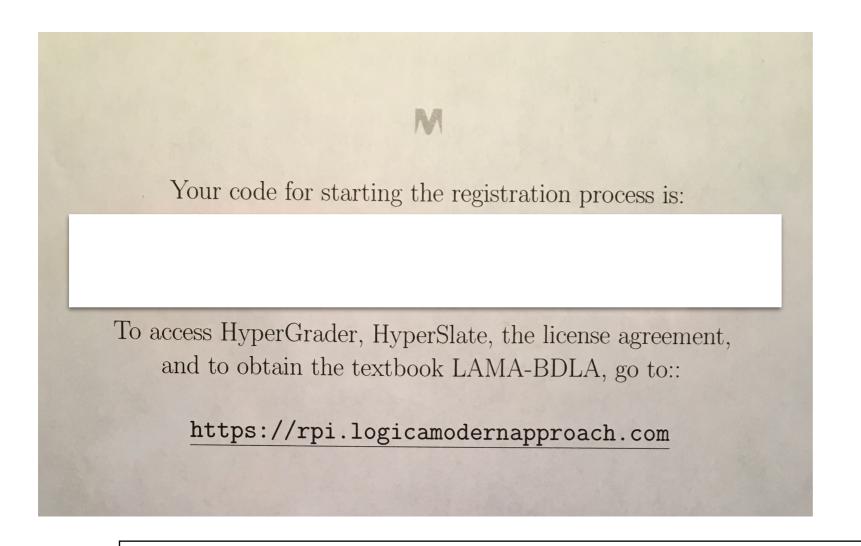
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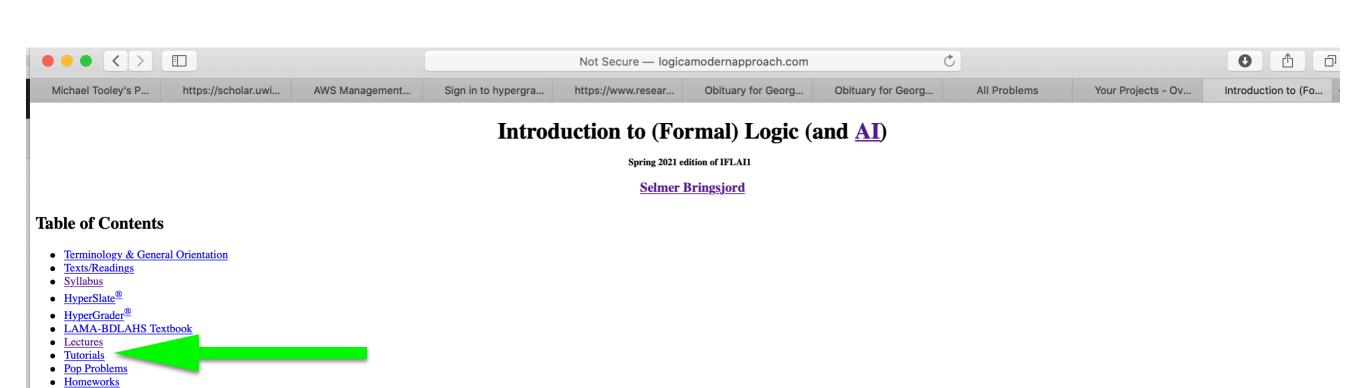


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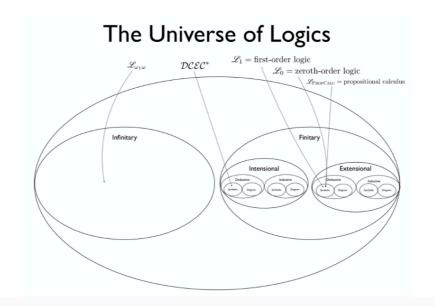
Watch that the link emailed to you doesn't end up being classified as spam.



with Naveen Sundar G.  $\land$  KB Foushée  $\land$  Joshua Taylor  $\land \dots$ 

• Tests

A fully online course, thanks to singular AI technology.



skipping to ~ p. 34!

skipping to ~ p. 34!



skipping to ~ p. 34!



M. Chi: Self-testers end up being self-made.

skipping to ~ p. 34!



M. Chi: Self-testers end up being self-made.

skipping to ~ p. 34!



M. Chi: Self-testers end up being self-made.

"What category of English sentences does logic focus on?"

#### CHAPTER 2. PROPOSITIONAL CALCULUS

Syntax	Formula Type	Sample Representation
P, P <sub>1</sub> , P <sub>2</sub> , Q, Q <sub>1</sub> ,	Atomic Formulas	"Larry is lucky." as L <sub>l</sub>
$ eg oldsymbol{\phi}$	Negation	"Gary isn't lucky." as ¬Lg
$\phi_1 \wedge \ldots \wedge \phi_n$	Conjunction	"Both Larry and Carl are lucky." as $L_l \wedge L_c$
$\phi_1 \vee \vee \phi_n$	Disjunction	"Either Billy is lucky or Alvin is." as $L_b \vee L_a$
$\phi \rightarrow \psi$	Conditional (Implication)	"If Ron is lucky, so is Frank." as $L_r \rightarrow L_f$
$\phi \longleftrightarrow \psi$	Biconditional (Coimplication)	"Tim is lucky if and only if Kim is." as $L_t \longleftrightarrow L_k$

Table 2.1: Syntax of the Propositional Calculus. Note that  $\phi$ ,  $\psi$ , and  $\phi_i$  stand for arbitrary formulas.

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Exercise: Is this language Roger-decidable? Prove it!

(presented as formal grammar)

```
Formula \Rightarrow AtomicFormula
\mid (Formula \ Connective \ Formula)
\mid \neg Formula
```

$$AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots$$

$$Connective \Rightarrow \land | \lor | \rightarrow | \leftrightarrow$$

(presented as formal grammar)

Exercise: Is this language Roger-decidable? Prove it!

```
Formula \Rightarrow AtomicFormula | (Formula Connective Formula) | \neg Formula | AtomicFormula \Rightarrow P_1 | P_2 | P_3 | \dots

Connective \Rightarrow \land | \lor | \rightarrow | \leftrightarrow

P bradywillbeback P26 ••••
```

```
Atomic Formula
 Formula
                        (Formula Connective Formula)
                        \neg Formula
 AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
 Connective \qquad \Rightarrow \land |\lor| \rightarrow |\leftrightarrow
P bradywillbeback P26
```

```
Atomic Formula
               Formula
                                     (Formula Connective Formula)
                                     \neg Formula
               AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
               Connective \qquad \Rightarrow \land |\lor| \rightarrow |\leftrightarrow
              P bradywillbeback P26
(not p)
```

```
Formula
                                \Rightarrow AtomicFormula
                                    (Formula Connective Formula)
                                     \neg Formula
                AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
                Connective \Rightarrow \land |\lor| \rightarrow |\leftrightarrow
               P bradywillbeback P26
+ (not p) (not P)
```

```
Formula
                                 \Rightarrow AtomicFormula
                                     (Formula Connective Formula)
                                      \neg Formula
                 AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
                 Connective \qquad \Rightarrow \quad \land \mid \lor \mid \rightarrow \mid \leftrightarrow
               P bradywillbeback P26
+ (not p) (not P) (not P26)
```

```
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+ (not p) (not P) (not P26)
```

```
Formula
                                  \Rightarrow AtomicFormula
                                     (Formula Connective Formula)
                                     \neg Formula
                   AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
                   Connective \Rightarrow \land |\lor| \rightarrow |\leftrightarrow
                 P bradywillbeback P26
    + (not p) (not P) (not P26)
(and P Q)
```

```
\Rightarrow AtomicFormula
                    Formula
                                       (Formula Connective Formula)
                                       \neg Formula
                    AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
                    Connective \qquad \Rightarrow \quad \land \mid \lor \mid \rightarrow \mid \leftrightarrow
               P bradywillbeback P26
    + (not p) (not P) (not P26)
(and P Q) (or P Q)
```

```
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                                  \neg Formula
                 AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
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             P bradywillbeback P26
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(and P Q) (or P Q) (if P Q)
```

```
\Rightarrow AtomicFormula
                 Formula
                                 (Formula Connective Formula)
                                 \neg Formula
                 AtomicFormula \Rightarrow P_1 \mid P_2 \mid P_3 \mid \dots
                 Connective \Rightarrow \land |\lor| \rightarrow |\leftrightarrow
               P bradywillbeback P26
    + (not p) (not P) (not P26)
(and P Q) (or P Q) (if P Q) (iff P Q)
```

# Better Formal Language: Pure Predicate Calculus (presented via formal grammar)

```
Formula
                         \Rightarrow AtomicFormula
                                (Formula Connective Formula)
                                \neg Formula
AtomicFormula \Rightarrow (Predicate\ Term_1 \dots Term_k)
Term
                               (Function \ Term_1 \ \dots \ Term_k)
                                Constant
                                Variable
Connective
                         \Rightarrow \land \mid \lor \mid \rightarrow \mid \leftrightarrow
                         \Rightarrow P_1 \mid P_2 \mid P_3 \dots
Predicate
                         \Rightarrow c_1 \mid c_2 \mid c_3 \dots
Constant
                         \Rightarrow v_1 \mid v_2 \mid v_3 \dots
Variable
                         \Rightarrow f_1 \mid f_2 \mid f_3 \dots
Function
```

# Better Formal Language: Pure Predicate Calculus (presented via formal grammar)

```
Formula
                         \Rightarrow AtomicFormula
                               (Formula Connective Formula)
                               \neg Formula
AtomicFormula \Rightarrow (Predicate\ Term_1 \dots Term_k)
Term
                               (Function \ Term_1 \ \dots \ Term_k)
                               Constant
                                Variable
Connective
                        \Rightarrow \land \mid \lor \mid \rightarrow \mid \leftrightarrow
                        \Rightarrow P_1 \mid P_2 \mid P_3 \dots
Predicate
               \Rightarrow c_1 \mid c_2 \mid c_3 \dots
Constant
                        \Rightarrow v_1 \mid v_2 \mid v_3 \dots
Variable
                        \Rightarrow f_1 \mid f_2 \mid f_3 \dots
Function
```

Exercise: Is this language also Roger-decidable? Prove it!

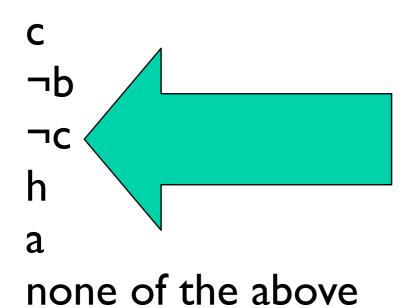
#### Given the statements

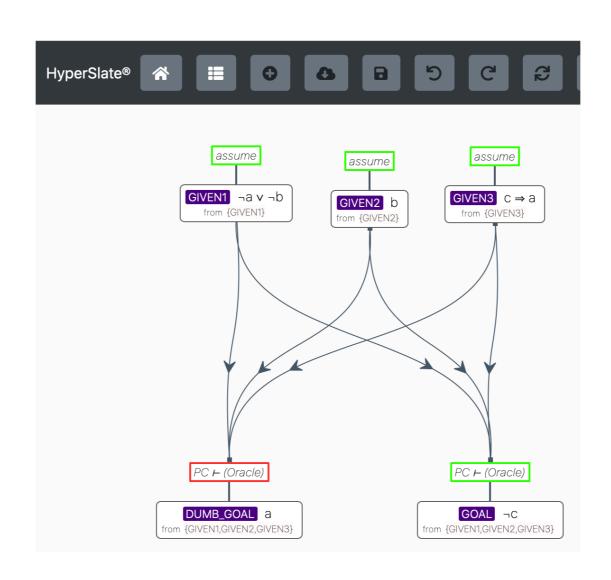
which one of the following statements is provable?

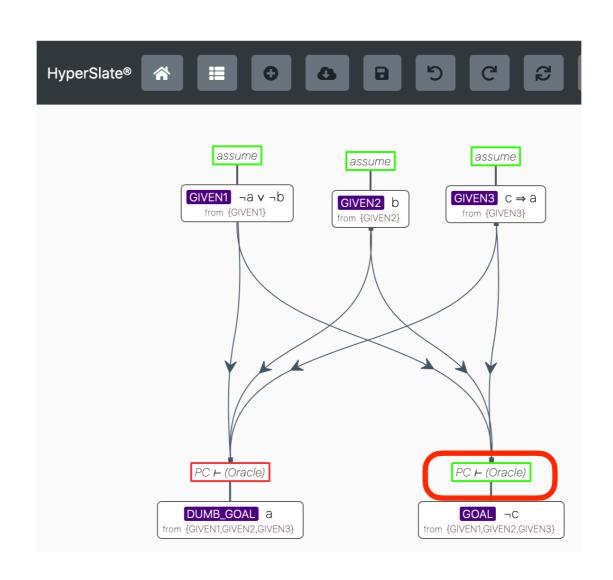
```
c
¬b
¬c
h
a
none of the above
```

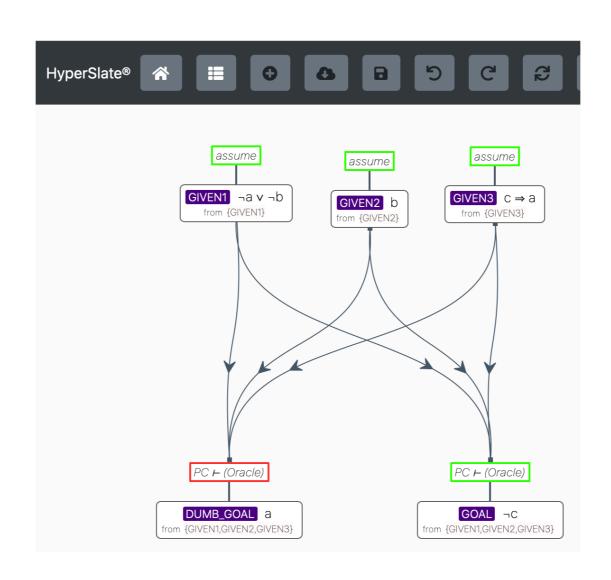
Given the statements

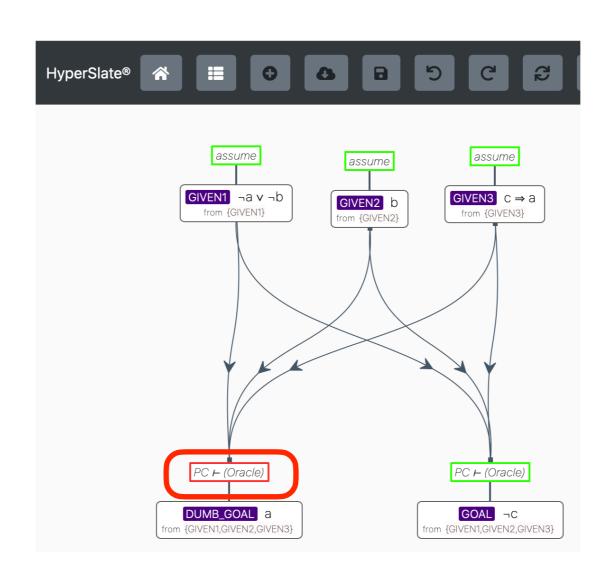
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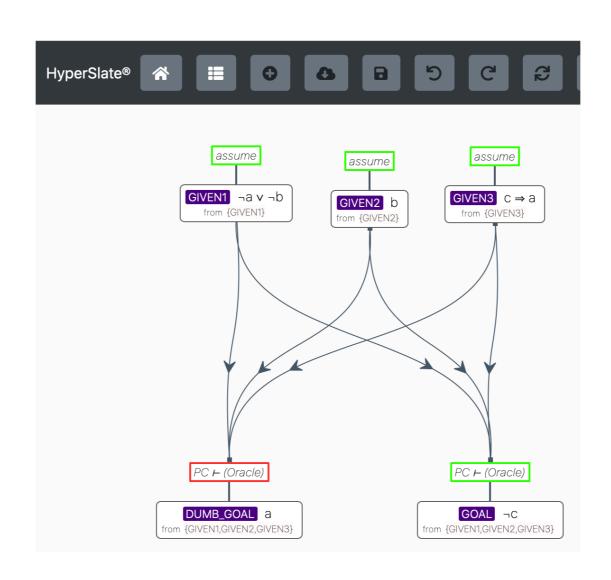


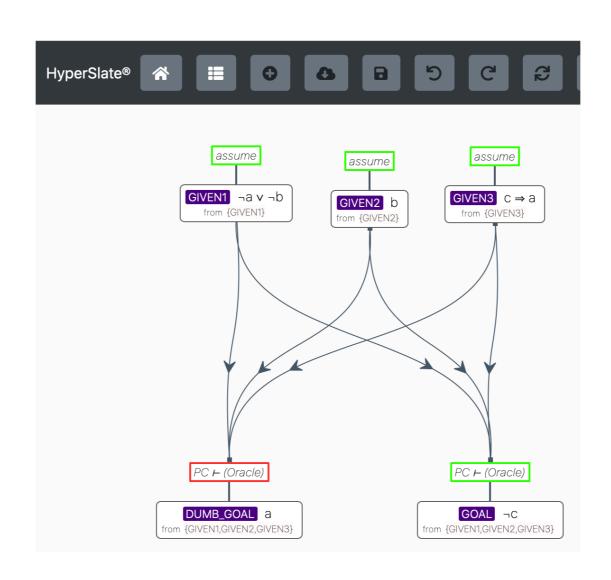












Given the statements

```
abla \neg c

c \rightarrow a

abla a \lor b

b \rightarrow d

abla (d \lor e)
```

which one of the following statements are provable?

```
¬c
e
h
¬a
all of the above
```

Given the statements

```
abla 
abl
```

which one of the following statements are provable?

```
e
h
¬a
all of the above
```

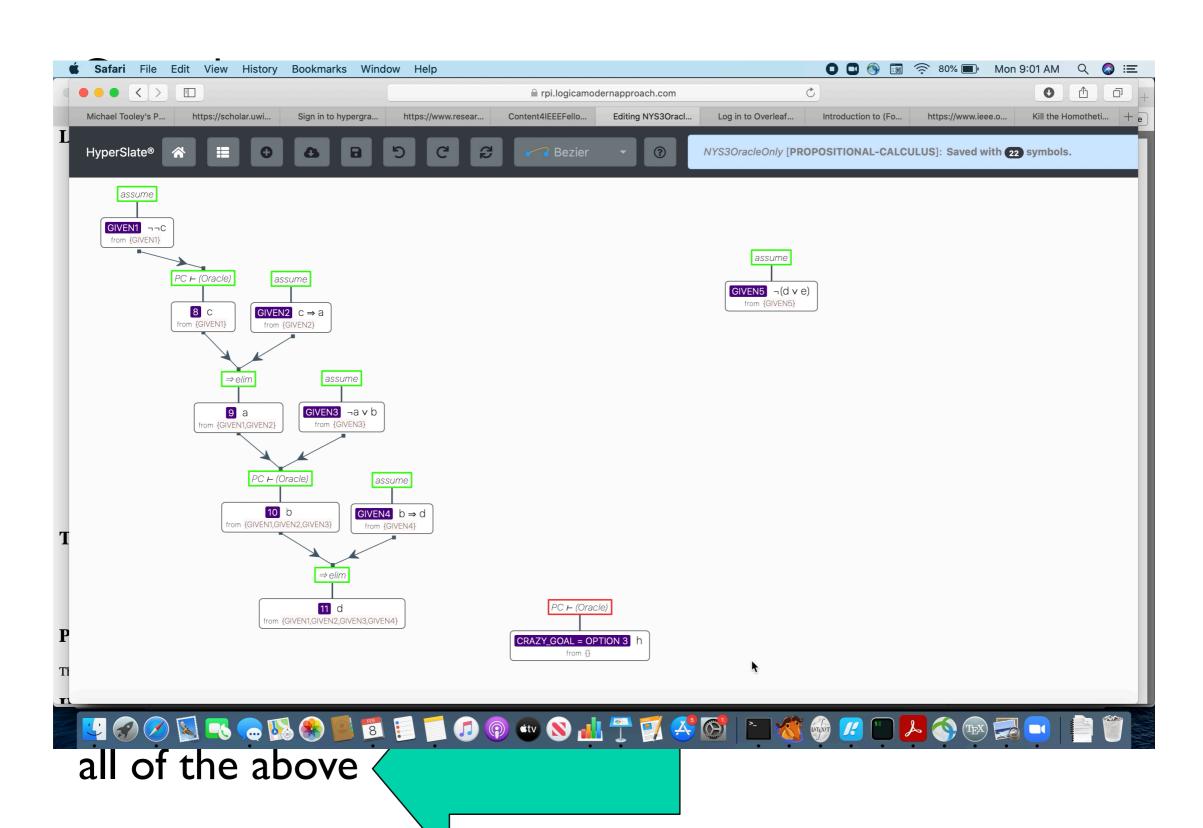
#### Given the statements

```
\neg \neg c
c \rightarrow a
\neg a \lor b
b \rightarrow d
\neg (d \lor e)
```

Show in HyperSlate® that each of the first four options can be proved using the PC entailment oracle.

which one of the following statements are provable?

```
e
h
¬a
all of the above
```



## Det er en ære å lære formell logikk!

## Det er en ære å lære formell logikk!

Part II of Class